

RT-11
Software Support Manual

DEC-11-ORPGA-B-D

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PREFACE

The RT-11 Software Support Manual covers the internal description of the RT-11 software system. Chapter 1 presents an overview of the system and discusses conventions used throughout the manual. Chapters 2 through 6 describe in detail various aspects of the monitor and system structure, including memory layout, monitor tables, file structures, file formats, system device structure, bootstrap operation, I/O queuing system, device handlers and F/B monitor description. Chapter 7 discusses the operation of the BATCH compiler and run-time handler.

The appendixes provide example handler listings, including a foreground terminal handler (Appendix B) and a sample foreground program (Appendix D). Complete flowcharts of both the Single-Job and Foreground/Background Monitors are shown in Appendix E.

The reader should be thoroughly familiar with the RT-11 system. Although the information in this manual is aimed at V02B users, it should be adequate for Version 2 users also; excluding a few minor alterations (to permit the addition of the new V02B devices), the construction of the monitors has changed very little between the two versions. A comprehensive list of differences between the V2 and V02B systems is included in Getting Started with RT-11 (V02B) (DEC-11-ORCPA-E-D).

It is assumed that the user has read the RT-11 System Reference Manual (DEC-11-ORUGA-B-D) or (DEC-11-ORUGA-C-D) and all other documentation included in the RT-11 kit, and is an experienced PDP-11 programmer. It is recommended that RT-11 monitor source listings be available for reference.

CHAPTER 1

RT-11 OVERVIEW

1.1 INTRODUCTION

RT-11 is a single-user programming and operating system designed for the PDP-11 series of computers. It permits the use of a wide range of peripherals and up to 28K of either solid state or core memory (hereafter referred to as memory).

RT-11 provides two operating environments: Single-Job (S/J) operation, and a powerful Foreground/Background (F/B) capability. Either environment is controlled by a single user from the console terminal keyboard by means of the appropriate monitor--S/J or F/B. The monitors are upwards compatible; features that are used only in a F/B environment are treated as no-ops under the S/J Monitor.

A feature common to both operating environments is the inclusion of a full complement of system development and utility programs to aid the programmer in the development of his own applications.

The normal use and operation of the monitors and system programs is discussed in detail in the RT-11 System Reference Manual. Concepts and applications that are specialized and useful to the more experienced programmer are included in this manual.

1.2 SYSTEM CONCEPTS AND TERMINOLOGY

The basic concepts necessary to use RT-11 effectively are defined in the RT-11 System Reference Manual. The user should be familiar with those concepts before proceeding to use this manual.

Abbreviations used throughout this document are:

<u>TERM</u>	<u>MEANING</u>
KMON	<p>Keyboard Monitor</p> <p>The console terminal interface to RT-11. KMON runs as a background job and allows the user to run programs, assign device names, and generally control the system.</p>
USR	<p>User Service Routines</p> <p>The nonresident (swapping) part of RT-11. The USR performs file-oriented operations.</p>
CSI	<p>Command String Interpreter</p> <p>The CSI is part of the USR. It accepts a string of characters from memory or from the console and performs specified file operations, or syntactically analyzes a command string and constructs a table from the information supplied.</p>
RMON	<p>Resident Monitor</p> <p>RT-11 provides a choice of two Resident Monitors: a Single-Job Monitor and a Foreground/Background Monitor. RMON specifically provides the following services:</p> <ul style="list-style-type: none">EMT dispatcherKeyboard (console) interrupt serviceTT: resident device handler (F/B only)Read/Write processorUSR swap routinesI/O queuing routinesSystem device handlerSystem I/O tablesMessage handler (F/B only)Job scheduler (F/B only)
\$CSW	<p>Channel Status Word</p> <p>Each bit in the CSW contains information relevant to the status of a channel; see Chapter 9 (.SAVESTATUS) of the <u>RT-11 System Reference Manual</u>.</p>

<u>TERM</u>	<u>MEANING</u>
JSW	Job Status Word The JSW contains information in bytes 44 and 45 about the job currently in memory.
F/B	The Foreground/Background version of the monitor
S/J	The Single-Job version of the monitor
B/G	The background job
F/G	The foreground job
<CR>	Carriage Return
<LF>	Line Feed

Various mnemonic names (e.g., BLIMIT, SYSLOW), referred to from within the text and in diagrams and flowcharts, represent the actual symbolic names as they appear in the monitor source listings.

To avoid confusion, underlining is used in most examples to designate computer printout; square brackets, [and], are used to enclose comments. Values for symbolic names used in examples can be found in Chapter 5 of Getting Started With RT-11 (V02B) (DEC-11-ORCPA-E-D).

CHAPTER 2
MEMORY LAYOUT

RT-11 operates properly in any configuration between 8K and 28K (words) of memory (16K to 28K for the F/B Monitor). No user intervention is required when programs are moved to a different size machine; i.e., programs correctly developed in one environment will work in any size environment (providing there is sufficient memory) with no relinking necessary.

Figure 2-1 shows a general diagram of the memory layout in an RT-11 system.

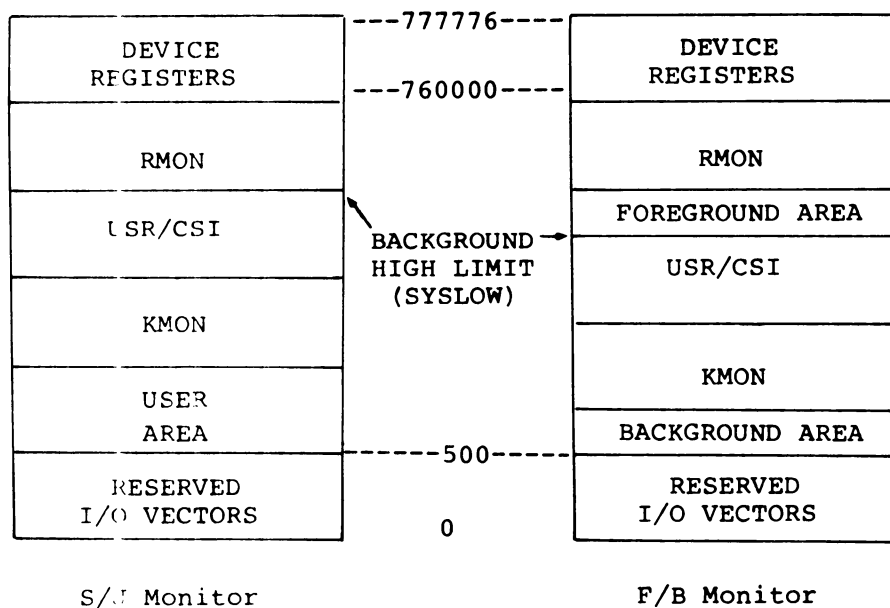


Figure 2-1
Monitor Memory Layout

The memory area diagrammed is arranged as follows:

<u>Memory Area</u>	<u>Use</u>
0-477	Reserved for I/O vectors, RT-11 system communication area.
500-SYSLOW	Space available for user (background) programs. (The high limit of memory for the background is contained in SYSLOW, a location in the monitor data base.)

Space for foreground programs and LOAded handlers is allocated as needed, reducing the amount of space available for a background job.

The areas marked KMON and USR/CSI are the areas that these units normally occupy when they are in memory. The amount of memory that a user program occupies is determined by:

1. The initial size of the program, or
2. The amount of memory the user program requests via a .SETTOP programmed request.

When a user program (background job) is executed (via the KMON commands R, RUN, or GET and START), the top of memory is set to correspond to the size of the program. If the top of user memory never exceeds KMON, both KMON and USR/CSI are resident. If all of memory (up to SYSLOW) is requested (via a .SETTOP), neither the KMON nor the USR is resident and swapping of the USR is required. Programs performing many file-oriented operations gain from having the USR resident, since no time is spent swapping the USR.

The KMON, USR, and RMON modules normally occupy the upper segment of memory. This implies that larger memory configurations automatically have more free memory available.

The area marked DEVICE REGISTERS is the top 4K of memory in any PDP-11 computer. This area is reserved for the status and control registers of peripheral devices.

2.1 FOREGROUND JOB AREA LAYOUT

The foreground job area is located above the KMON/USR, as shown in Figure 2-1, and is allocated by the FRUN command. The actual layout of the job within the foreground area is shown in Figure 2-2. The *impure area* (described in Section 2.5.3) occupies the lowest 207 words of the job area and contains terminal ring buffers, I/O channels, and other job-specific information.

The foreground stack is located immediately above the impure area with a default size of 128 words; this may be changed using the FRUN /S switch. The program may specify a different location for the stack by using an .ASECT into location 42, in which case the /S switch is ignored and the program itself must allocate stack space. Whenever the stack is located, stack overflow will most probably cause program malfunction before penetrating the task area boundary, since either the program itself or the impure area will be corrupted.

NOTE

Users must not use a relocatable symbol as the contents of location 42 when resetting the initial stack pointer via an .ASECT in a foreground job; such a symbol is not relocated when it occurs in an .ASECT in a foreground job. To set the stack to relative location 1000 in a foreground job, use:

```
.ASECT  
.=42  
.WORD 1000
```

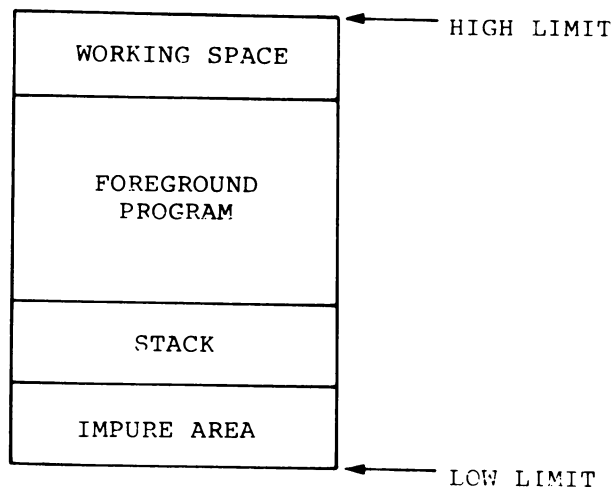


Figure 2-2
Foreground Job Area Layout

The space allocated for the foreground program is sufficient to contain the program code itself, as indicated by location 50 (in block 0 of the file); location 50 is set by the Linker and designates the program's high limit. If the foreground job requires working space, this space must either be reserved from within the program (e.g., using .BLKW) or allocated at run-time using the FRUN /N switch. Space allocated with the /N switch is located above the program as shown in Figure 2-2. Location 50 will point to the top of the program area and a .SETTOP will permit access to any working space.

2.2 JOB BOUNDARIES IN F/B

The actual job boundaries are stored (in RMON) in limit tables for both foreground and background jobs. The FLIMIT table contains high and low boundaries for the foreground, and the BLIMIT table contains boundaries for the background. .SETTOPs are permitted for any job up to its high limit. The SYSLOW pointer mentioned earlier is equivalent to the background BLIMIT high pointer entry. This is shown in Figure 2-3.

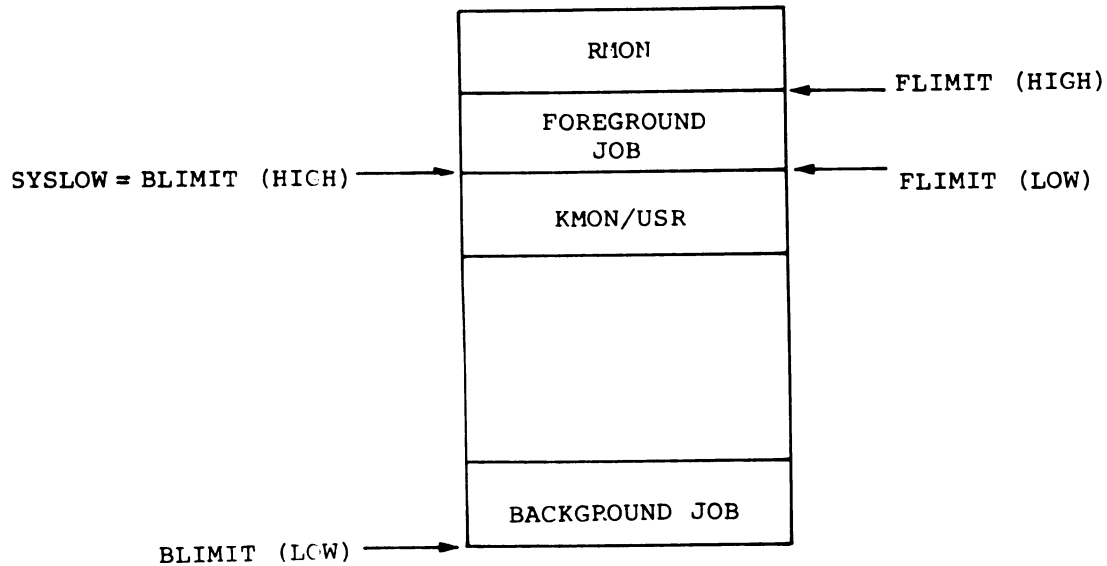


Figure 2-3
Job Limits

The limit pointers for a foreground job are fixed once the job has been loaded into memory. A program that requires working space and uses a .SETTOP will fail if the space is not allocated with the /N switch (a FORTRAN program is a typical case; see Appendix G, Section G.1, of the RT-11 System Reference Manual). The high limit pointer (SYSLOW) for the background, however, is not fixed and will change as space is allocated for LOADED handlers, the text scroller, and foreground jobs. In addition, if the USR is made permanently resident (using the SET USR NOSWAP command), SYSLOW (BLIMIT HIGH) will again change. This is shown in Figure 2-4.

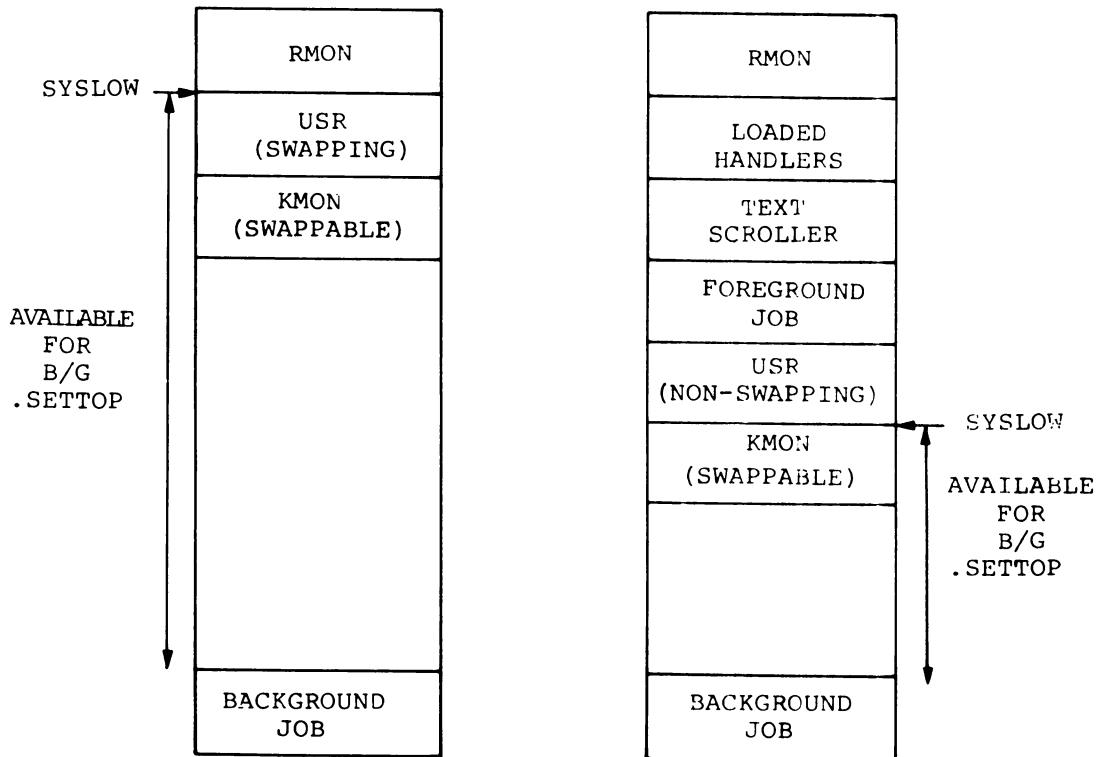


Figure 2-4
Background SYSLOW Examples

2.3 'FLOATING' USR POSITION

The RT-11 USR is normally located in the memory area directly below that pointed to by SYSLOW. For the Version 1 monitor, this was directly below the RMON. For the Version 2 and 2B monitors, the USR position varies as handlers, the scroller, and foreground jobs (in F/B) are loaded into memory; the SYSLOW pointer is corrected for each change in memory configuration. In any case, the SYSLOW position is considered the normal USR swapping position.

It is possible, however, to cause the USR to swap into another location in memory. This is done by setting location 46 (in the system communication area) to the address at which the USR is to swap; if the contents of location 46 are nonzero and even, the monitor loads the USR at the new address. Note, however, that if no swapping is required, the USR is *not* loaded at the address indicated in location 46. Location 46 is cleared by an exit to the Keyboard Monitor (via an .EXIT, .HRESET, .SRESET, or CTRL C).

It is possible to make the USR permanently resident (i.e., non-swapping). Using the SET USR NOSWAP Keyboard Monitor command makes the USR permanently resident at its *normal* position, that is, below the memory area pointed to by SYSLOW.

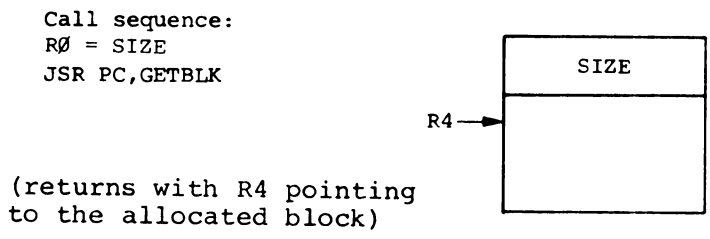
2.4 MONITOR MEMORY ALLOCATION

RT-11 uses a dynamic memory allocation scheme to provide memory space for LOADED handlers, foreground jobs (F/B Monitor only) and the display text scroller. Memory is allocated in the region above the KMON/USR and below RMON. If there is insufficient memory in this region (initially, after the system is bootstrapped, there is none), memory is taken from the background region by "sliding down" the KMON/USR the required number of words.

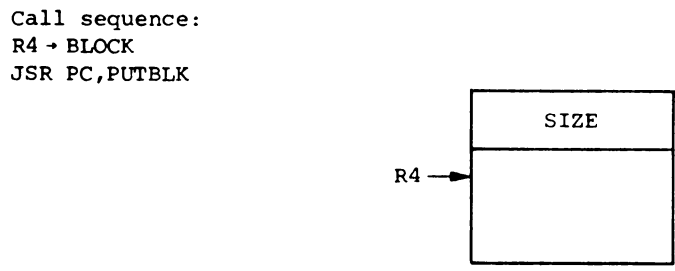
When memory allocated in this manner is released, the memory block is returned to a singly-linked free memory list, the list head of which is in RMON. Any contiguous blocks are concatenated into a single larger block. A block found to be contiguous with the KMON/USR is reclaimed by "sliding up" the KMON/USR, removing the block from the list.

Memory allocation and release is achieved by calls to the GETBLK and PUTBLK routines located in the KMON overlays (the GETBLK and PUTBLK routines are flowcharted in Appendix E). The requested number of words is passed to GETBLK in R0, and the address of the block is returned in R4. An extra word of memory is allocated by GETBLK, which then stores the size of the block in that word. R4 points to the first available word in the block (see Figure 2-5a). When releasing memory, R4 must point to the first available word, the same address returned by GETBLK during allocation (as shown in Figure 2-5b). The block will be linked into the free memory list (shown in Figure 2-5c).

a) Allocating a memory block



b) Releasing a memory block



c) Free memory list

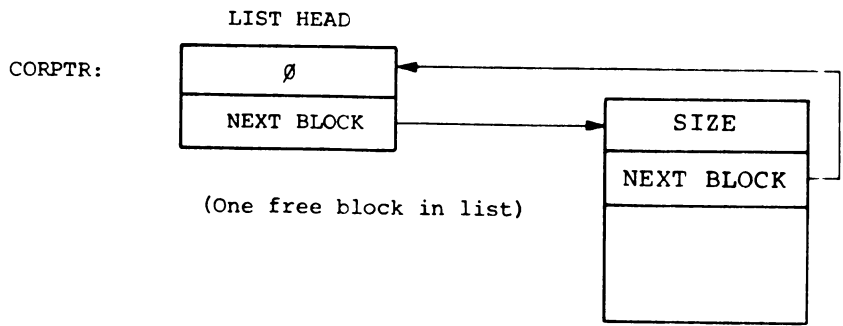


Figure 2-5
 Memory Allocation

When a block of memory of sufficient size is not available, GETBLK must create a hole in memory by sliding down the KMON/USR. This is achieved by a call to KUMOVE, a small routine located physically at the front of the KMON. KUMOVE does the actual work of moving the KMON/USR up in memory. For moves downward, an auxiliary subroutine, MOVEDN, located at the top of the USR, is used.

Whenever a request is made for a block of a certain number of words, the memory allocator searches memory for the first highest block that is large enough to satisfy the request (that is, equal to or larger than the requested number). The goal of the memory allocator is to minimize the amount of free (unused) memory in the foreground region, making the maximum amount of memory available to the background. Contiguous blocks of free memory are merged and reclaimed whenever possible. The search time of the singly-linked list is not a factor, since at any time there will be few nodes (free memory areas) in the list, and the allocator minimizes the number.

2.5 MEMORY AREAS OF INTEREST

This section describes memory areas of particular interest and indicates the contents of those locations. The areas covered are:

1. Monitor Fixed Offsets (F/B & S/J)
2. F/B Impure Area
3. Resident Bitmap (F/B & S/J)
4. Tables

2.5.1 Monitor Fixed Offsets

Certain values are maintained at fixed locations from the start of the Resident Monitor in both F/B and S/J; these quantities (listed in Table 2-1) may be accessed by user programs. The technique used to access these offsets is as follows:

OFFSET = the byte offset to the word desired
RMON = 54

```
MOV @#RMON,Rn          ;ANY GENERAL REGISTER
MOV OFFSET(Rn),Rn
```

Rn now contains the desired quantity. If a byte quantity is desired, a better method is:

```
CLR Rm
MOV @#RMON,Rn
BISB OFFSET(Rn),Rm
```

This ensures that the high-order bits of the register are not set by a MOV_B into the register.

Table 2-1
Fixed Offsets

Offset (from Start of RMON) Octal Decimal	Tag	Byte Length	Description
0	-	2	Serves as a link to interrupt entry code.
2	\$CSW	160 ₁₀	Default I/O channels for the background (16 ₁₀ @ 5 words each).
244 164	\$SYSCH	10 ₁₀	Internal I/O channel used for system functions.
256 174	BLKEY	2	Segment number of the directory now in memory. 0 implies no directory is there.
260 176	CHKEY	2	Device index and unit number of the device whose directory is in memory. Bits 1-5 are the device index, bits 8-10 are the unit number.
262 178	\$DATE	2	Current date value. (The format is shown in Chapter 3, section 3.1.2.5.)
264 180	DFLG	2	"Directory operation in progress" flag. Used to inhibit ^C from aborting a job until directory operation is finished.
266 182	\$USRLC	2	Normal location of USR.
270 184	QCOMP	2	Address of I/O completion manager, COMPLT.
272 186	SPUSR	2	Flag word used by MT/CT. If a USR function performed by MT or CT fails, this word is made non-zero.

(continued on next page)

Table 2-1 (Cont.)
Fixed Offsets

Offset (from Start of RMON)		Tag	Byte Length	Description																				
Octal	Decimal																							
274	188	SYUNIT	2	High-order byte contains the unit number of the current system device.																				
276	190	SYSVER	1	Monitor version number (2 in Version 2).																				
277	191	SYSUPD	1	Version update number.																				
300	192	CONFIG	2	System configuration word. A 16-bit series of flags whose meanings are: <table border="0" style="margin-left: 40px;"> <thead> <tr> <th style="text-align: left;"><u>Bit #</u></th> <th style="text-align: left;"><u>Meaning</u></th> </tr> </thead> <tbody> <tr> <td>0</td> <td>0 → S/J Monitor 1 → F/B Monitor</td> </tr> <tr> <td>2</td> <td>1 → VT11 hardware exists</td> </tr> <tr> <td>3</td> <td>1 → RT-11 BATCH controls the background</td> </tr> <tr> <td>5</td> <td>0 → 60-cycle KW11L clock 1 → 50-cycle clock</td> </tr> <tr> <td>6</td> <td>1 → 11/45 FPP present</td> </tr> <tr> <td>7</td> <td>0 → No foreground job present 1 → Foreground job is in memory</td> </tr> <tr> <td>8</td> <td>1 → User is linked to VT11 scroller</td> </tr> <tr> <td>9</td> <td>1 → USR is resident via SET USR</td> </tr> <tr> <td>15</td> <td>1 → KW11L clock is present</td> </tr> </tbody> </table> <p>Any bits not currently assigned are reserved by DIGITAL for future use and should not be used arbitrarily by user programs.</p>	<u>Bit #</u>	<u>Meaning</u>	0	0 → S/J Monitor 1 → F/B Monitor	2	1 → VT11 hardware exists	3	1 → RT-11 BATCH controls the background	5	0 → 60-cycle KW11L clock 1 → 50-cycle clock	6	1 → 11/45 FPP present	7	0 → No foreground job present 1 → Foreground job is in memory	8	1 → User is linked to VT11 scroller	9	1 → USR is resident via SET USR	15	1 → KW11L clock is present
<u>Bit #</u>	<u>Meaning</u>																							
0	0 → S/J Monitor 1 → F/B Monitor																							
2	1 → VT11 hardware exists																							
3	1 → RT-11 BATCH controls the background																							
5	0 → 60-cycle KW11L clock 1 → 50-cycle clock																							
6	1 → 11/45 FPP present																							
7	0 → No foreground job present 1 → Foreground job is in memory																							
8	1 → User is linked to VT11 scroller																							
9	1 → USR is resident via SET USR																							
15	1 → KW11L clock is present																							
302	194	SCROLL	2	Address of the VT11 scroller.																				
304	196	TTKS	2	Address of console keyboard status.																				
306	198	TTKB	2	Address of console keyboard buffer.																				

(continued on next page)

Table 2-1 (Cont.)
Fixed Offsets

Offset (from Start of RMON)		Tag	Byte Length	Description
Octal	Decimal			
310	200	TTPS	2	Address of console printer status.
312	202	TTPB	2	Address of console printer buffer.
				(See Section 2.6, Using Auxiliary Terminals as the Console Terminal.)
314	204	MAXBLK	2	Largest output file permitted with an indefinite length request (initially defined as -1, which implies that no limit is defined).
316	206	E16LST	2	Offset from start of RMON to the dispatch table for EMT's 340-357. (This is used by the BATCH processor.)
320	208	CNTXT	2	Pointer to the impure area for the current executing job.
		(F/B)		
322	210	JOBNUM	2	Executing job's number (0 = B/G, 2 = F/G).
320	208	\$TIME	4	Two words of time of day in the S/J Monitor.
322	210	(S/J)		
324	212	SYNCH	2	Address of monitor routine to handle .SYNCH request.
326	214	LOWMAP	20 ₁₀	Start of low memory protection map. (This map protects vectors at locations 0-476.)
352	234	USRLOC	2	Pointer to current entry point of USR.
354	236	GTVECT	2	Pointer to VT11 vector. The vector is initially positioned at 320.
356	238	ERRCNT	1	Error count byte (for future use by system programs).
357	239	FUTURE	5	Reserved by DIGITAL for future use.

2.5.2 Table Descriptions

The monitor device tables discussed in this section include:

```
$PNAME
$STAT
$ENTRY
$DVREC
$HSIZE
$DVSIZ
$UNAM1,$UNAM2
$OWNER
```

The size of these tables is fixed and is governed by the \$SLOT assignment; the default value is 14₁₀ entries per table. To alter this, it is necessary to first edit a new value of \$SLOT into the monitor source program, then reassemble and relink new monitors.

2.5.2.1 \$PNAME (Permanent Name Table) - \$PNAME is the central table around which all the others are constructed. There is an entry in \$PNAME for each device in the system. Each entry consists of a single word that contains the .RAD50 code for the two-character permanent device name for that device; for example the entry for DECTape is .RAD50 /DT/. The position of devices in this table is non-critical, but their relative position determines the general device index used in various places in the monitor; thus, all other tables must be organized in the same order as \$PNAME (the index into \$PNAME serves as the index into all the other tables for the equivalent device).

2.5.2.2 \$STAT (Device Status Table) - Each device in the system must have a status entry in its corresponding slot in \$STAT. The status word is broken down into two bytes as follows:

Even byte - contains a device identifier. Each unique type of device in the system has an identifying integer. Those defined are:

```
0 = RK05 Disk
1 = TC11 DECTape
2 = Reserved
3 = Line Printer (LP11, LS11, LV11)
4 = Console Terminal (LT33/35, LA30/36,
  VT05, VT50)
5,6 = Reserved
7 = PC11 High-speed Reader
10 = PC11 High-speed Punch
11 = Magtape (TM11, TU10)
12 = RF11 Disk
13 = TA11 Cassette
```

14 = Card Reader (CR11, CM11)
 15 = Reserved
 16 = RJS03/4 Fixed-head Disks
 17 = Reserved
 20 = TJU16 Magtape
 21 = RP11/RP02/RP03 Disk
 22 = RX11/RX01 Floppy Disk

Odd byte - Bit flags with the following meanings:

<u>Bit #</u>	<u>Meaning</u>
15	1 = File-structured device (disk, DECTape) 0 = Nonfile-structured device (PC11, LT33, etc.)
14	1 = Read-only device
13	1 = Write-only device
12	1 = Device whose directory is not a standard RT-11 directory (MT, CT); the device handlers for these devices are expected to be able to perform their own directory operations
11-8	Reserved

2.5.2.3 \$ENTRY (Handler Entry Point Table) - Whenever a handler is made resident, either by a .FETCH or with the LOAD command, the \$ENTRY slot for that device is made to point to the fourth word of the device handler. The entry is zeroed when the handler is .RELEASED or UNLOADED.

2.5.2.4 \$DVREC (Device Handler Block Table) - This table (filled in at system bootstrap time) reflects the absolute block position of each of the device handlers on the system device. Since handlers are treated as files under RT-11, their position on the system device is not necessarily fixed. Thus, each time the system is bootstrapped, the handlers are located and \$DVREC is updated with the value of the second block of the handler file. (Because the handlers are linked at 1000, the actual handler code starts in the second block of the file.) A zero entry in the \$DVREC table indicates that no handler for the device in that slot was found on the system device.

2.5.2.5 \$HSIZE (Handler Size Table) - This table contains the size, in bytes, of each device handler. The table is set up at assembly time with the correct values and is used when a .FETCH is executed to provide the size of the specified handler. This size is also returned to the user as one of the values returned in a .DSTAT request.

2.5.2.6 SDVSIZ (Device Directory Size Table) - Entries in this table are non-zero for file-structured devices only and reflect the number of 256₁₀-word blocks contained on the device. The current devices and their entries are:

<u>Device</u>	<u>Number of 256-Word Blocks</u>	<u>Device</u>	<u>Number of 256-Word Blocks</u>
RK11	11300 ₈	RP02	116300 ₈
TC11	1102 ₈	RJS03	2000 ₈
		RJS04	4000 ₈
RF11	2000 ₈ (1 platter) 4000 ₈ (2 platters) 6000 ₈ (3 platters) 10000 ₈ (4 platters)	RX01	752 ₈

The default for RF11 and RJS03/4 is one platter, or 2000₈ blocks. It is possible to alter the system to indicate the correct number of platters. Instructions are in Appendix A of the RT-11 System Reference Manual.

2.5.2.7 \$UNAM1, \$UNAM2 (User Name Tables) - These tables are used in conjunction with ASSIGN keyboard functions. The form of the ASSIGN command is:

```
.ASSIGN pnam:unam<CR>
```

where:

```
pnam - a system device name/unit number
unam - a user-assigned device name
```

A typical example is:

```
.ASSIGN DT1:DK
```

The default device name, DK, is now directed to DECTape unit 1. The user-assigned name is stored in an available slot in \$UNAM2, while the device's permanent name/unit is stored in the corresponding slot in \$UNAM1. The system uses a common device name lookup routine that maps any user-assigned name in the \$UNAM2 table into a physical device name to be used in an operation. The total number of ASSIGNS permitted is limited by the value of \$SLOT.

The command:

```
.ASSIGN<CR>
```

zeroes \$UNAM2, thus removing all user assignments.

2.5.2.8 \$OWNER (Device Ownership Table) - This table is used only under F/B to arbitrate device ownership. The table is \$SLOT*2 words in length and is divided into 2-word entries per device. Each 2-word entry is divided into eight 4-bit fields capable of holding a job number. Thus, each device is presumed to have up to eight units, each assigned independently of the others. However, if the device is nonfile-structured, the ownership is assigned to all units.

When a job attempts to access a particular unit of a device, the F/B Monitor checks to be sure the unit being accessed is either public or belongs to the requesting job. If the unit is owned by the other job, a fatal error is generated.

The device is assumed to be public if the 4-bit field is 0. If it is not public, the field contains a code equal to the job number plus one. Since job numbers are always even, the ownership code is odd. Bit 0 of the field being set is then used to indicate that the unit ownership is assigned to a job (1 for the background job and 3 for the foreground job).

2.5.2.9 DEVICE Macro - The DEVICE macro call is used in RMON to allow quick and easy insertion of new devices at assembly time. The form of the macro call is:

```
DEVICE NAME,SIZ,STAT,ENTRY
```


where:

NAME - two characters of the permanent device name
SIZ - the size of the device's directory in 256-word blocks; 0 means nonfile-structured or special
STAT - the sum of all \$STAT table entries that apply for this device plus the device id (from section 2.5.2.2):

FILST\$ = 100000 File-structured device
RONLY\$ = 40000 Read-only device
WONLY\$ = 20000 Write-only device
SPECL\$ = 10000 Non RT-11 file-structured device (including MT and CT)

ENTRY - the 2-character device name with SYS appended, if this is a system device.

Thus, a sample call is:

```
DEVICE TT,0,4
```

The SIZ entry is 0, since TT is a nonfile-structured device.

The entry for DECTape is:

```
DEVICE DT,1102,1+FILST$,DTSYS
```

The 1+FILST\$ indicates that the device code is 1 and FILST\$ is defined as 100000. The entry for DTSYS is present because DT can be a system device.

In addition to the DEVICE macro, another macro, HSIZE, is defined and sets the handler size for the \$HSIZE table. The format of the HSIZE macro call is:

```
HSIZE HAN,BYT,TYPE
```

where:

HAN - the 2-letter device name
BYT - the handler size in bytes
TYPE - SYS if the device can be a system device; blank otherwise

Chapter 5 shows the use of HSIZE in adding a handler to the RT-11 system. The KMON portion of the monitor source listing should be consulted for greater detail.

2.5.3 F/B Impure Area

An impure area is defined here as that area of memory where the monitor stores all job-dependent data. Thus, the impure area contains all information that the monitor requires to effectively run two independent jobs, both of which are memory-resident. This section details the contents and location of each word (byte) in the impure area.

A table that points to the impure area for a particular job is in the F/B monitor's data base. This table is at \$IMPUR and currently consists of two words: the first is a pointer to the background's impure area (which is permanently resident in RMON at location BKGND), the second is the foreground's pointer.

Under RT-11, a background job is always running and will be the KMON if no other background job exists. However, the foreground impure area pointer may be 0 if no foreground job is in memory. When an FRUN command is given, a foreground impure area is created for the job and the \$IMPUR entry for the foreground pointer is updated to point to the impure area.

Table 2-2 is a detailed breakdown of the contents of the impure area. The offset mentioned is the offset from the start of the impure area itself; thus, the first word in the area has a 0 offset.

Table 2-2
Impure Area

Offset	Mnemonic	Octal Length (Bytes)	Contents
0	I.JSTA	2	Job status.
2	I.QHDR	2	I/O Queue Header.
4	I.CMPE	2	Last entry in completion queue. I/O completion routines are queued for execution. This is the pointer to the last routine to be entered.
6	I.CMPL	2	Completion queue header.
10	I.CHWT	2	Pointer to channel during I/O wait. When a job is waiting for I/O, the address of the channel area in use goes here.
12	I.PCHW	2	Saved channel pointer during execution of a completion routine. The contents of I.PCHW are put in R0 when a completion routine is entered.
14	I.PERR	2	Error byte 52 and 53 saved during completion routines.
16	I.PTTI	2	Previous TT input character.
20	I.TTLC	2	Terminal input ring buffer line count.
22	I.TID	2	Pointer to job ID area.
24	I.JNUM	2	Job number of job that owns this impure area.
26	I.CNUM	2	Number of I/O channels defined. 16 ₁₀ is default, .CDFN can be used to define new ones.
30	I.CSW	2	Pointer to job's channel area.
32	I.IOCT	2	Count of total I/O operations outstanding.
34	I.SCTR	2	Suspension count. Zero means the number of .SPNDs = the number of .RSUMs.
36	I.SPLS	2	Address of the .DEVICE request list.

(continued on next page)

Table 2-2 (Cont.)
Impure Area

Offset	Mnemonic	Octal Length (Bytes)	Contents
40	I.TRAP	2	Address of user trap routine. Set by .TRPSET.
42	I.FPP	2	Address of FPP exception routine. Set by .SFPA.
44	I.SWAP	4	Address and number of extra words to be included in the context switch operation. Set by .CNTXSW request.
50	I.SP	2	Saved stack pointer. When this job is made inactive, the active value of SP is saved here.
52	I.BITM	24	Low memory protection bitmap. This map reflects the user's .PROTECT requests.
(76 through 332 concern the console terminal)			
76	I.IRNG	2	Input ring buffer low limit.
100	I.IPUT	2	Input "PUT" pointer for inter- rupts.
102	I.ICTR	2	Input character counter.
104	I.IGET	2	Input "GET" pointer for .TTYIN.
106	I.ITOP	2	Input ring buffer high limit.
110		144	Input ring buffer.
254	I.OPUT	2	Output "PUT" pointer for interrupts.
256	I.OCTR	2	Output character counter.
260	I.OGET	2	Output "GET" pointer for interrupts.
262	I.OTOP	2	Output ring buffer high limit.
264		50	Output ring buffer.
334	I.QUE	20	Initial I/O queue element.

(continued on next page)

Table 2-2 (Cont.)
Impure Area

Offset	Mnemonic	Octal Length (Bytes)	Contents
354	I.MSG	12	Message channel. Used by .RCVD and .SDAT. This channel is permanently open.
366		10	Job ID area. Contains (<CR><LF>)B(<CR><LF>) or (<CR><LF>)F(<CR><LF>) for terminal prompting. Space has been left for up to a 3-character job name.

2.5.4 Low Memory Bitmap (LOWMAP)

RT-11 maintains a bitmap which reflects the protection status of low memory, locations 0-476. This map is required in order to avoid conflicts in the use of the vectors. In F/B, the .PROTECT request allows a program to gain exclusive control of a vector or a set of vectors. When a vector is protected, the bitmap is updated to indicate which words are protected. If a word in low memory is protected, it will not be destroyed when a new background program is run.

The bitmap is a 20₁₀ byte table which starts 326 bytes from the beginning of the Resident Monitor. Table 2-3 lists the offset from RMON and the corresponding locations represented by that byte:

Table 2-3
Bitmap Byte Table

Offset	Locations (octal)	Offset	Locations (octal)
326	0-16	340	240-256
327	20-36	341	260-277
330	40-56	342	300-316
331	60-76	343	320-336
332	100-116	344	340-356
333	120-136	345	360-376
334	140-156	346	400-416
335	160-176	347	420-436
336	200-216	350	440-456
337	220-236	351	460-476

Each byte in the table reflects the status of 16_{10} words of memory. The first byte in the table controls locations 0-16, the second byte controls locations 20-36, and so on. The bytes are read from left to right. Thus, if locations 0-3 are protected, the first byte of the table contains:

11000000

Note that only individual words are protected, not bytes. Thus, protecting location 0 always implies that the word at location 0 is protected, meaning both locations 0 and 1. If locations 24 and 26 are protected, the second byte of the table contains:

00110000

since the leftmost bit represents location 20 and the rightmost bit represents location 36. To protect locations 300-306, the leftmost 4 bits of byte 342 must be set:

11110000

resulting in a value of 360 for that byte.

2.5.4.1 S/J Restrictions - The S/J Monitor does not support the .PROTECT request. If users wish to protect vectors, the protection must be done in one of two ways:

1. Manually, with PATCH, or
2. Dynamically (from within the user's program)

To protect locations 300-306 dynamically, the following instructions are used:

```
MOV @#54,R0
BISB #↑B11110000,342(R0)
```

Protecting locations with PATCH implies that the vector is permanently protected, even if the system is re-bootstrapped, while the second method provides a temporary measure and does not hold across bootstraps. However, users are cautioned that the second method involves storing directly into the monitor; for this reason it is recommended that S/J users use method 1.

2.6 USING AUXILIARY TERMINALS AS THE CONSOLE TERMINAL

This section describes how RT-11 can be modified to allow a terminal other than the standard console unit 0 to become the console terminal. This procedure is useful in cases where it is desirable to be able to use different console capabilities at different times (for example, at certain times the hard copy output of an LA30 is required, while at other times the speed of a VT05 is desirable). The only information required to make the alteration is:

- 1) the address of the auxiliary terminal's interrupt vectors, and
- 2) the I/O page addresses of the keyboard and printer status register and buffer.

RT-11 is designed so that all console references are done indirectly through centralized pointers. Thus, changing several system locations causes all operations to be transferred to a new terminal.

For this example, assume that the new terminal's interrupt vectors are at 300,302 and 304,306 and that its I/O page addresses are:

TKS at 177500
TKB at 177502
TPS at 177504
TPB at 177506

Also assume that the new terminal is a parallel interface so that no fill characters are required.

.R PATCH <CR>

PATCH Version number

FILE NAME--

*MONITR.SYS/M<CR>

*BASE;ØR<CR>

*6Ø/ VECTIN<LF>

62/ STATIN<LF>

64/ VECTOUT<LF>

66/ STATOUT<CR>

*3ØØ/ nnnnn VECTIN<LF>

3Ø2/ nnnnn STATIN<LF>

3Ø4/ nnnnn VECTOUT<LF>

3Ø6/ nnnnn STATOUT<CR>

*Ø,xx3Ø4/ 17756Ø 1775ØØ<LF>

Ø,xx3Ø6/ 177562 1775Ø2<LF>

Ø,xx31Ø/ 177564 1775Ø4<LF>

Ø,xx312/ 177566 1775Ø6<CR>

*Ø,xx342\ Ø 36Ø<CR>

*E

:

[The current values for the BASE address and for the input/output vectors and status are in Chapter 5 of Getting Started with RT-11. They must be copied into the new terminal's vectors.] [nnnnn are arbitrary numbers]

[xx = 16 for S/J, 17 for F/B. Modify monitor's central I/O page pointers]

[Protect new vectors]

The bootstrap must also be changed to relocate the new vector locations when the monitor is first loaded into memory. The bootstrap contains a list of items that must be relocated; the list is located at RELLST in the bootstrap code. The exact position of RELLST varies with each monitor and must be obtained from Chapter 5 of Getting Started With RT-11 (V02B). The patching procedure is:

.R PATCH <CR>

PATCH Version number

FILE NAME--

*MONITR.SYS/M<CR>

*RELLST+1Ø/ 6Ø 3ØØ<LF>

RELLST+12/ 64 3Ø4<CR>

*E

.R PIP<CR>

*A=MONITR.SYS/U<CR>

*SY:/O<CR>

[Bootstrap must be rewritten. Rebootstrap; system will appear on new terminal.]

It is also possible to write a user program that would perform this procedure dynamically at run-time. Such a program would modify the monitor's protection map and the central I/O page pointers, then set up locations 300-306 and exit. If done dynamically, the monitor file itself is unchanged; thus when the system is bootstrapped, the console terminal reverts to the usual unit.

2.7 MAKING TTY SET OPTIONS PERMANENT IN F/B MONITOR

The F/B Monitor may be configured for different console terminal requirements by use of the TTY options of the SET command. These changes are not permanent and must be made each time the monitor is bootstrapped. By using the patching procedures in this section, the various options required for the installation may be made a permanent part of the F/B Monitor.

Table 2-4 is a description of the TTY options and their default functions in the F/B Monitor as distributed.

Table 2-4
Default Functions for TTY Options

Option	Default	Description
TAB/NOTAB	NOTAB	Hardware tabs converted to spaces.
CRLF/NOCRLF	CRLF	<CR><LF> inserted if WIDTH reached.
FORM/NOFORM	NOFORM	Form Feed converted to Line Feeds.
FB/NOFB	FB	CTRL F/CTRL B cause context switch.
PAGE/NOPAGE	PAGE	CTRL S holds output, CTRL Q continues it.
SCOPE/NOSCOPE	NOSCOPE	VT05, VT50, VT11 is the console terminal (rubout produces backspace, space, backspace).
WIDTH	72(10)	Width of carriage.

The three options enabled are PAGE, CRLF, and FB. The carriage width is set to 72(10) characters (110 octal).

To permanently change these options, the words TTCNFG, TTWIDT and LISTFB in the F/B Monitor must be patched. The exact locations of these words and the BASE address are found in Chapter 5 of Getting Started with RT-11 (V02B). The numbers used in the following examples are for illustration purposes only and may not be correct for all systems.

2.7.1 Carriage Width

The carriage width is the line width at which the CTRL option generates a carriage return/line feed. This width is changed by patching the word TTWIDT, which for this example is assumed to be located at 21354.

```
.R PATCH <CR>
PATCH Version number
FILE NAME--
*MONITR.SYS/M<CR>           [The /M is necessary; set
*BASE;ØR<CR>                relocation registers; open
*Ø, 21354\_____ 11Ø_____ 2Ø4<CR> with backslash]
*E
:
```

In this example, the width is changed from 72₁₀ to 132₁₀ (204₈).

2.7.2 Other Options

Other options are changed by setting or clearing the appropriate bits in TTCNFG. To determine the new value to be inserted in TTCNFG, Table 2-5 is used. For each option, select the permanent value desired. Add together the octal bit patterns for each value selected to determine the new value of TTCNFG.

Table 2-5
TTCNFG Option Bits

Option	Bit Pattern
TAB	000001
CRLF	000002
FORM	000004
FB	000010
PAGE	000200
SCOPE	100000
Any NO option	000000

For example, the monitor default is PAGE, CRLF and FB. Adding together the bit patterns for PAGE, CRLF and FB produces the octal value 212 (= 200 + 10 + 2).

To change this to SCOPE, PAGE, FB, add together the numbers 100000, 200 and 10 to get 100210, the new value of TTCNFG. Using the location of TTCNFG obtained from Getting Started With RT-11, the procedure is:

```
.R PATCH <CR>
PATCH Version number
FILE NAME--
*MONITR.SYS/M<CR>
*BASE;ØR<CR>
*Ø, TTCNFG/ _____ 212 _____ 1ØØ21Ø<CR>
*E
:
```

If the FB option is changed, an additional step is necessary. Bit 15 of LISTFB must be changed to reflect the new FB option. Bit 15 must be 0 if the option is FB and must be 1 if the option is NOFB. For example, to change the monitor default to FORM, TAB, NOFB, the value of TTCNFG is 5 (4 + 1 + 0), and bit 15 of LISTFB must be a 1. The patch procedure is:

.R PATCH <CR>

PATCH Version number

FILE NAME--

*MONITR.SYS/M<CR>

*BASE;ØR<CR>

*Ø,TTCNFG/2125<CR>

*Ø,LISTFB/33161Ø3316<CR>

*E

:

[The /M is necessary;
set relocation register;
change TTCNFG;
set bit 15 in LISTFB.]

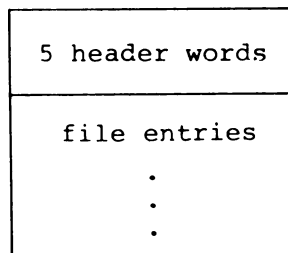
After making any of these patches, it is necessary to bootstrap the system to load the new version of the monitor.

CHAPTER 3
FILE STRUCTURES AND FILE FORMATS

3.1 DEVICE DIRECTORY SEGMENTS

The device directory begins with physical block 6 of any file-structured device and consists of a series of directory segments that contain the names and lengths of the files on that device. The directory area is variable in length, from 1 to 31 (decimal) directory segments. PIP allows specification of the number of segments when the directory is zeroed. The default value is four directory segments. Each directory segment is made up of two physical blocks; thus, a single directory segment is 512 words in length.

A directory segment has the following format:



3.1.1 Directory Header Format

Each directory segment contains a 5-word header block, leaving 507 (decimal) words for directory entries. The contents of the header words are described in Table 3-1.

Table 3-1
Directory Header Words

Word	Contents
1	The number of segments available for entries. This number is specified in PIP when the device is zeroed and must be in the range $1 \leq N \leq 31_{10}$.
2	Segment number of the next logical directory segment. The directory may, in certain cases, be a linked list. This word is the link word between logically contiguous segments; if equal to 0, there are no more segments in the list. Refer to Section 3.2.1, Directory Segment Extensions, for more details on the link word.
3	The highest segment currently open (each time a new segment is created, this number is incremented). This word is updated only in the first segment and is unused in any but the first segment.
4	The number of extra bytes per directory entry. This number can be specified when the device is zeroed with PIP. Currently, RT-11 does not allow direct manipulation of information in the extra bytes.
5	Block number where files in this segment begin.

3.1.2 Directory Entry Format

The remainder of the segment is filled with directory entries. An entry has the following format:

STATUS WORD	
NAME (CHARS 1-3)	
NAME (CHARS 4-6)	
EXTENSION	
TOTAL FILE LENGTH	
JOB #	CH #
DATE	
EXTRA WORDS	
.	
.	
.	

} IN RAD5Ø

} OPTIONAL

Figure 3-1
Directory Entry Format

3.1.2.1 Status Word - The Status Word is broken down into two bytes of data:

Even byte: Reserved for future use.

Odd byte: Indicates the type of entry. Currently RT-11 recognizes the file types listed in Table 3-2:

Table 3-2
File Types

Value	File Type
1	Tentative File, i.e., one that has been .ENTERed but not .CLOSEd. Files of this type are deleted if not eventually .CLOSEd and are listed by PIP as <UNUSED> files.
2	An empty file. The name, extension, and date fields are not used. PIP lists an empty file as <UNUSED> followed by the length of the unused area.

(continued on next page)

Table 3-2 (Cont.)
File Types

Value	File Type
4	A permanent entry. A tentative file that has been .CLOSEd is a permanent file. The name of a permanent file is unique; there can be only one file with a given name and extension. If another exists before the .CLOSE is done, it is deleted by the monitor as part of the .CLOSE operation.
10	End-of-segment marker. RT-11 uses this to determine when the end of the directory segment has been reached during a directory search.

3.1.2.2 Name and Extension - These three words (in .RAD50) represent the symbolic name and extension assigned to a file.

3.1.2.3 Total File Length - The file length consists of the number of blocks currently a part of the file. Attempts to read or write outside the limits of the file result in an End of File error.

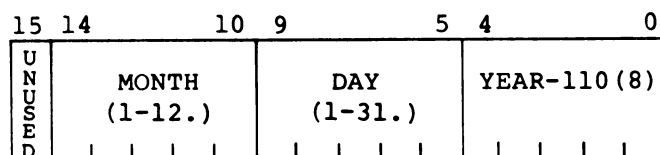
3.1.2.4 Job Number and Channel Number - A tentative file is associated with a job in one of two ways:

1. Under the S/J Monitor, the sixth word of the entry holds the channel number on which the file is open. This enables the monitor to locate the correct tentative entry for the channel when the .CLOSE is given. The channel number is loaded into the even byte of the sixth word.
2. In F/B, the channel number is put into the even byte of the sixth word; in addition, the number of the job that is opening the file is put into the odd byte of the word. This is required to uniquely identify the correct tentative file during the .CLOSE and is necessary because both jobs may have files open on their respective channels; the job number differentiates the tentative files.

NOTE

This sixth word is used only when the file is marked as tentative. Once it becomes permanent, the word becomes unused. Its function while permanent is reserved for future use.

3.1.2.5 Date - When a tentative file is created via .ENTER, the system date word is put into the creation date slot for the file. The date word is in the following format:



3.1.2.6 Extra Words - The number of extra words is determined by the number of extra bytes per entry in the header words. Although PIP provides for allocation and listing of extra words, RT-11 provides no direct facilities for manipulating this extra information. Any user program wishing to access these words must perform its own direct operations on the RT-11 directory.

Figure 3-2 shows a typical RT-11 directory segment:

HEADER BLOCK	4	FOUR SEGMENTS AVAILABLE
	0	NO NEXT SEGMENT
	1	HIGHEST OPEN IS #1
	0	NO EXTRA WORDS/ENTRY
	16	FILES START AT BLOCK 16 ₈
FILE ENTRIES	2000	PERMANENT ENTRY
	51646	RAD5Ø FOR "MON"
	35562	RAD5Ø FOR "ITR"
	75273	RAD5Ø FOR "SYS"
	42	FILE IS 34 ₁₀ (42 ₈) BLOCKS LONG
	0	
	0	NO CREATION DATE
	1000	AN EMPTY ENTRY
	0	(THE NAME AND EXTENSION OF AN
	0	EMPTY IS NOT IMPORTANT
	0	
	100	64 _{1Ø} (1ØØ ₈) BLOCKS LONG
	0	
	0	
	2000	PERMANENT
62570	RAD5Ø FOR "PIP"	
0		
50553	RAD5Ø FOR "MAC"	
11	FILE IS 9 ₁₀ (11 ₈) BLOCKS LONG	
0		
0	NO CREATION DATE	
4ØØ	TENTATIVE FILE ON CHANNEL 1	
62570	RAD5Ø FOR "PIP"	
0		
50553	RAD5Ø FOR "MAC"	
20		
1	JOB #, CHANNEL #	
0		
1000	EVERY TENTATIVE MUST BE FOLLOWED BY	
0	AN EMPTY ENTRY	
0		
0		
1020	FILE IS 528 _{1Ø} (1Ø2Ø ₈) BLOCKS LONG	
0		
0		
4000	END OF DIRECTORY SEGMENT	

Figure 3-2
Directory Segment

When the tentative file PIP.MAC is .CLOSED, the permanent file PIP.MAC is deleted.

To find the starting block of a particular file, first find the directory segment containing the entry for the desired file. Then take the starting block number given in the fifth word of that directory segment and add to it the length of each file in the directory before the desired file. For example, in Figure 3-2, the permanent file PIP.MAC will begin at block number 160 (octal).

3.2 SIZE AND NUMBER OF FILES

The number of files that can be stored on an RT-11 device depends on the number of segments in the device's directory and the number of extra words per entry. The maximum number of directory segments on any RT-11 device is 31_{10} . This theoretically leaves room for a maximum of:

$$31 \times \left[\frac{512-5}{7+N} \right]$$

directory entries, where N equals the number of extra information words per entry. If N=0, this indicates that the maximum is 2232_{10} entries.

If files are added sequentially (that is, one immediately after another) without deleting any files, roughly one half the total number of entries will fit on the device before a directory overflow occurs. This results from the way filled directory segments are handled.

When a directory segment becomes full and it is necessary to open a new segment, approximately one half the entries of the filled segment are moved to the newly-opened segment (this process is illustrated in Section 3.2.1); thus, when the final segment is full, all previous segments have approximately one half their total capacity. If this process were not done and a file was deleted from a full segment, the space from the deleted file could not be reclaimed. Every tentative file must be followed by an empty entry (for recovering unused blocks when the file is made permanent). Though only one file is deleted, two entries (tentative and empty) are needed to reclaim the space.

If files are continuously added to a device, the maximum number of entries will be:

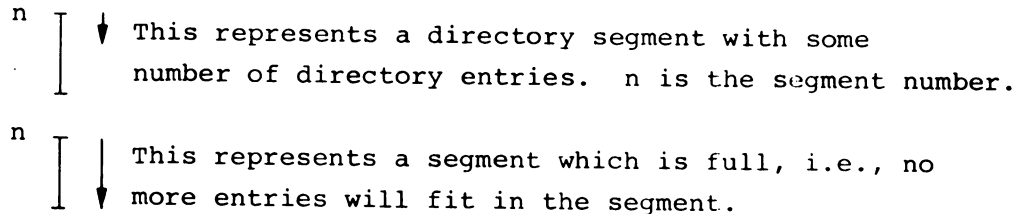
$$(M+1) \left[\frac{507}{2(7+N)} \right]$$

where M equals the number of segments available on the device and N equals the number of extra words.

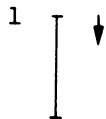
The theoretical total can be realized by compressing the device (using the PIP /S operation) when the directory fills up. PIP packs the directory segments as well as the physical device.

3.2.1 Directory Segment Extensions

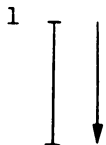
RT-11 allows a maximum of 31 (decimal) directory segments. This section covers the processing of a directory segment. For illustrative purposes, the following symbols are used:



Systems start out with entries entered into segment 1:



As entries are added, segment 1 fills:



When this occurs and an attempt is made to add another entry to the directory, the system must open another directory segment. If another segment is available, the following occurs:

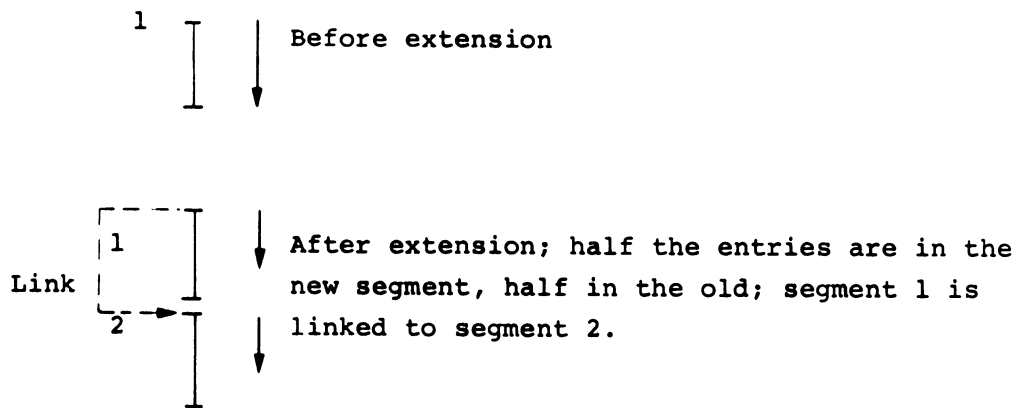
1. one half of the entries from the filled segment are put into the next available segment,
2. the shortened segment is re-written to the disk,
3. the directory segment links are set, and
4. the file is entered in the newly created segment.

NOTE

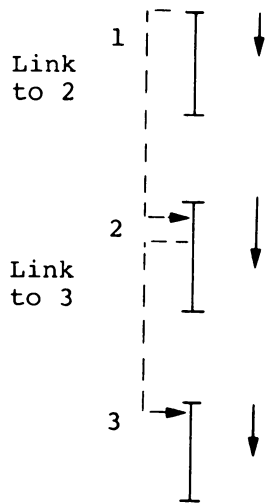
If the last segment becomes full and an attempt is made to enter another file, a fatal error occurs and an error message is generated:

?M-DIR OVFLO?

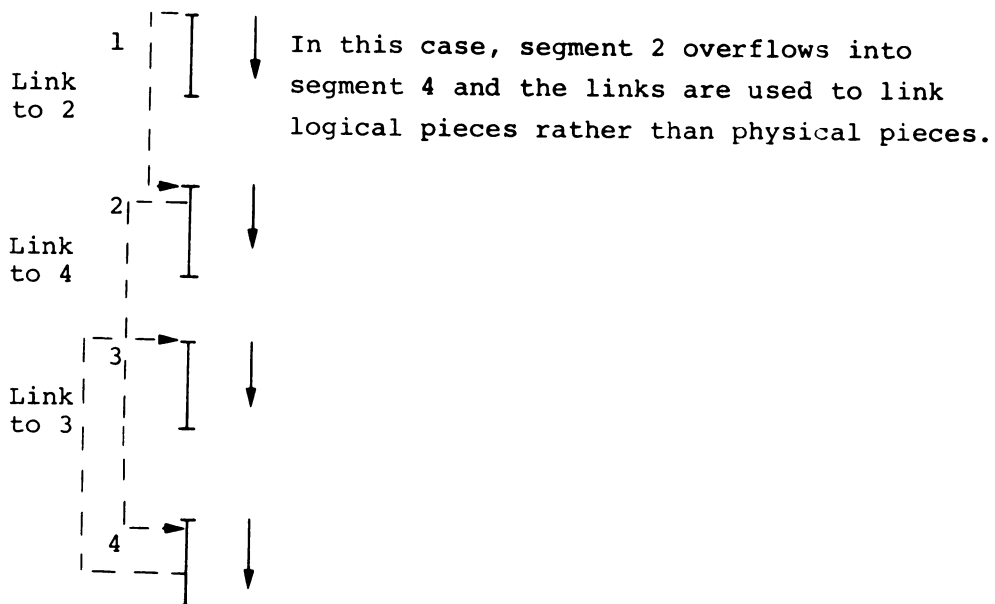
Thus, in the normal case, the segment appears as:



If many more files are entered, they fill up the second segment and overflow into the third segment, if it is available:



In this case, the links between the segments are not strictly necessary, as the segments are contiguous. However, the links do become necessary if a large file is deleted from segment 2 and many small files are entered, since it would then be possible to overflow segment 2 again. If this occurred and a fourth segment existed, the directory would appear:



3.3 MAGTAPE AND CASSETTE FILE STRUCTURE

3.3.1 Magtape File Structure

This section covers the magtape file structure as implemented in RT-11, Version 2B. The structure is slightly different from that of Version 2. However, RT-11 V02B can read magtapes written under Version 2.

RT-11 magtapes use a subset of the VOL1, HDR1, and EOF1 ANSI standard labels. Each magtape file has the format:

```
HDR1*---data---*EOF1*
```

where each asterisk represents a tape mark.

A volume containing a single file has the following format:

```
VOL1 HDR1*---data---*EOF1**
```

A volume containing two files has the following format:

```
VOL1 HDR1*---data---*EOF1*HDR1*---data---*EOF1**
```

A double tape mark following an EOF1 label indicates logical end of tape.

Each label occupies the first 80 bytes of a 256-word physical block, and each byte in the label contains an ASCII character (i.e., if the content of a byte is listed as '1', the byte contains the ASCII code for '1'). Table 3-3 shows the contents of the first 80 bytes in the three labels. Note that VOL1, HDR1, and EOF1 each occupy a full 256-word block, of which only the first 80 bytes are meaningful.

The meanings of the table headings are:

```
CP - character position in label  
Field Name - reference name of field  
L - length of field in bytes  
Content - content of field
```

Table 3-3
ANSI MT Labels Under RT-11

Volume-Header Label (VOL1)			
CP	Field Name	L	Content
1-3	Label identifier	3	VOL
4	Label number	1	1
5-10	Volume identifier	6	RT1101
11	Accessibility	1	Blank
12-37	(Reserved)	26	Blanks
38-51	Owner identifier	14	DD%% {used to indicate an RT-11 MT to RSX-11D
52-79	(Reserved)	28	Blanks
80	Label-Standard Version	1	1
First File Header Label (HDR1)			
CP	Field Name	L	Content
1-3	Label identifier	3	HDR
4	Label number	1	1
5-21	File identifier	17	6-character ASCII file name, followed by '.', followed by 3-character ASCII file extension; left justified, remainder of field is blanks
22-27	File Set identifier	6	RT1101
28-31	File Section Number	4	0001
32-35	File Sequence Number	4	0001
36-39	Generation Number	4	0001
40-41	Generation Vsn Number	2	00
42-47	Creation Date	6	Blank then year*1000+day of year in ASCII (YYDDD); e.g., 2/1/75=Δ75032
48-53	Expiration Date	6	blank then 00000
54	Accessibility	1	blank
55-60	Block Count	6	000000
61-73	System Code	13	RT11 left-justified followed by blanks
74-80	(Reserved)	7	blanks
First End-of-File Label (EOF1)			
Same as HDR1 except that the label identifier (CP 1-3) is <i>EOF</i> , not <i>HDR</i> , and the block count field (CP 55-60) contains the number of blocks in the file as a decimal value encoded in ASCII characters (for example, if the file was 12 blocks long, the block count field would be 00012).			

3.3.1.1 Moving MT to Other Industry-Compatible Environments - RT-11 V02B magtapes may be read by RSX-11D Version 6. RT-11 magtapes should be mounted, under RSX-11D, by using the /OVR switch of the MOUNT command, or by specifying a volume label of "RT1101". RSX-11D Version 6 will not allow the user to write on RT-11 V02B magtapes once they have been mounted. RT-11 V02B can read RSX-11D Version 6 magtapes, but RT-11 users should not attempt to write on tapes created by RSX-11D. Users should note that data structures differ between the two systems and these differences must be handled by the user.

RT-11 V02B magtapes may be read on IBM systems that support ANSI standard label processing. RT-11 V02B magtapes to be read by IBM systems should consist of single file volumes (one file per magtape). Important JCL parameters for reading RT-11 V02B tapes under an IBM OS system are as follows:

(In the DD statement of the Job Control Language)

DISP = OLD

LABEL = (01,AL,,IN)

VOL = (,RETAIN,SER=RT1101)

DSN = RTFILE.MAC

BLKSIZE = 512

DEN = 2 (for 800 bpi 7-track or 9-track tape)

The DSN parameter is the Data Set Name or the RT-11 filename and extension. Files to be moved to other systems should be created with full six-character filenames and three-character extensions.

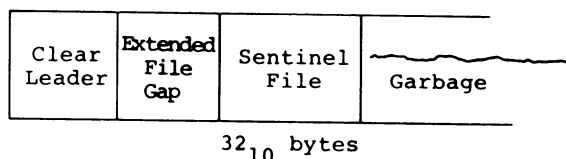
3.3.1.2 Recovering From Bad Tape Errors - When a bad tape error occurs on magtape, the magtape handler will retry the desired function, and, if the error persists, will attempt to save the tape's file structure. It does this on writes, for example, by retrying the write 10 times, using the write with extended file gap to space past the bad tape. If, after retrying, the error still exists, the file will be closed, containing all data written prior to the write on which the error occurred. The user should still be able to write additional files on the tape, since the bad portion of the tape will be within the area of the closed file.

If a bad tape error occurs when writing the file header during ENTER, and retry fails, the handler writes logical end of tape after the previous file on the tape. The remainder of the tape can be accessed only if the last complete file on the tape can be extended (or overwritten by a file of different length) so that the bad tape error does not occur on the file header when a subsequent file is ENTERed.

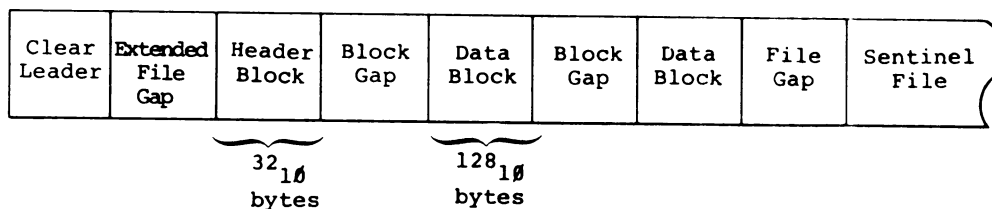
If a bad tape error occurs while writing the end of file label (EOF1) during CLOSE, the handler writes a triple tape mark to signify end of file and logical end of tape. Additional files can be added to the tape only if the last complete file can be extended (or overwritten by a file of different length) so that the bad tape error does not occur at the EOF1 label.

3.3.2 Cassette File Structure

A blank (newly initialized) cassette appears in the format:

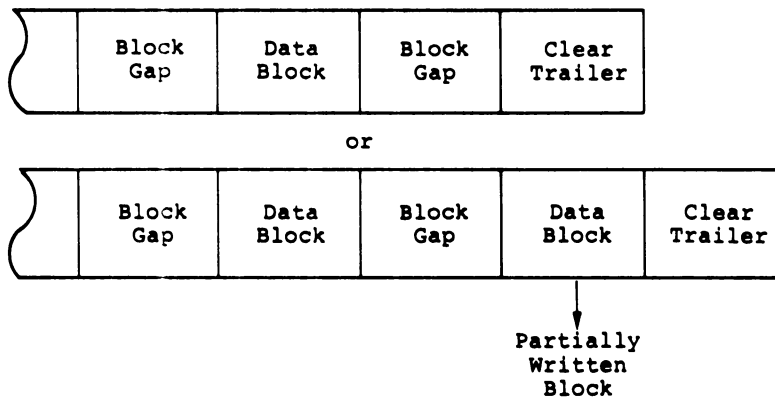


while a cassette with a file on it appears as:



Files normally have data written in 128₁₀-byte blocks. This can be altered by writing cassettes while in *hardware* mode. (In hardware mode, the user program must handle the processing of any headers and sentinel files; in *software* mode the handler automatically does this. Refer to Appendix H of the RT-11 System Reference Manual.)

The preceding diagram shows a file terminated in the usual manner (by a sentinel file). However, the physical end of cassette may occur before the actual end of the file. This format appears as:



In the latter case, for multi-volume processing the partially written block must be the first data block of the next volume.

3.3.2.1 File Header - The File Header is a 32₁₀-byte block that is the first block of any data file on a cassette. If the first byte of the header is null, the header is interpreted as a sentinel file, which is an indication of logical end of cassette. The format of the header is described in Table 3-4.

Table 3-4
CT File Header Format

Byte Number	Contents
0-5	File name in ASCII characters (ASCII is assumed to imply a 7-bit code)
6-8	Extension in ASCII characters
9	Data type (0 for RT-11)
10,11	Block length of 128_{10} (200_8); Note: byte 10=0 (high-order), byte 11= 200_8 (low-order)
12	File sequence number. (0 for a single-volume file or the first volume of a multi-volume file; successive numbers are used for continuations)
13	Level 1; this byte is a 1
14-19	Date of file creation (6 ASCII digits representing day (01-31); month (01-12), and last two digits of the year; 0 or 40_8 in first byte means no date present)
20,21	Zero
22	Record attributes (0 in RT-11 cassettes)
23-28	Reserved for future use
29-31	Reserved for user

3.4 RT-11 FILE FORMATS

3.4.1 Object Format (.OBJ)

An object module is a file containing a program or routine in a binary, relocatable form; object files normally have an .OBJ extension. Object modules are produced by language processors (such as MACRO or FORTRAN) and are processed by the Linker to become a runnable program (in SAV, LDA, or REL format, discussed later). Object files may also be processed by the Librarian to produce library .OBJ files, which are then used by the Linker. Figure 3-3 illustrates this process.

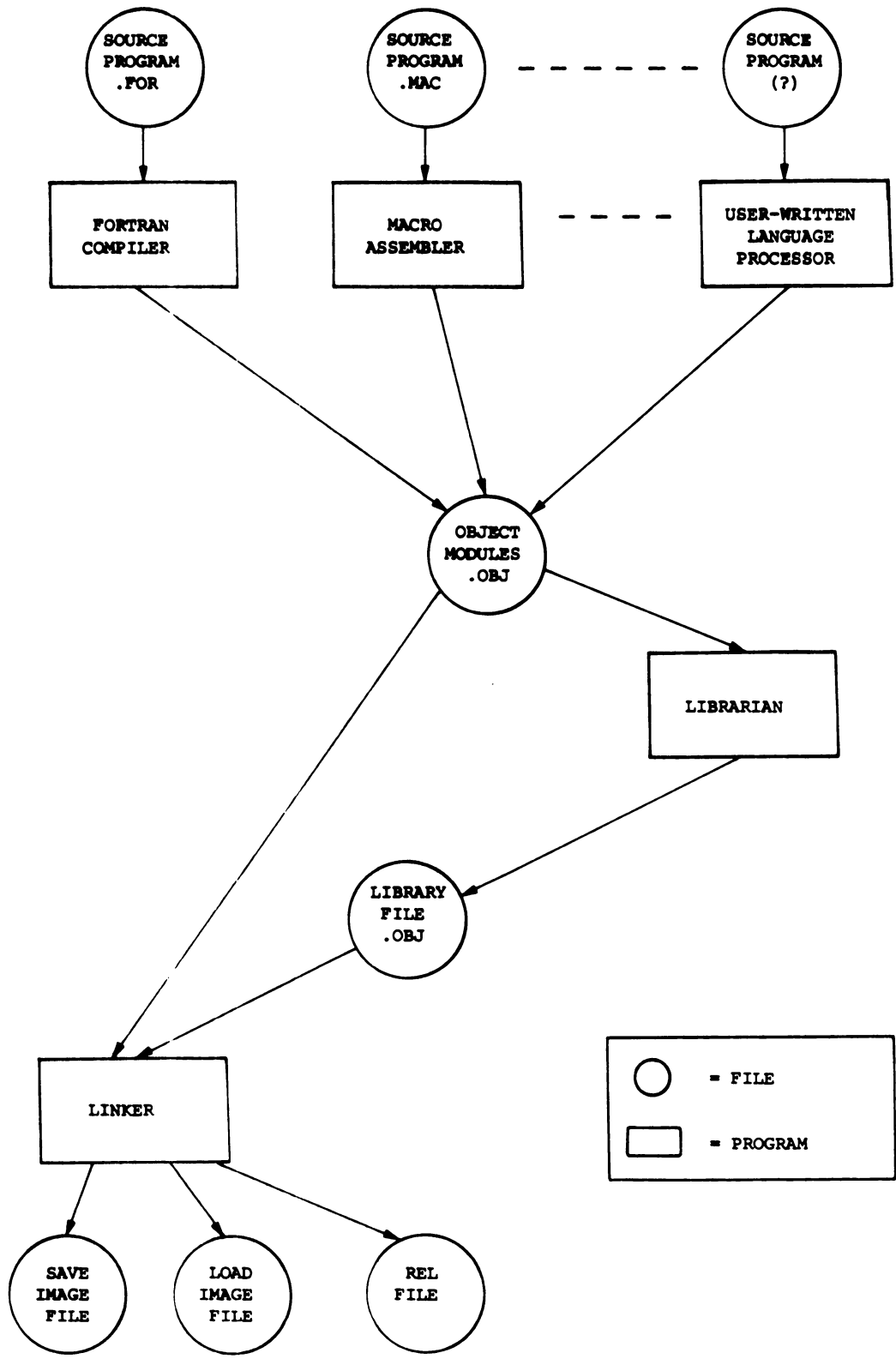


Figure 3-3
Object Module Processing

Many different object modules may be combined to form one file; each object module remains complete and independent. However, object modules combined into a library by the Librarian are no longer independent -- they become part of the library's structure.

Object modules are made up of formatted binary blocks. A formatted binary block is a sequence of 8-bit bytes (stored in an RT-11 file, on paper tape, or by some other means) and is arranged as illustrated in Figure 3-4.

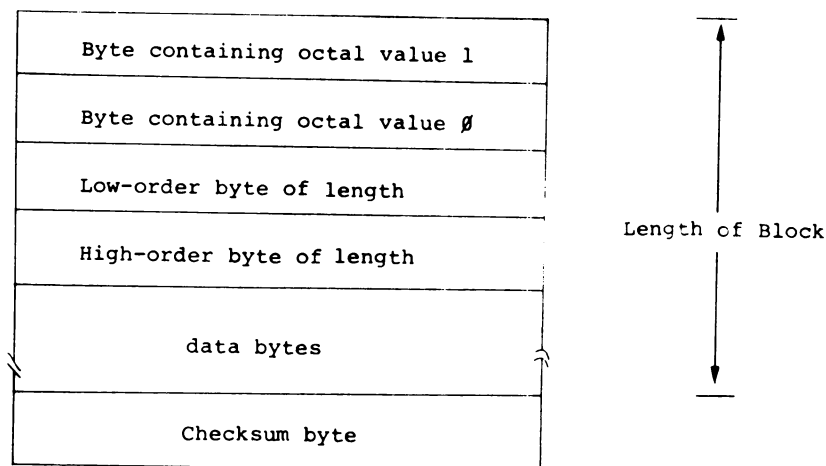


Figure 3-4
Formatted Binary Block

Each formatted binary block has its length stored within it; the length includes all bytes of the block except the checksum byte. The data portion of each formatted binary block contains the actual object module information (described later). The checksum byte is computed such that the sum of all bytes in the formatted binary block, including the checksum byte, is zero when the sum is masked to 8 bits.

Formatted binary blocks are used to hold various kinds of information in an object module; this information is always contained completely in the data portion of the block, surrounded by the formatted binary block structure.

Eight types of data blocks may be present in an object module:

<u>Identification Code</u>	<u>Type of Block</u>	<u>Function</u>	
1	GSD blocks	hold the Global Symbol Directory information	
2	ENDGSD block	signals the end of GSD blocks in a module	
3	TXT blocks	hold the actual binary "text" of the program	
4	RLD blocks	hold Relocation Directory information	
5	ISD blocks	hold Internal Symbol Directory - not supported by RT-11	
6	ENDMOD block	signals end of the object module	
7	Librarian Header Block	17 words holding the status of the library file	} Library File Only
10	Librarian End Block	signals the end of the library file	

The structure of object modules produced by a language processor will be described first, followed by details specific only to Library .OBJ files.

The first block of an object module must be a GSD block, and all GSD blocks must appear before the ENDGSD block. The ENDMOD block must be the last block of the module. Except for these three restrictions, blocks may appear in any order within an object module.

When a 16-bit word is stored as part of the data in a block, it is always stored as two consecutive 8-bit bytes, with the low-order byte first.

The first word (data word) of each type of block mentioned above contains the identification code of that block type (1 = GSD block, etc.) with any information present following the identification word.

3.4.1.1 Global Symbol Directory - The object module's global symbol directory contains the following information:

- 0 - Module Name
- 1 - Program Section (CSECT) Definitions
- 2 - Internal Symbol Table Name (not supported by RT-11)
- 3 - Transfer (Start) Address
- 4 - Global Symbol Definitions or References

Each piece of information in the GSD is contained in a *GSD item*, formatted as shown in Figure 3-5:

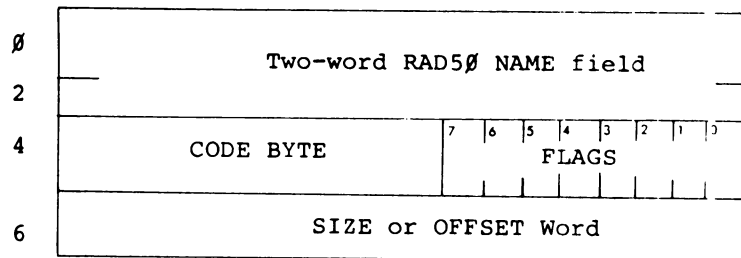


Figure 3-5
GSD Structure

The code byte identifies the information contained in a GSD item according to the codes listed above (0 = Module Name, 1 = Program Section Definition, etc.). The first GSD item of an object module must contain the Module Name information (FLAGS, CODE, and SIZE = 0). There may be no more than five GSD items per GSD block (i.e., per formatted binary block). As many GSD blocks as necessary may be present, but all must appear before the ENDGSD block. GSD blocks need not be contiguous.

Flags are coded as follows:

- Bits 0,1,2,4,7 unused
- Bit 3: 0 = undefined, 1 = defined (used only with Global Symbols)
- Bit 5: 0 = absolute, 1 = relocatable
- Bit 6: 0 = internal, 1 = global

All program sections (CSECTs) defined in a module must be declared in GSD items (code byte = 1). The size word of each program section definition should contain the size in bytes to be reserved for the section. Program sections may be declared more than once, in which case the largest declared size of the section will be used. All global symbols that are defined in a given program section must appear in the GSD items immediately following the definition item of that program section.

A special program section named ". ABS." (where represents a space) is called the absolute section. The absolute section has the special attribute that it is always allocated by the Linker beginning at location 0 of memory. All global symbols that contain absolute (non-relocatable) values should be declared immediately after the GSD item that defines the absolute section. If it is not desired to allocate any memory space to the absolute section, its size word may be specified as zero, even if absolute global symbol definitions occur after it. Flag bit 5 of each absolute global symbol is always set to zero. GSD items that contain the definitions of global symbols (code byte = 4) must immediately follow the program section declaration into which they are to be defined. Flag bit 3 is set to 1 to indicate a symbol definition, bit 5 is set if and only if the symbol is relocatable, and bit 6 is set to indicate that the symbol being defined is a global. In addition, the offset word is set to contain the defined value of the global symbol, relative to the base of the program section in which the global is defined. At link time, the Linker assigned section base is added to get the final value of the global symbol.

Global symbols that are referenced but not defined in the current object module must also appear in GSD items. These *global references* may appear in any GSD item except the very first (which contains the module name). Global references are recognized by code byte 4 with flag bit 3=0, bit 5 is undetermined, and bit 6=1. All global symbols used in the RLD of the object module (described later) must appear in at least one Global Symbol or Program Section GSD item.

If RT-11 is to begin execution of a program within a particular object module of that program, then the information on where to start is given in a Transfer Address (code=3) GSD item. The first even transfer address encountered by the Linker will be passed to RT-11 as the program start address. Whenever the resulting program is run (using R or RUN

for SAV images, FRUN for REL files, or the absolute loader for LDA files), the start address is used to indicate the first executable instruction. If no transfer address is present or if all are odd, the resulting program will not self-start when run. In a Transfer Address GSD item, the name field is used to specify a program section (or global name) and the offset word is used to indicate the offset from the base of that program section (or global) to the starting point of the program. The program section or global name referenced need not be defined in the current object module, but must be defined in some object module included at link time.

NOTE

Program Section and Global names must begin with an alphabetic or numeric character, except for the names `._ABS.` and `□□□□□□.`

3.4.1.2 ENDGSD Block - The ENDGSD block contains a single data word, and that is the identification code of the ENDGSD block (2). All GSD blocks in an object module must precede the ENDGSD block.

3.4.1.3 TXT Blocks and RLD Blocks - The first TXT block (3) in an object module (if present) must be preceded by an RLD block (4).

TXT blocks contain the actual binary form of the programs and are formatted as shown in Figure 3-6:

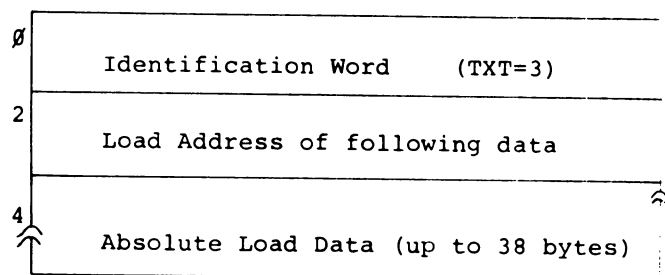


Figure 3-6
TXT Block Format

The load address of a TXT block gives the relative address of the first byte of the absolute load data. The address is relative to the base of the last program section given in a Location Counter Definition RLD command (explained later).

The Absolute Load Data contains the actual bytes that will be loaded into memory when the program is run (except for relocations, described later).

RLD blocks contain variable length RLD commands, used to modify and complete the information contained in TXT blocks. Except for the Location Counter commands, RLD information must appear in an RLD block immediately following the TXT block to be modified.

Available RLD commands are:

1. Internal Relocation
2. Global Relocation
3. Internal Displaced Relocation
4. Global Displaced Relocation
5. Global Additive Relocation
6. Global Additive Displaced Relocation
7. Location Counter Definition
8. Location Counter Modification (not used by RT-11)
9. Set Program Limits

The location counter commands (numbers 7 and 8) are the only two RLD commands that must appear in an RLD block preceding the text blocks modified. The first RLD block must precede the first TXT block and must contain only a location counter definition command (7) in order to declare a program section for loading the first text block. (The location counter modification command (8) is included for compatibility with other systems, but is not used by RT-11.)

The data portion of an RLD block must not be larger than 42_{10} bytes including the identification word (RLD=4) and all RLD commands.

All global names and program section names that appear in RLD commands must appear in GSD items in the same object module. Figure 3-7 shows the format of each RLD command (each part except the first word is optional and may not appear in some commands):

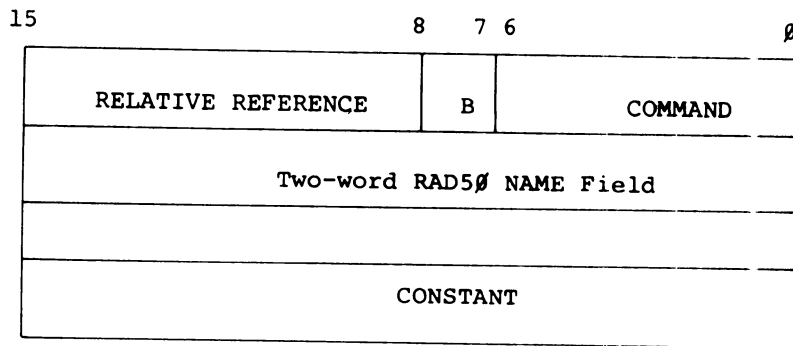


Figure 3-7
RLD Format

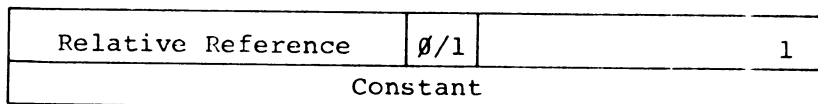
An RLD command may be 1, 2, 3, or 4 words long.

The Command Field contains the command code (1 = Internal Relocation, etc.). The Command Field occupies bits 0-6 of the first word of the command. The B field (bit 7) indicates a word command if 0 or a byte command if 1 (only valid for commands 1 through 6). The Relative Reference Field is a pointer into the preceding TXT block and is used with RLD commands that require text locations for modification (commands 1 through 6 and 9). This field specifies the displacement from the beginning of the preceding TXT block to the referenced text data byte (or word). The beginning of the TXT block is the identification word (the first word of the data portion of the block). Thus, the smallest relative reference will normally be 4 (the first byte (word) of the preceding TXT block).

The Name Field is used to hold a Global or Program Section name if the command requires it.

The Constant Field is used to hold a relative address or additive quantity if the command requires it. RLD commands are processed by the Linker as shown in the following situations:

1. Internal Relocation (code 1) - Add the current program section's base to the specified constant and place the result where indicated. This command relocates a direct pointer to an internal relocatable symbol.



Examples:

- a) .WORD LOCAL
- b) MOV #LOCAL,%Ø

2. Global Relocation (code 2) - Place the value of the specified global symbol where indicated. This command generates a direct pointer to an external symbol.

Relative Reference	Ø/1	2
Global Name		

Examples:

- a) .WORD GLOBAL
- b) MOV #GLOBAL,RØ

3. Internal Displaced Relocation (code 3) - Calculate the displacement from the position of the current location plus two to the specified absolute address, and store the result where indicated. This command occurs only when there is a reference to an absolute (non-relocatable) location from a relocatable section.

Relative Reference	Ø/1	3
Constant		

Examples:

- | | | |
|---------------|---|---|
| a) ABS=17755Ø | } | both addresses cause internal displaced relocation to occur |
| TST ABS | | |
| b) CLR 17755Ø | | |

4. Global Displaced Relocation (code 4) - Calculate the displacement from the current location plus two to the specified global address, and store the result where indicated.

Relative Reference	Ø/1	4
Global Name		

Example:

```
.GLOBL GLOBAL
MOV GLOBAL,RØ
```

5. Global Additive Relocation (code 5) - Add the value of the specified global symbol to the specified constant, and store the result where indicated.

Relative Reference	Ø/1	5
Global Name		
Constant		

Example:

```
.GLOBL GLOBAL
CMP #GLOBAL+6,RØ
```

6. Global Additive Displaced Relocation (code 6) - Calculate the displacement from the current location plus two to the address specified by the sum of the global symbol value and the given constant, and place the result where indicated.

Relative Reference	Ø/1	6
Global Name		
Constant		

Example:

```
.GLOBL GLOBAL
CLR GLOBAL+6
```

7. Location Counter Definition (code 7) - This command is used to specify the program section into which the following TXT blocks are to be loaded.

7
Program Section Name
Constant

This command is generated whenever .ASECT or .CSECT is used to initiate or continue a program section. The constant word is effectively ignored by RT-11 and may be used for diagnostic purposes to indicate the relative point at which a program section is being entered.

8. Location Counter Modification (code 10₈) - This command is used to enter the current program section at a different point. This command is effectively ignored by RT-11 and is used for diagnostic purposes only.

10
Constant

Examples:

- a) . = 100 ; IF WE ARE IN THE ASECT
 b) . = -20 ; IF WE ARE IN A RELOCATABLE SECTION

9. Set Program Limits (code 11₈) - This command (generated by the .LIMIT assembler directive) causes two words in the preceding TXT block to be modified. The first word is to be set to the lowest relocated address of the program. The second word is to be set to the address of the first free location following the relocated code. Note that both words to be modified must appear in the same TXT block.

Relative Reference	11
--------------------	----

In addition to the above commands, note that commands numbered 14₈, 15₈, and 16₈ can be generated by MACRO. These commands are identical to commands 4, 5, and 6 respectively, but are used when the *global* is really a program section name.

3.4.1.4 ISD Internal Symbol Directory - Not supported by RT-11.

3.4.1.5 ENDMOD Block - Every object module must end with an ENDMOD block. The ENDMOD block contains a single data word -- the identification code of the ENDMOD block (6).

3.4.1.6 Librarian Object Format - A library .OBJ file contains information additional to that previously defined. The object modules in a library file are preceded by a Library Header Block and Library Directory, and are followed by the Library End Block or trailer. This is illustrated in Figure 3-8.

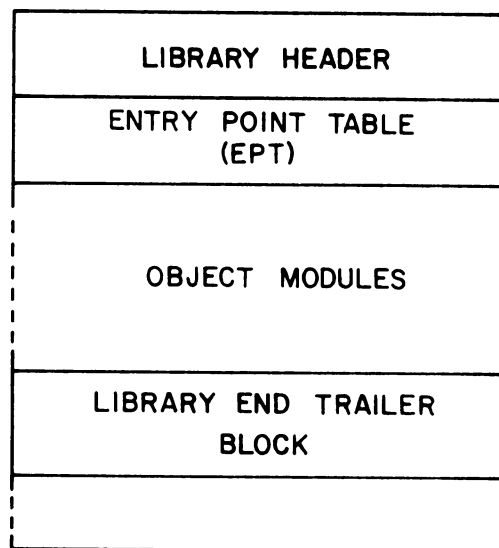


Figure 3-8
Library File Format

Diagrams of each component in the library file structure are included here, but Chapter 7 of the RT-11 System Reference Manual should be consulted for details.

The library header is composed of 17₁₀ words describing the status of the file. The contents of the 17 words are shown in Figure 3-9.

1	} FORMATTED BINARY BLOCK HEADER
56 ₈	
7	LIBRARIAN CODE
X	VERSION NUMBER
0	RESERVED
X	MONTH-DAY-YEAR (OR Ø IF NO DATE)
0	} RESERVED
0	
0	
0	
0	
12 ₈	EPT RELATIVE START ADDRESS
X1	EPT ENTRIES ALLOCATED IN BYTES
0	EPT ENTRIES AVAILABLE (NOT USED IN VI)
X2	NEXT INSERT RELATIVE BLOCK NUMBER
X3	NEXT BYTE WITHIN BLOCK
0	NOT USED (MUST BE ZERO)

Figure 3-9
Library Header Format

The Entry Point Table (EPT), Figure 3-10, is composed of four-word entries which contain information related to all object modules in the library file.

0	SYMBOL CHARS 1-3 (RAD5Ø)		
2	SYMBOL CHARS 4-6 (RAD5Ø)		
4		ADDRESS OF BLOCK	BIT 15=1-MODULE NAME Ø-CSECT OR ENTRY POINT NAME RELATIVE BYTE MAXIMUM=777 ₈ CSECTS MAXIMUM =177 ₈
6	# OF CSECTS IN OBJECT MODULE	RELATIVE BYTE IN BLOCK	

Figure 3-10
Entry Point Table Format

Object modules follow the Entry Point Table and consist of the types of data blocks already discussed: GSD, ENDGSD, TXT, RLD, and ENDMOD. The information in these blocks is used by the Linker during creation of the load module.

Following all object modules is a specially coded Library End Block (trailer), which signifies the end of the file, shown in Figure 3-11.

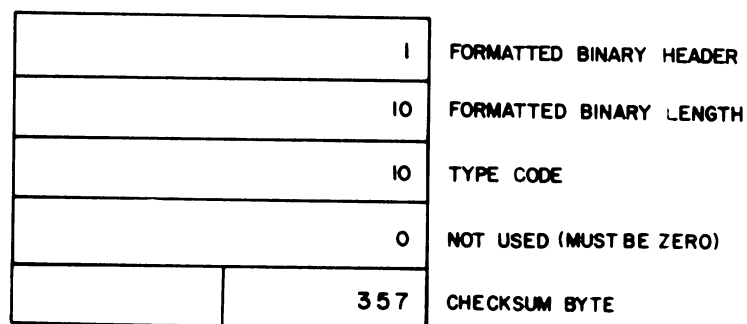


Figure 3-11
Library End Trailer

3.4.2 Formatted Binary Format (.LDA)

The Linker /L switch produces output files in a paper tape compatible binary format.

Paper tape format, shown in Figure 3-12, is a sequence of formatted binary blocks (as explained in Section 3.4.1 and in Figure 3-4). Each formatted binary block represents the data to be loaded into a specific portion of memory. The data portion of each formatted binary block consists of the absolute load address of the block followed by the absolute data bytes to be loaded into memory beginning at the load address. There may be as many formatted binary blocks as necessary in an LDA file. The last formatted binary block of the file is special; it contains only the program start address in its data portion. If this address is even, the loader passes control to the loaded program at this address. If it is odd, the loader halts upon completion of loading. The final block of the LDA file is recognized by the fact that its length is 6.

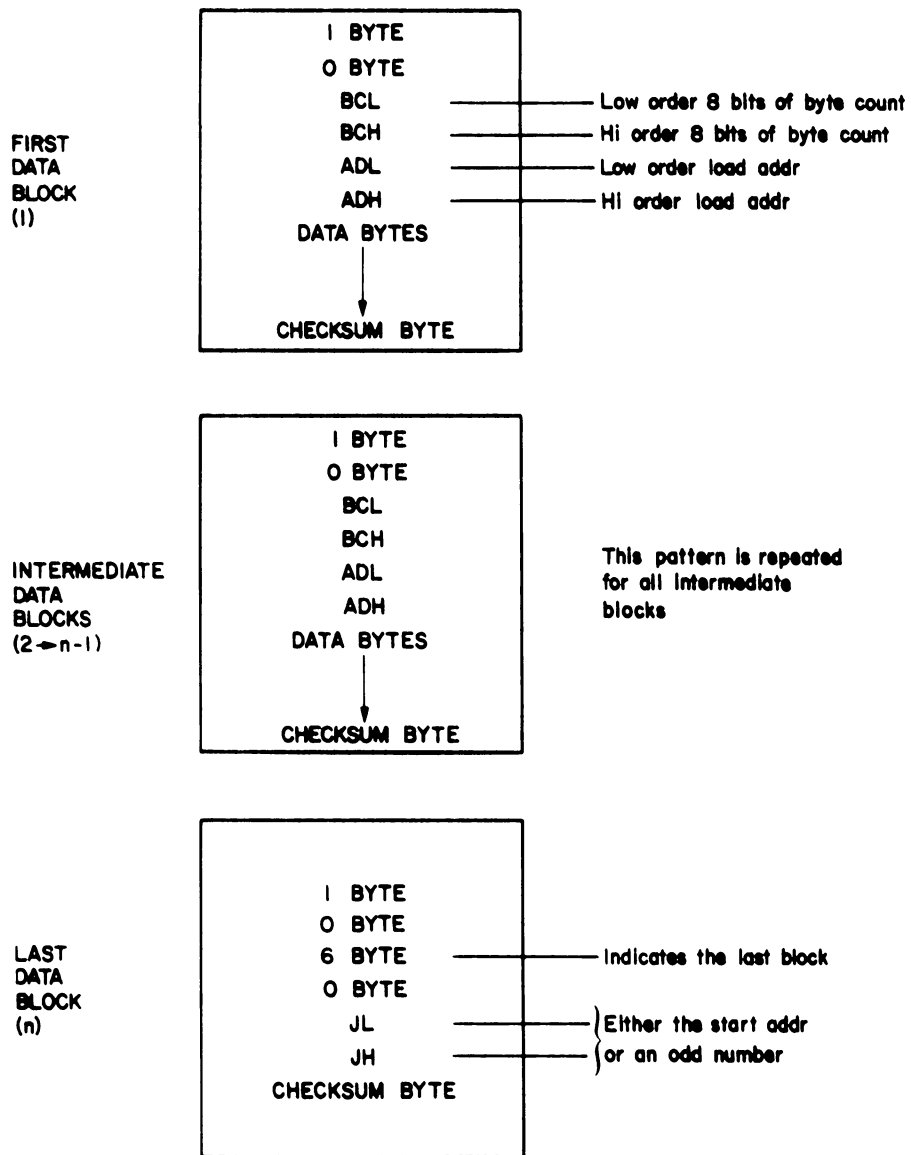


Figure 3-12
Formatted Binary Format

The load module's binary blocks contain only absolute binary load data and absolute load addresses; all global references have been resolved and the appropriate relocation has been performed by the Linker.

3.4.3 Save Image Format (.SAV)

Save image format is used for programs that are to be run in the background. This format is essentially an image of the program as it would appear in memory (block 0 of the file corresponds to memory locations 0-776, block 1 to locations 1000-1776, and so forth).

Locations 360-377 in block 0 of the file are restricted for use by the system. The Linker stores the program memory usage bits in these eight words. Each bit represents one 256-word block of memory and is set if the program occupies that block of memory. This information is used by the R, RUN, and GET commands when loading the program.

When loading a save image program into memory, KMON reads block 0 of the file to extract the memory usage bits. These bits are used to determine whether the program will overlay either the KMON or the USR. If these portions of the monitor will not be overlaid, the entire program is loaded; if the USR and KMON must swap, KMON loads the resident portion of the program, up to the start of KMON. It then puts the portion of the program that overlays KMON/USR into the system swap blocks. When the program starts, the monitor swaps in the virtual portion of the program, overlaying KMON.

When block 0 of a save image file is loaded, each word is checked against the protection bit map (LOWMAP), which is resident in RMON. Locations that are protected in the map, such as location 54 and the system device vectors, are not loaded.

3.4.4 Relocatable Format (.REL)

A foreground job is linked using the Linker /R switch. This causes the Linker to produce output in a linked, relocatable format, with a REL file extension.

The object modules used to create a REL file have been linked and all global references have been resolved. The REL file is not relocated, so it has an effective start address of 0, with relocation information included to be used at FRUN time. The relocation information in the file is used to determine which words in the program must be relocated when the job is installed in memory.

In order to determine if the code to be relocated (as indicated in the relocation information blocks) is to have positive or negative relocation (relative to the start address of the program), the following criteria from the text modification commands is used (R = relative address, G = global address, C = constant):

1. Internal Relocation (.WORD R) - always positive relocation (absolute)
2. Global Relocation (.WORD G) - positive relocation only if the global is not absolute

3. Internal Displaced Relocation - always negative relocation (MOV 54,R)
4. Global Displaced Relocation - negative relocation only where the global is defined as absolute elsewhere (MOV G,R)
5. Global Additive Relocation - same as 2 above (.WORD G + C)
6. Global Additive Displaced - same as 4 above (MOV G + C,R)
7. Program Counter Commands - not applicable
8. Set Program Limits - always positive relocation (requires 2 RELs; limit is two words)

There are two types of REL files to consider, those programs with overlay segments and those without.

3.4.4.1 Non-Overlay Programs - A REL file for a non-overlaid program appears as shown in Figure 3-13:

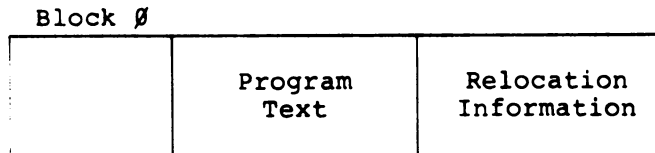


Figure 3-13
REL File Without Overlays

Block 0 (relative to start of the file) contains certain information required by the FRUN processor:

<u>Offset from Beginning of Block 0</u>	<u>Contents</u>
52	Size of the program root segment in bytes
54	Size of the overlay region in words; 0 if no overlays
56	REL file identification word, which must contain the RAD50 value of the characters 'REL'
60	Relative block number of relocation information

In addition, the system communication locations (34-50) contain the following information:

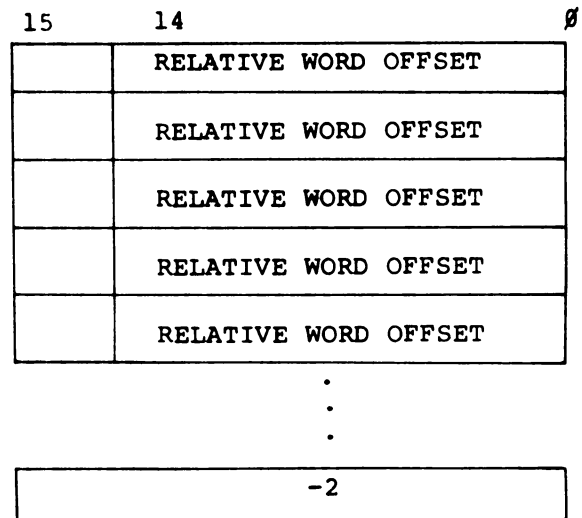
<u>Offset from Beginning of Block 0</u>	<u>Contents</u>
34,36	TRAP vector
40	Start address of program
42	Initial setting of stack pointer
44	Job Status Word
46	USR swap address
50	Highest memory address in user's program

In the case of non-overlaid programs, the FRUN processor performs the following general steps to install a foreground job.

1. Block 0 of the file is read into an internal monitor buffer.
2. The amount of memory required for the job is obtained from location 52 of block 0 of the file, and the space is allocated.
3. The program text is read into the space just allocated for it.
4. The relocation information is read into an internal buffer.
5. The locations indicated in the relocation information area are relocated by adding the relocation quantity, which is the starting address the job occupies in memory.

The relocation information consists of a list of addresses relative to the start of the user's program. This list is scanned, and the appropriate locations in the user's program area are updated with a constant. The job is then ready to be started.

The relocation information is in the following format:



Bits 0-14 represent the relative address to relocate divided by two. This implies that relocation is always done on a word boundary, which is the case. Bit 15 is used to indicate the type of relocation to perform, positive or negative. The relocation constant (which is the load address of the program) is added to or subtracted from the indicated location depending on the sense of bit 15; 0 implies addition, 1 implies subtraction. 17776 terminates the list of relocation information.

Following is an example of a simple, non-overlaid program linked to produce a REL file. A dump of the file follows the program.

```

.TITLE FTFST
.MCALL ..V2...REGDEF,.LOOKUP,.READW,.QSET,.PRINT,.EXIT
..V2..
.REGDEF
ST: .QSET #QLIST,#7
.LOOKUP #AREA,#0,#PTR
#CC 1$
.PRINT #LKFAIL
.EXIT
1$: .READW #AREA,#0,#RUFF,#254.,#R
#CC 2$
.PRINT #RDFAIL
.EXIT
2$: .PRINT #OK
.EXIT

QLIST: .BLKW 7*7
AREA: .BLKW 20.
PTR: .RAD50 /PR FILE12/
RUFF: .NLIST
.REPT 254.
.WORD 0
.ENDR
.LIST
LKFAIL: .ASCIZ /LOOKUP FAILED/
RDFAIL: .ASCIZ /READW FAILED/
OK: .ASCIZ /READW OK/
.EVEN
.NLIST
.REPT <ST+1776-.>/2
.WORD 0
.ENDR
.NLIST
.END ST

```



```

14 000142
000142 104350
17
18 000144
19 000306
20 000356 063320 023344 022070
24 001364 114 117 125 120
001367 113 125 106 101
001372 040 106 114 105
001375 111 114 100
001400 104 000
27 001402 122 105 101
001405 104 127 040
001410 106 101 111
001413 114 105 104
001416 000
28 001417 122 105 101 OK: .ASCYZ /READW OK/

```

.EXIT

FMT

0150

```

.OLTST: .BLKW 7*7
.ARFA: .BIKW 20.
.PTR: .RAD50 /PR FILE12/
.LKFATL: .ASCYZ /LOOKUP FATLFD/

```

```

.RPFATL: .ASCYZ /READW FAILED/

```

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```

001422 104 127 040
001425 117 113 000

```

29

.EVEN

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SYMBOL TABLE

```

ARFA 000306 RUFF 0003640 IKFATL 0013640 CK 001417P PC =X0000007
PTR 000356 .OLTST 0001440 RPFATL 001402P R0 =X0000000 P1 =X0000001
R2 =X000002 R3 =X000003 P4 =X000004 R5 =X000005 SP =X000006
ST 000000 ...V2 = 000001
. ABS. 000000 000
001776 001
ERRORS DETECTED: 0
FRFE CORR: 15895. WORDS

```

,LP:NIITM/LIMFB=FTFST

```

BLOCK NUMBER 0000
000/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
020/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
040/ 000000 000000 000000 000000 000000 001776 001776 000000 000000 *.....*
060/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
100/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
120/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
140/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
160/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
200/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
220/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
240/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
260/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
300/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
320/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
340/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
360/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
400/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
420/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
440/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
460/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
500/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
520/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
540/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
560/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
600/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
620/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
640/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
660/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
700/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
720/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
740/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
760/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*

```

In block 0, word 50 shows the highest, non-relocated, memory address in the user program. Word 52 shows the program size in bytes. Word 54 shows the size of the overlay region. The value is non-zero only for programs with overlays. Word 60 contains a 3, indicating that the relocation information begins at block 3 of the file.

```

BLOCK NUMBER 0001
000/ 112700 000007 012744 000144 104353 012700 000704 112740 * . . . F . N . K . P . P . *
020/ 000001 000001 105010 012760 000356 000002 005060 000004 * . . . . . P . N . . . . *
040/ 104375 105004 012700 001344 104351 104340 012700 000306 *) . . . . . T . Y . H . P . *
060/ 112760 000010 000001 105010 012760 000000 000000 012740 * P . . . . . P . . . . . *
100/ 000364 000004 012760 000400 000006 012740 000000 000010 * T . . . . . P . . . . . *
120/ 104375 105004 012700 001402 104351 104350 012700 001417 *) . . . . . T . Y . H . P . *
140/ 104351 104350 000000 000000 000000 000000 000000 000000 * T . Y . . . . . P . . . . . *
160/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
200/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
220/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
240/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
260/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
300/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
320/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
340/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
360/ 003364 022070 000000 000000 000000 000000 000000 067320 * T . Y . A . S . . . . . P . . . . . *
400/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
420/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
440/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
460/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
500/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
520/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
540/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
560/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
600/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
620/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
640/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
660/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
700/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
720/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
740/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *
760/ 000000 000000 000000 000000 000000 000000 000000 000000 * . . . . . P . . . . . P . . . . . *

```

This block corresponds to locations 0-776 in the assembly listing.

```

BLOCK NUMBER 0000
000/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
020/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
040/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
060/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
100/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
120/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
140/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
160/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
200/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
220/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
240/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
260/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
300/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
320/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
340/ 000000 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
360/ 000000 000000 047514 045517 050125 047040 044501 042514 *.....*
400/ 000104 042522 042101 020127 040506 046111 042105 051000 *D,READY FAILED,R
420/ 040505 053504 047440 000113 000000 000000 000000 000000 *EADW OK.....*
440/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
460/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
500/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
520/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
540/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
560/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
600/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
620/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
640/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
660/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
700/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
720/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
740/ 000000 000000 000000 000000 000000 000000 000000 000000 *.....*
760/ 000000 000000 000000 000000 000000 000000 000000 041056 *.....*

```

This block corresponds to locations 1000-1776 in the assembly listing.

```

BLOCK NUMBER 0003
000/ 000003 000006 000014 000023 000027 000040 000053 000067 *
020/ 177776 000000 000000 000000 000000 000000 000000 000000 *
040/ 000000 000000 000000 000000 000000 000000 000000 000000 *
060/ 000000 000000 000000 000000 000000 000000 000000 000000 *
100/ 000000 000000 000000 000000 000000 000000 000000 000000 *
120/ 000000 000000 000000 000000 000000 000000 000000 000000 *
140/ 000000 000000 000000 000000 000000 000000 000000 000000 *
160/ 000000 000000 000000 000000 000000 000000 000000 000000 *
200/ 000000 000000 000000 000000 000000 000000 000000 000000 *
220/ 000000 000000 000000 000000 000000 000000 000000 000000 *
240/ 000000 000000 000000 000000 000000 000000 000000 000000 *
260/ 000000 000000 000000 000000 000000 000000 000000 000000 *
300/ 000000 000000 000000 000000 000000 000000 000000 000000 *
320/ 000000 000000 000000 000000 000000 000000 000000 000000 *
340/ 000000 000000 000000 000000 000000 000000 000000 000000 *
360/ 000000 000000 000000 000000 000000 000000 000000 000000 *
400/ 000000 000000 000000 000000 000000 000000 000000 000000 *
420/ 000000 000000 000000 000000 000000 000000 000000 000000 *
440/ 000000 000000 000000 000000 000000 000000 000000 000000 *
460/ 000000 000000 000000 000000 000000 000000 000000 000000 *
500/ 000000 000000 000000 000000 000000 000000 000000 000000 *
520/ 000000 000000 000000 000000 000000 000000 000000 000000 *
540/ 000000 000000 000000 000000 000000 000000 000000 000000 *
560/ 000000 000000 000000 000000 000000 000000 000000 000000 *
600/ 000000 000000 000000 000000 000000 000000 000000 000000 *
620/ 000000 000000 000000 000000 000000 000000 000000 000000 *
640/ 000000 000000 000000 000000 000000 000000 000000 000000 *
660/ 000000 000000 000000 000000 000000 000000 000000 000000 *
700/ 000000 000000 000000 000000 000000 000000 000000 000000 *
720/ 000000 000000 000000 000000 000000 000000 000000 000000 *
740/ 000000 000000 000000 000000 000000 000000 000000 000000 *
760/ 000000 000000 000000 000000 000000 000000 000000 000000 *

```

This block shows the root relocation information. The first word of block 3 is a 3; since this is positive, positive relocation is indicated. Locations 6, 14, 30, 46, 56, 100, 126, and 136 must all be positively relocated at FRUN time. (On examination of the assembly listing, those locations marked with a ' need to be relocated.) The 177776 terminates the list.

Had negative relocation been indicated at relative location 6, block 3 would have shown 100003, 6, 14, 23, 27, 40, 53, 57, 177776.

3.4.4.2 RFL Files with Overlays - When overlays are included in a program, the file is similar to that of a non-overlaid program. However, the overlay segments must also be relocated. Since overlays are not permanently memory resident but are read in from the file as needed, they require an additional operation. Each overlay segment is relocated (by FRUN) and then rewritten into the file. Then, when the overlay is called in, it will be properly relocated. This process takes

place each time an overlaid file is run with FRUN. The relocation information for overlay files contains both the list of addresses to be modified and the original contents of each location. This allows the file to be FRUN after the first usage.

NOTE

.ASECTs are illegal above 1000₈ and restricted in an overlaid foreground job. Refer to Chapter 6 of the RT-11 System Reference Manual.

A REL file with overlays appears as shown in Figure 3-14:

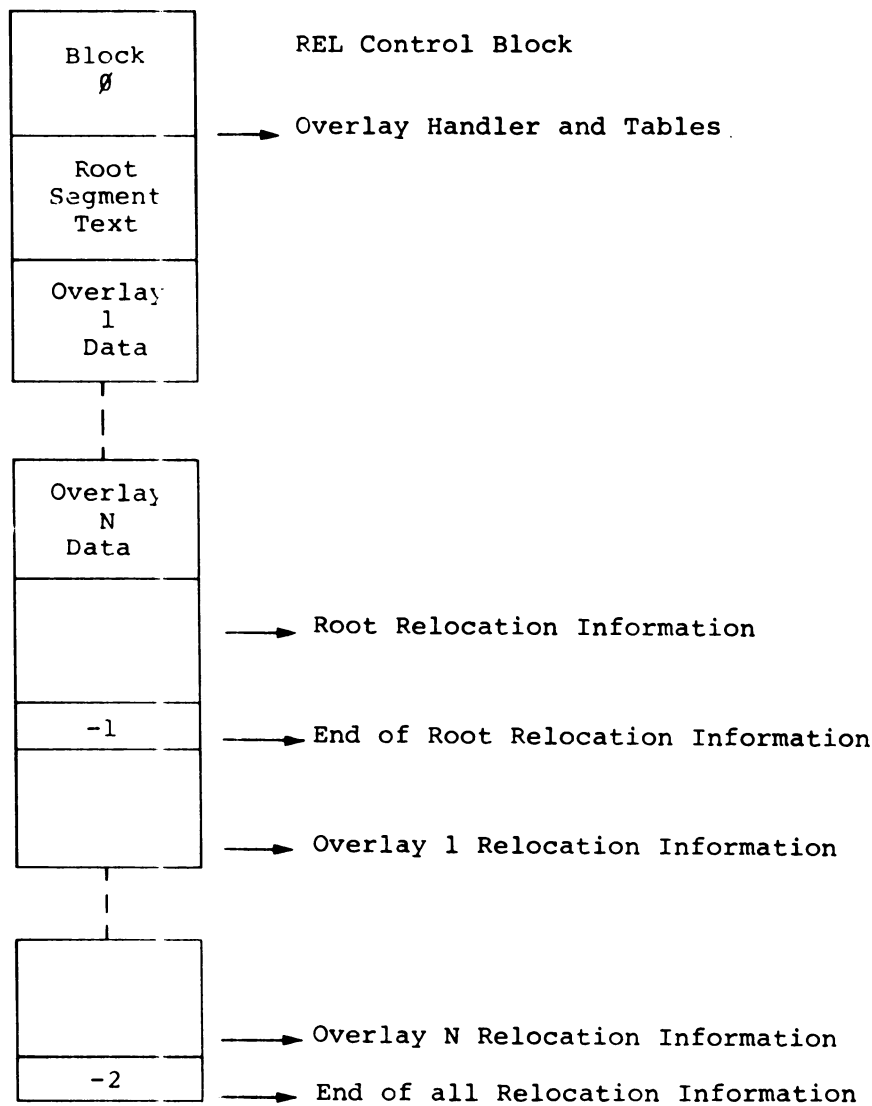


Figure 3-14
REL File With Overlays

In this case, location 54 of block 0 of the REL file contains the size of the overlay region, in words. This is used to allocate space for the job when added to the size of the program base segment in location 52.

After the program base (root) code has been relocated, each existing overlay is read into the program overlay region in memory, relocated via the overlay relocation information, and then written back into the file.

The root relocation information section is terminated with a -1. This -1 is also an indication that an overlay segment relocation block follows. The overlay segment relocation block is shown in Figure 3-15:

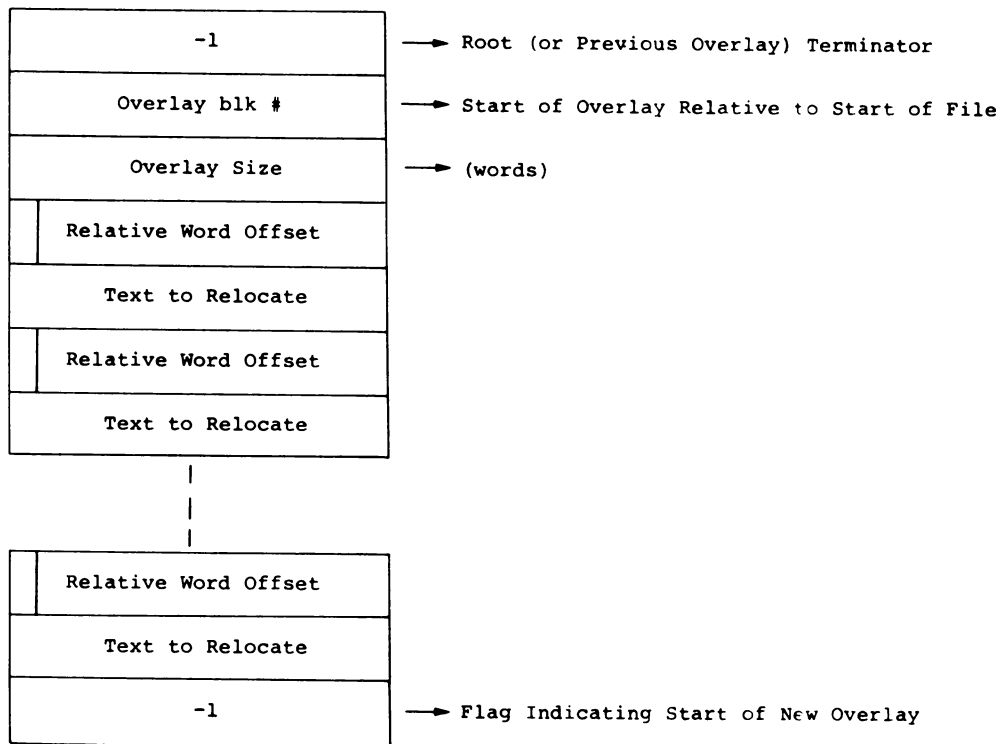


Figure 3-15
Overlay Segment Relocation Block

The displacement is relative to the start of the program and is interpreted as in the nonoverlaid file (i.e., bit 15 indicates the type of relocation, and the displacement is the true displacement divided by two). Encountering -1 indicates that a new overlay region begins here. A -2 indicates the termination of all relocation information.

CHAPTER 4
SYSTEM DEVICE

4.1 DETAILED STRUCTURE OF THE SYSTEM DEVICE

The RT-11 system device holds all the components of the system and is used by RT-11 to store device handlers and the monitor file. The layout of the system device is:

<u>Block #</u>	<u>Contents</u>
0	Bootstrap
1	Reserved for volume identification information
2	Bootstrap
3 to 5	Reserved for monitor or bootstrap expansion
6 to (N*2)+5	Directory segments; N is the number of directory segments
(N*2)+6 to end	File storage

All other system components, i.e., the monitor and device handlers, are files on the system device:

<u>File</u>	<u>Contains</u>
MONITR.SYS	The current RT-11 monitor; contains bootstrap, KMON, USR/CSI, RMON, KMON overlays, scratch blocks
SYSMAC.SML	System Macro Library
SYSMAC.S8K	8K System Macro Library
LP.SYS	Line printer handler

<u>File</u>	<u>Contains</u>
DT.SYS	DECTape handler
TT.SYS	Console handler (S/J only)
RK.SYS	RK disk handler
DS.SYS	RJS03/4 fixed-head disk handler
DX.SYS	RX01 flexible disk handler
DP.SYS	RP disk handler
PR.SYS	High-speed reader handler
PP.SYS	High-speed punch handler
CR.SYS	Card reader handler
RF.SYS	RF disk handler
CT.SYS	Cassette handler
MT.SYS	TM11 magtape handler
MM.SYS	TJU16 magtape handler
BA.SYS	BATCH run-time handler

In general, files with the .SYS extension are parts of the monitor system. The bootstrap records the block numbers of the relevant areas in the monitor tables at bootstrap time. Thus, RT-11 is extremely flexible with respect to the interchange and construction of systems.

4.2 CONTENTS OF MONITR.SYS

Following is the block layout of the RT-11 monitor file, MONITR.SYS. Block numbers are relative to the start of the file.

<u>F/B</u> <u>Monitor</u>	<u>Block # (decimal)</u>	<u>Contents</u>
	0-1	Copy of system bootstrap (blocks 0 and 2 of the system device)
	2-17	Swap blocks
	18-24	KMON (includes 2-block KMON overlay area)
	25-32	USR/CSI
	33-47	RMON
	48-57	KMON overlays

<u>S/J</u> <u>Monitor</u>	<u>Block # (decimal)</u>	<u>Contents</u>
	0-1	Copy of system bootstrap
	2-16	Swap blocks
	17-22	KMON (includes 1-block KMON overlay area)
	23-30	USR/CSI
	31-37	RMON
	38-44	KMON overlays

4.3 KMON OVERLAYS

The KMON overlays are one block in size in the S/J Monitor and two blocks in size in the F/B Monitor. The contents of each overlay are described in this list:

<u>Overlay #</u>	<u>S/J</u>	<u>F/B</u>
0	DATE, TIME	DATE, TIME, SAVE, ASSIGN
1	SAVE, ASSIGN	LOAD, UNLOAD, SUSPEND, RESUME, CLOSE, FRUN (Part 1)
2	LOAD, UNLOAD, CLOSE	FRUN (Part 2)
3	GT ON/OFF	GT ON/OFF, SET
4	SET	

4.4 DETAILED OPERATION OF THE BOOTSTRAP

Bootstrapping a system causes a fresh copy of that system to be installed in memory. In the RT-11 boot, certain system device resident tables are also updated. Following is a detailed description of the bootstrap.

<u>Action</u>	<u>Explanation</u>
1. User executes hardware bootstrap	This causes block 0 of the system device to be read into 0-777. Control then passes to location 0.
2. Second part of bootstrap is read	The first part of the boot reads the second half into 1000-1777.
3. Determine how much memory is available	Boot sets a trap at location 4 and then starts addressing memory. When the trap is taken, illegal memory has been addressed.
4. Look for special devices	Boot sets a trap at location 10 and then tries to address the clock, FPU, and VT11 display processor. Their presence or absence is indicated in the CONFIG word in RMON.
5. Check memory size	If memory is too small to read in the monitor, a message is printed and the boot halts.
6. Read in directory and find MONITR.SYS	The entire directory is searched. If MONITR.SYS is not found, a HALT occurs after the boot prints an error message.
7. Read the monitor into memory	The monitor file, MONITR.SYS, is read into the highest 4K of memory.
8. Put pointers to monitor file blocks into RMON	RMON references the monitor swap blocks directly. Thus, the position of the swap blocks varies as the placement of MONITR.SYS varies. The real position of the blocks is updated for each boot operation.
9. Update position-dependent areas in RMON	MONITR.SYS is initially linked at 8K. However, if more than 8K is available, RT-11 uses it. To do that, certain words must be updated to point to the actual areas of high memory where they will be. Boot contains a list of all words to be updated, located at RELLST in BSTRAP.MAC.
10. LOOKUP the device handlers in system and store their record numbers in \$DVREC	Boot looks at \$PNAME table to find the names of the devices in the system. The extension .SYS is appended. Thus, the PR handler is a file called PR.SYS. The location of the handler is then placed in \$DVREC. If the LOOKUP fails, the device gets a 0 in its \$DVREC entry. That implies that the device handler does not exist.

<u>Action</u>	<u>Explanation</u>
11. Print bootstrap header	Boot prints monitor identification message "RT-11" followed by monitor type ("FB" or "SJ") followed by version number.
12. Set up locations 0 and 2	Boot puts a "BIC R0,R0" in location zero and an .EXIT EMT in location 2.
13. Turn on KW11-L Clock	The bootstrap turns on the clock, if present in the configuration.
14. Exit to Keyboard Monitor	

4.5 FIXING THE SIZE OF A SYSTEM

RT-11 is designed to automatically operate from the top of the highest available 4K memory bank. However, it is possible to force the system to operate from a specified area that is not necessarily the highest. For instance, the following series of commands causes RT-11 to run in a 16K environment, even though the configuration actually has 28K of memory:

```
.R PATCH<CR>                                [Run RT-11 PATCH program.]

PATCH Version number

FILE NAME--
*MONITR.SYS/M<CR>
*BHALT/ 407 <CR>
*E
.R PIP
*A=MONITR.SYS/U<CR>
*SY:/O
```

[Specifying MONITR.SYS/M indicates it is a monitor file. Change location "BHALT" from a 407 to a 0 (HALT). The correct address of BHALT can be found in Chapter 5 of Getting Started With RT-11 (V02B). E causes an exit to the monitor. Now run PIP to update the bootstrap and reboot the system.]

When the bootstrap is performed, the computer halts. The halt allows the user to enter the desired size in the switch register. With this patch installed, the V2 bootstrap uses the top five bits (bits 11-15) of the switch register to determine memory size. If the switch register contains the number 160000 or greater (e.g., if the register is unchanged after booting the system), a normal memory determination is performed. Otherwise, the top five bits are taken to be a number representing the number of 1K word blocks of memory. Each bit has the following value:

<u>Switch Register</u>	<u>Memory Size</u>
4000	1K
10000	2K
20000	4K
40000	8K
100000	16K

A combination of the bits will produce the range of system sizes from 8K through 28K, in 1K increments.

Examples:

1. To boot a system into 24K on a 28K configuration, use the combination:

$$140000 = 100000 (16K) + 40000 (8K)$$

2. To boot the S/J Monitor into 11K, use the combination:

$$54000 = 40000 (8K) + 10000 (2K) + 4000 (1K)$$

When the switch register is set properly, press the CONTInue switch and the bootstrap will be executed.

If the CONTInue switch is pressed immediately following the halt without changing the switch settings, a normal memory determination is done. To change the bootstrap back to its original (non-halting) form, execute the same commands as above, but change the 0 at BHALT back to a 407.

This procedure allows the user to 'protect' memory areas, since RT-11 never accesses memory outside the bounds within which it runs.

Another useful procedure, when desiring to always boot a system into a specific memory size or when the console switch register is not available, is to determine the bit combination corresponding to the choice of memory size, as explained above. Then enter the following commands, where xxxxx represents the bit pattern just determined:

```

.R PATCH<CR>                                [Run RT-11 PATCH program.]

PATCH Version number

FILE NAME--
*MONITR.SYS/M<CR>
*BHALT/ 407 240<LF>                          [NOP the branch at BHALT
*BHALT+2/ 13702 12702<LF>                    Change MOV @#SR,R2 to
*BHALT+4/ 177570 xxxxxx<CR>                 MOV #VAL,R2. Address of
*E                                             switch register is replaced
:                                             with one of the bit combina-
:                                             tions described previously.]

```


For the patch addresses for other system devices, and for the address of BHALT, consult the Getting Started with RT-11 (V02B) manual.

CHAPTER 5

I/O SYSTEM, QUEUES, AND HANDLERS

I/O transfers in RT-11 are handled by the monitor through routines known as device handlers. Device handlers are resident on the system mass storage device and can be called into memory at a location specified by the user (via a .FETCH handler request or KMON LOAD command). Only the device handlers distributed with the system in use (V2 or V02B) may be used; the system will malfunction otherwise.

This chapter describes how to write a new device handler and add it to the system. A summary of differences between Version 1 and Version 2 Device Handler requirements is included for the user who wishes to update old device handlers. Instructions and examples for making a device the system device and for writing a new bootstrap for the device are also included.

5.1 QUEUED I/O IN RT-11

Once a device handler is in memory, any .READ/.WRITE requests for the corresponding devices are interpreted by the monitor and translated into a call to the I/O device handler. To facilitate overlapped I/O and computation, all I/O requests to RT-11 are done through an I/O queue. This section details the structure of the I/O queueing system.

5.1.1 I/O Queue Elements

The RT-11 I/O queue is made up of a linked list of queue elements. A single element has the structure shown in Figure 5-1:

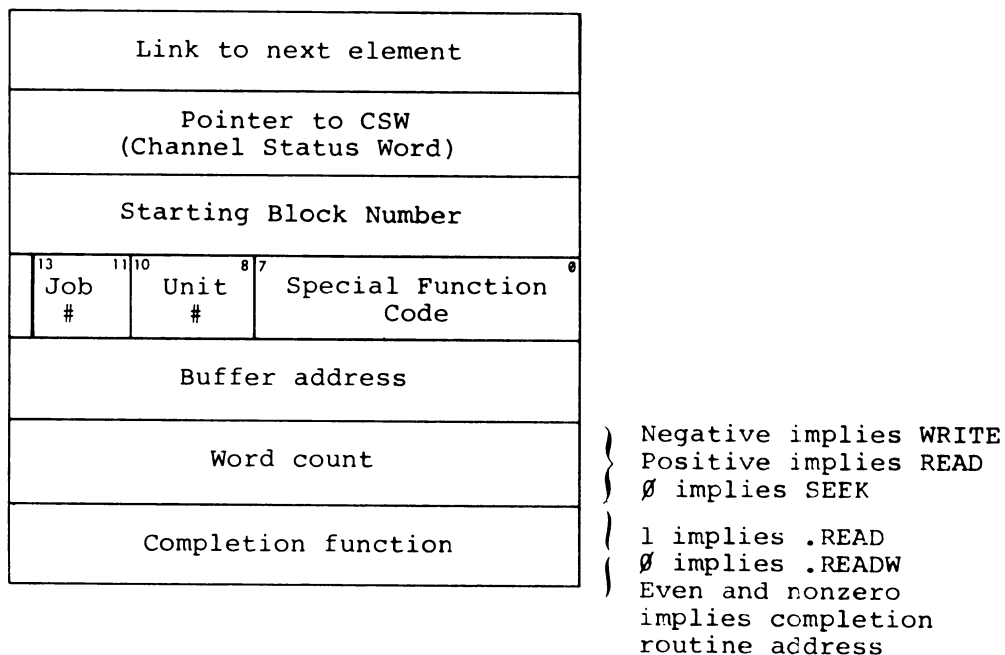
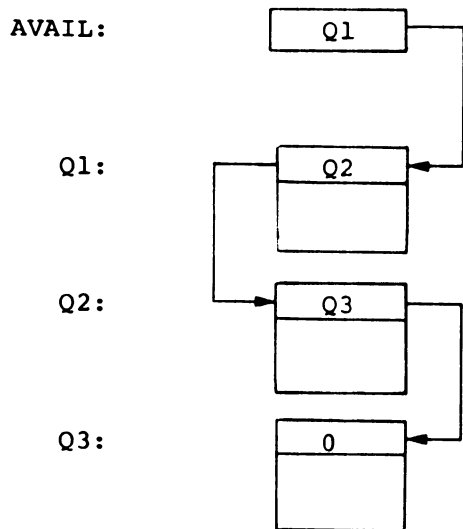


Figure 5-1
I/O Queue Element

RT-11 maintains one queue element in the Resident Monitor. (In F/B, one element per job is maintained in the job's impure area.) This is sufficient for any program that uses wait-mode I/O (.FEADW/.WRITW). However, for maximum throughput, the .QSET programmed request should be used to create additional queue elements.

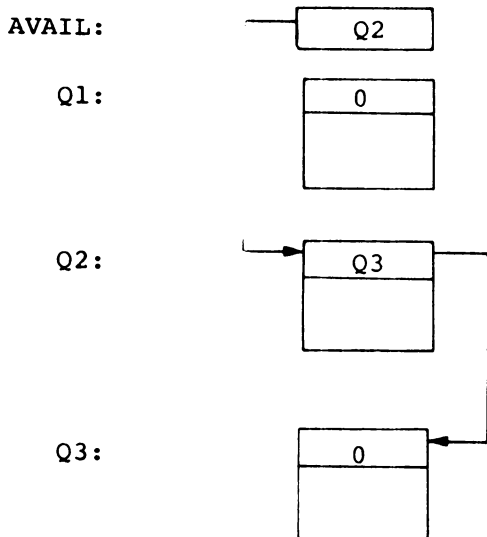
If an I/O operation is requested and a queue element is not available, RT-11 must wait until an element is free to queue the request. This obviously slows up program execution. If asynchronous I/O is desired, extra queue elements should be allocated. It is always sufficient to allocate N new queue elements, where N is the total number of pending requests that can be outstanding at one time in a particular program. This produces a total of N+1 available elements, since the Resident Monitor element is added to the list of available elements.

Diagrammatically, the I/O queue appears as follows:



AVAIL is the list header. It always contains a pointer to an available element. If AVAIL is 0, no elements are currently available.

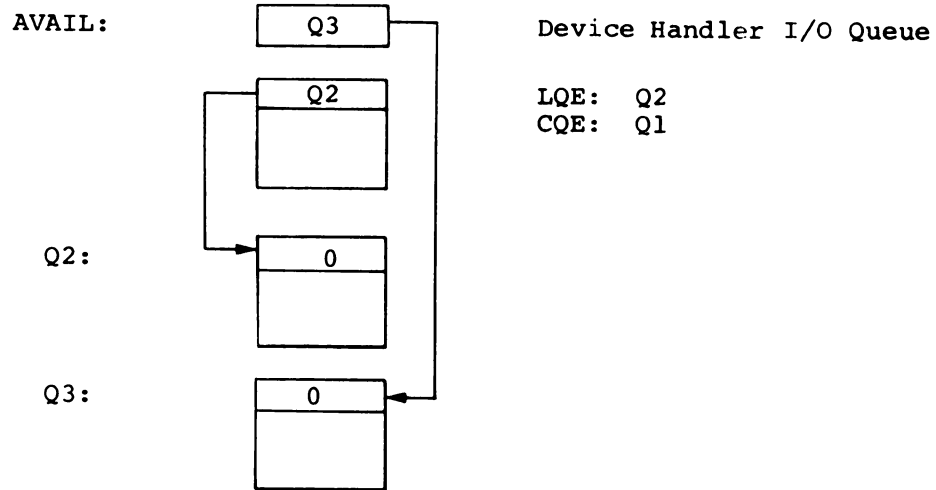
When an I/O request is initiated, an element is allocated (removed from the list of available elements) and is linked into the appropriate device handler's I/O queue. The handler's queue header consists of two pointers: the current queue element (CQE) pointer, pointing to the element at the top of the list, and the last queue element (LQE) pointer, pointing to the last element entered in the queue. The LQE pointer is used for fast insertion of new elements into the queue.



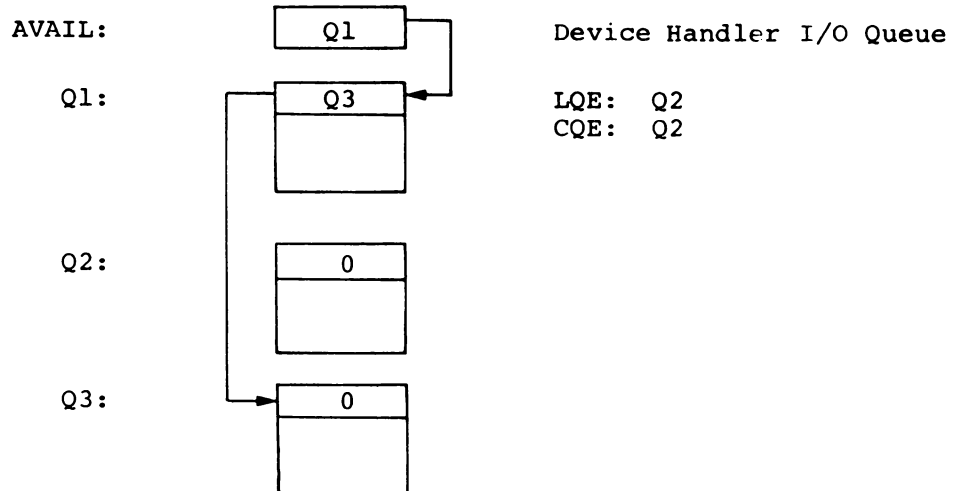
Device Handler I/O Queue

LQE: Q1 (Pointer to last queue element)
 CQE: Q1 (Pointer to current queue element)

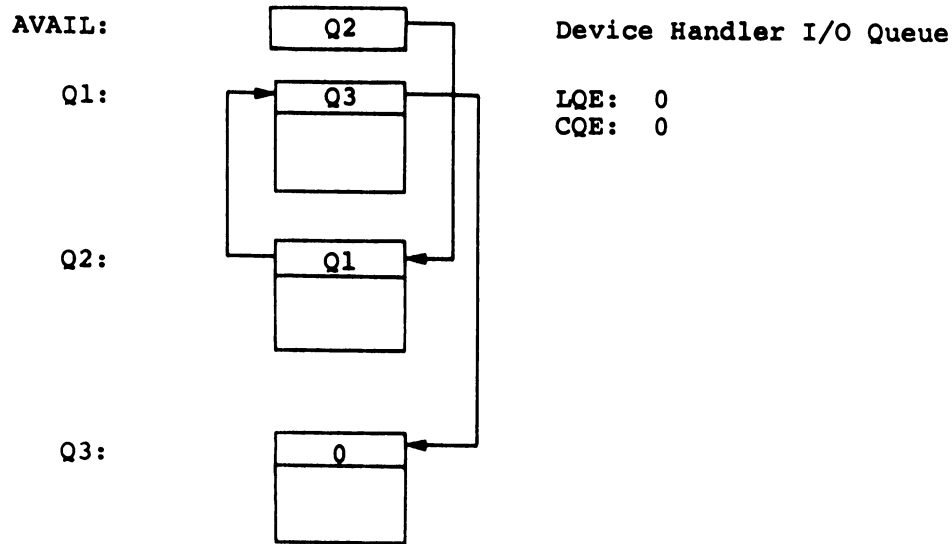
In this case, the device is associated with element Q1. If another request comes in for that same device before the first completes, a waiting queue is built up for that device.



When the I/O transfer in progress completes, Q1 is returned to the list of available elements, and the transfer indicated by Q2 will be initiated:



When Q2 is completed, it too is returned to the list of available elements.



Note that the order of the queue element linkages may be altered.

A distinction between S/J and F/B operation is that F/B maintains two separate queue structures, one for each active job. The queue headers (AVAIL) are words in the user's impure area. The centralized queue manager dispatches transfers in accordance with job priority. Thus, if two requests are queued waiting for a particular device, the foreground request is honored first. At no time, however, will an I/O request already in progress be aborted in favor of a higher priority request; the operation in progress will complete before the next transfer is initiated.

Another difference between S/J and F/B operation is that the F/B scheduler will suspend a job pending the availability of a free queue element and will try to run another job.

5.1.2 Completion Queue Elements

The F/B Monitor maintains, in addition to the queue of I/O transfer requests, a queue of I/O completion requests. When an I/O transfer completes and a completion routine has been specified in the request (i.e., the seventh word of the I/O queue element is even and non-zero), the queue completion logic in the F/B Monitor transfers the request node (element) to the completion queue, placing the channel

status word and channel offset in the node. This has the effect of serializing completion routines, rather than nesting them. Completion routines are called by the completion queue manager on a *first-in/first-out* basis, and the completion routines are entered at priority level 0 rather than at interrupt level.

The .SYNCH request also makes use of the completion queue. When the .SYNCH request is entered, the seven-word area supplied with the request is linked into the head of the completion queue, where it appears to be a request for a completion routine. The .SYNCH request then does an interrupt exit. The code following the .SYNCH request is next called at priority level 0 by the completion queue manager. To prevent the .SYNCH block from being linked into AVAIL (the queue of available elements), the word count is set to -1. The completion queue manager checks the word count before linking a queue element back into the list of available elements, and skips elements with the -1 word count.

Figures 5-2 and 5-3 show the format of the completion queue and .SYNCH elements.

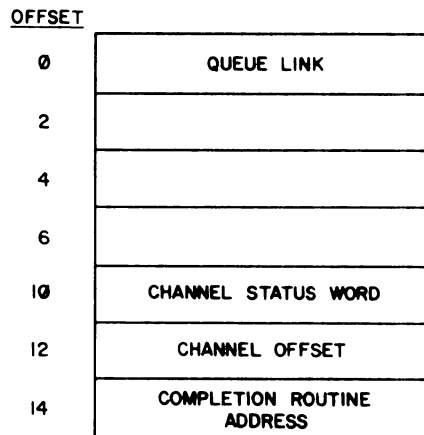


Figure 5-2
Completion Queue Element

OFFSET	
0	QUEUE LINK
2	JOB NUMBER
4	
6	
10	SYNCH I.D.
12	-1
14	SYNCH RETURN ADDRESS

Figure 5-3
.SYNCH Element

5.1.3 Timer Queue Elements

Another queue maintained by the F/B Monitor is the timer queue. This queue is used to implement the .MRKT request, which schedules a completion routine to be entered after a specified period of time. The first two words of the element are the high- and low-order time and the seventh word is the completion routine address. An optional sequence number can be added to the request to distinguish this timer request from others issued by the same job.

The F/B Monitor uses the timer queue internally to implement the .TWAIT request. The .TWAIT request causes the issuing job to be suspended and a timer request is placed in the queue with the .RSUM logic as the completion routine. Refer to Figure 5-4 for the format of the timer queue element.

OFFSET		
0	HIGH-ORDER TIME	C.HOT
2	LOW-ORDER TIME	C.LOT
4	LINK TO NEXT ELEMENT	C.LINK
6	JOB # OF OWNER	C.JNUM
10	OWNER'S SEQUENCE #	C.SEQ
12	0	
14	COMPLETION ADDRESS	C.COMP

Figure 5-4
Timer Queue Element

5.2 DEVICE HANDLERS

This section contains the information necessary to write an RT-11 device handler. It is illustrated with an example, a driver for the RS64 fixed-head disk (with RC11 controller). A source listing is included in Appendix A, Section A.1; portions of this listing are referenced throughout the remainder of this section and in future sections.

The user should refer to the PDP-11 Peripherals Handbook for details regarding the operation of any particular peripheral.

NOTE

All RT-11 handlers must be written in position independent code (PIC). Consult the PDP-11 Processor Handbook for information on writing PIC.

5.2.1 Device Handler Format

The first five words of any device handler are header words. The format is:

<u>Word #</u>	<u>Contents</u>
1	Address of first word of device's interrupt vector.
2	Offset from current PC to interrupt handler.
3	Processor status word to be used when interrupt occurs. Must be 340 (priority 7).
4,5	Zero. These are the queue pointers.

See area C in the example handler (Section A.1).

A word must be provided at the end of the handler. When the handler is .FETCHed, the monitor places a pointer to the monitor common interrupt entry code in the last word of the handler. This requires that the handler size in the monitor's \$HSIZE table be exact or the handler will malfunction. See area M in the example in Section A.1.

The word preceding the interrupt handler entry point must be an unconditional branch to the handler's abort code. The abort code is used by the F/B Monitor to stop I/O for the device. The abort entry point is shown at area G in the example and the abort code is at area K.

5.2.2 Entry Conditions

The device handler is entered directly from the monitor I/O queue manager, at which time it initiates the data transfer. The fifth word of the header contains a pointer into the queue element to be processed. This word (called CQE, for Current Queue Element) points to the third word of the queue element, which is the block number to be read or written. Referring to the example, location RCCQE contains the address of the third word of the queue element to be processed. It is generally advisable to put the pointer into a register, as that greatly facilitates picking up arguments to initiate the transfer. In the example, the entry point is at the location marked by E. Notice that registers need not be saved.

5.2.3 Data Transfer

Most handlers use the interrupt mechanism when transferring data. The handler initiates the transfer and then returns immediately to the monitor with an RTS PC, shown at area F. When the transfer is completed, the device interrupts. When the interrupt routine determines that I/O is complete or that an error has occurred, it jumps to the monitor completion routine in the manner shown at area J in the listing.

If the interrupt mechanism is not used, the data transfer must be completed before returning to the monitor. The handler must loop on a device flag with the interrupt disabled. When I/O is complete, the driver returns to the monitor with a jump to the monitor completion code, similar to that shown at area J in the example.

5.2.4 Interrupt Handler

Once the transfer has been initiated and control has passed back to the monitor, data interrupts will occur.

Information in the header of the handler causes the interrupt to be vectored to the interrupt handling code within the handler. The code at the interrupt location should keep the transfer going, determine when the transfer is complete, and detect errors.

When the transfer is done, control must be passed to the monitor's I/O queue manager, which performs a cleanup operation on the I/O queue.

Restrictions that apply to the interrupt code are:

1. The common interrupt entry into the monitor must be taken. Interrupt routines linked into a program use the `.INTEN` request described in Chapter 9 of the RT-11 System Reference Manual. Handlers made part of the system have a more efficient method of entry. The last word of the handler is set to point to the monitor common interrupt entry code when the handler is fetched. Upon reception of an interrupt, the handler must execute this code by performing a `JSR R5, @$INPTR`, where `$INPTR` is the tag commonly used by RT-11 handlers for the pointer word. See areas I and N in the example. The `JSR` instruction must be followed by the complement of the priority at which the handler will operate. See area I for an easy method to make the assembler compute the complement. On return from the monitor's interrupt entry code, R4 and R5 have been saved and may be used by the handler. Other registers must be saved and restored if they are to be used.
2. A check must be made to determine if the transfer is complete. However, with nonfile-oriented devices, such as paper tape, line printer, etc., an interrupt occurs whenever a character has been processed. For these devices, the byte count, which is in the queue element, is used as a character count.

Nonfile-structured input devices should be able to detect an end of file condition, and pass that on to the monitor.

NOTE

The queue element contains a word count, not a byte count. The initial entry to the handler should change the word count to a byte count if the device interrupts at each character. The transfer is complete when the byte count decrements to 0.

Before the conversion to bytes is made, the sign of the word count must be determined since it specifies whether this transfer is a Read or Write. A negative word count implies a Write and should be complemented before converting to bytes.

3. Check for occurrence of an error. If a hardware error occurred, the hard error bit in the channel status word (CSW) should be set, the transfer should be aborted, and the monitor completion code executed. The address of the channel status word is in word 2 of the queue element. The error bit is bit 0 of the CSW. Generally, it is advisable to retry a certain number of times if an error occurs. RT-11 currently retries up to eight times before deciding an error has occurred. (Note that this is true for file-structured devices only.) It is

desirable, in case an error occurs, to do a drive or control reset, where appropriate, to clear the error condition before a retry is initiated. See the area between I and H in the example.

4. If the transfer is not complete and no error has occurred, registers used should be restored, and an RTS PC executed.

To pass an EOF (End of File) to the monitor, the 2000 bit in the CSW should be set. Refer to the sample handler in Appendix A for an example of setting the EOF bit. When EOF is detected on non-file structured devices, the remainder of the input buffer must be zeroed.

5. When the transfer is complete, whether an error occurred or not, the monitor I/O completion code must be entered to terminate activity and/or enter a completion routine. When return is made to the monitor, R4 must point to the fifth word of the handler (RCCQE in the example). See area J in the example for the method of returning to the monitor completion routine.

Handlers should check for special error conditions that can be detected on the initial entry to the handler. For example, trying to write on a read-only device should produce a hard error. It must be emphasized that the user handlers should interface to the system in substantially the same way as the handler in Section A.1. This handler is included as a guide and an example.

5.3 ADDING A HANDLER TO THE SYSTEM

When the handler has been written and debugged, it may be installed in the system by following the procedures in this section. The process consists of inserting information about the handler into the monitor tables listed below.

<u>Table to be Changed</u>	<u>Contents</u>
\$HSIZE	Size of handler (in bytes).
\$DVSIZ	Size of device in 256-word blocks. If nonfile device, entry = 0.
\$PNAME	Permanent name of the device (should be two alphanumeric characters entered in .RAD50 notation, left-justified).
\$STAT	Device status table. Refer to Section 2.5.2.2 for the format of \$STAT table.
LOWMAP	Low memory protection map; refer to Section 2.5.4.

There is no restriction on handler names; any 2-letter combination not currently in use may be chosen for the new handler and the name may be inserted in any unused slot in the \$PNAME table, or in a slot occupied by a nonexistent device (i.e., a device not installed on the user's system). Note that the name must be entered in .RAD50. Since PATCH does not have a .RAD50 interpretation switch, the name must be entered to PATCH in its numerical form. Appendix C of the RT-11 System Reference Manual contains a .RAD50 conversion table; ODT can also be used to perform .RAD50 conversions.

As an example, assume again the handler for the RC11/RS64 disk (the sample handler in Section A.1) is to be inserted in the system. First, the values of the table entries for this device are determined (the addresses used in the example are for illustrative purposes only; consult Getting Started With RT-11 (V02B) for the correct table addresses for the version in use):

\$HSIZE:	316	After assembly, the handler was found to take up 316 bytes. See area 0 in the example listing.
\$DVSIZ:	2000	The disk has 1024 (decimal) 256-word blocks for storage.
\$PNAME:	.RAD50 /RC/ or 70370	The name assigned is RC. The .RAD50 value of RC is 70370.
\$STAT:	100023	The device is file-structured, is a read/write device, and uses the standard RT-11 file structure. The identifier (selected by the user) is 23. Refer to Section 2.5.2.2 for the format of the \$STAT table.
LOWMAP:	14	Protect RC vector 210,212 at byte 336 of LOWMAP (refer to Section 2.5.4.).

Once these values have been decided, the steps for inserting the device handler are:

1. Assemble the handler, using either MACRO or ASEMBL.
2. Link the handler at 1000. The name of the handler should be whatever the \$PNAME entry is, with the .SYS extension appended:

```
.R LINK
*RC.SYS=RC   where RC.OBJ is the handler object
UNDEF GLBLS module. The default link address is
              1000.
```

NOTE

If the handler being linked is one that could also be a system device handler, the user can expect one undefined global, \$INTEN.

3. Run PATCH to modify the tables and protect the interrupt vectors.

For this example, assume that the table addresses are found to be:

<u>Table</u>	<u>S/J Address</u>	<u>F/B Address</u>
\$HSIZE	13674	14602
\$DVSIZ	13730	14636
\$PNAME	16516	17640
\$STAT	16552	17674

NOTE

The addresses above are for illustration only. Consult the Getting Started With RT-11 (V02B) manual for current table addresses and for the address of the monitor base location, BASE. ~ 17000

The tables have room for fourteen (decimal) device entries; all are already assigned by the monitor. Assuming that a given configuration never has all supported devices, however, at least one slot should be available to be overlaid. For example, assume the twelfth slot is occupied by a device not installed on the system, and therefore available for change. The octal offset is 26, which, added to the table addresses above, gives the address of the empty slot:

S/J Monitor:

.R PATCH<CR>

PATCH Version number

FILE NAME--

*MONITR.SYS/M<CR>

*BASE;0R<CR>

*0,13674/ 4000 316<CR>

*0,13730/ 0 2000<CR>

*0,16516/ 6250 7037<CR>

*0,16552/ 4 100023<CR>

*0,16336\ 77 <CR>

*E

.

[/M is necessary;

Monitor base;

\$HSIZE table;

\$DVSIZ table;

\$PNAME table;

\$STAT table;

Check that vectors in

permanent map are protected;

Exit to monitor]

F/B Monitor

.R PATCH

PATCH Version number

FILE NAME--			
*MONITR.SYS/M<CR>			[/M is necessary;
*BASE;ØR<CR>			Monitor base;
*Ø,146Ø2/	4ØØØ	316<CR>	\$HSIZE table;
*Ø,14636/	Ø	2ØØØ<CR>	\$DVSIZ table;
*Ø,1764Ø/	625Ø	7Ø37Ø<CR>	\$PNAME table;
*Ø,17674/	4	1ØØØ23<CR>	\$STAT table;
*Ø,17336	77	<CR>	Check that vectors in
*E			permanent map are protected;
.			Exit to monitor]
.			

At this point, the system should be re-bootstrapped to make the modified monitor resident. The device RC will then be available for use.

5.4 WRITING A SYSTEM DEVICE HANDLER

This section describes the procedures for writing a new system device handler. A system device is the device on which the monitor and handlers are resident. RT-11 currently supports the RK, RF, DP, DS, and DX disks, and DEctape as system devices. The procedures for writing the handler and creating a new monitor are explained, illustrated by the example in Section A.1, the RC11/RS64 handler.

The basic requirements for a system device are random access and read/write capability. These requirements are met by the RC11 disk, which is a multiple platter, fixed-head disk. When writing the driver, the procedures in Section 5.2 should be followed. Because the system handler is linked with the monitor, the additional tagging and global conventions described here must also be followed.

5.4.1 The Device Handler

The following conditions must be observed when writing a system handler. Refer to the example listing in Section A.1.

1. The handler entry point must be tagged xxSYS, where xx is the 2-letter device name. For the RC disk, this is RCSYS. See area D in the listing.
Important: Note that the tag is placed after the third word of the header block.
2. The entry points of all current system devices must be referenced in a global statement. These

currently include RKSYS, RFSYS, DSSYS, DXSYS, DPSYS, DTSYS and RCSYS. See Area A.

3. The entry point tags of all other system devices must be equated to zero. See area B in the listing.
4. A .CSECT SYSHND must be included at the top of the handler code. It is located above area C in the example.
5. The last word of the handler is used for the common interrupt entry address. This should have the tag \$INPTR and should be set to the value \$INTEN. See areas M and N in the example listing. These tags should be global. See area A.
6. The interrupt entry point should have the tag xxINT, or RCINT for this example, and this must be a global. See areas A and H.
7. The handler size must be global, with the symbolic name xxSIZE, or RCSIZE. See area A. This step is not necessary if the monitor sources are available and are being reassembled, since the global will be generated by the HSIZE macro. See Step 3 in Section 5.4.3.

5.4.2 The Bootstrap

This section describes the procedure for modifying the system bootstrap to operate with a new system device. Either the bootstrap source must be acquired, or the listing in Section A.2 may be used. Again, the RC11/RS64 disk is used for an example. The references in this section, however, are to the bootstrap listing found in Section A.2 of Appendix A.

The following changes must be made to the bootstrap to support a new system device:

1. Add a new conditional, \$xxSYS, to the list at point AA. Here xx is the 2-letter device name, and in this case the conditional is \$RCSYS.
2. Add a simple device driver for the device inside a \$xxSYS conditional. This is shown at area CC. Because the RC11 is similar to the other disks, it is possible to share code with the other device drivers, reducing the implementation effort. To do this, the \$RCSYS conditional is added at area BB and the device specific code is at area FF. This code merges with the common code at area GG.

3. The device driver has these characteristics:
 - a. The SYSDEV macro must be invoked for the device. The macro arguments are the 2-letter device name and the interrupt vector address. For this example, the arguments are "RC" and "210", shown at area DD on the listing.
 - b. The device driver entry point must have the tag READ. See area EE.
 - c. When the driver is entered:
 - RØ = Physical Block Number
 - R1 = Word Count
 - R2 = Buffer Address
 - R3,R4,R5 = are available for use by the driver routine
 - d. The driver must branch to BIOERR if a fatal I/O error occurs.

5.4.3 Building the New System

This section describes the procedure for building a new monitor using the system device handler and bootstrap just developed. Again, the example used is the RC11/RS64 disk, and the appropriate listings are those in Sections A.1 and A.2.

The procedure is:

1. Assemble the handler, producing an object module with the name xx.OBJ, where xx is the 2-letter device name. In this example, the name is RC.

```
.R MACRO
*RC.OBJ=RC.MAC
```

2. Assemble the bootstrap, defining the conditional \$xxSYS (where xx is again the device name; e.g., \$RCSYS). Define the conditional BF if an F/B bootstrap is desired. Let BF be undefined for an S/J bootstrap. For the S/J bootstrap:

```
.R MACRO <CR>
*RCBTSJ=TT:,DK:BSTRAP<CR>
^$RCSYS=1<CR>
^Z^$RCSYS=1<CR>
^ZERRORS DETECTED: Ø
FREE CORE: 156Ø8. WORDS
```

For the F/B bootstrap:

```
.R MACRO<CR>
*RCBTFB=TT:,DK:BSTRAP<CR>
^$RCSYS=1<CR>
BF=1<CR>
^Z^$RCSYS=1<CR>
BF=1<CR>
^ZERRORS DETECTED: 0
FREE CORE: 15580. WORDS
```

3. If the monitor sources are available, the DEVICE macro described in Section 2.5.2.9 can be invoked for the new device by editing the macro call into RMONFB.MAC and RMONSJ.MAC and reassembling the monitor. For the RC device, the macro would be:

```
DEVICE RC 2000 100020 RCSYS
```

The HSIZE macro, described in the same section, must also be invoked. For the RC device, the macro would be:

```
HSIZE RC,316,SYS
```

Monitor assembly instructions are in Appendix A of the RT-11 System Reference Manual. If this approach is used, the table patching procedure in step 5 is not necessary.

4. Link the monitor with the new bootstrap and device handler.

For S/J:

```
.R LINK
*RCMNSJ.SYS,MAP=RCBTSJ,RT11SJ,RC
```

For F/B:

```
.R LINK
*RCMNFBSYS,MAP=RCBTFB,RT11FB,RC
```

5. If step 3 was not done and step 4 used the current monitor object modules, then the monitor tables must be patched to enter the device information. The monitor device tables are located using the procedure in Section 5.3. An additional table, the \$ENTRY entry point table, must also be patched. For this example, assume the table addresses are:

<u>Table</u>	<u>S/J Address</u>	<u>F/B Address</u>
\$HSIZE	13674	14602
\$DVSIZ	13730	14636
\$PNAME	16516	17640
\$STAT	16552	17674
\$ENTRY	16612	17612

NOTE

These table addresses are for illustration only. Consult the Getting Started with RT-11 (V02B) manual for the table addresses of the current monitor release and for the address of BASE.

A link map was made during the linking sequence in Step 4. Locate the value of the system handler entry point, xxSYS. For this example, the tag is RCSYS and its value is found to be 56266 for F/B. This value is put in the \$ENTRY table. The other values were determined in Section 5.3:

```
$HSIZE = 316
$DVSIZ = 20000
$PNAME = 70370
$STAT  = 100023
$ENTRY = 56266 (F/B) 45056 (S/J)
```

The patch procedure for the S/J monitor, using the twelfth slot, would then be:

```
.R PATCH<CR>
PATCH Version number
FILE NAME--
*RCMNSJ.SYS/M<CR>           [The /M is necessary;
*BASE;0R<CR>               Monitor base;
*0,13674/ 4000 316<CR>     $HSIZE table;
*0,13730/ 0 20000<CR>      $DVSIZ table;
*0,16516/ 6250 70370<CR>   $PNAME table;
*0,16552/ 4 100023<CR>    $STAT table;
*0,16612/ 0 45056<CR>     $ENTRY table;
*E                           Exit to monitor]
-
:
```

For the F/B monitor:

```
.R PATCH<CR>
PATCH Version number
FILE NAME--
*RCMNF.B.SYS/M<CR>
*BASE;0R<CR>
*0,14602/ 4000 316<CR>     [$HSIZE table;
*0,14636/ 0 20000<CR>     $DVSIZ table;
*0,17640/ 6250 70370<CR>   $PNAME table;
*0,17674/ 4 100023<CR>    $STAT table;
*0,17564/ 0 56266<CR>     $ENTRY table;
*E                           Exit to monitor]
-
:
```

The new monitor is now complete and may be used by transferring it to an RC disk and renaming it to MONITR.SYS.

5.5 DEVICES WITH SPECIAL DIRECTORIES

The RT-11 monitor can interface to devices having nonstandard (that is, non RT-11) directories. This section discusses the interface to this type of device.

5.5.1 Special Devices

Special devices are file-structured devices that do not use an RT-11 directory format. Examples are magtape and cassette as supported under RT-11. They are identified by setting bit 14 in the device status word. The USR processes directory operations for standard RT-11 devices; for special devices, the handler must process directory operations, as well as data transfers.

5.5.1.1 Interfacing to Special Device Handlers - There are three types of processes that a special device handler must perform:

1. Directory operations (.LOOKUP, .ENTER, etc.)
2. Data transfer operations (.READ, .WRITE)
3. Special operations (rewind, backspace, etc.)

The particular process required is passed to the handler in the form of a function code, located in the even byte of the fourth word of the I/O queue element (see Section 5.1.1). The function code may be positive or negative. Positive codes are used for processes of types 1 and 2 above; negative codes indicate device-dependent special functions.

The positive function codes are standard for all devices and include:

<u>Code</u>	<u>Function</u>
Ø	Read/Write
1	Close
2	Delete
3	Lookup
4	Enter

These functions correspond to the programmed requests .READ/.WRITE, .CLOSE, .DELETE, .LOOKUP, and .ENTER, described in Chapter 9 of the RT-11 System Reference Manual. The .RENAME request is not supported for special devices.

A queue element for a special handler will look identical to an element for a standard RT-11 handler when the function is a .READ/.WRITE (negative word count implies a .WRITE). For the remaining positive functions, word 5 of the queue element (the buffer address word discussed in Section 5.1.1) will contain a pointer to the file descriptor block, containing the device name, file name, and file extension in .RAD5Ø format.

Negative function codes are used for device-dependent special functions. Examples of these are backspace and rewind for magtape. Because these functions are characteristic of each device type, no standard definition of negative codes is made; they are defined uniquely for each device.

Software errors (for example, file not found or directory full) occurring in special device handlers during directory operations are returned to the monitor through the procedure described next. A unique error code is chosen for each type of error. This error code is directly returned by placing it in SPUSR (special device USR error), located at a fixed offset (272) into RMON. (Section 2.5.1 discusses monitor fixed offsets.) Hardware errors are returned in the usual manner by setting bit Ø in the channel status word pointed to by the second word of the queue element.

5.5.1.2 Programmed Requests to Special Devices - Programmed requests for directory operations and data transfers to special devices are handled by the standard programmed requests. When a .LOOKUP is done, for example, the monitor checks the device status word for the special device bit. If the device has a special directory structure, the proper function code is inserted into the queue element and the element is directly queued to the handler, by-passing any processing by the RT-11 USR. Device independence is maintained, since .FEAD, .WRITE, .LOOKUP, .ENTER, .CLOSE, and .DELETE operations are transparent to the user.

Requests for device-dependent special functions having negative function codes, must be issued by using the .SPFUN special function programmed request, described in Chapter 9 of the RT-11 System Reference Manual.

5.6 ADDING A SET OPTION

The Keyboard Monitor SET command permits certain device handler parameters to be changed from the keyboard. For example, the width of the line printer on a system can be SET with a command such as:

```
SET LP WIDTH=80
```

This is an example of a SET command that requires a numeric argument. Another type of SET command is used to indicate the presence or absence of a particular function. An example of this is a SET command to specify whether an initial form feed should be generated by the LP handler:

```
SET LP FORM           (generate initial form feed)
SET LP NOFORM        (suppress initial form feed)
```

In this case, the FORM option may be negated by appending the NO prefix.

The SET command is entirely driven by tables contained in the device handler itself. Making additions to the list of SET options for a device is easy, requiring changes only to the handler, and not to the monitor. This section describes the method of creating or extending the list of SET options for a handler. The example handler used is the LP/LS11 line printer handler, listed in Appendix A in Section A.3. The SET command is described in Chapter 2 of the RT-11 System Reference Manual.

Device handlers have a file name in the form xx.SYS, where xx is the 2-letter device name; e.g., LP.SYS. Handler files are linked in save image format at a base address of 1000, in which a portion of block 0 of the file is used for system parameters. The rest of the block is unused, and block 0 is never FETCHed into memory. The SET command uses the area in block 0 of a handler from 400 to 776 (octal) as the SET command parameter table. The first argument of a SET command must always be the device name; e.g., LP in the previous example command lines. SET looks for a file named xx.SYS (in this case LP.SYS) and reads the first two blocks into the USR buffer area. The first block contains the SET parameter table, and the second block contains handler code to be modified. When the modification is made, the two blocks are written out to the handler file, effectively changing the handler.

The SET parameter table consists of a sequence of 4-word entries. The table is terminated with a zero word; if there are no options available, location 400 must be zero. Each table entry has the form:

```

.WORD      value
.RAD50     /option/           [2 words of RAD50]
.BYTE      <routine-4000>/2
.BYTE      mode

```

where:

value is a parameter passed to the routine in register 3.

option is the name of the SET option; e.g., WIDTH or FORM.

routine is the name of a routine following the SET table that does the actual handler modification.

mode indicates the type of SET parameter:

- a. Numeric argument - byte value of 100
- b. NO prefix valid - byte value of 200

The SET command scans the table until it finds an option name matching the input argument (stripped of any NO prefix). For the first example command string, the WIDTH entry would be found (area 2 in the listing in Section A.3). The information in this table entry tells the SET processor that O.WIDTH is the routine to call, that the prefix NO is illegal and that a numeric argument is required. Routine O.WIDTH is located at area 4 on the listing. It uses the numeric argument passed to it to modify the column count constant in the handler. The value passed to it in R3 from the table is the minimum width and is used for error checking.

The following conventions should be observed when adding SET options to a handler:

1. The SET parameter tables must be located in block 0 of the handler file and should start at location 400. This is done by using an .ASECT 400 (area 1 on the listing).
2. Each table entry is four words long, as described previously. The option name may be up to six .RAD50 characters long, and must be left-justified and filled with spaces if necessary. The table terminates with a zero (area 3 on the listing).

3. The routine that does the modification must follow the SET table in block 0 (area 4 on the listing). It is called as a subroutine and terminates with an RTS PC instruction. If the NO prefix was present and valid, the routine is entered at entry point +4. An error is returned by setting the C bit before exit. If a numeric argument is required, it is converted from decimal to octal and passed in R0. The first word of the option table entry is passed in R3.
4. The code in the handler that is modified must be in block 1 of the handler file, i.e., in the first 256 words of the handler. See areas 6 and 7 on the listing for code modified by the WIDTH option.
5. Since an .ASECT 400 was used to start the SET table, the handler must start with an .ASECT 1000. See area 5 on the listing.
6. The SET option should not be used with system device handlers, since the .ASECT will destroy the bootstrap and cause the system to malfunction.

5.7 CONVERTING USER-WRITTEN HANDLERS

User-written device handlers must, in all cases, conform to the standard practices for Version 2 (2B). General programming information is discussed in Appendix H of the RT-11 System Reference Manual. Points to consider when converting user-written device handlers (written under Version 1 of the RT-11 System) follow; the details of these procedures have already been discussed.

1. The last word of a device handler is used by the monitor, thus the user must be sure to include one extra word at the end of his program when indicating the handler size.
2. The third header word of the handler should be 340, indicating that the interrupt should be taken at level 3.
3. It is not necessary to save/restore registers when the handler is first entered, although to do so is not harmful.
4. When an interrupt occurs, the handler must execute an .INTEN request or its equivalent. On return from .INTEN, R4 and R5 may be used as scratch registers. Device handlers may not do EMT requests without executing a .SYNCH request.
5. The handler must return from an interrupt via an RTS PC.
6. When the transfer is complete, the handler must exit to the monitor to terminate the transfer or enter a completion routine. When return is made to the monitor, R4 should point to the fifth word of the handler.

7. The handler should contain an abort entry point (located at INTERRUPT SERVICE -2) to which control is transferred on forced exit. The abort entry point should contain a BR instruction to code that will perform the necessary operations (stop device action and exit to monitor completion code).

CHAPTER 6

F/B MONITOR DESCRIPTION

The RT-11 Foreground/Background Monitor permits two jobs to simultaneously share memory and other system resources. The foreground job has priority and executes until it is blocked (i.e., execution is suspended pending satisfaction of some condition, such as I/O completion). When the foreground job is blocked, the background job is activated and executes until it finishes or until the foreground blocking condition is removed.

6.1 INTERRUPT MECHANISM AND .INTEN ACTION

All interrupt handlers must be entered at priority level 7 and must execute a .INTEN request on entry. The handler will then be called (as a co-routine of the monitor in system state) at its normal priority level. This is essential to the operation of RT11 for two reasons:

1. As a co-routine of the monitor, the interrupt handler exits to the monitor, which then does job scheduling.
2. Because of the above condition, there is a danger that interrupt processing may be postponed due to a context switch. For example, if a disk interrupts a lower priority device handler and goes to I/O completion, the monitor may switch to the foreground job and delay the lower priority interrupt until the foreground job is again blocked. By requiring the .INTEN request of all interrupt handlers, the monitor can assure that all interrupts are processed before the context switch is made.

The .INTEN request is implemented as a JSR R5 to the first fixed-offset location of RMON, which contains a jump to the interrupt entry code. This code saves R4 (R5 was saved by the JSR) and increments the system state counter. If the interrupt occurred on a job stack, the stack pointer is switched to use the system stack. The priority is lowered

to the handler's requested priority and control returns to the handler via another JSR instruction.

The handler interrupt code now executes in system state, with several results: any further interrupts are handled on the system stack, preventing their loss by a context switch to another job's stack; a context switch or completion routine cannot occur until all pending interrupts are processed; any error occurring in the handler occurs in system state, causing a fatal halt. When the handler exits via an RTS PC instruction, control returns to the monitor, which can now enter the scheduling loop if all interrupts have been processed.

6.2 CONTEXT SWITCH

When passing control from one job to another, the F/B Monitor does a complete context switch, changing the machine environment to that of the new job. The current context is saved on the stack of the current job and is replaced by the context of the new job.

The information saved on the stack includes:

1. The general registers (R0-R5)
2. The system communication area (memory locations 34-52)
3. The FPP registers, if used
4. The list of special locations supplied by the job (via .CNTXSW), if any

In addition, the stack pointer (R6) is saved in the job's impure area at offset I.SP (=50). The switch requires a minimum of 23_{10} words of stack, not including the special swap list.

The following are the minimum calculated times to context switch between jobs. The assumptions are that the F/G job is waiting for I/O completion, the handler completes an I/O request, and there are no user I/O completion routines.

Processor	$\frac{11}{20}$	$\frac{11}{40}$	$\frac{11}{45}$
(core memory)	.66 ms	.36 ms	.28 ms

6.3 BLOCKING A JOB

The F/B Monitor gives priority to the foreground job, which runs until it is blocked by some condition. In this case, the background, if runnable (i.e., not blocked itself), is scheduled. The conditions which may block a job are flagged in the I.JSTA word, which is located in the job's impure area:

<u>Tag</u>	<u>Bit in I.JSTA Word</u>	<u>Condition</u>
TTIWT\$	14	Waiting for terminal input
TTOWT\$	13	Waiting for room in output buffer
CHNWT\$	11	Waiting for channel to complete
SPND\$	10	Suspended
NORUN\$	9	Not loaded
EXIT\$	8	Waiting for all I/O to stop
KSPND\$	6	Suspended from KMON
USRWT\$	4	Waiting for theUSR

6.4 JOB SCHEDULING AND USE OF .SYNCH REQUEST

The F/B Monitor uses a scheduling algorithm to share system facilities between two jobs. The goal of the scheduler is to maximize system utilization, with priority given to the foreground job. The scheduler is generalized to use job numbers for scheduling, the higher job number having the higher priority. The background job is assigned job number 0 and the foreground job number 2. Job numbers must be even.

The foreground job runs until it is blocked by some condition (see Section 6.3), at which point the scheduler is initiated. The job list is scanned top down (from highest to lowest priority) for the highest priority job that is runnable. A job is runnable if it is not blocked, or if it is only blocked pending completion and is not suspended. If no jobs are currently runnable, the idle loop is entered.

If the new job is runnable, a context switch is made. The context switch routine tests for the completion pending condition (i.e., I/O is finished and a user completion routine was queued). In this case, a pseudo-interrupt is placed on the job's stack to call the completion queue manager when the scheduler exits to the job.

The scheduler is event driven and is entered from the common interrupt exit path whenever an event has occurred which requires action by the scheduler. The set of such events include:

1. An .EXIT or .CHAIN request
2. A job abort from the console, or an error abort
3. I/O transfer completed
4. Expiration of timed wait
5. A blocking condition encountered:
 - a. .TWAIT request or SUSPEND command
 - b. .TTYIN or .CSI waiting for end of line
 - c. .TTYOUT or .PRINT waiting for room in output buffer
 - d. Attempt to use busy channel
6. A blocking condition removed
7. No queue elements available
8. .SYNCH request (see below).

The .SYNCH request is used in interrupt routines to permit the issuing of other programmed requests. The .SYNCH macro is expanded as a JSR R5 to the .SYNCH code in the F/B resident monitor. The .SYNCH routine uses the associated 7-word block as a queue element for the completion queue.

If the .SYNCH block is not in use, register R5 is incremented to the successful return address and placed in the block as the completion address. The word count is set to -1 to prevent the block from being linked into the AVAIL queue. The block is placed in the completion queue, at its head, and the job associated with the .SYNCH request is flagged to have a completion routine pending. A request for a job switch is entered before the .SYNCH logic exits with an interrupt return.

On exit from the interrupt with a job switch pending, the scheduler is entered and the completion queue manager is called. When control finally returns to the code following the .SYNCH request, it is executing as a completion routine at priority level 0. It can now issue programmed requests without fear of being interrupted. If another interrupt comes in, it will be queued pending return of the current

completion routine, since the .SYNCH block is freed before calling the completion routine. Further interrupts will be rejected by the .SYNCH code, unless provision is made for supplying extra .SYNCH blocks.

6.5 USR CONTENTION

The directory operations handled by the USR are not re-entrant, particularly since the directory segment is buffered within the USR. Therefore, to use the USR in F/B, a job must have ownership of the USR. To facilitate this, the F/B monitor maintains a USR queuing mechanism.

Before issuing a USR request, a job must request ownership of the USR. If the USR is in use by another job, even of lower priority, the requesting job is blocked and must wait for the USR. The USRWT\$ flag is set in the I.JSTA word (see Section 6.3) and the job cannot continue until the USR is released and the blocking bit cleared. When the USR is released, the job list is scanned for jobs waiting for the USR, starting with the job having highest priority.

Because of the impact this may have on system performance, CSI requests are handled differently in the F/B system than in the S/J Monitor. If the command string is to come from the console keyboard, the prompting asterisk is printed and then the USR is released, pending completion of command line input. This prevents a job doing a CSI request from locking up the USR and blocking another, perhaps higher priority, job from executing. A job can determine if the USR is available by doing a .TLOCK request (see Chapter 9 of the RT-11 System Reference Manual).

6.6 I/O TERMINATION

Because of the multi-job capabilities of RT-11 F/B, termination of I/O on job exit or abort must be handled differently than in the S/J Monitor. The use of the RESET instruction is unacceptable, and a form of I/O rundown must be used. This is done by the IORSET routine, called when doing an abort or hard exit.

The IORSET routine searches the queue of every resident handler for elements belonging to the aborted job. If a handler is found to be resident and active (i.e., there are elements on its queue), the IORSET routine "holds" the handler from initiating a new transfer by setting bit 15 of the LQE word (entry point) in the handler. The

current transfer may complete, but the hold bit will prevent the queue manager from initiating a new transfer.

While it is held, the handler's queue is examined for the current request. If it belongs to the aborted job, the handler's abort entry point is called to stop the transfer. The queue of pending I/O requests is then examined and any elements belonging to the aborted job are discarded. The hold flag is cleared and a test is made to see if the current transfer completed while the handler was held. If it did, the completion queue manager, COMPLT, is again called to return the completed element and initiate the next transfer. At this point, any elements belonging to the aborted job will have been removed from the queue.

After the device handlers are purged, the internal message handler is examined for waiting messages that were originated by the aborted job. All such messages are discarded. Finally, all mark time requests belonging to the aborted job are cancelled.

CHAPTER 7
RT-11 BATCH

The RT-11 BATCH system is composed of a BATCH compiler and a run-time handler. The BATCH compiler converts BATCH Job Control language into a format comprehensible to the BATCH run-time handler. The compiler creates a control (CTL) file (from the BATCH language statements) which is then scanned by the handler; the CTL format is a versatile programming language in its own right. The result is a BATCH system that is simple to use, and yet easily customized to handle different situations.

7.1 CTL FORMAT

The BATCH run-time handler uses a unique language format that includes many programming features, such as labels, variables, and conditional branches. The directives are explained in detail in Chapter 12 of the RT-11 System Reference Manual.

Each directive consists of a backslash character followed by one or more other characters. For example, to run PIP and generate a listing, the CTL directives \E (execute) and \D (data line) are used:

```
\ER PIP  
\DLP:=/L
```

Messages are sent to the console device by using the \@ directive:

```
\@ PLEASE MOUNT DT2
```

Labels and unconditional branches are implemented with the \L (label) and \J (jump) directives:

```
\JEND 1
      .
      .
      .
\LEND
```

Each BATCH command is sent to the log as it is executed, using the \C (comment) directive:

```
\C
$JOB
```

In this case, every character up to the next backslash is sent to the log.

7.2 BATCH RUN-TIME HANDLER

The BATCH run-time handler (BA.SYS) is constructed as a standard RT-11 device handler. To use the handler, it must be made permanently resident via the monitor LOAD command. The handler links itself into the monitor, intercepting certain EMTs described later.

The linking occurs the first time the BATCH compiler is run after the BA handler is loaded. The compiler does a .READW to the BA handler, which then links itself to the monitor and returns a table of addresses to the BATCH compiler. The linking is achieved by replacing the addresses of monitor EMT routines with corresponding addresses in the BATCH handler. Those EMTs that are diverted include:

<u>EMT</u>	<u>BATCH Handler Routine</u>
.TTYIN	B\$TIN
.TTYOUT	B\$TOT
.EXIT	B\$EXT
.PRINT	B\$PRN

Once the link is established, the BATCH handler cannot be unloaded. The links must first be undone by again running the BATCH compiler and specifying the /U switch. The compiler removes the links and prints a prompting message, after which the UNI. BA command can be issued.

With the BA handler linked to the monitor, all console terminal communication is diverted to BA, along with program exits. The BA handler then dispatches the program request to the monitor routine or diverts it to a routine in BA, depending on the values of switches in BATSW1. The switches are:

<u>TAG</u>	<u>BIT</u>	<u>DESCRIPTION</u>
HELP	0	0 = Do not log terminal input (.TTYIN) 1 = Log terminal input
DESTON	1	0 = EMT is going directly to monitor 1 = BA intercepts the EMT
SOURCE	2	0 = Character input by monitor from console terminal 1 = Character input comes from BATCH stream
COMWAT	3	0 = No command 1 = Command is waiting
ACTIVE	4	0 = Console terminal inactive 1 = Console terminal is active; i.e., BA is waiting for input from console terminal
DATA	5	0 = Characters are going to KMON; i.e., KMON is active in B/G 1 = Characters are going to B/G programs
BDESTN	6	0 = Output characters are going to console terminal 1 = Output characters are going to LOG
BGET	7	0 = Normal mode 1 = Get mode (\G); input comes from console terminal until <CR><LF> is encountered
NOTTY	8	0 = Log terminal output 1 = Do not log terminal output (.TTYOUT, .PRINT)
	9-13	Reserved
BSOURC	14	0 = BA directives come from console terminal 1 = BA directives come from CTL file
BEXIT	15	1 = A program has done an .EXIT while DATA switch was set

The BATSW1 word, located six bytes past the handler entry point, determines the state of the system at any given moment. If the word is zero, RT-11 operates normally. When the DESTON bit is set, EMTs are diverted to routines in BA for action, but the specific action taken by those routines is determined by the other switch bits.

For example, if the BDESTN bit is set, output from .TTYOUT and .PRINT is diverted from the console terminal to the log device. If SOURCE is set, the characters for the .TTYIN request are taken from the BATCH stream rather than from the console terminal via the monitor ring buffer. Directives for the BA handler itself may come from either the CTL file or the console terminal, depending on the state of the BSOURC bit.

The state of the background is reflected in the DATA bit. Either the KMON is active (DATA=0) or a program is active (DATA=1). If a program issues an .EXIT request while in DATA mode, the BEXIT state is entered until the BA handler encounters the next KMON directive (\E) in the BATCH stream, causing any unused \D lines to be ignored. A program can be aborted by diverting any of the .TTYIN, .TTYOUT or .PRINT requests to the .EXIT code in the monitor.

7.3 BATCH COMPILER

The obvious function of the BATCH compiler is to convert BATCH Standard Commands into the BA handler directives mentioned in Section 7.1, creating a control (CTL) file. BATCH jobs entered from a card reader or a file-structured device are compiled into a CTL file stored on a file-structured device for execution by the BA handler. However, the BATCH Compiler has other important functions; these are described in this section along with details on the initiation and termination of BATCH jobs.

7.3.1 BATCH Job Initiation

The following sequence of actions is performed by the BATCH Compiler when setting up a job for execution:

1. A check is made to ensure that LOG and BA device handlers are loaded and assigned properly. The LOG handler must be assigned the logical name LOG;; the BATCH Compiler may be run several times during the course of a job to do special tasks for the BA handler, and it will reference LOG:.
2. A nonfile-structured .LOOKUP is done on BA and a .READW is issued. If this is the first time BATCH has been run since BA was loaded, the handler links itself to the monitor (see Section 7.2). BA returns a list of

eleven pointers to important parameters within BA. These include:

- BA state word (BATSW1)
- CTL file savestatus area (INDATA)
- LOG file savestatus area (ODATA)
- Output (LOG) buffer (OUTBUF)
- Output buffer pointer (BATOPT)
- Output character counter (BATOCT)
- Input character counter (BATICI)
- Monitor EMT dispatch address save areas

3. A command string is collected from the console terminal and is processed by .CSISPC. An input file must be specified.
4. If the input file is a .BAT file to be compiled, a .CTL file is entered. If the LOG: device is file-structured, a fixed-size enter is done and then the file is initialized by writing zeroes in all blocks.
5. A .LOOKUP is done on all input files.
6. The .LOG file is .CLOSED so that a .LOOKUP and .SAVE-STATUS may be done. The savestatus data is placed in the ODATA area in BA.
7. If the input file is a .BAT file, it is now compiled, with output going into the .CTL file.
8. The .CTL file is closed, again so that a .LOOKUP and .SAVESTATUS may be done. The .SAVESTATUS data is transferred to the INDATA area in BA. Buffer pointers and counters in BA are initialized.
9. The BA handler is activated by setting the SOURCE, DESTON, BSOURC and BDESTN bits in the BATSW1 state word in BA. Control passes to BA when the compiler does an .EXIT, assuming an abort is not requested.
10. If an abort is requested (an error occurred during compilation or the /N switch was used), the .LOG file is .REOPENed and all \$ command lines are logged out with any error diagnostics. The BATSW1 word is then cleared before exiting, preventing the execution of the job.

The following switches are used by the BATCH system during job initiation and continuation, and should not be typed by the user:

- /B BATCH continuation of jobs in input stream
- /D Print the physical device name assigned a logical device name in a \$DISMOUNT command
- /M Make a temporary source file
- /R Return from \$CALL
- /S \$CALL subroutine

7.3.2 BATCH Job Termination

Every BATCH job must be terminated with an \$EOJ statement. The \$EOJ statement causes the compiler to insert the CTL directives:

```
\R BATCH
\D/R
```

The /R switch for the BATCH compiler, which is legal only when entered from a BATCH stream, is used to terminate a BATCH job. This switch causes the compiler to pop the BATCH stack up a level. If the stack was empty, the stream is finished and the compiler cleans up, clears the BATSWL word in BA, and exits. If the stack is not empty, the /R switch implies a return from a \$CALL. The stack contents are used to restore parameters in the BA handler so that control will return to the calling BATCH stream at the next statement after the \$CALL.

7.3.3 BATCH Compiler Construction

The BATCH Compiler is constructed in two pieces: a data area and a program area. The data area is located in low memory, in a .CSECT named UNPURE. The contents are described in the accompanying table (Table 7-1). The program section, located in the .CSECT named PROGRAM, starts at the symbol START. The general register R4 always points to UNPURE and all references to the data base are made as indexed references relative to R4.

Locations in the data base are created with the ENTRLO macro. For example,

```
ENTRLO BOTLCT,0
```

allocates one word in the data base and initializes it to zero. The symbol BOTLCT is an offset into the data base, so that references to BOTLCT are made in the form BOTLCT(R4).

Table 7-1
 BATCH Compiler Data Base Description

Tag	Byte Offset	Description
BATSWT	0	<p>BATCH Control Switches</p> <p>ABORT = 100000 ABORT after compile DATDOL = 40000 DATA or DOLLARS set NO = 20000 "NO" prefix on switch CTYOUB = 10000 Output to CTY (\@) LOGOUB = 4000 Output to LOG (\C) DATOUB = 2000 Output to user prog (\D) COMOUB = 1000 Output to monitor (\E) JOB = 400 \$JOB encountered MAKEB = 200 /B switch on command COMMA = 100 Comma terminates command BFORLI = 40 Next link requires FORTRAN library UNIQUE = 20 UNIQUE command option set BANNER = 10 Print BANNER on \$JOB, \$EOJ RT11 = 4 RT11 default on NO '\$' in Column 1 TIME = 2 Print time of day MAKE = 1 Create a source file</p>
BATSW2	2	<p>More BATCH Control Switches</p> <p>ABORT =100000 Second time through ABORT FIRST = 10000 First card processed SBIT = 4000 /S switch on command SEQ = 2000 \$SEQ card processed LSTBIT = 1000 Request temporary listing file COMSWB = 400 Command switches MAKEB = 200 Same as BATSWT STARFD = 100 Asterisk in FD field STAROK = 40 Wild card option is valid BNOEOJ = 20 \$JOB or \$SEQ before \$EOJ LSTDAT = 10 List DATA sections BEOF = 4 EOF encountered on .BAT file XSWT = 2 /X switch set EOJ = 1 \$EOJ encountered</p>
TMPSWT	4	Temporary command switches
COMSWT	6	Current command switches
LINSIZ	10	Input line buffer size
BINLCT	12	Last buffer character count
INSTAT	14	Input buffer status (see OTSTAT)
ICHRPT	16	Input character pointer
BINCTR	20	Input buffer counter

(continued on next page)

Table 7-1 (Cont.)
 BATCH Compiler Data Base Description

Tag	Byte Offset	Description
BINARG	22	Input file EMT argument list
BATIBK	24	Input file block number
BATIBP	26	Input buffer address
	30	Input buffer size
	32	Wait I/O
BOTLCT	34	Last output buffer character count
OTSTAT	36	Output buffer status BFREE = 1 0 → Buffer is free BWAIT = 2 In I/O wait BEOF = 4 End of file
OCHRPT	40	Output character pointer
BOTCTR	42	Output character count
BOTARG	44	Output file EMT argument list
BATOBK	46	Output file block number
BATOBP	50	Output buffer address
	52	Output buffer size
	54	Wait I/O
STACK	56	Compiler stack pointer save area These are the arguments passed between BATCH and BA:
BATSW1	60	Pointer to BATSW1 in BA.SYS
INDATA	62	Pointer to INDATA
ODATA	64	Pointer to ODATA
OUTBUF	66	Pointer to BATCH handler output buffer
BATOPT	70	Pointer to output character pointer
BATOCT	72	Pointer to output character counter
BATICT	74	Pointer to input character counter

(continued on next page)

Table 7-1 (Cont.)
 BATCH Compiler Data Base Description

Tag	Byte Offset	Description
		Pointers to EMT intercept pointers:
O\$EXT	76	.EXIT
O\$TIN	100	.TTYIN
O\$TOT	102	.TTYOUT
O\$PRN	104	.PRINT
		CSI Buffer:
SPC0	106	Channel 0
SPC1	120	1
SPC2	132	2
SPC3	144	3
SPC4	154	4
SPC5	164	5
SPC6	174	6
SPC7	204	7
SPC8	214	10
LINIMP	224	Pointer to command line buffer (LINIMM)
LINIMM	226	Command line input buffer
LINIMS	350	Command line buffer save area
LIBLST	470	ASCIZ name of FORTRAN default library plus a line buffer
BATIBF	610	BATCH Compiler input buffers (INBSIZ * 2)
BATOBF	2610	BATCH Compiler output buffers (OTBSIZ * 2)
QSET	4610	Seven I/O queue elements for double/buffering
SOUTMP	4700	Source temporary file descriptor
OBJTMP	4714	Object temporary file descriptor
LOGTYP	4730	LOG device status word (word 0 of .DSTATUS)
ARGARG	4732	EMT argument list for BA handler initialization

(concluded on next page)

Table 7-1 (Cont.)
 BATCH Compiler Data Base Description

Tag	Byte Offset	Description
STKBLK	4744	EMT argument list for READ/WRITE of BATCH stack
DEFCHN	4756	Default channel numbers
DEVSPC	4770	Pointer to device handler space
WDBLK2	4772	Two-word EMT argument block
WDBLK5	5000	Five-word EMT argument block
FTLPC	5012	Contents of PC on BATCH fatal error
AREA0	5014	Pointer to impure area
LSTTMP	5016	Listing temporary file descriptor
SWTMSK	5026	Switch mask for this BATCH directive
FD0	5030	File descriptor 0 for BATCH directive
FD1	5034	1
FD2	5040	2
FD3	5044	3
FD4	5050	4
FD5	5054	5

7.4 BATCH EXAMPLE

The following example demonstrates how the compiler converts BATCH Standard Commands into RT-11 BATCH handler directives. The example consists of a main BATCH stream, EXAMPL.BAT, and a BATCH subroutine file, EDITIT.BAT. EXAMPL creates a program, assembles and runs it. The program, called FILE.MAC, prints a message that is diverted to the log. The listing file from the assembly is printed and then deleted. The BATCH variable S is then tested and, if it is zero, the BATCH subroutine EDITIT is called. The EDITIT stream uses EDIT to edit the file FILE.MAC, changing the message to be printed. After return from EDITIT, the stream branches unconditionally to label L1, repeating the assembly and execution of FILE.MAC. EDITIT increments the variable S before returning, so that the BATCH stream, on encountering the IF statement again, now branches to label L2, skipping the call to EDITIT. \$DIRECTORY and \$DELETE operations are performed before finally exiting from BATCH.

Note the following about the .CTL files created:

1. The \$JOB command produces a comment for the log (the \C directive, but no action directives). Its function is to initialize the BATCH compiler.
2. The \$CREATE command produces directives that run the BATCH compiler, using the file name to be created with a /M switch. This is a special function of the BATCH compiler used to create data files. The compiler will enter the data that follows in the CTL file into the newly created file, until an EOF (CTRL/Z) is encountered. The data is fed to the compiler by the BATCH handler through the .TTYIN programmed request. After the EOF character is encountered, the BATCH compiler closes the new file and exits, returning control to the BATCH handler through the .EXIT request. In this example, the file created is called FILE.MAC.
3. The \$MACRO command has the /RUN switch appended, which forces the compiler to generate a series of assembly, link and execute instructions. A temporary execution file, 000000.SAV, is created from the assembled object module, FILE.OBJ. After execution with the monitor R command, the temporary execution file is deleted with PIP.
4. PIP is used to implement \$PRINT, \$DELETE, \$COPY, and \$DIRECTORY. The compiler translates these commands into the appropriate PIP command strings.

5. The variable S is defined to be zero with the LET statement. This translates into the BATCH handler directive,

```
\KS1<null>
```

which instructs the BATCH handler to set variable S to the value in the byte following the character 1.

6. Labels are implemented by inserting a \L directive followed by the 6-character label name into the CTL stream where the label was declared. The label is also logged out with the \C directive so that the labels will appear in the log.
7. The unconditional branch, or GOTO command, is implemented with the \J directive immediately followed by the label. Note that the BATCH programmer must indicate whether the branch is forward or reverse. In this case, the branch is a backward reference and a minus sign is prefixed to the label:

```
GOTO -L1
```

There is no error checking done by the compiler. If an error is made (e.g., the minus sign is left off the L1), the BATCH handler searches forward in the CTL stream until it finds the label. Since an error was made, the label will not be found. The search (and consequently the BATCH job) terminates when the label stopper (\L\$\$\$\$) is encountered at the end of the CTL file.

8. The IF conditional branch is implemented with the \I directive. The \I directive is followed by the name of the variable to be tested, the value to be tested against, and three label fields. Each label field consists of the 6-character label name with a reference character appended. The character 1 indicates the label is a forward reference, a 0 indicates a backward reference. The test value is subtracted from the current value of the variable and the appropriate branch is taken. If no label is specified for a field, it is filled with spaces and causes the BATCH stream to fall through to the next command if that branch is elected.
9. The \$CALL command is very useful and permits a BATCH stream to call another BATCH file as a subroutine, with control returning to the command following the \$CALL. The \$CALL is implemented by simply running the BATCH compiler, passing it the name of the \$CALLED routine with a /S switch appended. Another BATCH compile/execute sequence will follow, but the /S switch will cause the compiler to save certain locations in the BATCH handler in an internal stack in the BA.SYS file. In this example, the \$CALL EDITIT statement causes the file EDITIT.BAT to be compiled and executed.

10. BATCH variables may be used to enter ASCII values into a job stream. In the file EDITIT, the variable A is set equal to the value of the ESC (or ALT MODE) character. The variable A is inserted into a string of EDIT commands in place of the ALT MODE character.
11. The \$EOJ must terminate every BATCH job. The \$EOJ command generates the stopper label, \L\$\$\$\$\$, and then produces directives to run the BATCH compiler again, this time with a /R switch. The compiler, when given a /R switch, checks the BATCH stack. If it is empty, the compiler exits. Otherwise, the stack is popped to restore conditions in the BATCH handler prior to the \$CALL causing the push, and the BATCH stream continues. The \$EOJ finally generates a \E to bring in the KMON and a \F<CR> to terminate the BATCH stream.

EXAMPL.BAT

```
SJOB
SMESSAGE EXAMPLE BATCH STREAM
SCREATF FILE,MAC
      .MCALL .RPGDEF,.PRINT,.FXTT
      .REGDEF

START: .PRINT #MSG
      .EXIT
      .NLIST REY
MSG:   .ASCIZ /THIS MESSAGE COMES FROM THE BATCH STREAM/
      .EVEN
      .LIST REY
      .END START

SEOD
SRT11
      LET S=0

L1:
SMACRO/RUN FILE,LST/LIST FILE,MAC/INPUT FILE/OBJECT
SPRINT FILE,LST
SDELETEF FILE,LST
SRT11
      IF(S=0) ,,L2
SCALL EDITIT      !CALL EDITIT TO EDIT FILE,MAC
SRT11
      GOTO -L1

L2:
SDIRECTORY FILE.*
SDELETEF FILE.*
SEOJ
```

EDITIT.BAT

```
SJOB/RT11
S! JOB TO EDIT FILE,MAC
      XS          !INCREMENT S TO PREVENT RECURSION
      LET A=33    !A IS ALT MODE

.R EDIT
*ERFILE,MAC'A'R'A'A'
*GMSG:'A'KI      .ASCIZ /MODIFIED BY EDITOR RUN BY BATCH/
*'A'FX'A'A'
SEOJ
```

EXAMPL.CTL

```
\C
SJOB

SMESSAGE EXAMPLE BATCH STREAM
\E\O EXAMPLE BATCH STREAM
\C
SCREATF FILE.MAC
\ER BATCH
\DFILE.MAC/M=
    .MCALL .RFGDEF,.PRINT,.FXTT
    .REGDEF
START: .PRINT #MSG
    .EXIT
MSG:   .NLIST REY
    .ASCIZ /THIS MESSAGE COMES FROM THE BATCH STREAM/
    .EVEN
    .LIST REY
    .END START

\C
SEOD

SRT11
    LET S=0
\KS1 \LL1 \CL1:

SMACRO/RUN FILE.LST/LIST FILE.MAC/INPUT FILE/OBJECT
\ER MACRO
\DFILE,FILE.LST=FILE.MAC
\F\D\ER LINK
\O000000=FILE
\ER \O000000
\ER PIP
\O000000.SAV/O
\C
SPRINT FILE.LST
\ER PIP
\DLST:*,*/X=FILE.LST
\F\D\C
SDELEFT FILE.LST
\ER PIP
\DFILE.LST/O
\C
SRT11
    IF(S=0) ,,L2
\IS 1 1L2 1\L \C
SCALI EDITIT !CALL EDITIT TO EDIT FILE.MAC
\F\ER BATCH
\OEDITIT/S
\C
SRT11
    GOTO -L1
\JL1 0\LL2 \CL2:
```

EXAMPL.CTL (Cont)

```
$DIRECTORY FILE.*
\ER PIP
\DFILE.* /L
\F\D\C
$DELETE FILE.*
\ER PIP
\DFILE.* /D
\C
$EQU
\LS$SS$S\F\ER BATCH
\D/R
\E\F
```

EDITIT.CTL

```
\C
$JOB/RT11

$! JOB TO EDIT FILE.MAC
      XS          !INCREMENT S TO PREVENT RECURSION
\KS0\C LET A=33  !A IS ALT MODE
\KA1 \FR EDIT
\DFBFILE.MAC\KA2R\KA2\KA2
\DGMSG:\KA2KT .ASCIZ /MODIFIED BY EDIT RUN BY BATCH/
\D\KA2FX\KA2\KA2
\C
$EQU
\LS$SS$S\F\ER BATCH
\D/R
\E\F
```


EXAMPL.LOG

SJOB

SMESSAGE EXAMPLE BATCH STREAM

SCREATE FILE.MAC

SEOD

SRT11

LET S=0

L1 L11

SMACRO/RUN FILE.LST/LIST FILE.MAC/INPUT FILE/OBJECT

*ERRORS DETECTED: 0
FREE CORE: 15100. WORDS

*

EXAMPL.LOG (Cont.)

.MAIN. RT-11 MACRO VM02-10 10-APR-75 10:33:45 PAGE 1

```
1
2 000000          .MCALL .RFGDEF,.PRINT,.EXIT
3 000000          .REGDEF
4 000006          START: .PRINT #MSG
5                .EXIT
6 000010          .NLIST REY
7                .ASCIZ /THIS MESSAGE COMES FROM THE BATCH STREAM/
8                .EVEN
9                .LIST REY
                .END START
```

EXAMPL.LOG (Cont.)

,MAIN, RT=11 MACRO VM02-10 10-APR-75 10133:45 PAGE 1+
SYMBOL TABLE

MSG	000010R	PC	=X000007	R0	=X000000
R1	=X000001	R2	=X000002	R3	=X000003
R4	=X000004	R5	=X000005	SP	=X000006

START 000000R
. ABS. 000000 000
000062 001

ERRORS DETECTED: 0
FREE CORE: 15100. WORDS

FILE,FILE.LST=FILE.MAC

THIS MESSAGE COMES FROM THE BATCH STREAM

SPRINT FILE.LST

SDELFTF FILE.LST

SRT11
IF(S=0) ,,12

SCALL EDITIT ICALL EDITIT TO EDIT FILE.MAC

SJOB/RT11

S! JOB TO EDIT FILE.MAC
XS IINCREMENT S TO PREVENT RECURSTON
LET A=33 IA IS ALT MODE

*ERFILE.MACSRS
*
GMSGISKI .ASCIZ /MODIFIED BY EDITOR RUN BY BATCH/
SEXSS

SE0J
S88888

SRT11
GOTO -L1
L1:

SMACRO/RUN FILE.LST/LIST FILE.MAC/INPUT FILE/OBJECT

*ERRORS DETECTED: 0
FREE CORE: 15136. WORDS

*

EXAMPL.LOG (Cont.)

.MAIN. RT-11 MACRO VM02-10 10-APR-75 10:34:00 PAGE 1

```
1
2 000000          .MCALL  .REGDEF,.PRINT,.EXIT
3 000000          .REGDEF
4 000006          .PRINT  #MSG
5                .EXIT
6 000010          .NLIST  REX
7                .ASCIZ  /MODIFIED BY EDITOR RUN BY BATCH/
8                .EVEN
9                .LIST   REX
000000'          .END   START
```

EXAMPL.LOG (Cont.)

.MAIN. RT-11 MACRO VM02-10 10-APR-75 10:34:08 PAGE 1+
SYMBOL TABLE

MSG	000010R	PC	=X000007	R0	=X000000
R1	=X000001	R2	=X000002	R3	=X000003
R4	=X000004	R5	=X000005	SP	=X000006

START 000000R

. ABS. 000000 000
000050 001

ERRORS DETECTED: 0

FREE CORE: 15136. WORDS

FILE,FILE.LST=FILE,MAC

MODIFIED BY EDITOR RUN BY BATCH

\$PRINT FILE.LST

\$DDELETE FILE.LST

\$RT11

IF(S=0) ,,12

L2:

\$DIRCTORY FILE.*

10-APR-75

FILE .BAK 1 10-APR-75

FILE .MAC 1 10-APR-75

FILE .OBJ 1 10-APR-75

3 FILES, 3 BLOCKS

417 FREE BLOCKS

\$DDELETE FILE.*

\$EOJ

7.5 CTT TEMPORARY FILES

In certain cases the BATCH compiler will produce temporary files with the extension CTT and the file name of the BAT file being compiled. These files occur when a multiple input file command string is issued, or when an unexpected \$JOB or \$SEQ statement occurs in a BATCH stream, or when multiple jobs are run from the card reader or a .BAT file.

The CTT file is actually a CTL file used to link together execution of several BATCH jobs. Each CTT file contains the BA directives:

```
\ER BATCH  
\D/B
```

which execute the BATCH compiler, passing it the /B switch.

The CTT file also contains the following information:

1. Current input channel number (range is 3-10₈)
2. Current input file block number
3. The CTL file descriptor block (device, file name and file size)
4. The LOG file descriptor block (device, file name, and file size)
5. The set of input (BAT) file descriptor blocks (device and file name)

When the CTT file is executed, the compiler restores the input channel number and block number and the entire set of file descriptor blocks from the CTT file. If, for example, the input channel number is 4, the second of a string of .BAT files is compiled and executed.

APPENDIX A

SAMPLE HANDLER LISTINGS

A.1 RC11/RS64 DEVICE HANDLER

RC11 V01-01 (FIXED HEAD DISK) RT-11 MACRO V02-00 A-APR-75 12104126 PAGE 1

1
2
3
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35

TITLE RC11 V01-01 (FIXED HEAD DISK)
RT-11 RC11/RS64 DEVICE HANDLER

IDPC-11-XXXXX-A

IJPO

1 OCTOBER 1974

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```

37 000164 100397
38 000166 100705
39 000170 001755
40 000172 005405
41 000174 010506
42
43
44
45
46 000176 012746
47 000202 000205
48 000204 100066
49 000210 100316
50 000212 001373
51 000214 000726
52 000216 001401
53 000220 000205
54 000222
55
56 000222 000524
57 000224 010605

          RPL      RCHOME
          TSTB
          REG
          NEG
          MOV      R5,=(SP)
          I CALCULATE THP # OF SECTORS IN
          I TSP WORDS = ONP (SECTOR)

          MOV      R5,=(SP)
          ASB
          RORB   1(RP)
          NEFB   PSP
          RNF    18
          TST    (SP)+
          REG    PB
          TNP    PS
          PSI
          I PND OF SECTOR CALCULATION
          ADD
          MOV     (SP)+,PS

          ICURRENT QIE ELEMENT (RCC002)
          INPILL FOR READS
          I EVEN # BLOCKS WRITTEN?
          I VPS = FILL NOT NECESSARY.
          I WRITE WORD COUNTS ARE NEGATIVE IN
          I THE QIE.
          I THE CURRENT OPERATION

          I FINISH REPEAT COUNT ONTO STACK
          I DVTOP THP # WORDS
          I CHECK FOR SECTOR OVERFLOW
          I INCREMENT REPEAT COUNT
          I SPECTOR OVERFLOW ?
          INN
          I INCLUDE NEXT SECTOR

          ICALCULATE CURRENT DISK ADDR. (RCC0A)

```

```

RC11 V01-01 (PTXED HFAN DISK) RT=11 MACRO VM02=00 A=APP=75 I2104126 PARE 4+
RC=11 INTERRUPT ROUTINP

```

```

58 000226 050705 17700
59
60
61 000232 000724
62 000234 010724 000103
63 000240 010524
64 000242 010714
65 000244 060714 177500
66 000250 000207

          R18      017400,PS
          TST
          MOV     (R4)+
          MOV     R0R,(R4)+
          MOV     R5,(R4)+
          MOV     PC,0RA
          ADD     02PRO-..0RA
          RTS
          PC

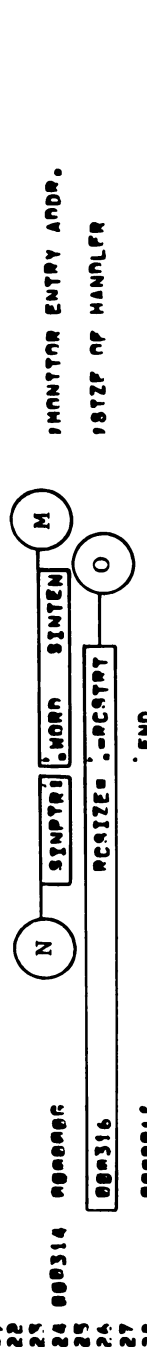
```



```

1  .SRTYL COMMON SUBROUTINE
2  IRPCOMN SUBROUTINE USED BY INTERRUPT
3  IAND ENTRY ROUTINES
4
5  000252 012704 177486 RC0M11 MOV 0R00,R4
6  000256 031014 RC0M11 RIT 0R0,0R4
7  000260 001402 RC0M11 RER 10
8  000262 012600 MOV (0R3)+,R0
9  000264 000717 RR R0,R0ME
10 000266 014705 MOV R0,R000,R5
11 000272 012306 MOV (R0)+,-(R0)
12 000274 004316 ASL 00P
13 000276 004316 ASL 00P
14 000300 004316 ASL (R0)+,0R4
15 000302 052014 R10 -(R0),-(R0)
16 000304 024404 CMP (R0)+,0R4
17 000306 012014 MOV (R0)+,
18 000310 004725 TST
19 000312 000000 RTS

```



MONITOR ENTRY ADDR.
START OF HANDLER

A.2 RC11/RS64 BOOTSTRAP

ROOT VM2R-01 RT-11 BOOTSTRAP RT-11 MACRO VM2R-00 A-APR-75 11149104
 TABLE OF CONTENTS

2-	2	MACROS, GLOBALS
3-	1	ASPCY
7-	29	BOOTSTRAP I/O DRIVER - RC11
10-	1	BOOTSTRAP CORE DETERMINATION
11-	1	READ MONITOR, LOOKUP HANDLERS
12-	1	RELOCATION LIST

ROOT VM2R-01 RT-11 BOOTSTRAP RT-11 MACRO VM2R-00 A-APR-75 11149104 PAGE 1

```

1 000001
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27

SRCVSB01
.TITLE ROOT VM2R-01 RT-11 BOOTSTRAP
; RT-11 BOOTSTRAP
;
; DEC-11-ORBITA-0
; COPYRIGHT (C) 1975
; DIGITAL EQUIPMENT CORPORATION
; MAYNARD, MASSACHUSETTS 01754
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; OR RELIABILITY OF ITS SOFTWARE ON EQUIPMENT
; WHICH IS NOT SUPPLIED BY DIGITAL.

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.SMTTL MACROS, GLOBALS
.MCALL .:V1:.
.P.V1:.
.MCALL .:EXIT, .:LOOKUP, .:PRINT, .:SAVESTATUS
;*****
; CONDITIONAL ASSEMBLY OF ROOT FOR SINGLE USER OR RP SYSTEM
;.IF NDP 0P=0
; GLOBAL REFERENCES TO MONITOR:
;GLOBL SOVREC, SENTRY, SINPTR, AKMLOC, SHONBL, SPNAME, SBLOT
;GLOBL SWPRL, SUSRLC, SPNAME
;GLOBL RSTRNG, CORPTR, AKASSG, FILLR, MHPPUS, MHQSPS, KMLOC
;GLOBL KMON, KMONSZ, KWILLS, MAPOFF, GCQMP, RT11SZ
;GLOBL RTLEN, RTSIZE, SWAPSZ, SVENTO, SYINDO, SYNCH, SYABSS
;GLOBL SYSLOW, TTIBUP, YTABUP, USRLOC, USRSZ, MAXSYN
;GLOBL RELLST
; FOLLOWING ARE GLOBALS FOR EITHER 0P OR SU SYSTEM, BUT NOT BOTH
;.IF NE 0P
;GLOBL RCNTXT, RKEND1, RKEND2, RKEND3, CNTXT, PUDBE1, PUDBE2
;GLOBL MSBENT, RMNSP, SHPTR, SHPTR, YTOUSR, YTOUSR, .SCRN
;.IF
;GLOBL AVAIL, I.CSW, PFPADD, PPIREN, MONLOC, TRAPLC, TRAPER
;.ENDC
PERM = 2000
ENDBLK = 4000
JSM = 44
SR = 177570
; REGISTER DEFINITIONS
R0=00
R1=01
R2=02
R3=03
R4=04
R5=05
R6=06
R7=07
; MONITOR OFFSET CONSTANTS
000000
000001
000002
000003
000004
000005
000006
000007

```


49 000300 HARDWARE CONFIGURATION WORD
 50 000274 SYSTEM UNIT #
 51
 52 177546 CLOCK STATUS REGISTER
 53 172000 67400 LOCATION
 54 177560 KEYBOARD STATUS
 55 177562 " " BUFFER
 56 177564 PRINTER STATUS
 57 177546 " " BUFFER

ROOT VR2R=01 RT-11 BOOTSTRAP RT-11 MACRO VM02-00 A-APR-75 11:49:04 PAGE 3

1 .START ASBCT
 2
 3 .IF NDP SRPSYS
 4 .IF NDP SDTSSYS
 5 .IF NDP SOPSYS
 6 .IF NDP SDSSYS
 7 .IF NDP SRCSYS
 8 .IF NDP SDXSYS
 9 SRKSYS= N
 10 .PENDC
 11 .PENDC
 12 .PENDC
 13 .PENDC
 14 .PENDC
 15 .PENDC
 16
 17 .ARECT
 18 . = N
 19 .PAR
 20 RR ROOT1
 21
 22 .IF NDP SDXSYS
 23 . = N
 24 JMP ROOT
 25
 26 .IFF
 27 CS00= 1
 28 CS01= 2
 29 CS02= 4
 30 CSUNTY= 20
 31 CSONE= 40
 32 CSTR= 200
 33 CSFR= 10000
 34
 35 RXCS= 177170



18
 19 00000
 20 00002 000016
 21
 22
 23 000034 000107 000040
 24
 25
 26
 27
 28
 29
 30
 31
 32
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 34
 35

RETURN ON SRKSYS IF ALL OTHERS ARE OFF
 ROOT VALIDATION PATTERN
 BRANCH TO REAL ROOT
 PUT THE JUMP ROOT IN TRAP VECTOR
 START THE BOOTSTRAP
 START FUNCTION
 EMPTY BUFFER
 PAR SECTOR
 UNIT 1 SELECTION
 TRY DONE
 TRY TRANSFER READY
 TRY PROR
 RXCS STATUS REGISTER


```

01      . = 200
02      CMPB (R3)+, (R3)+
03      RPT
04      CMPB (R3)+, (R3)+
05      RPT
06      RPT
07      CLR
08      MOV RB, (PC)+
09      RTI
10      STRWAIT, #200
11      JMP
12      .ENDC
13
14      .SECTOR 2 OF PX ROOT
15      .BUMP TO SECTOR 5
16      .CALL READS SURROUTINE
17      .BUMP TO SECTOR 7
18      .CALL READS SURROUTINE
19      .CHECK UNIT ID
20      .BRANCH IF BOOTING UNIT I, R0=1
21      .SET TO UNIT 0
22      .SAVE UNIT BOOTED FROM FOR LATER
23      .SAVE THE UNIT MEMP
24      .LPTS HANDLE ERRORS DIFFERENTLY
25      .NOW WP ARE READY TO DO THE REAL BOOT

```

ROOT VR2R=01 RT-11 ROOTSTRAP RT-11 MACRO VM02=09 A=APR-75 11149104 PAGE 4

```

1  . FOLLOWING ARE THE ROOTSTRAP I/O DRIVERS FOR EACH VALID
2  . SYSTEM DEVICE.
3  . CALLING SEQUENCE:
4  . R0 = PHYSICAL BLOCK TO READ/WRITE
5  . R1 = WORD COUNT
6  . R2 = BUFFER ADDRESS
7  . R3,R4,R5 ARE AVAILABLE AND MAY BE DESTROYED BY THE DRIVER
8  . IF THE DRIVER MUST GO TO HIGHERR IF A FATAL I/O ERROR OCCURS.
9  . IT MUST ALSO INVOKE THE MACRO SYSDEV
10 . MACRO SYSDEV NAME, VECTOR
11 . GLOB NAME, INT, NAME'STZP
12 . SYMNAME S R
13 . IRPC X, <NAME>
14 . SYMNAME S <SYMNAME>, 'X=1000'+S
15 . ENDR
16 . SYMVEC S VECTOR
17 . S SYMVEC. WORD NAME, INT, 340
18 . S SYMSIZE. WORD NAME'STZP
19 . S 400
20 . SYMTO S VECTOR / 20
21 . SYMITS S #1100000
22 . RPT <VECTOR & 17> / 4
23 . SYMITS S SYMITS / 4
24 . ENDR
25 . ENDM
26 . SYMDFV
27

```

```

1  .IF DF S0SSYS IRS SYSTEM
2  .SRTL ROOTSTRAP T/O DRIVER = R911
3
4
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32
33
34
35
    .IF DF M0USSC
    SYRDEV 08,150
    RSC92  0 174310
    .ENDC
    .IF M0P M0USSC
    SYSDPV 09,204
    RSC92=172050
    .ENDC

    READI  MOV  R0,R4
    MOV  RSC92,R5
    MOV  (R5),-(SP)
    MOV  R0,R5
    RIT  R2,16(R5)
    RNE  18
    ABL  R4
    RLC  R4
    MOV  (SP)+,(R5)
    MOV  R4,-(R5)
    MOV  R2,-(R5)
    MOV  R1,-(R5)
    RER  R5
    MOV  R71,-(R5)
    RIT  R10200,R5
    RER  28
    RMT  R10ERR
    RTS  PC

    COPY BLOCK NUMBER
    IPOINT TO REGISTERS
    ISAVE UNIT 0
    ICONTROLLER CLEAR
    IWHAT IS IT?
    IIT'S AN R804
    IIT'S AN R803
    ICONVERT TO TRACK/SECTOR
    ISTRIP TO UNIT BITS
    ISPT BLOCK

    IGO, READ, NO INTERRUPT
    IWAIT FOR DONE OR ERROR
    IBOOT ERROR

    .ENDC

```

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;IP OF 80858
;SMTL ROOTSTRAP I/O DRIVER - RP11
; RP11 DISK DRIVER

SYSDPV NP,254
RPCB= 176714
RPNB= 176710
RPNB= 176724
CB,GB= 000001
CB,RR= 000004
CB,ORV= 003400
DB,ATT= 000377

READ: MOV R0,R3
JMP R2,DIV
.MORN
MOV R4,=(0P)
MOV R5,R3
JMP R2,DIV
.MORN
SWAB
RIS (0P)+,R4
MOV R0PDA,R3
MOV R4,0R3
MOV R5,=(R3)
MOV R2,=(R3)
MOV R1,=(R3)
NEG R1
RIS #C<CB,ORV>,=(R3) ;CLEAR ALL BUT UNIT #
RIS #CS,RD+CS,RO,0R3 ; AND START READ
TSTB 0R3
RPL 10
TST 0R3
RMT RIGERR
MOVB 0DS,ATT,0RPNB
CLRB 0RPNB
RIS PC

; DIVIDE ROUTINE FOR RP HANDLER.
; R5 = R3 / 0RP, RPHANDLER IN R4

DIVE CLR R5
CLR R4
TST R3
REG 43
COM R5
ROL R3

; QUOT. = 0
; RPM = 0
; IS DIVIDEND 0?
; YES = JUST RETURN
; QUOT. = -1 & SET CARRY
; NORMALIZE

IR3 = BLOCK #
IGET SECTOR NUMBER
IBY DIVIDING BY 10
ISAVE SECTOR
ISPT NPM DIVIDEND
IAND COMPUTE CYL & TRACK
IBY DIVIDING BY 20
IPOSITION TRACK IN HIGH BYTE
IAND INSTALL SECTOR
IR3 -> DISK ADRS RPB
ISPT TRACK & SECTOR
IAND CYLINDER
IAND BUS ADDRESS
IAND WORD COUNT
IMAKE NEGATIVE
ICLEAR ALL BUT UNIT #
IAND START READ
IWAIT UNTIL TRANSFER COMPLETE

; ANY ERRORS?
IYPS
ICLEAR UNIT ATTN FOR BOTH
I OLD & ECOD CONTROLLERS
IELSP JUST RETURN

; QUOT. = 0
; RPM = 0
; IS DIVIDEND 0?
; YES = JUST RETURN
; QUOT. = -1 & SET CARRY
; NORMALIZE

```

```

49          18          /SHIFT & SUBTRACT
50          R4
51          R4,OR2
52          78
53          SRP,R4
54          R5
55          R3
56          R2
57          R5

```

```

ROOT V02R=01 RT-11 ROOTSTRAP RT-11 MACRO VM02=09 A-APR-75 11149104 PAGE 6+
ASFCY

```

```

58          49i      T8T      (R2)+
59          RTS      R2
60
61          ;ENDC

```

```

ROOT V02R=01 RT-11 ROOTSTRAP RT-11 MACRO VM02=09 A-APR-75 11149104 PAGE 7
ASFCY

```

```

1          ;IP DP SRMSYS|SRPSYS|SRCSYS
2
3          ;IPDP SRPSYS
4          ;SATTL ROOTSTRAP I/O DRVPR = RP11
5
6          / RP11 DISK HANDLER
7
8          SYSDEV RP,204
9          RPHC = 177460
10         RPHC = 177462
11         RPHA = 177464
12         RPHA = 177466
13         RPHB = 177470
14         RPHB = 177472
15
16         READI  MOV SRP0A,R3
17         MOV   R0,R5
18         SWAB  R5
19         MOV   R5,R4
20         CLRB  R5
21         MOV   R5,(R3)+
22         RLC   #17740,R4
23         MOV   R4,(R3)
24         T8T  -(R3)
25
26         ;DEVICE IS RP. IT VECTORS TO 204.
27         ;CONTROL & STATUS REGISTER
28         ;MARR COUNT
29         ;MEMORY ADDRESS
30         ;DISK ADDRESS
31         ;DISK ADDRESS EXTENSION
32         ;DATA BUFFER
33
34         ;POINT TO DISK ADDRESS
35         ;COPY BLOCK NUMBER
36         ;MULTIPLY BY 256 TO GET WORD # ON DISK
37         ;SAVE HIGH ORDER DISK ADDRESS
38         ;MAKE RA AN EVEN BLOCK NUMBER
39         ;PUT LOW ORDER ADDRESS IN CONTROLLER
40         ;ISOLATE HIGH ORDER ADDRESS
41         ;PUT IT IN CONTROLLER
42         ;RESET POINTER

```

```

26 .ENDC
27 .IFDP SRC8V8
28 .CONDITIONAL FOR RC (R864) DISK
29 .SATTL ROOTSTRAP T/O DRIVER = RC11
30
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32 DD
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```



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177440
177442
177444
177446
177450
177452
177454
177456
177460
177462
177464
177466
177450
177452
177454
177456
177460
177462
177464
177466
177450
177452
177454
177456
177460
177462
177464
177466
177450
177452
177454
177456
177460
177462
177464
177466

```

```

RCDA= 177440
RCNA= 177442
RCPR= 177444
RCCB= 177446
RCMC= 177450
RCCA= 177452
RCMN= 177454
RCOB= 177456
RCDA,R3
R0,R5
R5
R5
R5,R3
@12,R3

```

```

.SATTL ROOTSTRAP T/O DRIVER = RC11
)DPVICE IS RC. IT VECTORS TO 210
)LOOK AHEAD REGISTER
)DISK ADDRESS REGISTER
)DISK ERROR STATUS REGISTER
)DISK CONTROL & STATUS REGISTER
)WARD COUNT REGISTER
)CURRENT ADDRESS REGISTER
)MAINTENANCE REGISTER
)DATA BUFFER REGISTER
)PT TO DISK ADR REGISTER
)OPT BLOCK NUMBER
)CALCULATE DISK ADR FOR RCDA
)UNIT, TRACK # & SECTOR ADR)
)132 * 8=256)
)IND PROPER DISK ADR
)PT TO CURRENT ADR REG + 12
)/(INTERFACE TO COMMON CODE)
)
)ENDC
)IFDP
)SATTL ROOTSTRAP T/O DRIVER = RK05
) RK05 DISK HANDLER

```

```

BOOT V020-01 RT-11 ROOTSTRAP RT-11 MACRO VM02-09 R-APP-75 11149104 PAGE 7+
ROOTSTRAP I/O DRIVER = RC11

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```

```

RKDA
READ1
151
251
551
SYSDEV RK.220
= 177412
MOV @14,R3
RR R3
ADD @20,R3
SUB @14,R0
RPL 18
ADR R3,R0
MOV @RKDA,R3
R1C @1777,R3
R1S @0,(R3)
)DPVICE IS RK. IT VECTORS TO 220
)RK DISK ADDRESS
)PHYSICAL BLOCK TO RK DISK ADR.
)ENTER BLOCK # COMPUTATION
)CONVERT DISK ADDRESS
)RK HAS DISK ADDRESS
)POINT TO HARDWARE DISK ADR REGISTER
)LPAVE THE UNIT NUMBER
)PUT DISK ADDRESS INTO CONTROLLER

```

```

71
72
73
74 000424 010243
75 000426 010143
76 000430 005413
77 000432 012743
78 000436 105713
79 000440 100376
80 000442 003713
81 000444 100401
82 000446 000207
83
84

```

GG

```

.ENDC
I THIS CODE IS COMMON TO RK05,RC11 AND RP11 HANDLERS
IBUFFER ADD.
IWORD COUNT
I(NEGATIVE)
ISTART DISK READ
IWAIT UNTIL COMPLETE
IANY ERRORS?
IWARD HALT ON ERROR

```

```

MOV R2,=(R3)
RI,=(R4)
NEG (R3)
MOV #5,=(R3)
TSYB Y8
RPL (R3)
TSY (R3)
RMI R10ERR
RTS PC

```

```

3si

```

```

.ENDC

```

ROOT VM2R=01 RT=11 BOOTSTRAP RT=11 MACRO VMR2=00 8-APR-75 11149104 PAGE 8
 BOOTSTRAP I/O DRIVER = RC11

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```

```

.IF DF 80T8YS
.SATTI BOOTSTRAP I/O DRIVER = DFCTAPE
I DECTAPE BOOTSTRAP HANDLER
SYSDV DT,214
TCCH = 177342
TCNT = 177340
TCST = 177340
READI MOV #TCCH,R4
MOV #TCNT,R3
NTSRCH: MOV #0,R5
SUI #0003,0R4
RIY #100200,0R4
REG #0
RMT NTER
CMP #5,0R3
ALT NTSRCH
#3,0R4
DTPWRD: MOV #100200,0R4
48i REG #0
RMT NTER
CMP #0,0R3
RBT DTPWRD
ALT NTSRCH
MOV #2,=(R3)
R1
NEB R1
MOV RI,=(R3)

```

```

I DEVICE IS DT. IT VECTORS TO 214.
ICOMMAND REGISTER
IDATA REGISTER
ISTATUS REGISTER
IRA => COMMAND REG
IRS => DATA REG
ICOPY BLOCK NUMBER
ISEARCH FOR 2 EARLIER
IINVERSE,RNUM
IWAIT TILL BLOCK FOUND

```

```

IIS IT THE DESIRED BLOCK
INO,CONTINUE SEARCHING
ISEARCH FORWARD (RNUM)
IWAIT

```

```

IOPSRPD BLOCK
INO=SEARCH FORWARD
INO=SEARCH REVERSE
IBUFFER ADDRESS
IWORD COUNT

```



```

31      MOV     05,0R4
32      RIT     0100200,0R4
33      REG     014
34      RMT     RIGERR
35      CLR     0R4
36      RTS
37      TST     00YCSY
38      RPL     RIGERR
39      RIT     00000,0R4
40      RNE     0YPRD
41      RR
42
43      .ENDC

```

ROOT VM02=01 RT-11 ROOTSTRAP RT-11 MACRO VM02=09 A-APR-75 11109184 PAGE 9
 BOOTSTRAP I/O NATVR = RC11

```

1      .IF OP 30X86S
2      .SRTTL  ROOTSTRAP T/D DRIVER = FLOPPY
3
4      SYSDEV  DX,264
5
6      READI   00
7      ABL     00
8      ABL     00
9      ABL     R1
10     MOV     00,=(0P)
11     MOV     00,R3
12     CLR
13     RR
14     SUR     #23,,R3
15     INC
16     SUR     #26,,R4
17     RPL
18     CMP     #-14,,R4
19     POL
20     SUR     #26,,R3
21     RPL
22     ADD
23     RPT
24     MOV     (0P)+,R0
25     INC
26     TST
27     RGT
28     RTS
29     TST
30     REG
31     RPL

```

```

IREAD
IWAIT FOR COMPLETION
IRPAD ERROR
ISTOP NT
IWHAT KIND OF ERROR ?
INOT END ZONE
I REVERSE?
ITHEN GO SEARCH FORWARD
IELSP SEARCH REVERSE

```

```

.IF OP 30X86S
.SRTTL  ROOTSTRAP T/D DRIVER = FLOPPY
SYSDEV  DX,264

```



```

1 001186 011602
2 001130 002700
3 001134 010046
4 001136 012701
5 001142 104701
6 001146 002701
7 001152 006201
8 001154 004001
9 001156 002701
10 001162 002700
11 001166 004767
12 001172 012700
13 001176 012601
14 001200 012604
15 001202 102704
16 001206 010164
17 001212 002701
18 001216 002701
19 001222 010164
20 001226 000430
21 001230 000027
22 001234 101774
23 001236 012005
24 001240 000405
25 001242 000415
26 001244 012005
27 001246 001374
28 001250 011700 000054
29
30
31
32
33
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48

:STARTL READ MONITOR, LOOKUP HANDLERS
MONPND: MOV 00P,R2
        JRECALL LOAD LOCATION
        JSUMP RB OVER BOOT RECORDS
        JSAVE SWAP BLOCK POINTER
        J00 GLOBAL ARITHMETIC HERE
        JRI = MAXSYM-SYSIZE (BYTES)
        JADD AMOUNT OF EXTRA STUFF
        J(WORDS)
        J(YO SUBTRACT)
        JLPNSYM TO LOAD (WORDS)
        JPOINT TO BLOCK WITH XMON
        JREAD THE MONITOR INTO PLACE
        JPOINT TO LIST OF THINGS TO RELOCATE
        JRI = SWAP BLOCK NUMBER
        JRA =P XMON IN CORP
        JSUBTRACT LOCATION XMON WAS LINKED TO
        JRA = OIAB. SET UP SWAP BLOCK #

        JSET SWR BLOCK #
        JRELOCATE A POINTER IN THE ASBCT
        JDONE YET ?
        JND
        JGET POINTER TO THING IN MONITOR
        JBIAS THE POINTER
        JNDM RELOCATE THE WORD
        JGET NEXT POINTER
        JPOINT TO MONITOR

        JCODE FOR RK
        JGET THE RK UNIT NUMBER

        JEXTRACT IT
        JDP SRKSYS
        JCODE FOR RJS03/4
        JUNIT # INTO R1
        JSTRIP TO 3 BITS

        JRP11
        JGET CONTROLLER STATUS REG INTO R1
        JSTRIP TO UNIT NUMBER
        JUNIT # INTO BITS 2-0
    
```

```

49          .ENDC          IDP SDPSYS
50
51          .IF NP          IFLOPPY
52          MOV            RTUNIT,R1          IGPT BOOTED UNIT (STORED BY ROOT2)
53          .ENDC          IDP SDPSYS
54
55          ADD            R1,DKAS8G(R0)      IFIX PERMANENT PSEUDO-ASSIGNMENTS
56          ADD            R1,SVAS8G(R0)
57          MOVB           R1,SVUNIT+1(R0)    ISPT UNIT NUMBER WP ROOTED

```

ROOT V02R=01 RT=11 ROOTSTRAP RT=11 MACRO VM02=00 A=APR=75 11149104 PAGE 11+

```

58          .ENABL          .ENDC          IDP SRKSYS:SDXSYS:IDPSYS
59          LBR
60
61          62 001254 054760 000322 000300          RCNFG,CONFTR(R0) ISET HARDWARE CONFIGURATION
62          63 001262 004003          R3          ICMUNT DEVICE SLOTS
63          64 001264 012701 00000000          #SENTRY,R1          IPOINT TO SENTRY TABLE IN RMON
64          65 001270 060401          R4,R1
65          66 001272 009721          (R1)+          IRESIDENT DEVICE ?
66          67 001274 001414          R0,=(R1)          INC, SKIP IT
67          68 001276 069441          IYF0, FIX HANDLER POINTER
68          69 001300 024127 00000000 #PNAM0(R1),SSYNAM IIS THIS THE SYSTEM DEVICE?
69          70 001306 001006          INN
70          71 001310 013360 00000000          R3,SYIND0(R0) ISPT SYSTEM INDEX NUMBER
71          72 001314 060360 00000000          R3,SYIND0(R0) I(DOUBLED)
72          73 001320 011100 00000000          #R1,SYENT0(R0) IAND SET UP SYSTEM ENTRY POINTER
73          74 001324 009721          (R1)+          IANY MORE ?
74          75 001326 009203          R3
75          76 001330 022703 00000000          #SLOT,R3
76          77 001334 001356          48
77          78 001336 062700 000010          #BYBITO,R0          IAND IN OFFSET TO SYSTEM VECTOR IN MAP
78          79 001342 152760 000014 00000000          #BYBITS,MAPOFF(R0) IAND PROTECT IT
79          80 001350 012701 00000000          #SPNAME,R1          IPOINT TO PERM NAMP TABLE
80          81 001354 060401          R4,R1
81          82 001356 062704 00000000          #SVREC,R4          IPOINT R4 TO SVREC IN RMON
82          83 001362 012703 00000000          #SLOT,R3          INUMBER TO LOOK UP
83          84 001366 012167 000222          (R1)+,PNAMP          IPTLL IN NAME IN LOOKUP
84          85 001372          .LOOKUP #,SBLOCK          ILOOKUP SYTHM.SYS
85          86 001400 102002          CLR          IGO IF THERE
86          87 001402 009024          (R4)+          ICLEAR RECORD NUMBER
87          88 001404 000406          RR          ISAVE STATUS OF THING
88          89 001406          .SAVEST 0,SBLOCK          ISPT STARTING RECORD
89          90 001414 016714 000204          MOV          CBLCK+2,RR4          IFIX IT
90          91 001420 009224          TNC          (R4)+
91          92 001422 004303          R5          DEC

```

```

93 001424 001360 RNE 68
94
95 001426 012737 100000 000044 MOV #10000,00JSH
96 001434 000000 .PRINT #STRNG
97 001442 000000 CLR R0
98 001444 012720 MOV (PC)+,(R0)+
99 001446 040000 BIC R0,R0
100 001450 012720 MOV (PC)+,(R0)+
101 001452 .EXIT
102 001454 032767 000000 000120 RIT #KWILL,RCNFG
103 001462 001403 REG 103
104 001464 012737 000100 177546 MOV #100,0LKCS
105 001472 000000 CLR R0
106 001474 1001 .EXIT
107 .DSABL L0R
108
109 001476 000011 RCLR: CLR #R1
110 001500 000002 RTI RTI

```

```

;TRAP MEANS THIS CONFIGURATION NYET
;UNTRAP

```

BOOT V020-01 RT-11 BOOTSTRAP RT-11 MACRO V002-09 8-APR-75 11009104 PAGE 12

RELOCATION LIST

```

1 001502 000004
2 001504 000010
3 001506 000030
4 001508 000054
5 001510 000060
6 001512 000064
7 001514 000100
8 001516 000210
9 001520 000244
10 001522 000000
11 001524 000000
12 001526 000000
13 001530 000000
14 001532 000000
15 001534 000020
16 001536 000060
17 001540 000100
18 001542 000000
19 001544 000000
20 001546 000000
21 001550 000060
22 001552 000020
23 001554 000000
24 001556 000000
25 001560 000000

;SYTL RELOCATION LIST
;R0: 4
;R1: 10
;R2: 30
;R3: 54
;R4: 60
;R5: 64
;R6: 100
;R7: SYVEC
;R8: 204
;R9: USRLOC
;R10: SUSRLOC
;R11: RCOMP
;R12: SKALOC
;R13: TTIBUF
;R14: TTIBUF+2
;R15: TTIBUF+6
;R16: TTIBUF+10
;R17: TTIBUF
;R18: TTIBUF+4
;R19: TTIBUF+6
;R20: SYSL0W
;R21: COMPTR+2
;R22: SINPTR
;R23: SYNC

;ILLEGAL MEM AND INST. TRAPS
;EMT
;ADDRESS OF RMON
;TTY VECTORS
;CLOCK VECTOR
;SYSTEM DEVICE VECTOR
;LOCATION OF FPU TRAP
;LOCATION OF USR NOW
;ADDRESS OF 'NORMAL' USR
;ADDRESS OF RMON
;TTY RING BUFFER--INPUT
;TTY RING BUFFER--OUTPUT
;LOWEST USFD LOCATION
;FREE CORE LIST
;POINTER TO SINTEN IN RESIDENT HANDLER
;SYNCHRONIZATION ADDRESS

```


A.3 LP/LS11 DEVICE HANDLER

LP V02-03 25-JUN-74 07-11 MACRO VM02-10 14-APR-75 10:05:11 PAGE 1

```

1  .TITLE LP V02-03 25-JUN-74
2
3  / DT-11 LINE PRINTER (LP/LS11) HANDLER
4
5  / DEC-11-ORTIA-D
6
7  / PGR/FP/ARC/EF
8
9  / MARCH 1973/FEBRUARY 1974
10
11 / COPYRIGHT (C) 1970,1975
12
13 / DIGITAL EQUIPMENT CORPORATION
14 / MAYNARD, MASSACHUSETTS 01754
15
16 / THIS SOFTWARE IS FURNISHED UNDER A LICENSE FOR USE ONLY
17 / ON A SINGLE COMPUTER SYSTEM AND MAY BE COPIED ONLY WITH
18 / THE INCLUSION OF THE ABOVE COPYRIGHT NOTICE. THIS SOFTWARE,
19 / OR ANY OTHER COPIES THEREOF, MAY NOT BE PROVIDED OR OTHERWISE MADE
20 / AVAILABLE TO ANY OTHER PERSON EXCEPT FOR USE ON SUCH SYSTEM AND TO
21 / ONE WHO AGREES TO THESE LICENSE TERMS. TITLE TO AND OWNERSHIP OF THE
22 / SOFTWARE SHALL AT ALL TIMES REMAIN IN DIGITAL.
23
24 / THE INFORMATION IN THIS DOCUMENT IS SUBJECT TO
25 / CHANGE WITHOUT NOTICE AND SHOULD NOT BE CONSTRUED
26 / AS A COMMITMENT BY DIGITAL EQUIPMENT CORPORATION.
27
28 / DIGITAL ASSUMES NO RESPONSIBILITY FOR THE USE
29 / OR RELIABILITY OF ITS SOFTWARE OR EQUIPMENT
30 / WHICH IS NOT SUPPLIED BY DIGITAL.
31

```


49 001144 104327
 50 001146 000001
 51 001150 001423
 52 001152 110537 177516
 53 001156 000742
 54 001160 000207
 55 001162 120527 000011
 56 001166 001420
 57 001170 120527 000012

TARCNT: .WARR
 PC1: MOVR
 RET: RTA
 PRDTST: CMDB
 REO: TARSFT
 CMDB: PS,#LF

ASLH (PF)+
 .WARR
 PS,#CR
 PS,#LPR
 LPNEXT
 PC
 PS,#HT
 TARSFT
 PS,#LF

UPDATE TAB COUNT
 PRESFT TAB
 IPRINT THE CHAR
 ITRY FOR NEXT CHAR
 IIS CHAR A TAB?
 IYES-PRESFT TAB
 IIS IT LF?

LP 002=03 25-JUN-74 0T-11 MACRO VM02-10 14-APR-75 10:05:11 PAGE 4+

58 001174 001406
 59 001176 120527 000015
 60 001202 000240
 61 001204 120527 000014
 62 001210 001333
 63 001212 012747 000204
 64 001220 012767 000001
 65 001226 000751
 66 001230 014747 177712
 67 001236 012705 000000
 68 001242 000735
 69

REO RSTC
 CMDB PS,#CR
 NOP
 CMDB PS,#FF
 RNF IGNORE
 MOV #COLSIZ,COLCNT
 MOV #I,TABCNT
 RR PC1
 MOV TARCNT,TARFLG
 MOV #40,R5
 RR PCHAR

70 001244 000240
 71 001246 022424
 72 001250 012705 000014
 73 001254 000754
 74
 75 001256 052754 000001
 76
 77

IYES-PRESFTOPE COLUMN COUNT
 IIS IT CR?
 IIGNORF UNLESS MODIFIED
 IIS IT A FF?
 INC-CHAR IS NON-PRINTING
 IRE-INITIALIZE COLUMN COUNTER
 I PRESFT TAB COUNTER
 IPRINT THE CHAR
 I SFT UP TAB
 IPRINT SPACES

78 001262 005037 177514
 79 001266 010704
 80 001270 062704 177520
 81 001274 012705 000054
 82 001300 000175 000270
 83 001304 000000
 84 001306 000306
 85 000001

INP
 (R4)+,(R4)+
 #FF,R5
 R5TC
 #MDEPR,0-(04)
 #I,PS
 PC,R4
 #LPCRE-,R4
 #MONLOW,R4
 #OFFSET,PS1
 #

86 001304 000000
 87 000306
 88 000001

I MAKE SURE WF ONLY COMP HERE ONCF
 I POINT TO ADPS OF NEXT CHAR
 I PRINT INITIAL FF
 I SFT HARD ERROR BIT
 I TURN OFF INTERRUPT
 I ADD OF COE IN R3
 I JUMP TO O MANAGER

INTEN: 0
 IPRITE = .LOADPT
 .END

A.4 CR11 DEVICE HANDLER

CR.SYS RT-11 MACRO VM02-10 28-APR-75 16:00:38 PAGE 1

```

1 .TITLE CR.SYS
2 RT-11 CARD READER (CP11) HANDLER
3
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1 .TITLE CR.SYS
2 RT-11 CARD READER (CP11) HANDLER
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1 .TITLE CR.SYS
2 RT-11 CARD READER (CP11) HANDLER
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```

CR.SYS RT-11 MACRO VM02-10 28-APR-75 16:00:38 PAGE 2

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000000
010001
020002
030003
040004
050005
060006
070007

```

```

1 CARD READER CONTROL

```



```

13 000270      INTERRUPT VECTOR
14 177160      ICARD READER STATUS REGISTER
15 177162      IDATA BUFFER 1
16 177164      IDATA BUFFER 2
17
18
19
20 000001      I CONSTANTS FOR MONITOR COMMUNICATIONS
21 000054      IHAND FERR00
22 000270      IBASE ADDRESS OF MONITOR
23 177776      IOFFSET TO HANDLER RETURN
24 000300      IPROGRAM STATUS WORD
25 000300      IPRIVTY 7
26 000300      IPRIVTY 4
27
28
29
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41
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44
45
00015      I ASCII CONSTANTS
00012      ICARriage RETURN
00000      ILYNE FEED
00000      ISPACE
00001      IEND OF FILE
00001
00001      I CARD READER CONTROL AND STATUS BITS
00001      IREAD
00002      IJECT02
000100      INTERRUPT ENARLF
000200      ICOLLIM NONE
00000      IREADY00
000100      IBUSY
00000      ICNIN000
00000      IONLYNF
00000      IATLATE000
00000      IOTIN01000
00000      IOPCK02000
00000      ICARDNONE
00000      IERR00

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44
    .SRTTL      CONFIGURATION SECTION
    ; THE FOLLOWING CODE IS EXECUTED WHEN A "SET CP" CONSOLE COMMAND IS
    ; GIVEN.
    .BARR      ;CONFIGURATION AREA
    ; SET CR (NO) CRIF
    ; APPEND/DO NOT APPEND CARRIAGE RETURN/LINE FEED TO EACH CARD IMAGE
    BR          .+LXCRLF-YCRIF
    .RAD50     /CRLF/
    .WARR      <CRLF-4000>/2+100000
    ; SET CR (NO) TRIM
    ; TRIM/DO NOT TRIM TRAILING BLANKS FROM CARD IMAGES
    BR          .+LXTRIM-YTRIM
    .RAD50     /TRIM/
    .WARR      <TRIM-4000>/2+100000
    ; SET CR (NO) HANG
    ; HANG/RETURN HARD ERROR IF READER NOT READY AT START OF TRANSFER
    BR          .+LEPRRD-YHANG
    .RAD50     /HANG/
    .WARR      <HANG-4000>/2+100000
    ; SET CR CODE IS 024 (029) MODE
    ; SET TRANSLATION TO 026 (029) MODE
    .WARR      024.
    .RAD50     /CODE/
    .WARR      <CODE-4000>/2+00000
    ; SET CR (NO) IMAGE
    ; TRANSMIT EACH COLUMN AS 12 BITS (ONE WORD/COLUMN)
    .WARR      NOTMAG-IMAGEF
    .RAD50     /IMAGE/
    .WARR      <IMAGE-4000>/2+100000
    .WARR      0      ;END-OF-OPTIONS FLAG
  
```

```

1  .SRTTL  CONFIGURATION SUBROUTINES
2
3  000452 012703      INOP IF POSITIVE
4  000454 000200      MOV
5  000456 010347      MOV
6  000458 000206      MOV
7  000462 000207      PLS
8  000464 012703      TRM:
9  000466 000200      MOV
10 000470 010347      MOV
11 000474 000207      PLS
12
13 000476 012703      HANG:
14 000500 000200      MOV
15 000502 010347      MOV
16 000506 000207      PLS
17
18 000510 010701      CODE:
19 000512 062701      ADD
20 000516 160300      SUB
21 000520 100407      MUL
22 000522 001407      DIV
23 000524 062701      ADD
24 000530 022700      CMP
25 000534 001402      JEQ
26 000536 000201      SET
27 000540 000207      CONEVT: PLS
28 000542 010703      SETCND: MOV
29 000544 062703      SCDF: CLR
30 000550 000000      BISH
31 000552 152100      MOV
32 000554 001771      MOV
33 000556 060300      ADD
34 000560 112110      MOV
35 000562 000772      RR
36
37 000564 062703      THAGF: ADD
38 000570 060703      ADD
39 000572 012367      TRMASE: MOV
40 000576 010347      MOV
41 000602 012367      MOV
42 000606 012367      MOV
43 000612 000207      PLS
44
45 000614 001402      NOTMAG: REN
46 000616 114537      CHRTRL: XTM2(P5),0(PC)+
47 000622 000001      .MORN
48 000624 000402      YATMAG: RR
49 000626 012737      MOV
50 000632 000002      .MORN

```

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:BATTL 024, 029 CONVERSION TABLES

I 024 CONVERSION TABLE
I MODIFIS CHARACTER TABLE TO ACCEPT 026 KEYPUNCH CODES

SET026: .BYTE 012,137          TRACK ARROW      (0-2)
         .BYTE 013,075          EQUAL             (A-3)
         .BYTE 015,134          UP ARROW         (0-5)
         .BYTE 016,007          APOSTROPH       (0-6)
         .BYTE 017,134          BACKSLASH       (0-7)
         .BYTE 052,073          SFTCOLN         (0-8-2)
         .BYTE 054,050          LEFT PAREN      (0-8-4)
         .BYTE 055,042          QUOTES          (0-8-5)
         .BYTE 054,043          FLR. STGN      (0-8-6)
         .BYTE 057,045          PERCENT         (0-8-7)
         .BYTE 112,072          COLN            (11-A-2)
         .BYTE 115,133          L BRACKET       (11-A-5)
         .BYTE 114,074          GREATER THAN   (11-A-6)
         .BYTE 117,044          AMPERSAND      (11-A-7)
         .BYTE 200,053          PLUS           (12)
         .BYTE 212,077          HESTON         (12-A-2)
         .BYTE 214,051          RIGHT PAREN    (12-A-4)
         .BYTE 215,135          R BRACKET      (12-A-5)
         .BYTE 214,074          LESS THAN      (12-A-6)
         .WORD 0              END OF TABLE **
  
```

```

I 029 CONVERSION TABLE
I MODIFIS CHARACTER TABLE TO ACCEPT 029 KEYPUNCH CODES

SET029: .BYTE 012,072          COLN            (A-2)
         .BYTE 013,043          FLR./STGN      (A-3)
         .BYTE 015,007          APOSTROPH     (A-5)
         .BYTE 014,075          EQUAL          (A-6)
         .BYTE 017,042          QUOTES        (A-7)
         .BYTE 052,134          BACKSLASH     (0-A-2)
         .BYTE 054,045          PERCENT       (0-8-4)
         .BYTE 055,137          TRACK ARROW   (0-8-5)
         .BYTE 054,074          GREATER THAN (0-A-6)
         .BYTE 057,077          HESTON        (0-8-7)
         .BYTE 112,134          R BRACKET     (11-A-2)
         .BYTE 115,051          RIGHT PAREN   (11-A-5)
         .BYTE 114,073          SFTCOLN      (11-A-6)
         .BYTE 117,134          UP ARROW      (11-A-7)
         .BYTE 200,044          AMPERSAND    (12)
         .BYTE 212,133          L BRACKET     (12-A-2)
         .BYTE 214,074          LESS THAN    (12-A-4)
         .BYTE 215,050          LEFT PAREN   (12-A-5)
         .BYTE 214,053          PLUS         (12-A-6)
         .WORD 0              END OF TABLE **
  
```



```

40 000150 000000
50 000152 062785
51 000154 114537
52 000160 000000
53 000162 062787
54 000170 052737
55 000176 000207
56
57

```

FNDPTR: .WARR 0
 TNFOLT: ADD PC,R5
 YIM2: MOV CHPTRL=,(R5),0(PC)+
 CHPTRL: .WARR 0
 YIM3: ADD #1,CHPTRL
 TNTRPT: BIZ #INTFB,0PCOST
 BIZ PC

MAKE A PIC POINTER TO TRANS TABLE
 INPUT TRANSLATED CHAR IN LINE
 PUSH BUFFER POINTER BY 1 OR 2
 REARLF ONE LINE INTO IF NECESSARY

CR.SYS 01-11 MACRO VM02-1M 20-APR-75 16:00:33 PAGE 44
 HANDLER PROPR

```

58 000200 032785
59 000204 001400
60 000206 012704
61 000210 000000
62 000212 014705
63 000216 025225
64 000220 023727
65 000224 000000
66 000226 001471
67 000230 010106
68 000232 014701
69 000236 000200
70 000240 014701
71 000244 000201
72 000246
73 000250 000240
74 000254 112721
75 000258 010107
76 000262 010107
77 000266 000245
78 000270 024704
79 000274 001371
80 000278 032737
81 000282 001325
82 000286 010106
83 000290 062704
84 000294
85 000298
86 000302
87 000306
88 000310
89 000314
90 000318
91 000322
92 000326
93 000330
94 000334
95 000338
96 000342
97 000346
98 000350
99 000354

```

I CARD NAME OR ERROR
 CARD: BIT #CARDN,R5
 MOV FRR1
 MOV (PC)+,R4
 RUPPTR: .WARR 0
 MOV RRCOP,R5
 CMP (R5)+,(R5)+
 CMP #7817,(PC)+
 CHAR12: .WARR 0
 REC FNDFL
 MOV R1,(SP)
 MOV CHPTR,R1
 YTRIM: NOP
 MOV FNDPTR,R1
 TNC R1
 I XTRIM:
 YCOLF: NOP
 MOVR #CR,(R1)+
 MOVR #LF,(R1)+
 I XPRIF: MOV R1,ENDPTD
 MOV (SP)+,R1
 RR CONT
 FI RUIF: MOVR (R4)+,0(R5)+
 REC ORS
 REC DETHAN
 TNC -(RS)
 CMP FNDPTR,R4
 RPT FI RUIF
 BIT #READY+RUIV,0CRRT
 RNF TNTRPT
 MOV PC,R4
 JDN #CHRRUF=..04

ICARD DONE?
 INDF -- SPURIOUS INTERRUPT - IGNORE
 IRA -- CHRRUF
 IPRINTR TO CHRRUF
 IPRINT TO ELEMENT
 IPRINT OVER BUFFER POINTER
 IEND OF FILE?
 I12-BIT FOR FIRST CARD COL.
 IYES
 ISAVE R1
 IPRINT TO PART END OF AN CARD
 I PATCH HERE TO SUPPRESS TRIMMING
 IN4 GIVE IT TO HIM
 I PATCH HERE TO SUPPRESS CR/LF
 IJA, GIVE IT HIM
 ITHIS IS NOW THE END
 IRESTORE R1
 IENTER FILLING LOOP
 IPUT A BYTE IN MYS BUFFER
 IIS WE FULL ?
 IYFF
 IPRINT BUFFER POINTER
 IEND OF OUR CARD ?
 IINT YET
 IREADY+RUIV,0CRRT OKAY TO INT READ?
 INDF - GO HAND UNTIL READY
 IPRINT TO CHRRUF

```

91 000324 010447 177630          ;START FTLING FROM NEW CARD
92 000330 010447 177614          ;WHICH IS AS YET EMPTY
93 000334 010447 177650          ;ESTABLISH BUFFER POINTERS
94 000340 005047 177640          ;CLEAR END FLAG
95 000344 012727 000120          ;SET COLUMN POINT
96 000350 000000 000000          ;COUNT OF COLUMNS REMAINING IN CARD
97 000352 012737 000101 177140 ;PREPARED, PREPARED START A CARD GOING
98 000354 000207 000000          ;SAVE
99
100
101
102 000362 014705 177422          ;POINT TO BUFFER ELEMENT
103 000366 052755 000001          ;YOU CAN'T WRITE ON A READER
104 000372 000044 000000          ;
105
106 000374 032705 000000          ;DATA RATE IS ONLY NOT CURABLE
107 000400 001370 000000          ;ISSUE HARD FROM YF SN
108 000402 000747 177742          ;NONE WITH DATA COLUMNS?
109 000406 1000337 000000          ;INPF -- MUST BE PCK CMFCK, ETC.
110
111 000410 000073 000000          ; START A NEW READ TO CORRECT CONDITION.
112
113
114
; END OF FILE CARD FOUND

```

CR.SYS RT-11 MACRO VM02-10 28-APR-75 16:00:34 PAGE 4*

```

115 000412 012504          ;NOFFL: MOV (RS)+,R4
116 000414 105020          ;LDR:FL: CLRB (R4)+
117 000416 005315          ;DEC BR
118 000420 001375          ;RNF CLEARIF
119 000422 052775 020000 177770 ;RIS #2000,0-10(R5) ;SET END BYT IN CHANNEL
120 000430 005047 177514          ;ARRT: CLR FNPTR
121
122
123
124 000434 010447 177520          ; RETURN TO MONITOR (REQUEST DONE, END, OR ERROR)
125 000440 005037 177140          ;PETHAN: MOV P4,CHRPD
126 000444 010704 000000          ;CLB #PRST
127 000446 062704 177342          ;MOV PC,R4
128 000452 013705 000050          ;ADD #COCOE-1,R4
129 000456 000175 000270          ;MOV #MONLOW,R5
130

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1  .SRTTL CHARACTER TABLE
2
3  ! THE FOLLOWING MACRO TAKES AS ARGUMENTS THE ASCII TRANSLATION
4  ! DESIRED AND THE LIST OF PUNCH COMBINATIONS FOR THAT CHARACTER.
5
6  .MACRO P $LIST
7      TER
8      X<SLIST>
9      .IF NE Y'
10     .IF LE Y'-7
11     TBT+Y'.
12
13     .IF
14     .IF 10
15     X'-A.
16     .IF 11
17     TBT+11
18
19     .FNDF
20     .IF
21     .FNDF
22     .ENDR
23
24     .SCHARL+T
25     .RYTE SCHAR
26     SCHAR = SCHAR + 1
27     P
28
29 ! THE FOLLOWING TABLE TRANSLATES Q29 KEYPUNCH CODES TO ASCII.
30
31 .SCHARL:
32 .REPT 254.
33 .RYTE 130
34 .ENDR
35
36 .SCHARL
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34	000000	START AT OCTAL 000
37	000442	INITI
38	000754	ICTRI-A
39	000764	ICTRI-B
40	000785	ICTRI-C
41	000786	ICTRI-D
42	000512	ICTRI-E
43	000540	ICTRI-F
44	000541	ICTRI-G
45	000542	ICTRI-H
46	000611	ICTRI-I
47	000710	ICTRI-J
48	000540	ICTRI-K
49	000716	ICTRI-L
50	000717	ICTRI-M
51	000720	ICTRI-N
52	000721	ICTRI-O
53	000722	ICTRI-P
54	001014	ICTRI-Q
55	000604	ICTRI-R
56	000605	ICTRI-S

PR. SVS	RT-11	MACRO	VM02-10	28-APR-75	16:00:3A	PAGE	7+
58	000604	ICTRI-T					
59	000517	ICTRI-U					
60	000520	ICTRI-V					
61	000505	ICTRI-W					
62	000541	ICTRI-X					
63	000613	ICTRI-Y					
64	000614	ICTRI-Z					
65	000522	ICTRI-7					
66	000542	ICTRI-1					
67	000617	ICTRI-1					
68	000620	ICTRI-1					
69	000621	ICTRI-1					
70	000622	ICTRI-1					
71	000443	ICTRI-1					
72	000702	ICTRI-1					
73	000502	ICTRI-1					
74	000476	ICTRI-1					
75	000576	ICTRI-1					
76	000537	ICTRI-1					
77	000643	ICTRI-1					
78	000500	ICTRI-1					
79	000700	ICTRI-1					
80	000600	ICTRI-1					

81 000577	C	<12,R,4>	I*
82 000701	C	<0,R,3>	I.
83 000536	C	<11>	I-
84 000563	C	<12,R,3>	I.
85 000676	C	<0,1>	I/
86 000524	C	<0>	I0
87 000523	C	<1>	I1
88 000444	C	<2>	I2
89 000445	C	<3>	I3
90 000446	C	<4>	I4
91 000447	C	<5>	I5
92 000470	C	<6>	I6
93 000471	C	<7>	I7
94 000472	C	<8>	I8
95 000473	C	<9>	I9
96 000503	C	<0,2>	I1
97 000475	C	<11,R,4>	I1
98 000601	C	<12,R,0>	I4
99 000677	C	<0,6>	I8
100 000501	C	<0,R,6>	I2
101 000541	C	<0,R,7>	I2
102 000542	C	<0,4>	I0
103 000477	C	<12,1>	I4
104 000644	C	<12,2>	I8
105 000645	C	<12,3>	I0
106 000646	C	<12,4>	I0
107 000647	C	<12,5>	I0
108 000670	C	<12,4>	I0
109 000671	C	<12,3>	I0
110 000672	C	<12,4>	I0
111 000673	C	<12,0>	I1
112 000703	C	<11,1>	I1
113 000544	C	<11,2>	I8
114 000545	C	<11,3>	I1

CR.SYS PT-11 MACRO VM02-01M PA-APR-75 14:00:38 PAGE 74
CHARACTER TABLE

115 000546	C	<11,0>	I4
116 000547	C	<11,5>	I4
117 000570	C	<11,4>	I0
118 000571	C	<11,7>	I0
119 000572	C	<11,8>	I0
120 000573	C	<11,0>	I8
121 000603	C	<0,2>	I5
122 000525	C	<0,3>	I7

123	000526	L	<R,0>	III
124	000527	L	<R,5>	IV
125	000528	L	<R,4>	IX
126	000529	L	<R,7>	IX
127	000530	L	<R,8>	IX
128	000531	L	<R,9>	IX
129	000532	L	<R,0>	IX
130	000533	L	<R,1>	IX
131	000534	L	<R,2>	IX
132	000535	L	<R,3>	IX
133	000536	L	<R,4>	IX
134	000537	L	<R,5>	IX
135	000538	L	<R,6>	IX
136	000539	L	<R,7>	IX
137	000540	L	<R,8>	IX
138	000541	L	<R,9>	IX
139	000542	L	<R,0>	IX
140	000543	L	<R,1>	IX
141	000544	L	<R,2>	IX
142	000545	L	<R,3>	IX
143	000546	L	<R,4>	IX
144	000547	L	<R,5>	IX
145	000548	L	<R,6>	IX
146	000549	L	<R,7>	IX
147	000550	L	<R,8>	IX
148	000551	L	<R,9>	IX
149	000552	L	<R,0>	IX
150	000553	L	<R,1>	IX
151	000554	L	<R,2>	IX
152	000555	L	<R,3>	IX
153	000556	L	<R,4>	IX
154	000557	L	<R,5>	IX
155	000558	L	<R,6>	IX
156	000559	L	<R,7>	IX
157	000560	L	<R,8>	IX
158	000561	L	<R,9>	IX
159	000562	L	<R,0>	IX
160	000563	L	<R,1>	IX
161	000564	L	<R,2>	IX
162	000565	L	<R,3>	IX
163	000566	L	<R,4>	IX
164	000567	L	<R,5>	IX
165	000568	L	<R,6>	IX
166	000569	L	<R,7>	IX
167	000570	L	<R,8>	IX
168	000571	L	<R,9>	IX
169	000572	L	<R,0>	IX
170	000573	L	<R,1>	IX
171	000574	L	<R,2>	IX

PARCENT GRAVE

PLP A
PLP B
PLP C
PLP D
PLP E
PLP F
PLP G
PLP H
PLP I
PLP J
PLP K
PLP L
PLP M
PLP N
PLP O
PLP P
PLP Q
PLP R
PLP S
PLP T
PLP U
PLP V
PLP W
PLP X
PLP Y
PLP Z

INDEX SPACE
VERTICAL BAR
DIAGONAL BAR
TILDE
DOLLAR

... 8 CHARACTERS + 256.

CHARACTER: RIKEW 81.
SYMBOL: WORN 8

PLUGGED TO POINT TO COMMON ENTRY

172 001324 CRRITE=-LNDNPT
 173 000001 .END
 174

ABORT	0004300	002	RUEPTR	0002100	002	BUSY	0010000	002	CARD	0002000	002	CARDN	0000000	002
CHAR12	0002200	002	CH00IF	0010020	002	CH0PTR	0001000	002	CH0TBL	0004020	002	CH0UIF	0000100	002
CODE	0005100	002	CODEXT	0005000	002	COLCNT	0003500	002	COLD	0002000	002	CONT	0003000	002
CR	0000150	002	CR01	0011020	002	CR02	0011000	002	CR0F	0000100	002	CRHAND	0000120	002
CRINT	0000020	002	CRIF	0000050	002	CRIOF	0000000	002	CRITE	0013240	002	CRRT	0017100	002
CRVECT	0002100	002	CATLAT	0000000	002	FJ00	0000000	002	FANFTL	0000120	002	FNDPTR	0001000	002
FOF	0000001	002	FR0	0000000	002	FR00	0003700	002	FR01	0003000	002	FIL0IF	0002700	002
HANG	0000000	002	HDFR	0000000	002	H0PCK	0000000	002	HAGF	0005000	002	YMRK	0000000	002
INCOLT	0001000	002	INTFR	0001000	002	INTFT	0001700	002	IF	0000000	002	IF	0000000	002
LOADPT	0000000	002	IXTRIF	0002000	002	IXTRM	0002000	002	LEFRNR	0003000	002	MOTIN	0010000	002
NOYHAG	0000100	002	NYTCHR	0001000	002	NFSFT	0002100	002	MONLW	0000000	002	PC	0000000	002
PR0	0003000	002	PR7	0003000	002	PS	0017700	002	PR0	0000000	002	PC	0000000	002
READY	0000000	002	REYMON	0000000	002	RE	0000000	002	READ	0000000	002	READ	0003100	002
R3	0000000	002	R0	0000000	002	R5	0000000	002	R1	0000000	002	R2	0000000	002
SET026	0000000	002	SET029	0007000	002	SP	0000000	002	SCDF	0005000	002	SETCD	0000000	002
TRM	0000000	002	TRM00	0001000	002	TR	0000000	002	TRCF	0000000	002	TR	0000000	002
YIM1	0001000	002	YIM2	0001000	002	YIM3	0001000	002	YOLF	0002000	002	YHANG	0000000	002
YIM2	0001000	002	YIM3	0001000	002	YIM3	0001000	002	YTRM	0002000	002	YATHAG	0000000	002
YIM3	0001000	002	YIM3	0001000	002	YIM3	0001000	002	YTRM	0002000	002	YATHAG	0000000	002

.LPT:/N/STH/C/SC0

A.5 TC11 DEVICE HANDLER

NT VR2-87 12-APR-74 RT-11 MAPRO VM82-18 28-APR-74 16188124 PAGE 1

1 .TITLE NT VR2-87 12-APR-74
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1 000000
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3 000000
4 000000 000214
5 000002 000072
6 000004 000340
7 000006
8 000006 000000
9 000010 000000
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11 000012 012727 000010
12 000016 000000
13 000020 017746 177776
14 000024 000747 000004
15 000030 000207
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20 000032 000510
21 000034 000577 000262
22 000040 000000
23 000042 010006
24 000044 014700 177740
25 000048 013704 177302
26 000054 012006
27 000056 012005 170377
28 000060 003705 100100
29 000064 032714
30 000070 100452
31 000072 001502
32
33 000074 032714 000002
34 000100 001043
35 000102 023767 177350 000206
36 000110 001422
37 000112 003407
38 000114 052705 000000
39 000120 162716 000002
40
41
42 000124 000002
43 000126 052705 010000
44
45 000132 052705 000103
46 000136 012667 000150
47 000142 010514
48 000144 012600

```

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REGS
I LOAD POINT
.WARN TCVEC
.WARN INTNT-.
.WARN PR7
NTSYS
NTIOFS .WARN 0
NTIOF: .WARN 0
I ENTRY POINT
MOV #0.,(PC)+
NTTRY: .WARN 0
MOV #005.-(SP)
.TSP PC.DTINT
PTA PC
INTERRUPT SERVICE
RR NTSTOP
TSP #5.0SINSTR
.WARN ZC<300>BPR7
MOV #0.-(SP)
MOV NTIOF,RR
MOV STIOF,RA
MOV (RR)+,-(SP)
MOV (RR)+,RS
R1C #7<300>,05
R1T #100100,(RA)
R1T #TFRP
REC RETRY
R1T #2.(R4)
REC #TNAME
PHD #ATCNT,AWANT
REC #LKEND
DIRECT: #LF
REVERSF: #RS
CLR #2.(R4)
RR FORWARD
R1C #10000,RS
FORWARD: #RS
FORWARD: #RS
MOV #M3,RS
RETN: #M3
MOV #5.(R4)
MOV (SP)+,RR
MOV (SP)+,RR

```

```

ADDRESS OF INTERRUPT VECTOR
INFER TO INTERRUPT SERVICE
PRIORITY 7
I POINTER TO LAST 0 ENTRY
I POINTER TO CURRENT 0 ENTRY
INIT THE RETRY COUNT
IRETRY COUNTER
IFAKE AN INTERRUPT
IBACK TO MONITOR
IABORT CALL FROM AF SYSTEM
IABORT J9R TO COMMON CODE
IABORT AT LEVEL 4
IR0 POINTS TO 0 ELEMENT
IR4 POINTS TO CONTROL REGISTER
IDESTROY BLOCK 0 ONTO STACK
IUNIT 0 INTO RR
IISOLATE UNIT NUMBER
IFERRR RTT ON?
IYES
IF INTERRUPT IS OFF, WE ARE INITIATING
IA REQUEST
ISARCHING0
IN0-A READ OR WRITE JUST COMPLETED
ICOMPARE ACTUAL BLOCK TO DESTROYED BLOCK
IFOUND IT
ISFARCH IN THE FORWARD DIRECTION
ISFT REVERSE RTT
ISFARCH FOR TWO BLOCKS BEFORE ONE
IPARTIALLY DESIRED (TO ALLOW
ISPACE FOR THE TURN-AROUND
IDONT SFT DELAY INHIBIT
IDONE IS ALREADY MOVING FORWARD
ISD INHIBIT HARDWARE DELAY
IFREMPERUP ENARLF,ONUM,AND GO
ITFLI CONTROLLER TO CO
IFRSTORE RR

```

49 000146 000207
 50
 51 000150 032714 0001000
 52 000154 001757 PFLVDFSE
 53 000156 032714 0000000
 54 000162 001343
 55
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TRACK INTO MONITOR

WERE WE IN REVERSE?
 INC-REVERSE TAPE
 WERE WE GOING FORWARD?
 INC-WE HAVE TO TURN AROUND

INITIATE READ/WRITE REQUEST

DT V02-07 12-APR-74 07-11 WAFHO VM02-10 2A-APR-75 16100134 PAGE 3*

58 000144 052705 010115
 59 000170 012037 177306
 60 000174 010114
 61 000176 100400
 62 000200 001423
 63 000202 000416
 64 000204 022705
 65 000210 012037 177304
 66 000214 000752
 67
 68
 69 000216 032737 104000
 70 000224 100003
 71 000226 032714 000002
 72 000232 001546
 73 000234 000347 177556
 74 000240 001417
 75 000242 052770 000001 177772
 76
 77

ASSUME WRITE
 CORP ADDRESS
 WORD COUNT (OVER BLOCK #)
 WRITE WAS A GOOD GUESS
 IF ZERO, SEEK
 REAP-NEGATE WORD COUNT
 SET REAP FUNCTION
 SET WORD COUNT

END7 ERROR?
 INT END7
 WERE WE SEARCHING?
 YES-REVERSE TAPE
 MORE TAPER LEFT?
 YES
 INC-SET HAD ERROR BIT

OPERATION FINISHED

POP BLOCK
 RESTORE RM
 STOP SELECTED DRIVE
 JUMP OF ONE IN 04

78 000250 000724
 79 000252 012400
 80 000254 112737
 81 000262 010700
 82 000264 062700
 83 000270 013705
 84 000274 000175
 85
 86
 87
 88
 89 000300 100737 177300
 90 000304 100710
 91 000306 062747 000000 000002
 92 000314 022716

TAPE UP TO SPEED?
 YES-AVOID STOPPING TAPE
 INT-T TAKES 4 BLOCKS TO START AND STOP
 MAKE AN ATTEMPT TO START IN THE

93 000316 000000
 94 000320 000074
 95
 96 000322 000000
 97
 98 000324
 99
 100 000001
 .END

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 SYMBOL TABLE

REC	000000	VM2	BLKEND	000150	002	000316	002	DIRFCT	000112	002	000000	002	000000	002
NSYS	= 000000	G	NTQF	000100	002	000250	002	NTFRD	000210	002	000000	002	000000	002
NTLOF	000000	VM2	NTSIZ	= 000324	002	000250	002	NTSYS	000060	002	000000	002	000000	002
NSYS	= 000000	G	FNZP	000150	002	000120	002	FORM1	000120	002	000000	002	000000	002
MONLW	= 000000	G	NOTE7	000210	002	000270	002	PC	= 000007	002	000000	002	000000	002
PS	= 17776		DETRN1	000120	002	000300	002	PFVRS	000110	002	000000	002	000000	002
RKSYS	= 000000	G	OP	= 000000	002	000000	002	P2	= 000002	002	000000	002	000000	002
RA	= 000000	G	PS	= 000005	002	000006	002	TCRA	= 177346	002	000000	002	000000	002
TCNT	= 177350		TCST	= 177300	002	000210	002	TCMC	= 177300	002	000000	002	000000	002
\$INTFN	= 000000	G												
.ABR.	000000	VM2												
SYSHND	000324	VM2												
ERRORS DETECTED:	0													
FREE CORP:	10070	WORDS												

.LP:/N:TYM/C=DT

APPENDIX B

FOREGROUND TERMINAL HANDLER

The following listing is a terminal handler for the foreground. The user can write his own handler using this code as an example, or use the copy provided in the software kit. Instructions for its use are found on the second and third pages of the listing.

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```

.TITLE KB.MAC V01-01
 ; RT-11 V2 DEVICE INDEPENDENT TERMINAL HANDLER, KB.
 ; REC-11-ORRA-D
 ; COPYRIGHT (C) 1975
 ;
 ; DIGITAL EQUIPMENT CORPORATION
 ; MAYNARD, MASSACHUSETTS 01754
 ;
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 ;
 ; MARCH 1975
 ; RGR

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; THIS HANDLER HAS THE FOLLOWING CHARACTERISTICS:
 ; 1) PARITY RETURN CAUSES THE REPAIRMAN
 ; OF THE INPUT BUFFER FOR THE CALLING READ REQUEST TO BE
 ; ZERO-FILLED, AND THE READ IS COMPLETED. THUS, THE HANDLER
 ;
 ; THIS HANDLER CAN BE USED BY EITHER THE FOREGROUND OR BACKGROUND (BUT NOT
 ; BOTH SIMULTANEOUSLY) TO READ AND WRITE TO ANY DL-11A OR KL-11A
 ; CONTROLLER TERMINAL.

```

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```

TRANSFERS ONE LINE AT A TIME, NO MATTER HOW LONG THE
 INPUT BUFFER IS FOR THE READ REQUEST, THE UNUSED PORTION
 OF THE BUFFER IS ZERO-FILLED. CARRIAGE RETURN ECHOS
 CARRIAGE RETURN, LYME-PEEP, AND INSERTS CR AND LF CHARACTERS
 IN THE BUFFER IF THERE IS ROOM, ELSE ONLY CR IS PLACED IN THE BUFFER.
 FROM EEP REMOVES Y LINE PEPS, AND INSERTS A PP CHARACTER IN
 BUFFER.
 3) RUBOUT ECHOS "\ " AND DELETES THE LAST CHARACTER IN THE BUFFER.
 IF THERE ARE NO CHARACTERS IN THE BUFFER, RUBOUT DOES NOT ECHO
 AND IS IGNORED.
 4) YR ECHOS ENOUGH SPACES TO POSITION THE PRINT HEAD AT THE
 NEXT TAB STOP, AND INSERTS A YR CHARACTER IN THE BUFFER.
 5) CTRL U ECHOS "U" AND ERASES THE CURRENT LINE.
 6) CTRL Z ECHOS "Z" AND CAUSES THE HANDLER TO REPORT END-OF-FILE.
 THE CTRL Z CHARACTER IS NOT INSERTED IN THE BUFFER.
 7) THE LOW-SPEED READER WILL RUN IF IT IS TURNED ON WHILE A READ
 REQUEST IS PENDING TO THE HANDLER. IF THE TAPE BEING READ HAS
 MANY TABS, HOWEVER, THE TIME NECESSARY TO ECHO THE TABS WILL
 CAUSE CHARACTERS FOLLOWING THE TABS TO BE LOST. TO DISABLE THE
 ENDING OF TABS, THE "SET" COMMAND CAN BE USED AS FOLLOWS:
 "SET RR L00" WILL DISABLE TAB ECHOS, ALLOWING A TAPE
 "SET RR L00" TO BE READ WITHOUT CHARACTER LOSS.
 "SET RR N000" WILL ENABLE TAB ECHOS, PER NORMAL KEYBOARD
 INPUT, THIS IS THE DEFAULT.
 8) WHEN THE HANDLER RECEIVES A READ REQUEST, A "C" CHARACTER IS
 PRINTED IN THE LEFT MARGIN OF THE TERMINAL TO DENOTE THAT THE
 HANDLER IS READY FOR INPUT. THIS CHARACTER CAN BE CHANGED, OR THE
 PROMPT FEATURE CAN BE REMOVED, BY RE-ASSIGNING THE SYMBOL
 "PROMPT" TO THE ASCII VALUE OF THE
 DESIRED CHARACTER. SETTING PROMPT TO "0" WILL CAUSE NO CHARACTER
 TO BE PRINTED.
 9) IF NO READ REQUEST IS ACTIVE, THE HANDLER WILL NOT ACCEPT INPUT,
 AND THE KEYS WILL NOT ECHO. IF IT DOES ECHO, THE HANDLER IS
 ACCEPTING INPUT.

THIS HANDLER CONTAINS CONDITIONAL CODE TO SUPPORT TERMINALS THAT
 REQUIRE FILLER CHARACTERS AFTER A PARTICULAR CHARACTER. TO ENABLE THE
 FILLER FUNCTION, DEFINE THE SYMBOL "FILLER" EQUAL TO THE ASCII
 VALUE FOR THE CHARACTER TO BE FILLED AFTER, AND THE SYMBOL "FILLCNT"
 TO BE THE OCTAL NUMBER OF NULLS TO BE ISSUED AFTER EACH OCCURRENCE
 OF THE CHARACTER DEFINED BY "FILLCHR". FOR EXAMPLE, TO PROVIDE
 FIFTEEN FILLER CHARACTERS AFTER A CARRIAGE RETURN, SET "FILLCHR" TO "15"
 AND "FILLCNT" TO "15".

```

1  / THE HANDLER IS INSTALLED VIA THE FOLLOWING PROCEDURE:
2  /
3  / 1) ASSEMBLE IT AS FOLLOWS:
4  /   .R MACRO
5  /   .R FILLER
6  /   .R MACRO
7  / 2) LINK IT AS FOLLOWS:
8  /   .R LINK
9  /   .R SYSERR
10 / 3) INSTALL IT AS DEVICE "KB", AS DESCRIBED IN SECTION XXXXX
11 / OF THE RT-11 V2 SOFTWARE SUPPORT MANUAL. REMEMBER THAT
12 / THE VECTORS FOR THE TERMINAL MUST BE PROTECTED IN THE BIT MAP
13 / AS DESCRIBED IN THAT SECTION.
14 / THE VALUES FOR THE VARIOUS TABLE ENTRIES SHOULD BE
15 /   HORIZEVALUP OF SYMBOL "KRBSTR" ON LAST LINE OF LISTING
16 /   DIVISOR (NON-FILE STRUCTURED DEVICE)
17 /   NAME "KBSTR" (R4000 FOR "KB ")
18 /   STATE HIGH ORDER BYTES, LOW ORDER BYTES AND DEVICE NUMBER
19 /   AVAILABLE. NOTE THAT IT CANNOT BE 0, A VALUE >50
20 /   IS RECOMMENDED.
21 / 4) ONCE INSTALLED, "KB" WILL BE AVAILABLE WHEN THE SYSTEM IS REBOOTED.
22 /
23 / THE HANDLER ITSELF IS ACTIVATED WITH READ AND WRITE REQUESTS, AS ARE ALL
24 / RT-11 DEVICE HANDLERS. WHEN USING SYSTEM PROGRAMS WHICH OPERATE ON
25 / LARGE BUFFERS, SEVERAL LINES MAY ACCUMULATE IN THE BUFFER BEFORE
26 / THEY APPEAR ON THE TERMINAL, AND THEN ALL AT ONCE. TO AVOID THIS PROBLEM,
27 / EACH OUTPUT BUFFER CAN BE ZERO-FILLED AND SENT TO THE TERMINAL TO PRINT.
28 / EACH LINE-THE HANDLER WILL IGNORE NULLS ON OUTPUT.
29 / IN FORTRAN, EACH LINE CAN BE FORCED IN OR OUT BY USING A REMIND
30 / FOLLOWING EACH READ OR WRITE TO THE DEVICE. FOR EXAMPLE:
31 /   LOGICAL I INPUT, (OR)
32 /   CALL ASSIGN (7, 'KB/C')
33 /   WRITE (7, I)
34 /   REMIND 7
35 /   WRITE (7, 2)
36 /   REMIND 7
37 /   INPUT
38 /   REMIND 7
39 /   .
40 /   .
41 /   .
42 /   .
43 /   .
44 /   .
45 /   .
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51 /   .
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54 /   .
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100 /   .

```



```

48 000216 001743          TPAUT2          IBRANCH IF NULL
49          .IFDF          .ICONDITIONAL CODE FOR FILLPR
50          CMPB          IDMES THIS CHAR NEED TO BE FILLED AFTERS
51          ANP          IBRANCH IF NOT
52          MOVB          IYPS=SET UP COUNT OF FILLS NEEDED
53          .ENDC
54 000220 01737 000100 174504 TPAUT3I          IENABLE PRINTER INTERRUPT
55 000226 110437 174506          IPRINT CHARACTER
56 000232 000207          RTS          IRETURN TO MONITOR
57

```

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```

58          IREQUEST TERMINATION AND ABORT COOP
59          ITHIS ROUTINE IS ENTERED WHEN THE I/O TRANSFER IS
60          ICOMPLETED OR ABORTED. THE DEVICE INTERRUPTS ARE DISABLED, AND
61          ISTANDARD MONITOR COMPLETION ENTRY CODE IS EXECUTED.
62 000234 000037 174504          ABORTI CLR 00TPCSR          IENABLE OUTPUT INTERRUPTS
63 000240 000037 174506          DONEI CLR 00KBCSR          IENABLE INPUT INTERRUPTS
64 000244 010704          MOV PC,R4          ISTANDARD MONITOR
65 000246 06P704 177502          ADI 00KCODE-.R4          ICOMPLETION ENTRY
66 000252 013705 000054          MOV 00R4,R5          ICODEP
67 000256 000175 000270          JMP 00PPSET(R5)

```

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```

1  IKEYBOARD INTERRUPT SERVICE
2  ITHIS IS THE KEYBOARD INTERRUPT SERVICE ROUTINE. AFTER ENTERING
3  ISYSTEM STATE, IT GETS THE TYPED CHARACTER INTO R4, THEN
4  IPROCEEDS DOWN A CHAIN OF CHECKS FOR THE SPECIAL CASE CHARACTERS
5  I(RUBOUT, CTRL U, CTRL Z, CR, FF). IF IT IS ONE OF THE SPECIAL
6  ICHARACTERS, THE ROUTINE "ECHO" IS CALLED TO ECHO APPROPRIATE
7  ICHARACTERS ON THE TERMINAL, THEN APPROPRIATE ACTION FOR THE SPECIAL
8  IIS TAKEN. IF A NORMAL CHARACTER IS TYPED, IT IS ECHOED AND PLACED
9  IIN THE USER BUFFER BEFORE THE INTERRUPT IS RESUMED.
10 K0TNYI JBR 00SINPYR          IENTER SYSTEM STATE
11 000262 000577 000460          .WORD 00CPR4+0BRT          IGET CHAR
12 000266 000140          MOVB 00KBRUP,R4          ISTRTP TO SEVEN BITS
13 000270 113704 174502          BIC 017760,R4          IIS THIS CHARACTER A RUBOUT?
14 000274 042704 177600          CMPB R4,#DELETE          IBRANCH IF NOT
15 000300 120427 000177          ANE 113          IANY CHARS LEFT TO RUB OUT?
16 000304 001020          CMP UBPTR,UBPTR1
17 000306 026767 000430

```

18	000314	001520	REG	K0TN	IND=IGNORE RUBOUT
19	000316	004367	REC	UBPTR	IBACK UP POINTER INTO USPR BUFFER
20	000322	004167	JOB	R1,ECHO	
21	000326	134	.BYTF	'\,377	
22	000330	104267	TNCB	TARCNT	IBUMP TAR COUNTER FOR "\
23	000334	104077	CLPB	SURPTR	IZERO RUBOPO OUT CHAR
24	000340	004267	TNC	RYCNT	IAND INCREASE TRANSFER COUNT TO REFLECT LOAT CHAR
25	000344	000304	RR	K0TN	IR=ENABLE INTERRUPTS AND PXTY
26	000346	120427	CHPB	R6,0FF	IIS THIS CHAR A FORM FEED?
27	000352	001006	ANP	AS	IBRANCH IF NOT
28	000354	004167	JOB	R1,ECHO	IYPS=ECHO Y LINE FEEDS
29	000360	012	.BYTF	LF,LF,LF,LF,LF,LF,LF,LF,377	
30	000363	012			
31	000366	012			
32	000370	120427	CHPB	R4,0CR	IIS THIS CHAR A CR?
33	000374	001017	ANP	Y8	IBRANCH IF NOT
34	000376	004167	JOB	R1,ECHO	IYPS=ECHO CR,LP
35	000402	015	.BYTF	CR,LF,0,377	
36	000406	110477	MOV8	R4,0UBPTR	IFUT CR IN USER BUFFER
37	000412	004267	TNC	UBPTR	IBUMP USER BUFFER POINTER
38	000416	004367	REC	RYCNT	IROOM IN BUFFER FOR LP TMO?
39	000422	001776	REG	NONE	IDON'T INSERT IT IF NOT
40	000424	117777	MOV8	ALP,0URPTR	IELSE ADD LP TO BUFFER
41	000432	000702	RR	NONE	

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1	000434	120427	0000P5	78i	CHPB	R4,#CTRLU	IIS CHAR CTRL UP?
2	000440	001007	ANP		R8	AS	IBRANCH IF NOT
3	000442	004167	000120		R1,ECHO	R1,ECHO	IECHO "SU"
4	000446	136	1P5	015	'S,U,CR,LF,0,377		
5	000451	012	000	377	.BYTF		
6	000454	000167	177352		JMP	RETRY	IAND RSTART READ
7	000460	120427	000072	08i	CHPB	R4,#CTRL7	IIS CHAR CTRL Z?
8	000464	001013	000074		ANP	Q8	IBRANCH IF NOT
9	000466	004167	000074		JOB	R1,ECHO	IECHO "SZ"
10	000472	136	132	015	.BYTF	'S,Z,CR,LF,0,377	
11	000475	012	000	377	MOV	K0C0E,R5	IPPOINT R5 TO 0 ELEMENT
12	000500	014705	177304		RIS	SECF,0=2(R4)	IAND SET ENF FLAG IN CSM
13	000504	053775	020000	177776	RR	NONE	ISTOP TRANSFER
14	000512	000652					
15	000514	120427	000040	09i	CHPB	R4,#40	IIS THIS A PRINTING CHAR?
16	000520	004402			ALT	P18	IBRANCH IF NOT
17	000522	104267	000213		TNCB	YARCNT	IYPS-INCREASE TAR POSITION

```

1A 000526 110467 000004          ISPT UP TO ECHO CHAR
19 000532 000167 000030          R4,200
20 000536 000000 000377          R1,ECHO
21 000540 110477 000176          R4,0UBPTR
22 000544 000267 000172          JUMP RUPPTR
23 000550 000367 000162          JANY MORE TO TRANSFER
24 000554 000163 000163          BRANCH IF NOT
25 000556 010737 000101          SENABLE KEYBOARD INTERRUPT
2A 000564 000207 000207          RRETURN TO MONITOR

```

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```

1 1SUBROUTINE PCHO
2 THIS SUBROUTINE SERVES TO PLACE THE SPECIFIED CHARACTERS IN THE
3 ECHO BUFFER, AND START THE PRINTER IN CASE IT IS IDLE.
4 THE CALLING SEQUENCE IS
5
6 JBR R1,ECHO
7
8 .BYTE CHAR1,CHAR2,CHAR3,...,CHARN,377
9
10 INATE THAT THERE MUST BE AN EVEN NUMBER OF BYTES IN THE ARGUMENT LIST
11 AND THEREFORE THE NUMBER OF CHARACTERS EXCLUDING THE 377
12 MUST BE ODD.
13
14 WHEN ENTERED, PCHO SCANS THE ECHO BUFFER TO FIND THE END OF THE
15 ECHO LIST, WHICH IS MARKED BY A NULL BYTE. WHEN THE END OF THE LIST
16 IS FOUND, IT IS DETERMINED IF THERE ARE AT LEAST 8 FREE SLOTS IN THE LIST
17 TO ACCOMMODATE A POSSIBLE LINE FEED OR FORM FEED. IF NOT, THE
18 CHARACTER JUST TYPED IS IGNORED. IF SO, THE CHARACTERS FROM THE
19 ARGUMENT LIST FOLLOWING THE CALL ARE INSERTED IN THE BUFFER,
20 AND THE PRINTER IS STARTED IF IT IS IDLE, AND THE ROUTINE RETURNS.
21 INATE THAT TAB IS A SPECIAL CASE: IF R4 CONTAINS A TAB CHARACTER
22 WHEN THIS ROUTINE IS ENTERED, THE ARGUMENT LIST IS NOT USED. RATHER,
23 AN APPROPRIATE NUMBER OF SPACES TO MOVE THE PRINT HEAD TO THE
24 INPUT TAB STOP ARE PLACED IN THE ECHO BUFFER, AND THE ROUTINE RETURNS
25
26 .GENERAL LBR
27 PC,R5          ICALC ABSOLUTE ADDRESS
28 MOV R4,R5     JOP PCHO BUFFER
29 ADD R4,R5     ISAVE ADDRESS OF ECHO BUFFER
30 MOV R5,TEMP   SERLENGTH-I, TEMP POINTS TO END OF ECHO BUFFER
31 ADD R4,TEMP   IS THIS END OF ECHO LIST?
32 AND R4,R5     BRANCH IF NOT
33 JZ R5,TEMP    JY20-R5 POINTS TO FIRST FREE SLOT IN ECHO LIST
34 JZ R5,TEMP    IF END NUMBER OF FREE SLOTS IN ECHO LIST
35 JZ R5,TEMP    IS THERE ENOUGH ROOM TO ECHO TAB OR FF?
36 JZ R5,TEMP    BRANCH IF YES
37 JZ R5,TEMP    IND-IGNORE THIS CHAR THEN
38 JZ R5,TEMP    JY20MISS INTERRUPT
39
40 R4,R5
41 R4,R5
42 R4,R5
43 R4,R5
44 R4,R5
45 R4,R5
46 R4,R5
47 R4,R5
48 R4,R5
49 R4,R5
50 R4,R5
51 R4,R5
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93 R4,R5
94 R4,R5
95 R4,R5
96 R4,R5
97 R4,R5
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99 R4,R5
100 R4,R5

```

```

36 000634 010406
37 000636 120427
38 000640 000011
39 000642 001013
40 000644 110728
41 000650 100207
42 000654 130767
43 000662 001300
44 000664 000721
45 000666 100015
46 000670 000403
47 000672 110129
48 000674 100376
49 000676 100005
50 000700 000707
51 000704 010004
52 000706 000001
53
54

```

SAVE CHAR
IF THIS CHAR A TAB?
IF THIS COMPARE OPERAND CAN BE CHANGED BY SET LBR
IF BRANCH IF NOT
IF A SPACE
IF BUMP POSITION COUNTER
IF TAB STOP YES?
IF BRANCH IF NOT
IF YES-ARTIFICIALLY BUMP RETURN
IF END ECHO LIST
IF AND START ECHO
IF MOV CHAR INTO PCMD LIST
IF BRANCH IF END-OP-LIST NOT SEEN
IF LBR USE A TO MARK END OF ECHO LIST
IF PRINT A CHAR TO START PRINTER
IF STORE CHAR
IF RETURN

```

38i      MOV      R4,=(SP)
        CMPB    R4,(PC)+
        LSRPT:  HT
58i      ANP     @SPACE,(RS)+
        INCB    TARGNT
        AIB    @7,TABCNT
        ANP     @8
        TBT    (R1)+
        CLRB   (R4)
        RR     @8
        MOVB   (R1)+,(R4)+
        RPL    @8
        CLRB   @8
        JBR    PC,TPOUTP
        MOV     (SP)+,R4
        AYS    R1
        .DARL  LBR

```

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```

1 1
2 2
3 3
4 000710 000
5
6
7
8
9
10
11
12
13 000734 000000
14 000736 000000
15 000740 000
16 000741 000
17 000742 000000
18 000744 000000
19
20
21 000746 000000
22
23 000750 000001
24 000001 .END

```

DATA AREA
ECHO RING_BUFFER=PBLENGTH CHARACTERS LONG
R0STR1 ,0BYTE 0
.OLKR PBLENGTH=1
IVARTARLP ARPA
FILCHR1 ,0BYTE FILCHR
FILCHR11 ,0BYTE 0
TEMP1 ,0WORD 0
RYCNT1 ,0WORD 0
READFL1 ,0BYTE 0
TARGNT1 ,0BYTE 0
UBPTR1 ,0WORD 0
UBPTR11 ,0WORD 0
MONITOR ,SYSTEM STATE ENTRY LINK
SINPTR1 ,0WORD 0
KBRTZ= ,KRSTRY
.END

IF FILLER CONDITIONAL
IF CHARACTER TO BE FILLED AFTER
IF NUMBER OF FILLS REMAINING
IF TEMPORARY
IF USER TRANSFER COUNT
IF FLAG FOR "READ IN PROGRESS"
IF TAB POSITION COUNTER
IF POINTER INTO USER BUFFER
IF POINTER TO START OF USER BUFFER

SYMBOL TABLE

```

ABORT 000230R
DELET 000177
PF 000014
KBTN 000556R
KOVEC 000300
PC -X000007
READFL 000740R
R2 -X000002
SPACP 000040
TPINT 000120R
TPVEC 000304

. ABS. 000424 000
000750 001
ERRORS DETECTED: 0
FREE CORE: 14460. WORDS
KB,LP1/NI/TM/C=KR

BYTCNT 000736R
NONE 000240R
MT 000011
K01NT 000262R
LP 000012
PRAMP 000076
RETRY 000032R
R3 -X000003
TARCNT 000741R
TPDUT 000176R
UBPTR 000742R

CR 000015
PBLNG= 000024
K0RUP 000011
K0LOP 000000R
L0ROPT 000640R
PR4 000200
RTSPP 000232R
R4 -X000004
TEMP 000734R
TPAUT1 000214R
UBPTR1 000744R

CTRLZ = 000032
POF = 020000
KBC0R = 176500
K0STRY 000000R
OPLAR 000412
R0STRY 000710R
R1 -X000001
SP -X000006
YPC0R = 176504
TPPUT3 000200R
...V2 = 000001

CTRLU = 000025
PCMO 000566R
KBC0E 000010R
K0SIZE= 000750
RFP0ET= 000270
RR7 = 000340
R0 -X000000
R5 -X000005
TPAUP = 176506
TPAUT2 000126R
SINPTR 000740R

```

CORRECTIONS AND ENHANCEMENTS TO THE KB HANDLER AND INSTALLING HANDLERS (JD)

1. A CTRL U will double the byte count used by the handler.
2. The write interrupt is not disabled upon completion of a write.
3. The handler will not transfer a file to the KB terminal if it is 32K words or larger. This is due to the byte count overflowing a 16 bit word.
4. Execution time of both read and writes can be diminished with a minimum number of changes.

The following Edit changes correct the above problems. After the changes are made, KB.MAC must be reassembled, relinked, and reinstalled.

```

,R EDIT
*EBKB,MACSR$$
*FBYTCNTSASS
*I      MOV      (R5),R4 <CR>
        ROR      (R5) <CR>
$$
*GDONES=CFIN$$
*GPCSS=CBMISS
*G2S=C4FS
*AKSF#EPLESF
*QAKS2<AK>$$
*GBGTS=CBEQ$$
*GDONES=CABORT$$
*GPC;SASS
*ITPOUT4;      DEC      BYTCNT <CR>
        BR      TPOUT2 <CR>
$$
*GDONE;SASS
*IFIN;S$
*F#4Q$4ASS
*I      BEQ      PUT <CR>
$$
*4ASIPUT;SS
*FTEMP$=C-(SP)SS
*3<GTEMP$=C(SP)S>SS
*GTEMP$=C(SP)+SS
*FTEMP;SS
*QAKS$
*EXSS
    
```

It should also be noted that the slots occupied by the TT and BA handlers (slots 1 and 12) must not be used. On the page 5-13 and 5-14 of the RT-11 Software Support Manual (DEC-11-ORPCA-B-D) the aforementioned handler slots are used in the example patches. Replacement pages will be generated to correct this. Replacement pages will also be generated for pages 4-23 and 4-26 of the RT-11 System Generation Manual (DEC-11-ORGMA-A-D). The instructions on these pages should read:

The table entries to be patched are:

SHSIZE+OCTAL OFFSET
SDVSIZE+OCTAL OFFSET
SPNAME+OCTAL OFFSET
SSTAT+OCTAL OFFSET

ERRORS IN KB.MAC (SPR 11-10433,11-10434,11-10595 JM)

The following errors involve the operation of KB.MAC:

1. The output interrupt is not disabled when the echoing of characters on input is done.
2. Characters, upon requesting output, are sometimes lost.
3. Nulls decrement the character count but are not placed in the character buffer.
4. When KB is used for a second terminal on a LSI-11, the processor will trap to ODT micro-code upon enabling the printer interrupt.

The following edit changes will correct the above problems. Previous changes to KB.MAC (Article Seq. #4, published in the September, 1976, DIGITAL SOFTWARE NEWS) must be installed before applying the changes below. After the changes are made, KB.MAC must be reassembled, relinked, and reinstalled.

```
.R EDIT <CR>
*EBKB.MAC$$$
*FTPOUT:$$-A$$$
*I      BEQ      TPOUT <CR>
        CLR      J#TPCSR <CR>
        RTS      PC <CR>

$$$
*GTPOUT3:$$
*SKAU$$
*FPUT:$$-4D$$$
*-5A$GPUT$=CKBIN$$
*EX$$$
```


APPENDIX C

VERSION 1 EMT SUMMARY

Although Version 1 programmed requests are supported by Versions 2 and 2B of RT-11, it is strongly recommended that the Version 1 formats not be used. For purposes of compatibility, however, this section provides a brief review of the V1 format. The V2/V2B format is covered in detail in Chapter 9 of the RT-11 System Reference Manual.

In brief, the major distinctions between V1 and V2/V2B formats are:

1. V1 format has arguments pushed on the stack and in R0. V2/V2B requests generally accept a set of arguments, or an argument in R0.
2. V1 channel numbers are restricted to 16_{10} . Also, the channel number in V1 is not a legal assembler argument; it is merely an integer in the range 0 to 15_{10} .
3. V1 requests are non-reentrant because the channel number and function code are embedded within the EMT instruction.

Table C-1 lists all the Version 1 macro calls. Those in the left column have the same format as the corresponding Version 2/V2B request; those in the right column have a different format, shown after the table. The operations performed by the requests are the same in both versions.

Table C-1
V1 Programmed Requests

V1 - Format Same as V2/V2B	V1 - Format Different from V2/V2B
.CSIGN	.CLOSE
.CSISPC	.DELETE
.DATE	.ENTER
.DSTAT	.LOOKUP
.EXIT	.READ
.FETCH	.READC

(continued on next page)


```
.WAIT .chan
```

```
.WRITE }  
.WRITC } .chan,.buff,.wcnt,.crtn,.blk [ .crtn is required ]  
.WRITW } [ only for .WRITC ]
```

The system macro library (SYSMAC.SML) can be used with Versions 2 and 2B to generate Version 1 programmed requests.

Under Version 2, the `..V2..` macro is capable of handling V1 expansions. `..V2..` normally expands as:

```
.MCALL ...CM1,...CM2,...CM3,...CM4  
...V2=1
```

This causes Version 2 expansions in all cases. To allow expansion of all V1 requests in their V1 format (and all new Version 2 requests in V2 format) the `..V2..` macro should not be called, but the utility macros must still be defined:

```
.MCALL ...CM1,...CM2,...CM3,...CM4
```

Omitting both `..V2..` and the utility macros causes all old V1 requests to be expanded in V1 format; no V2 requests can be used.

Under Version 2B, the `..V1..` macro call enables expansion of all macros in Version 1 format. `..V1..` expands as:

```
...V1=1
```

To enable expansion of all Version 1 macros in V1 format and all new Version 2 macros in V2 format, these statements must be included:

```
.MCALL ..V1...CM1,...CM2,...CM3,...CM4  
..V1..
```

A listing of SYSMAC.SML is provided in the RT-11 System Reference Manual.

APPENDIX D
FOREGROUND SPOOLER EXAMPLE

The following program is an example of a line printer spooler for the foreground. Instructions for its use follow.

1. Create the program using the Editor and store it on the system device under the name LSPOOL.MAC.
2. Next assemble it under MACRO and then link it to create the REL format output file:

```
.R MACRO  
*LSPOOL=LSPOOL
```

```
.R LINK  
*LSPOOL=LSPOOL/R
```

3. Load the necessary handlers (in this case, LP and RF) and run the program. All files on device RF with the extension .LST are listed on the line printer and then deleted from RF:

```
.LOA LP,RF<CR>
```

```
.FRU LSPOOL<CR>
```

```
F>  
DEVICE TO SPOOL?
```

```
B>
```

```
.
```

[Control must be redirected
to the foreground via ^F.]

```
F>  
RF:*.LST<CR>
```

This program assumes device DK: and extension .LPT unless otherwise indicated.

```

1  TITLE LSP00L = LINF PRINTER SPOOLER
2  .SRTTL A USEFUL FORGROUND PROGRAM
3
4  THIS PROGRAM FOR THE FOREGROUND IS A LINF PRINTER SPOOLER.
5  IT SPUNCHES A SPECIFIED DEVICE FOR FILES WITH A PARTICULAR
6  EXTENSION (THE DEFAULT IS LPT) AND PRINTS THEM, DELETING
7  AFTER PRINTING. IF NONP ARE FOUND, IT WILL GO TO SLEEP FOR
8  HALF A MINUTE, PERMITTING THE BACKGROUND TO RUN.
9
10 TO RUN LSP00L, FIRST LOAD LP HANDLER AND INPUT DEVICE HANDLER
11 IF IT IS NOT THE SYSTEM DEVICE TYPE.
12
13 P.G..
14
15 LDA LP,RP
16 .PRU LSP00L
17
18 LSP00L WILL TYPE: "DEVICE TO SPOOL?"
19 TYPE INPUT DEVICE AND FILE DECRYPTION, P.G.i
20
21 .PFI,LIST
22
23
24 .MFAIL .V2...REDOFF.
25 .MFAIL .REARM..WRSTM..LOOKUP..DPLTF..CSISPC..TTYIN
26 .MFAIL .PRINT..TYOUT..SRPSPY..CTRLC..CLOSE..EXIT
27
28 .V2..
29 .RPGNER
30
31
32
33
34
35
36
37
38
39
40
41
42
43
44
45
46
47

```



```

4R 000044      01A706      177732      000000
40 000052      000000      000000      000000
50
51 000056      000000      000000      000000
52 000070      000000      000000      000000
53 000072      000000      000000      000000
54 000100      001102      000000      000000
55 000104      01A201      000000      000000
56 000106      000000      000000      000000
57 000112      022700      000015      000000

```

```

PRINT ERROR MESSAGE
IMPORT STACK, FALL THRU TO BPOIN

IMP WILL USE CH 0, SO CLEAN IT UP
IMPORT CTRL/D FLAG 00
IMPORTING MSG WILL PRINT.
POINT TO COMMAND STRING BUFFER
ICOPY THE POINTER AND INPUT COMMAND
IA CHARACTER AT A TIME.
ICARRIAGE RETURN

```

```

RARC0M1 .PRINT 0M02
MOV      STK0AV,0P

RERZM1  .CLOSE 00
.PRTLN
.PRINT  0M01
MOV      0C01LK,02
MOV      02,R1
.TTYTN
CMP      0C0,00

```

LSPOOL - LINE PRINTER SPOLLER 07-11 MACRO VM02-10 A-APR-75 1P03115 PAGE 10

```

58 000116      001773      000000      000000
59 000120      110022      000000      000000
60 000122      127000      000012      000000
61 000126      001347      000000      000000
62 000130      100042      000000      000000
63 000132      010000      000000      000000
64 000144      001336      000000      000000
65 000146      104735      000000      000000
66 000148      000000      000000      000000
67 000152      01A700      000776      000000
68 000204      001002      000000      000000
69 000210      012700      000000      000000
70 000212      040024      000244      000752
71 000216      010067      000000      000000
72 000220      01A700      000000      000000
73 000222      001002      000000      000000
74 000226      012700      000000      000000
75 000230      01A700      000000      000000
76 000232      01A700      000000      000000
77 000234      010011      000000      000000
78 000236      000001      000000      000000
79
80 000242      101003      000532      000000
81 000252      001767      000000      000000
82 000254      001004      000000      000000
83 000260      000000      000000      000000
84 000262      000000      000000      000000
85 000270      000000      000000      000000
86 000272      101007      000000      000000
87 000320      000000      000000      000000
88 000322      012700      001200      000000

```

```

IYPS, IGNORE IT.
IMOVE IT INTO RUPFR AND
IYPSY FOR PND OF LINE.
IMN, GET ANOTHER CHARACTER.
IYPS, CLEAR OUT TMP LINE FEED
IPROCESS THE COMMAND
IYPSY 0 OF SWITCHES IS UNCHANGED.
IMN SWITCHES ALL000
ISYNTAX ERROR

IGPT FILE EXTENSION TO PRINT.
IBRANCH IF USER SPECIFIED,
IELSP USE LPT EXTENSION

ISAVE TMP EXTENSION FOR LATER.
IGPT THE INPUT DEVICE NAME
IBRANCH IF USER SPECIFIED.
IELSE USE THE
IDFAULT DEVICE.
IGPT DEVICE NAME IN FILE
IOPSCRYPTOR BLOCK, CLEARING
IOUT ANY FILE NAME,
IINPUT DEVICE HANDLER MUST BE RESIDENT
IILLEGAL DEVICE
IYPSY ENTRY POINT
IBRANCH IF O.K.
IELSP PRINT MESSAGE
IOPEN CHANNEL TO SPAN DIRECTORY

IRP -> BUFFER

```

```

REG      10
MOV      00,(02)+
RMP      01P,00
RMP      10
FLAG     -(02)
.CSIRP  01,0NEPEXT,R1
MOV      (0P)+,00
RARC0M1 RARC0M
RCA      RARC0M
.LOOKUP 010,01,0LP
MOV      010LK+3A,00
RMP      00
MOV      (PC)+,00
.RAD00  /LPT/
MOV      00,LPT...
MOV      010LK+3P,00
RMP      00
MOV      (PC)+,00
.RAD00  /DR0/
MOV      00,001
CLR      2(01)

.DSTATUS 0100,R1
RCA      00
YBY      Y0R+4
RMP      00
.PRINT  0M02
RR      RERZM1
.LOOKUP 0100,00,R1
RCA      RANER
.ENARL  LSR
PINOLP: MOV      00,BUFF,R2

```



```

135 000730
136 000742
137 000772 000794
138 000774 044676
139 000776
140 001010
141 001016
142 001018
143 001037
144 001038
145 001051
146 001061
147 001111
148
149
150 001126
151 001132 044674
152 001134 000000
153 001142
154 001260
155

```

LSPOOL - LINE PRINTER SPOOLER RT-11 MACRO VM02-10 A-APR-75 12:03:19 PAGE 10

```

SYMBOL TABLE
RANCM 000000
CR = 00015
FINDLP 000320
MSG0 001010
PC =X00007
R4 =X00004
TIMBLK 001126
. ABS. 00000 000
. P2126 001
ERRORS OFTECEN: 0
PRPE CORP1 10055. WORDS
LSPOOL,LP1/NITTH=LSPOOL.

```

```

:CLOSE #2
:DPLOT #IND,03,R1
:WHEN CONTINUE
:SPOLLER OUTPUT DEVICE
:SENT ARGUMENT BLOCK

```

```

:CLOSE #2
:DPLOT #IND,03,R1
:WHEN CONTINUE
:SPOLLER OUTPUT DEVICE
:SENT ARGUMENT BLOCK

```

LSPOOL - LINE PRINTER SPOOLER RT-11 MACRO VM02-10 A-APR-75 12:03:19 PAGE 10

```

SYMBOL TABLE
RANCM 000000
CR = 00015
FINDLP 000320
MSG0 001010
PC =X00007
R4 =X00004
TIMBLK 001126
. ABS. 00000 000
. P2126 001
ERRORS OFTECEN: 0
PRPE CORP1 10055. WORDS
LSPOOL,LP1/NITTH=LSPOOL.

```

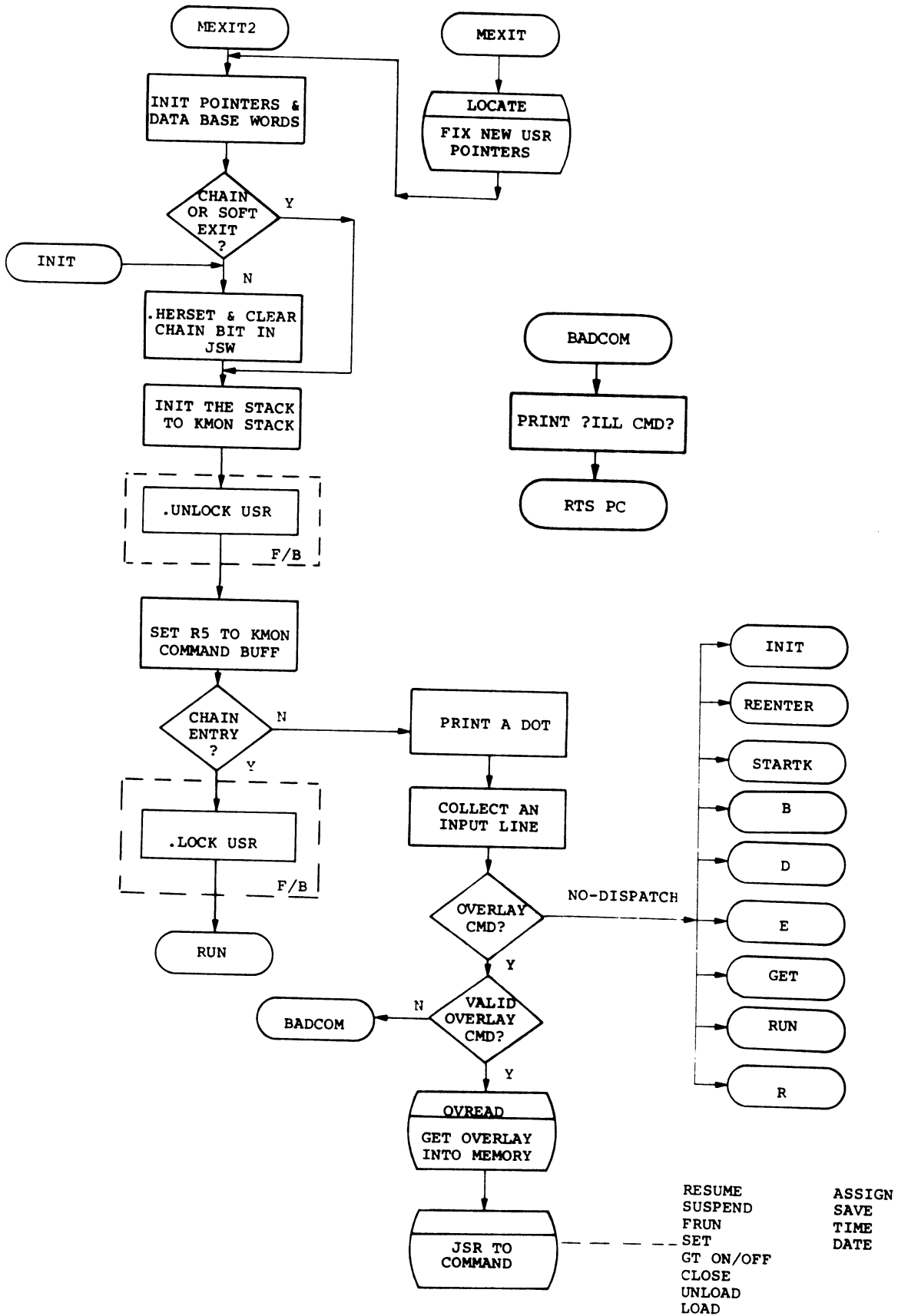

APPENDIX E
S/J AND F/B MONITOR FLOWCHARTS

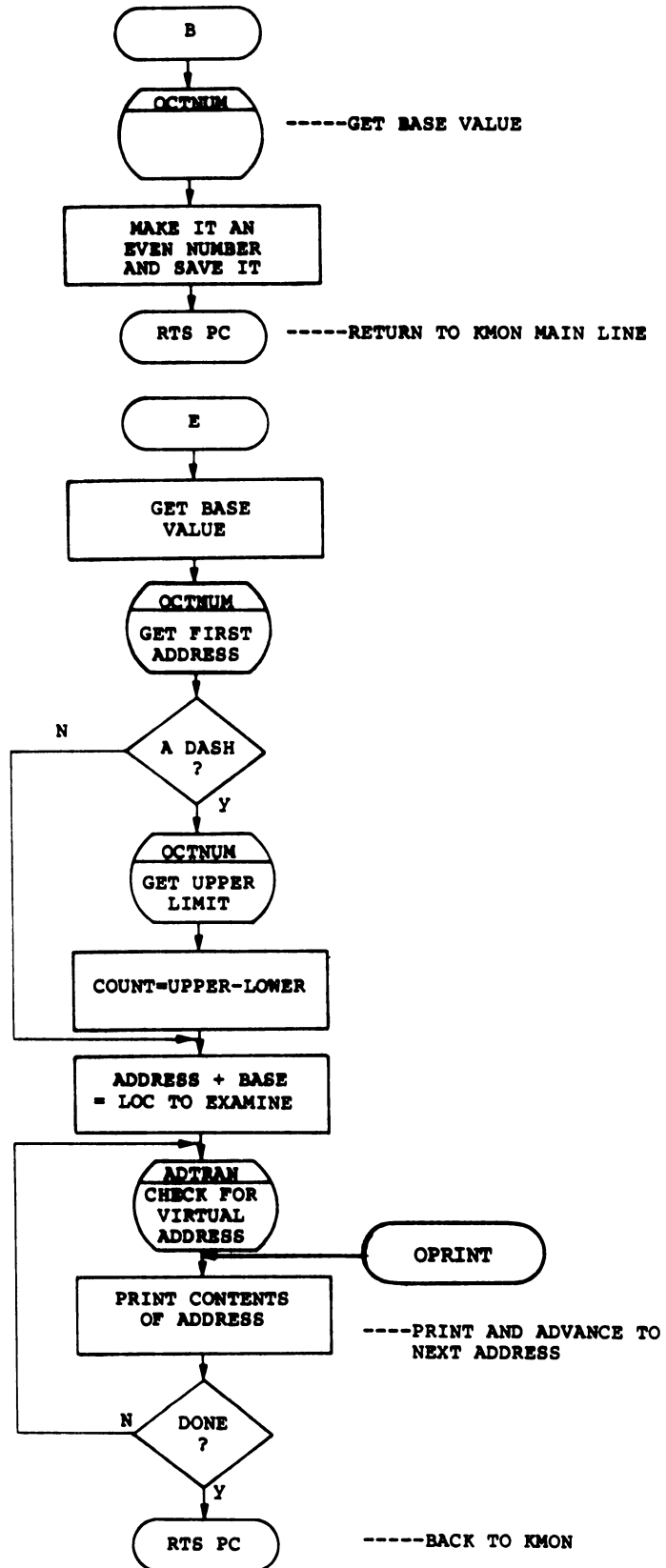
The following flowcharts are of the Single-Job and Foreground/
Background Monitors. It is recommended that the reader have source
listings available for reference. Steps inside are per-
formed only in the F/B or S/J Monitor, as noted.

An index of all entry points appears at the end of the appendix.

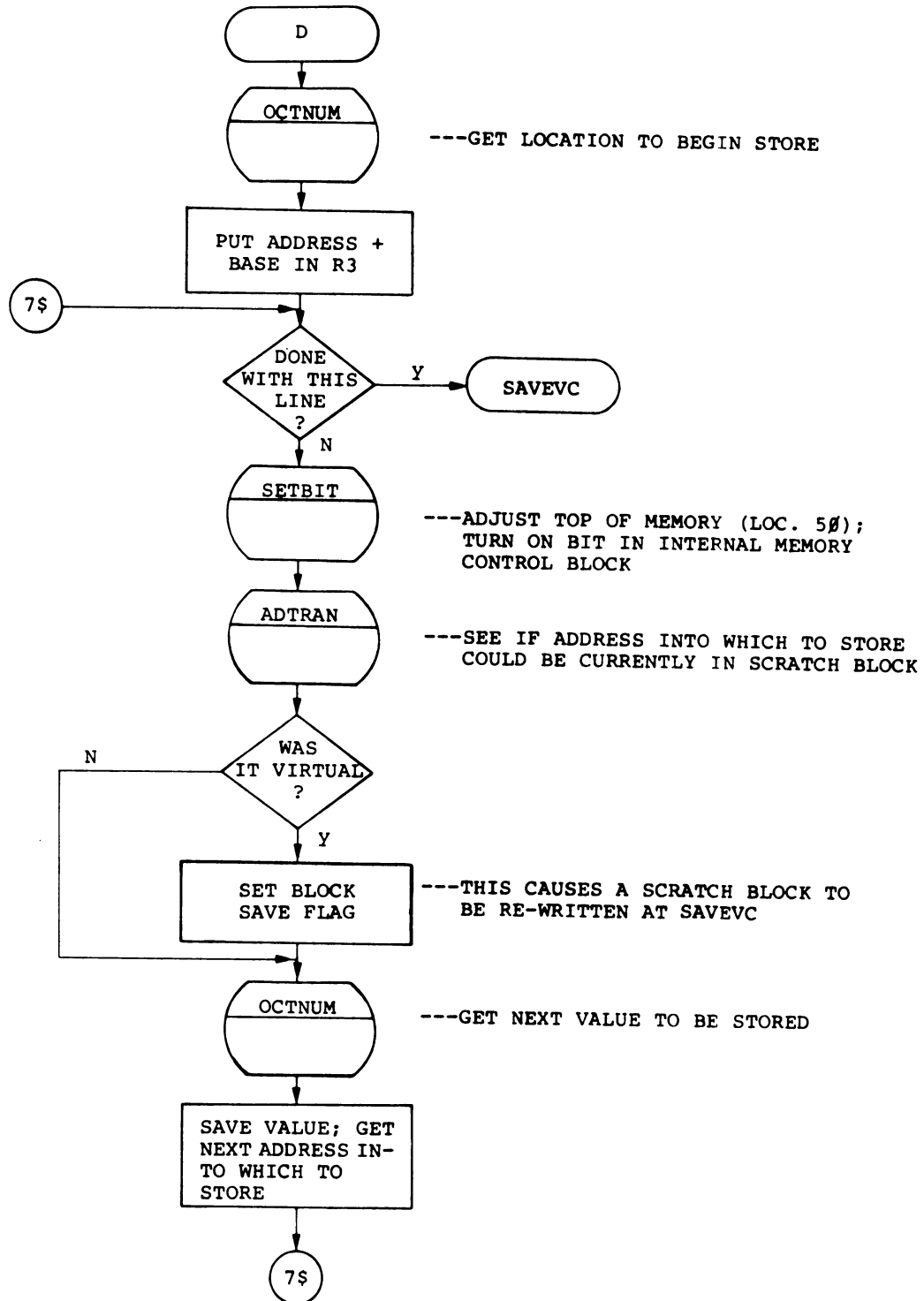
E.1 KMON (KEYBOARD MONITOR) FLOWCHARTS

KMON

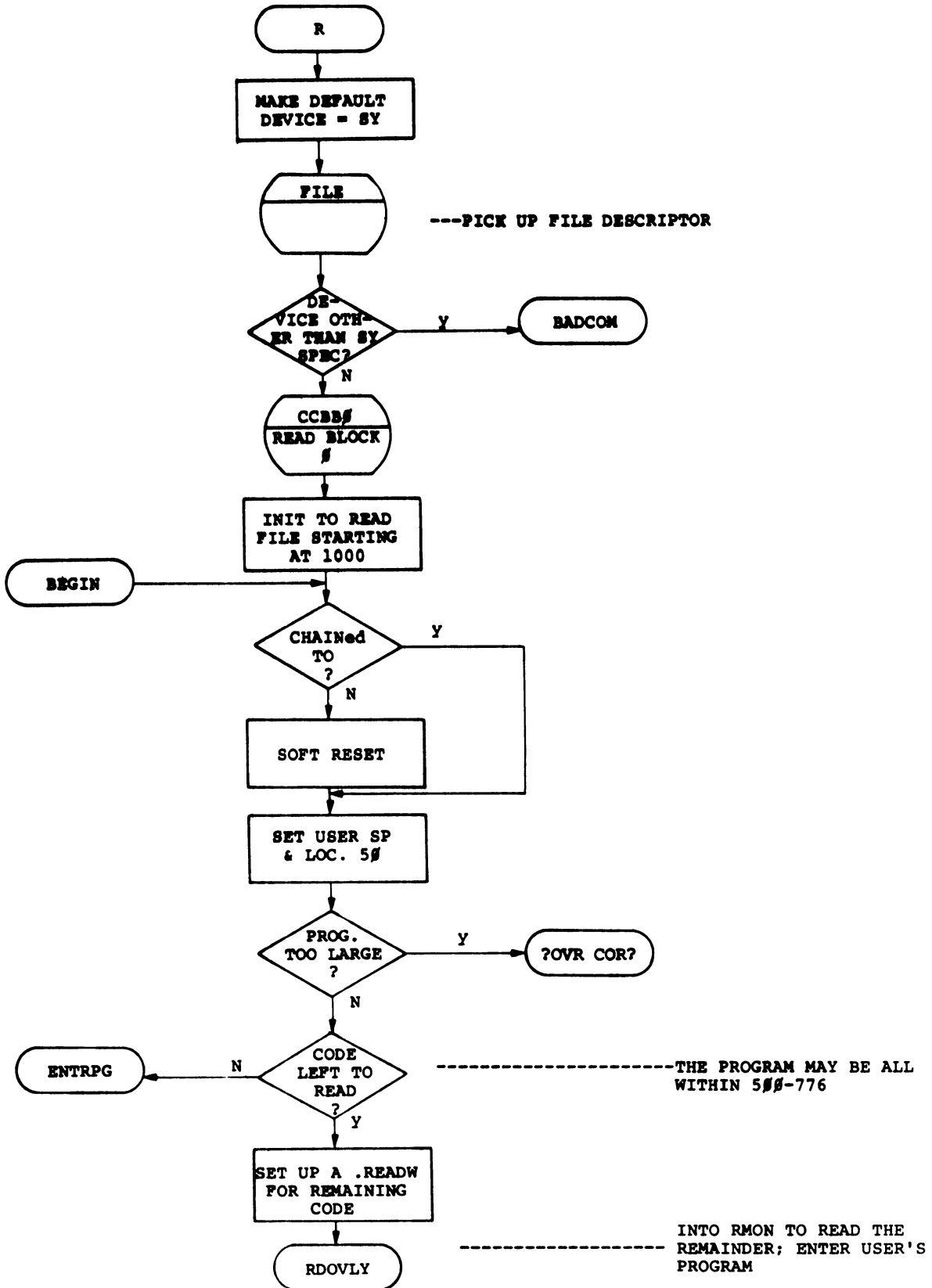




DEPOSIT

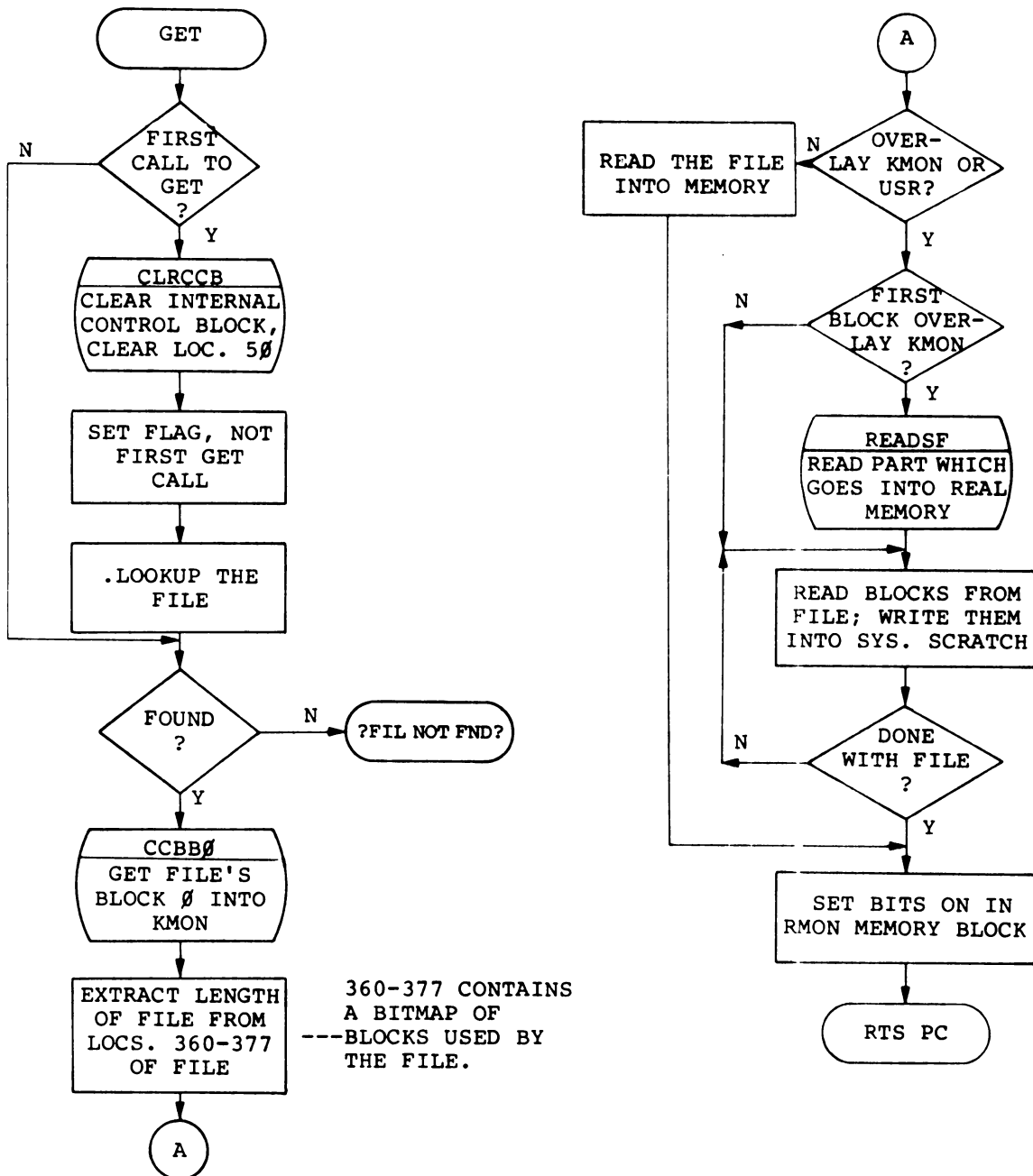


SAVEVC - Entered to rewrite the current virtual block back into the system scratch area. It also acts as the exit point for Deposit; The RTS PC will return control to KMON.

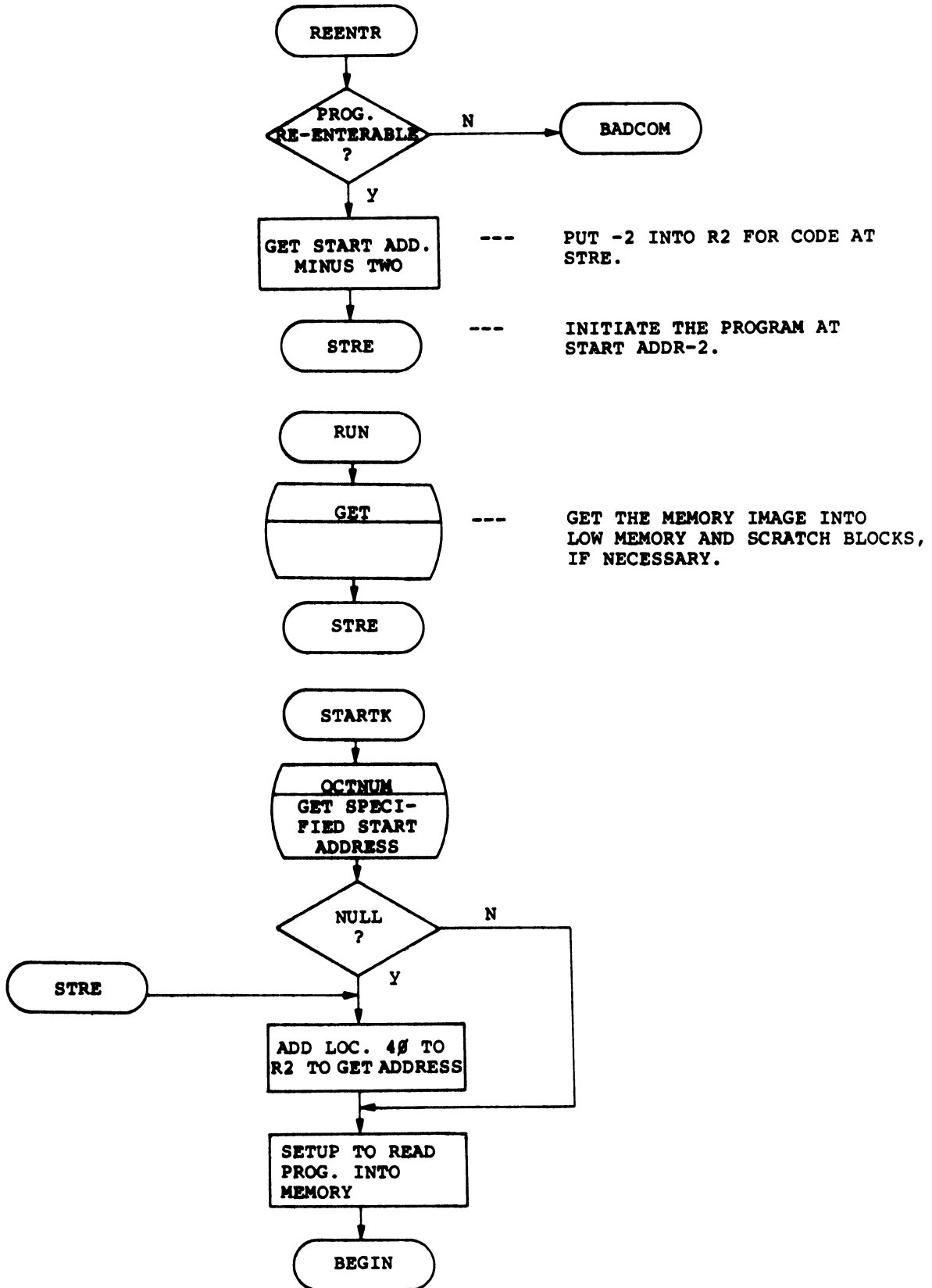


GET

GET - Used to load a .SAV image into memory. If parts of the file overlay KMON/USR, those parts are placed into system scratch blocks.



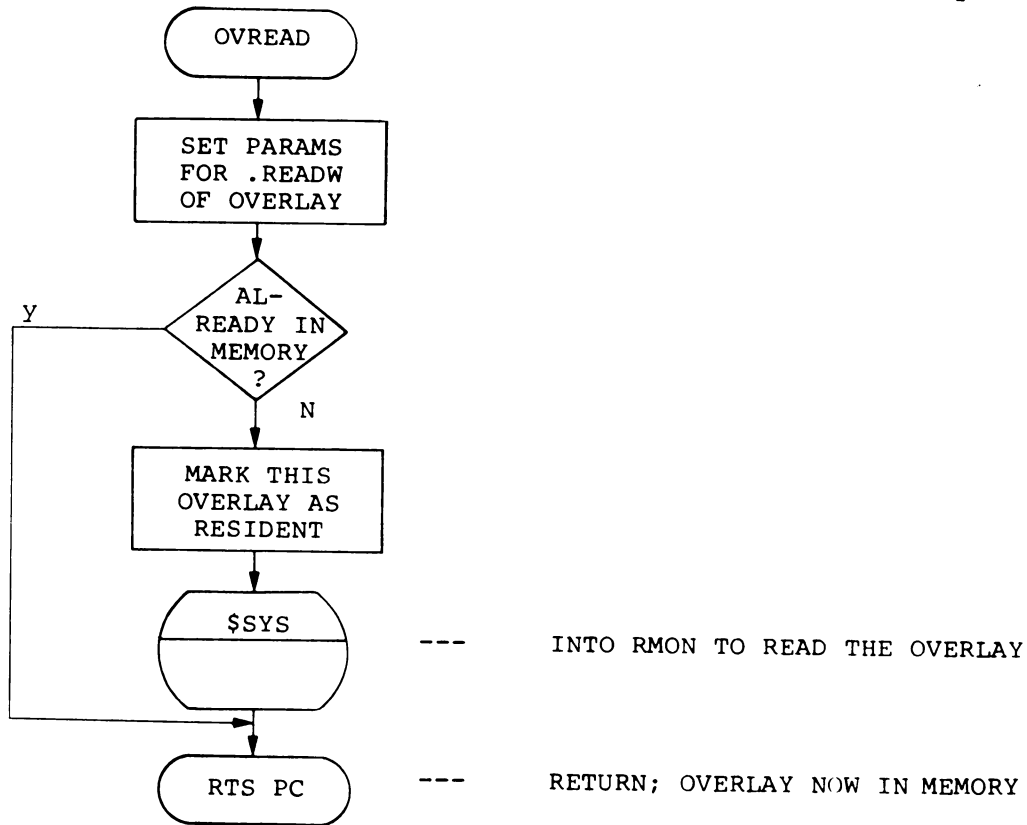
REENTER/RUN/START



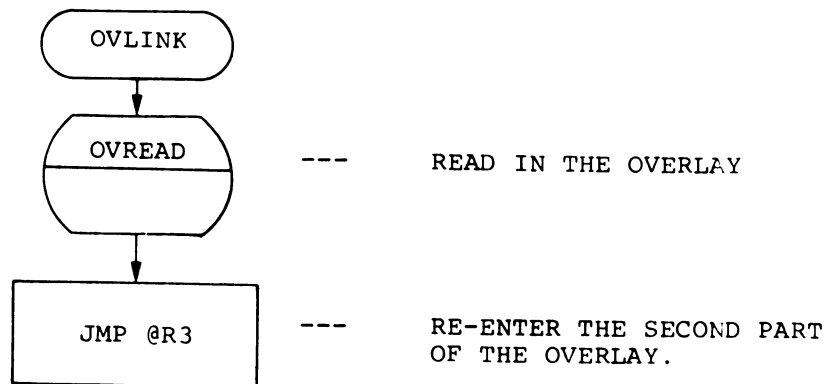
E.1.1 KMON Subroutines

OVREAD/OVLINK

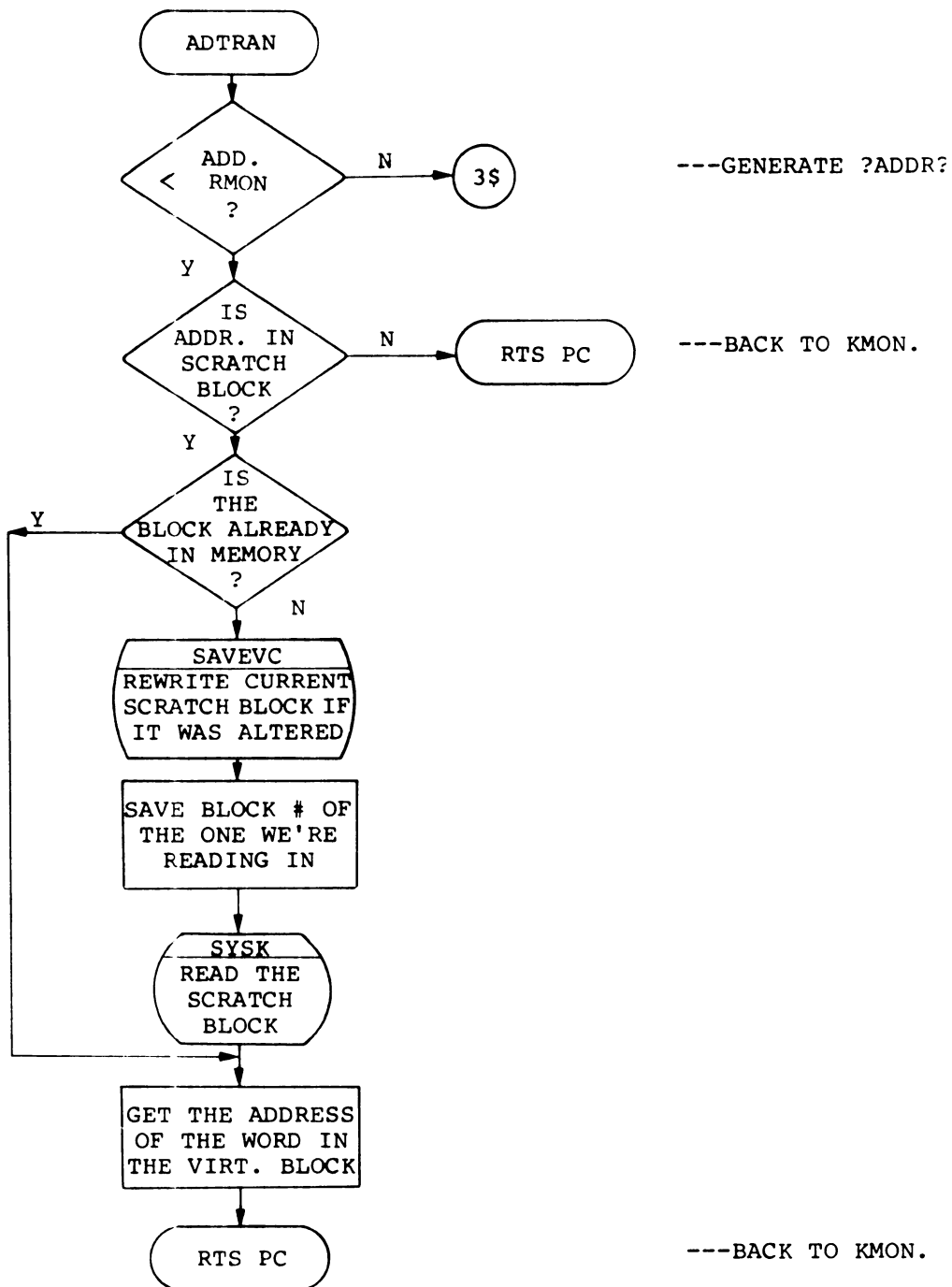
OVREAD - Used to read overlay command processors into memory.



OVLINK - Called from overlay processors to allow linking from one overlay to the other.

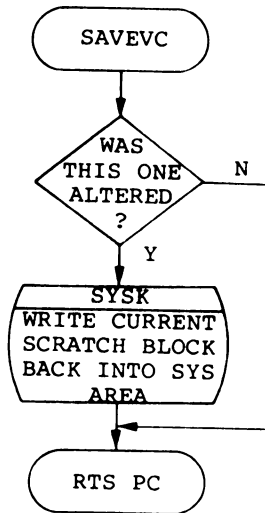


ADTRAN - Used to determine if a user-typed address is a) legal (i.e., address of RMON), b) in scratch blocks on system device.

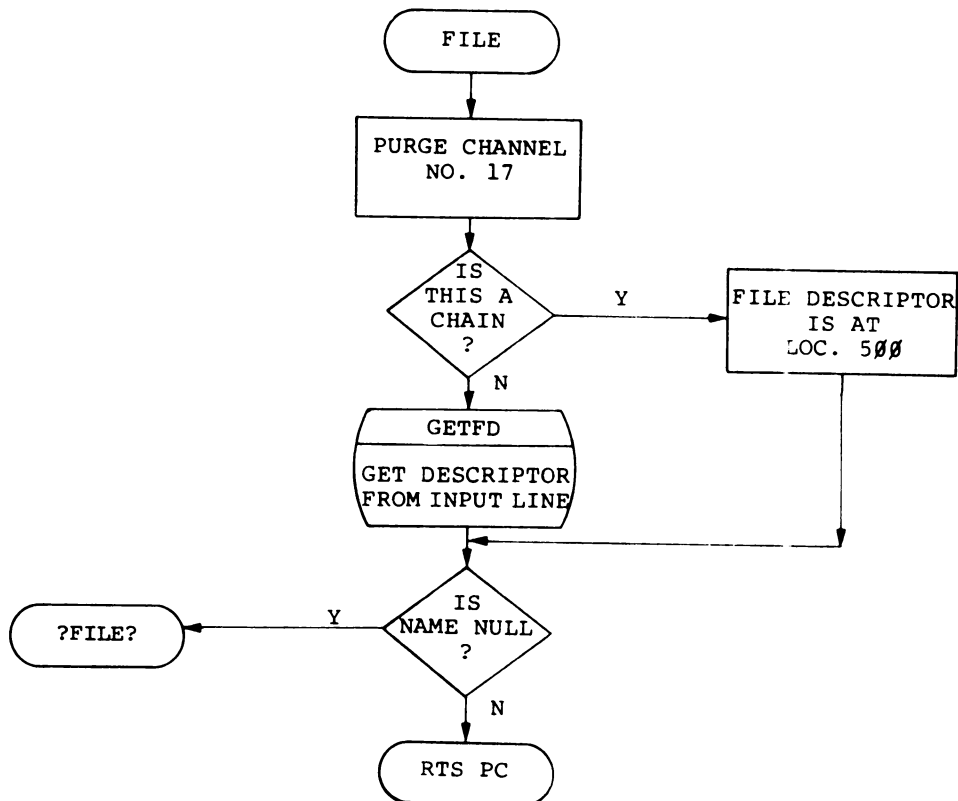


SAVEVC/FILE

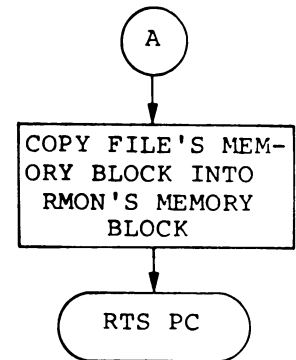
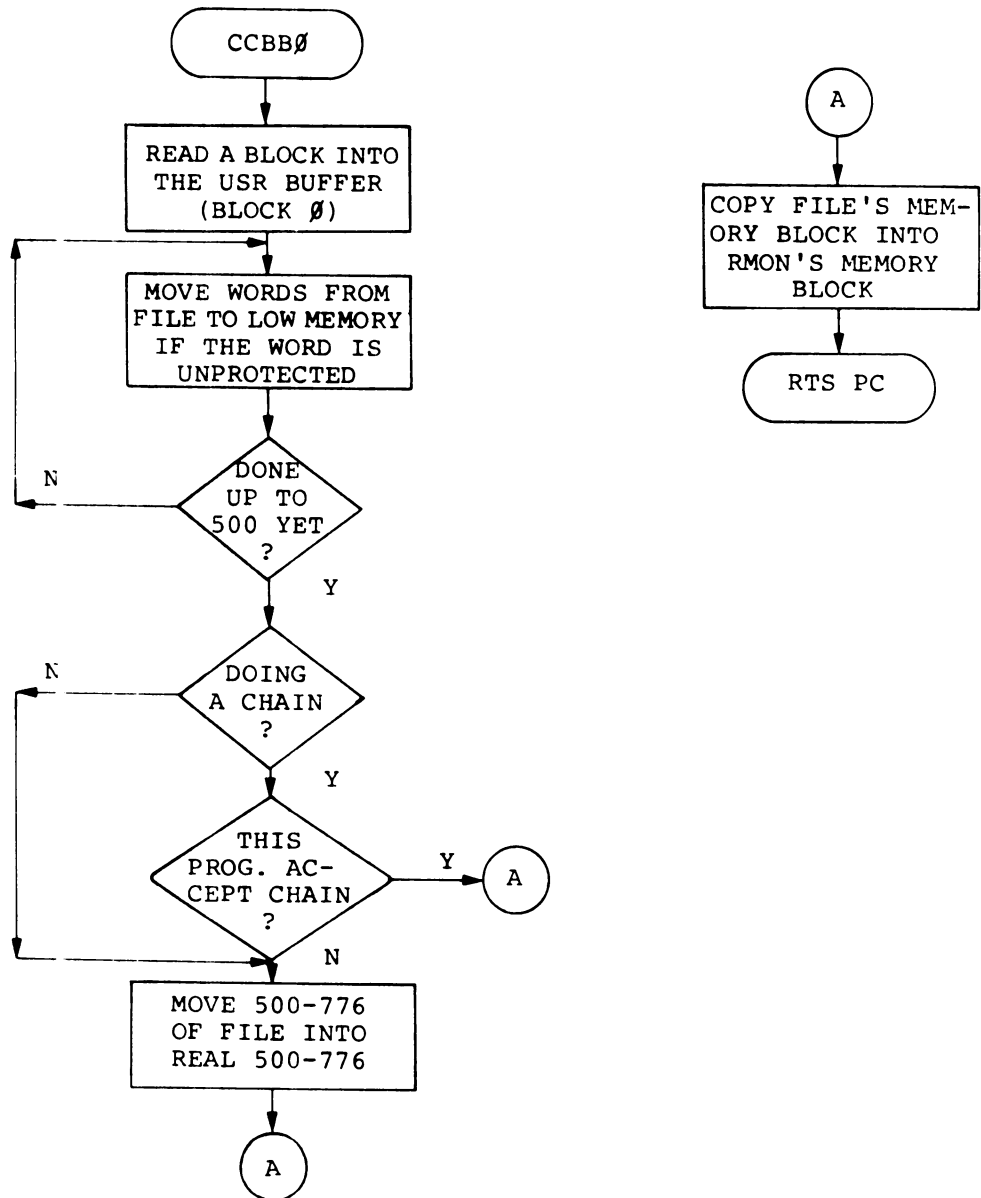
SAVEVC - Rewrites a block of memory back to the system scratch area if the block's contents were altered with a Deposit.



FILE - Called to pick up the .RAD50 representation of DEV:FILE.EXT. It will assume a default extension of .SAV.

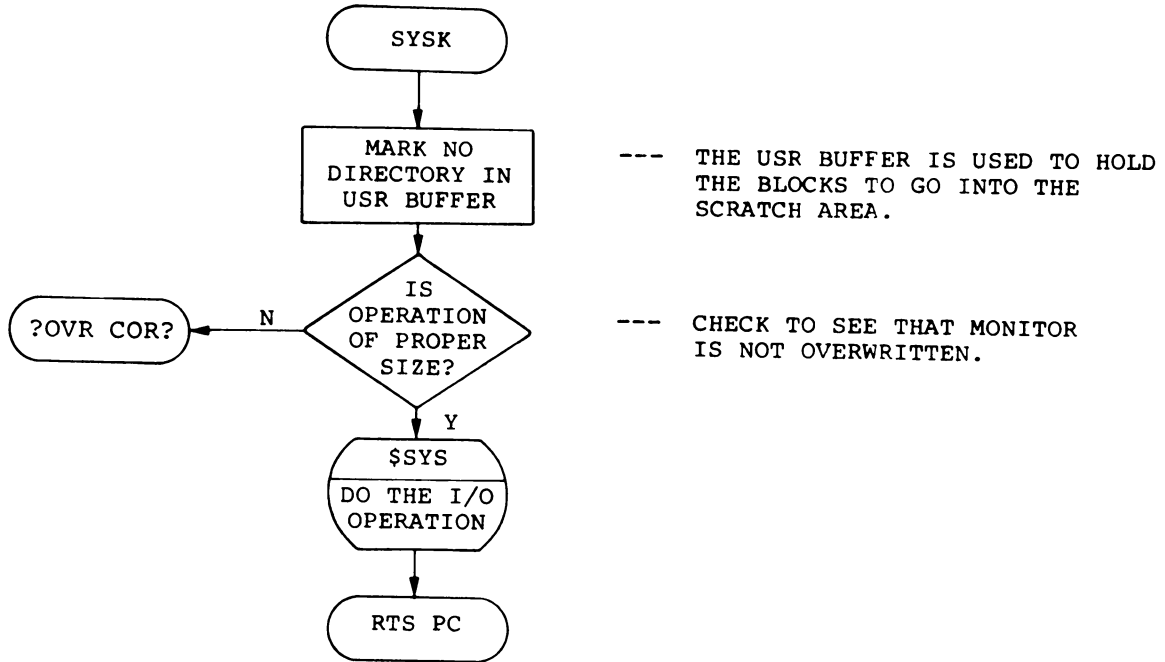


CCBBØ - The CCBBØ routine reads the first block of a .SAV file into the USR buffer, then moves selected locations from that block into the corresponding physical memory locations. The words moved are those marked with Ø's in the RMON bitmap. This procedure protects the system from having its vectors overlaid. If a chain is being done to a program which does not accept a CHAIN, 5ØØ-776 will be loaded with the contents of the file.



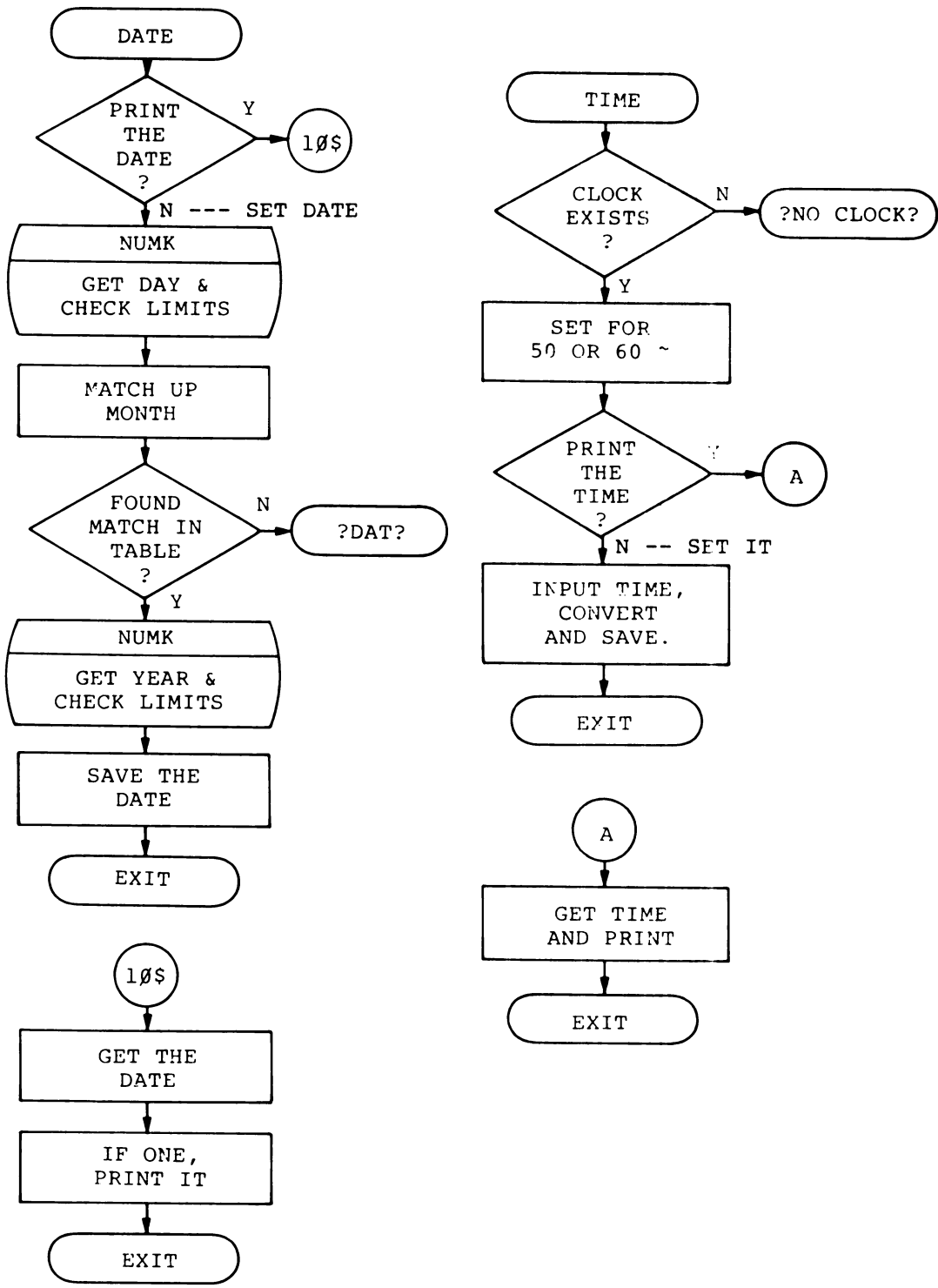
SYSK

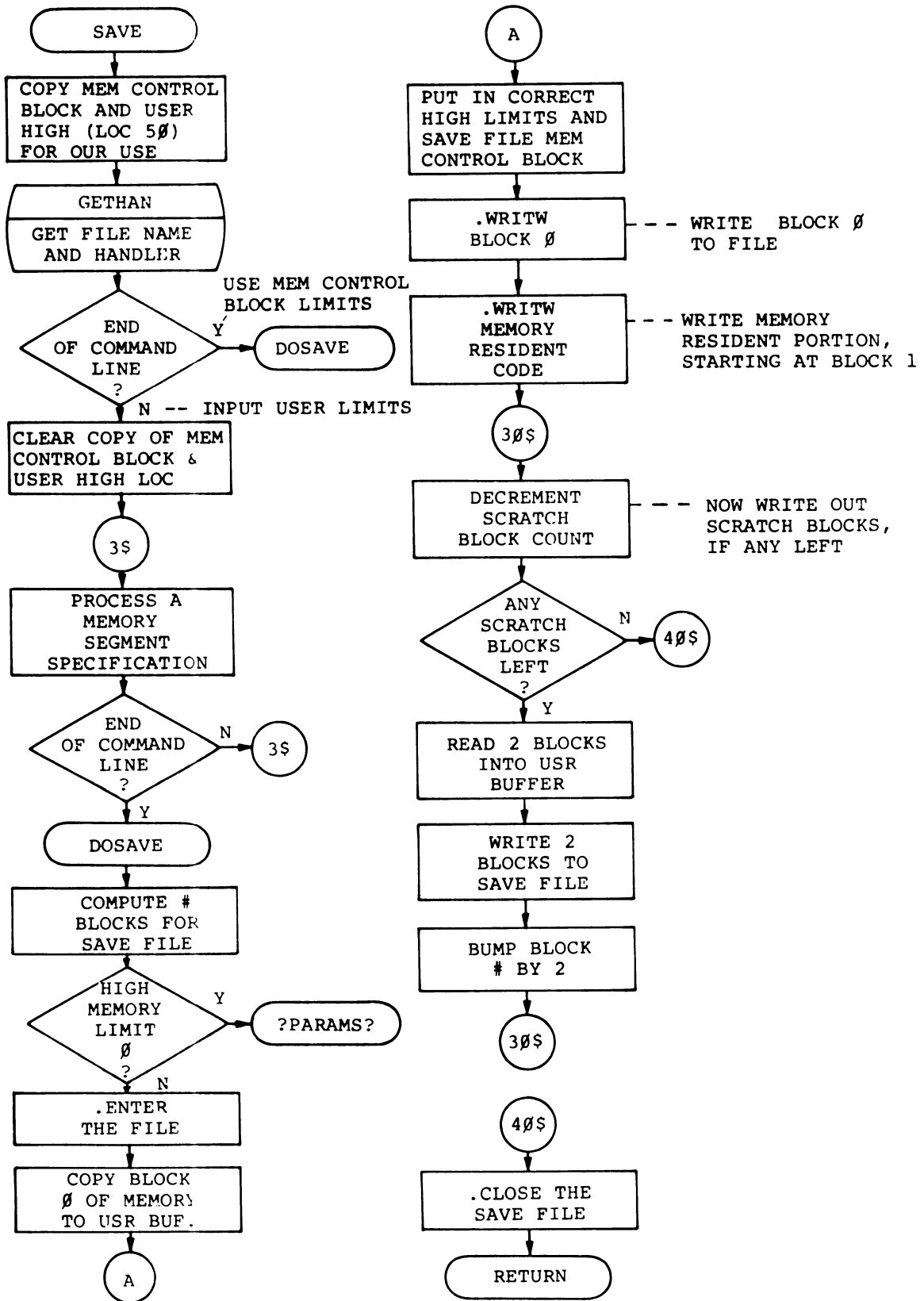
SYSK - Used to read/write blocks into and out of the system scratch area.



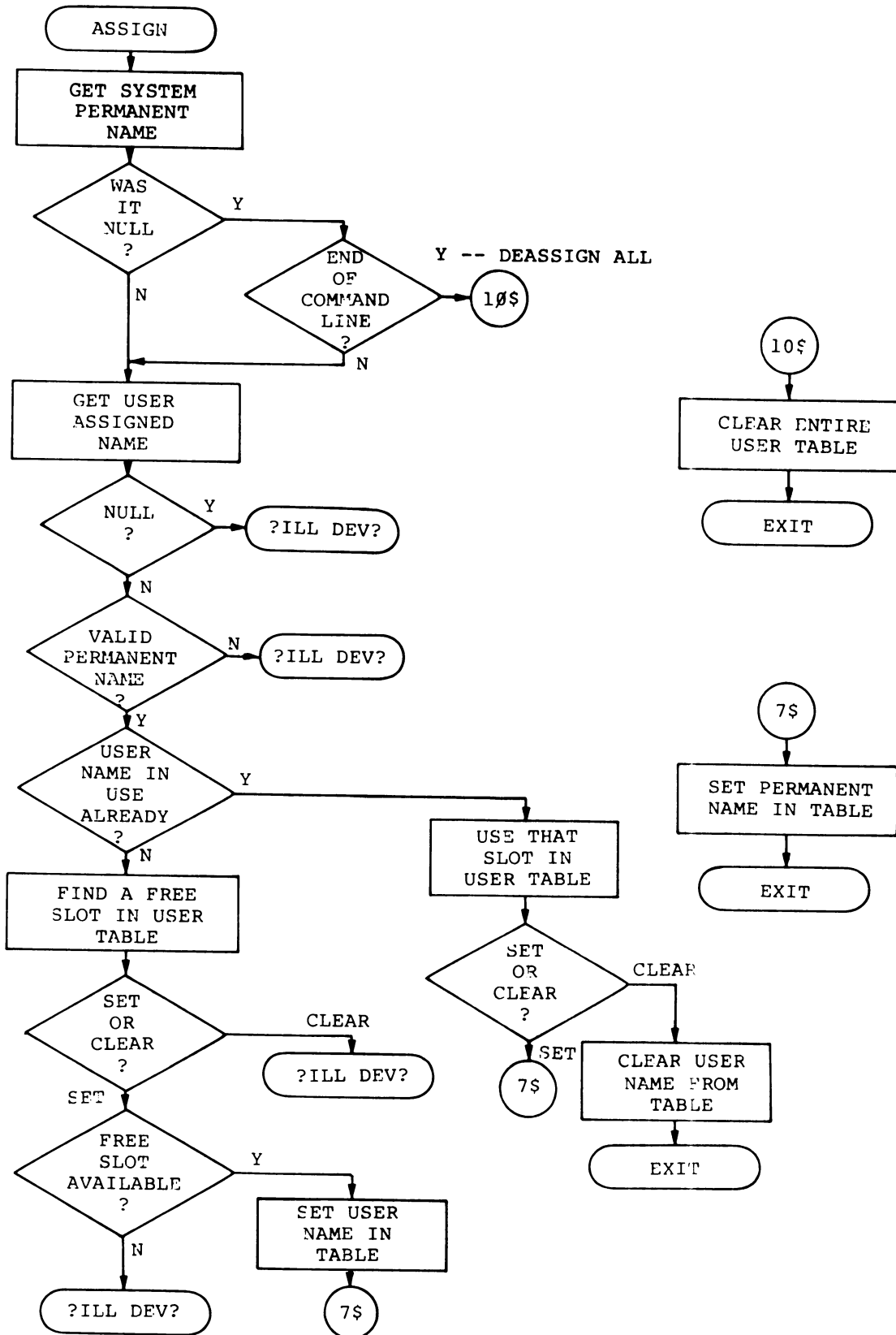
E.1.2 KMON Overlays

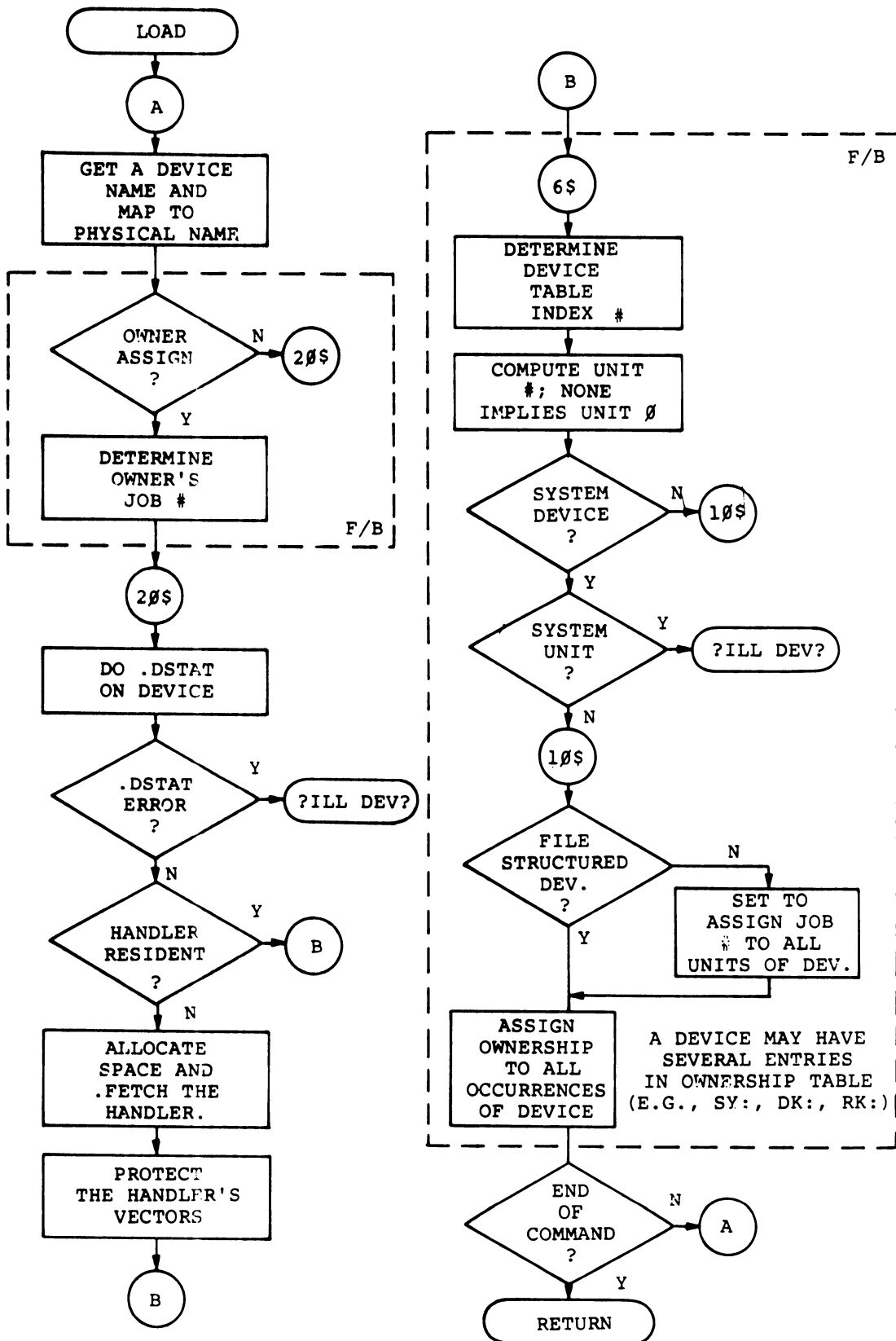
DATE/TIME



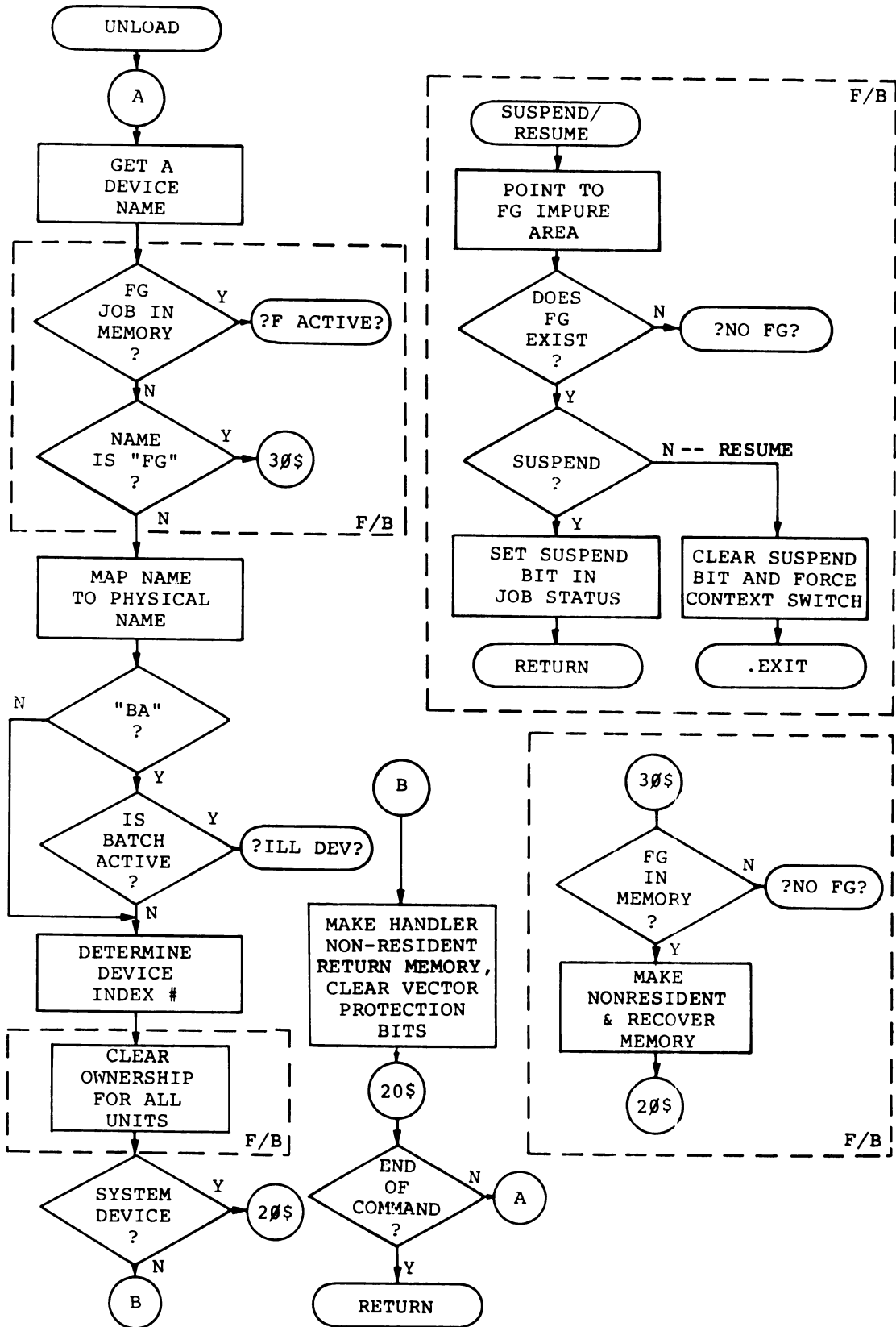


ASSIGN

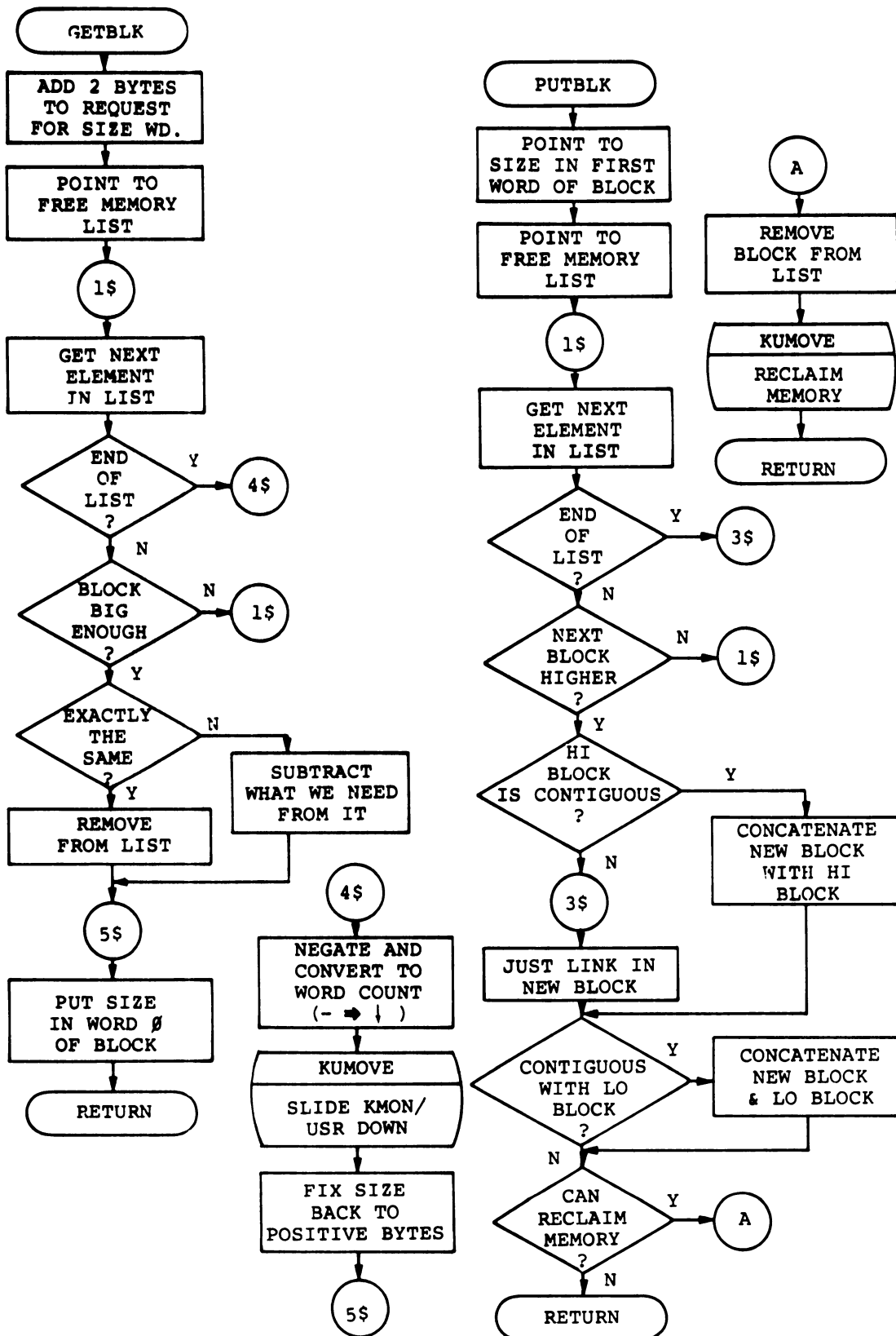




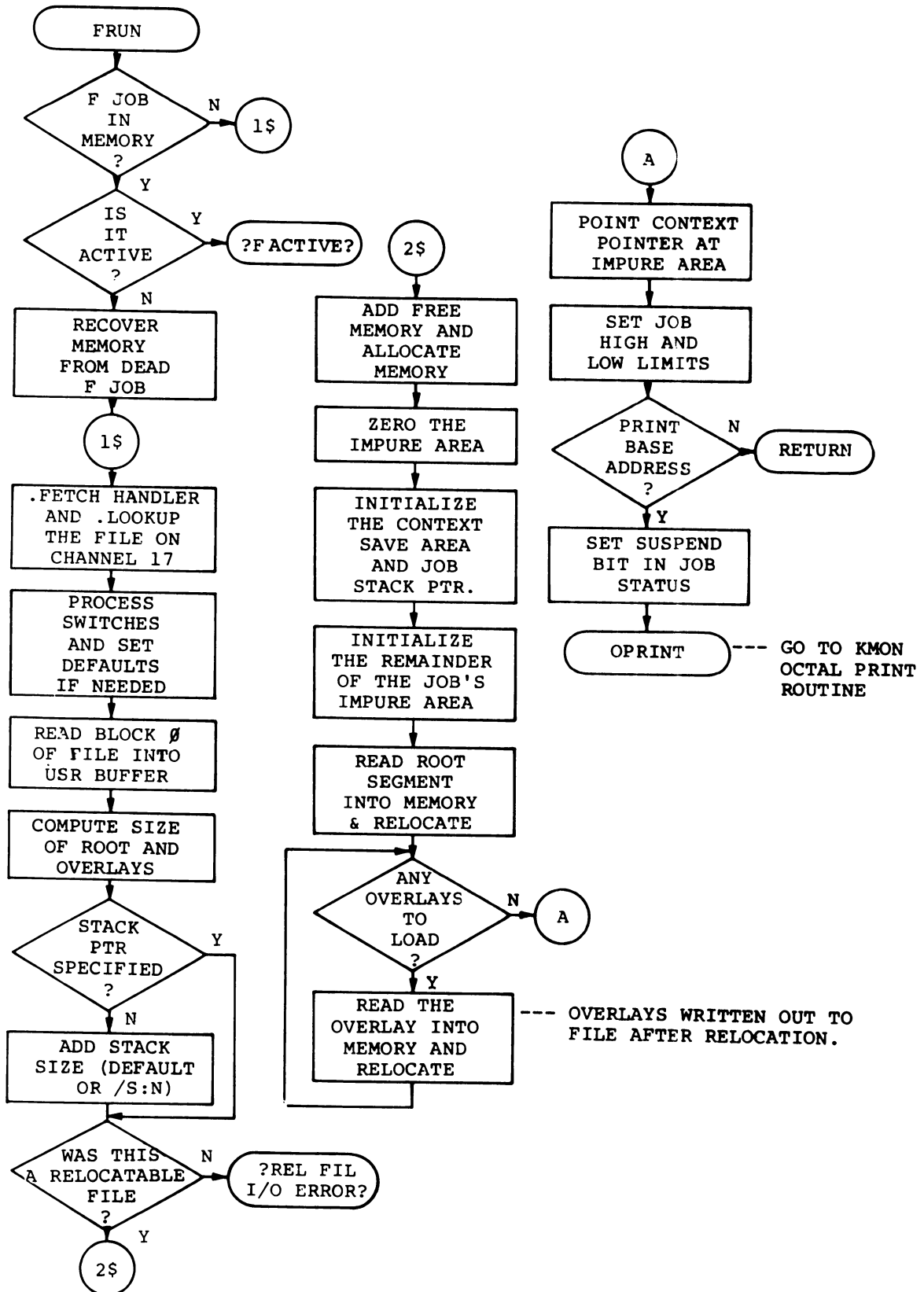
UNLOAD/SUSPEND/RESUME

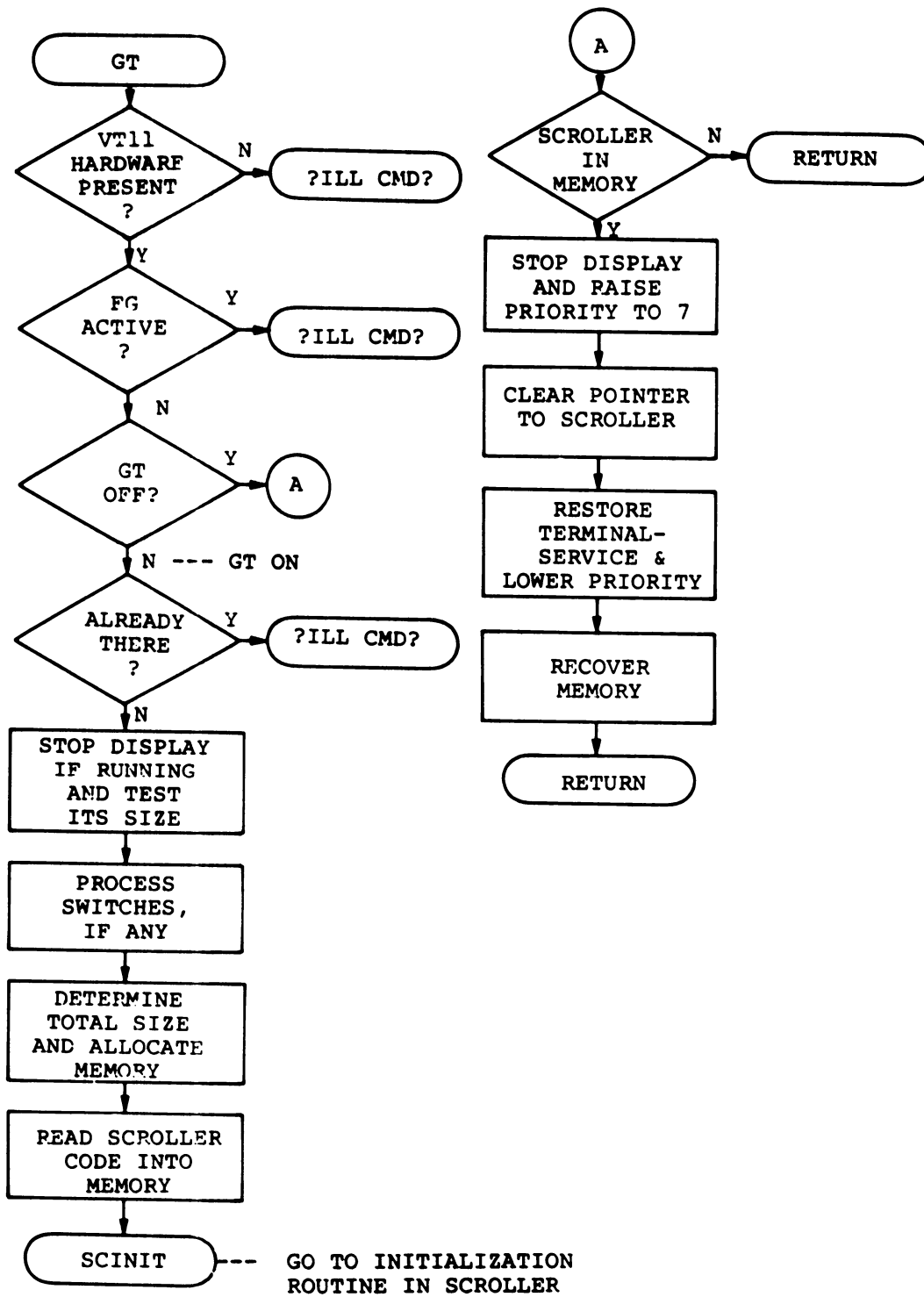


GET/PUT A BLOCK OF MEMORY



FRUN





E.2 USR (USER SERVICE ROUTINES) FLOWCHARTS

USRBUF/FATAL/CDFN

The first 2 blocks of the USR are used by the USR for directory operations. They are also used by the KMON at various points for a 2-block general purpose buffer. There is, however, executable code in the buffer that can be executed every time a fresh copy of the USR is read from the system device. The functions included in the buffer are:

1. USR Relocation

This code is executed whenever the USR is newly read into memory. It serves to make certain pointers into RMON absolute.

2. Fatal error processor and fatal error messages (S/J only)

3. CDFN (channel define) EMT (S/J only)

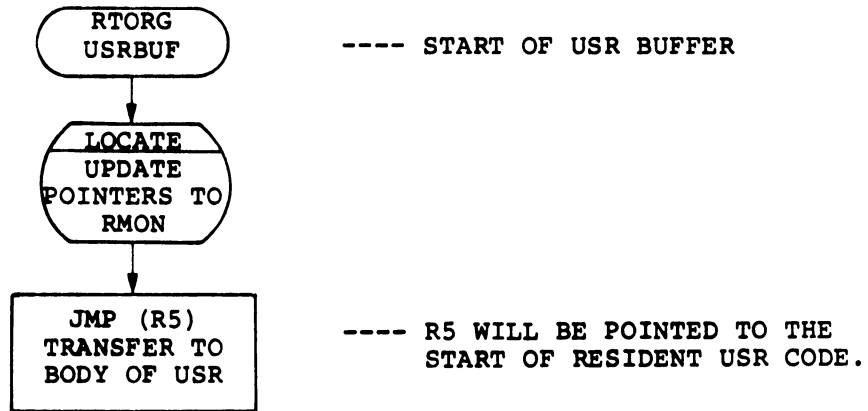
The CDFN EMT call forces a new copy of the USR into memory to guarantee the presence of the EMT processor.

The flows for these functions follow.

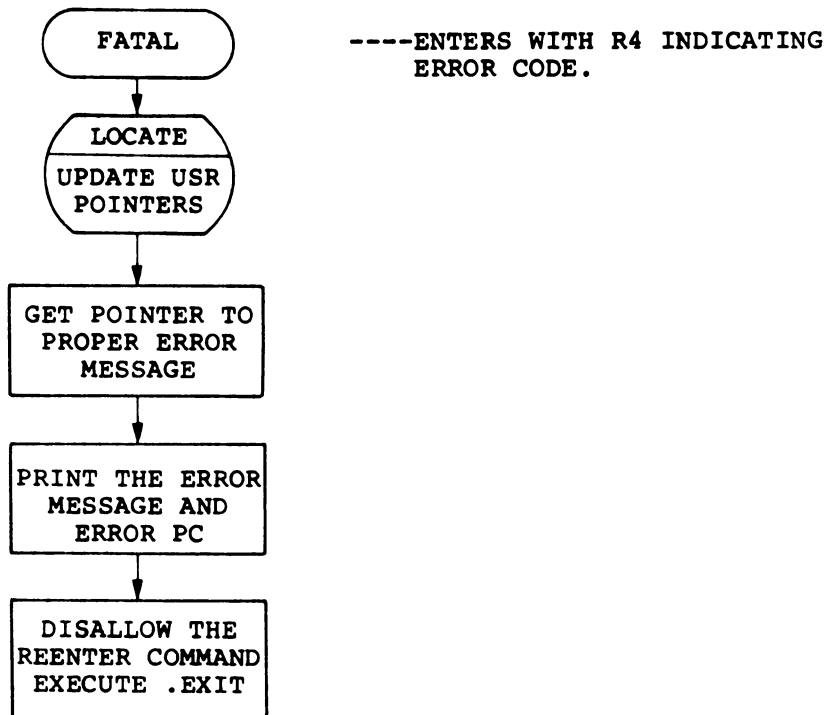
NOTE

Fatal error handler and CDFN processor are RMON functions in the F/B Monitor. The only code in the buffer in the F/B system will be the USR relocation code.

USRBUF is the initial entry point for USR calls when the USR has just been read into memory. LOCATE sets up pointers into RMON.

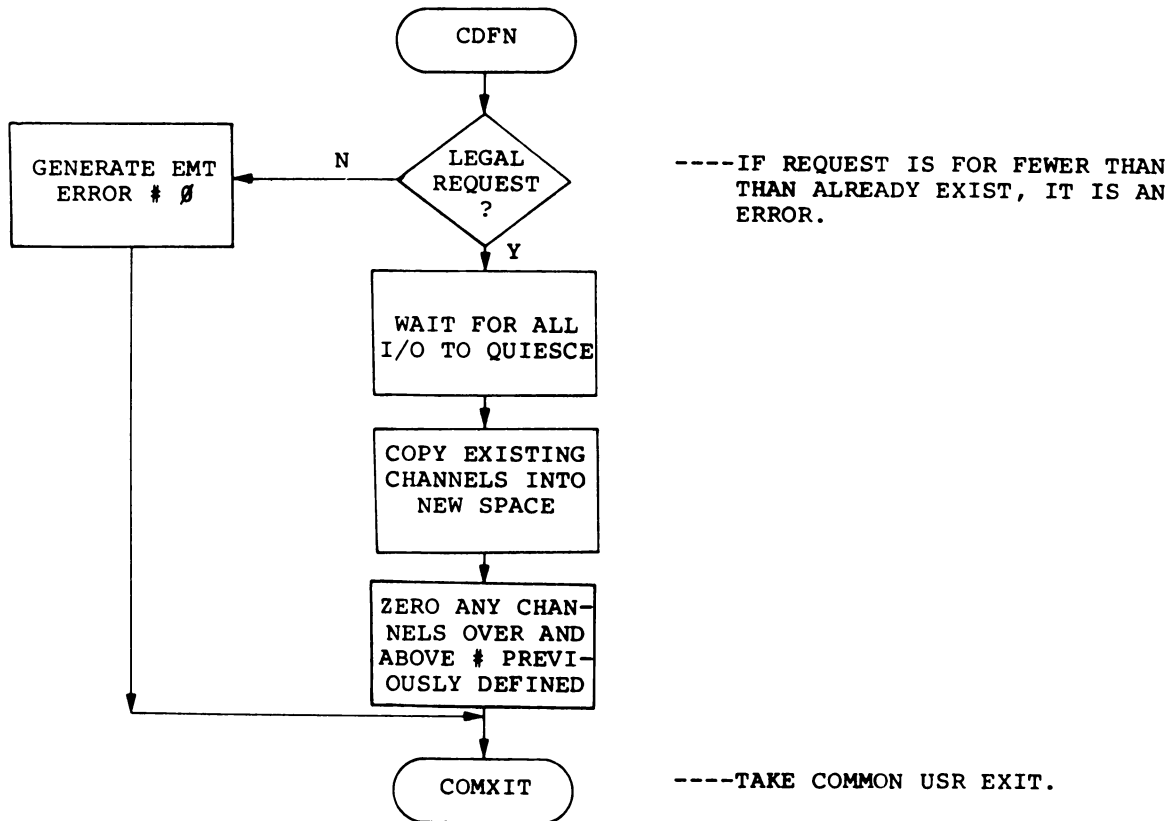


The LOCATE routine is called to update the list of pointers at RELIST. The list is initially a list of address differences (i.e., VALUE-\$RMON where VALUE is the desired location and \$RMON is the address of the start of RMON). LOCATE then makes all the differences into absolute addresses. Any errors which would generate a ?M-error use the FATAL error processor code to generate the message in the S/J system. This is a resident function in F/B.



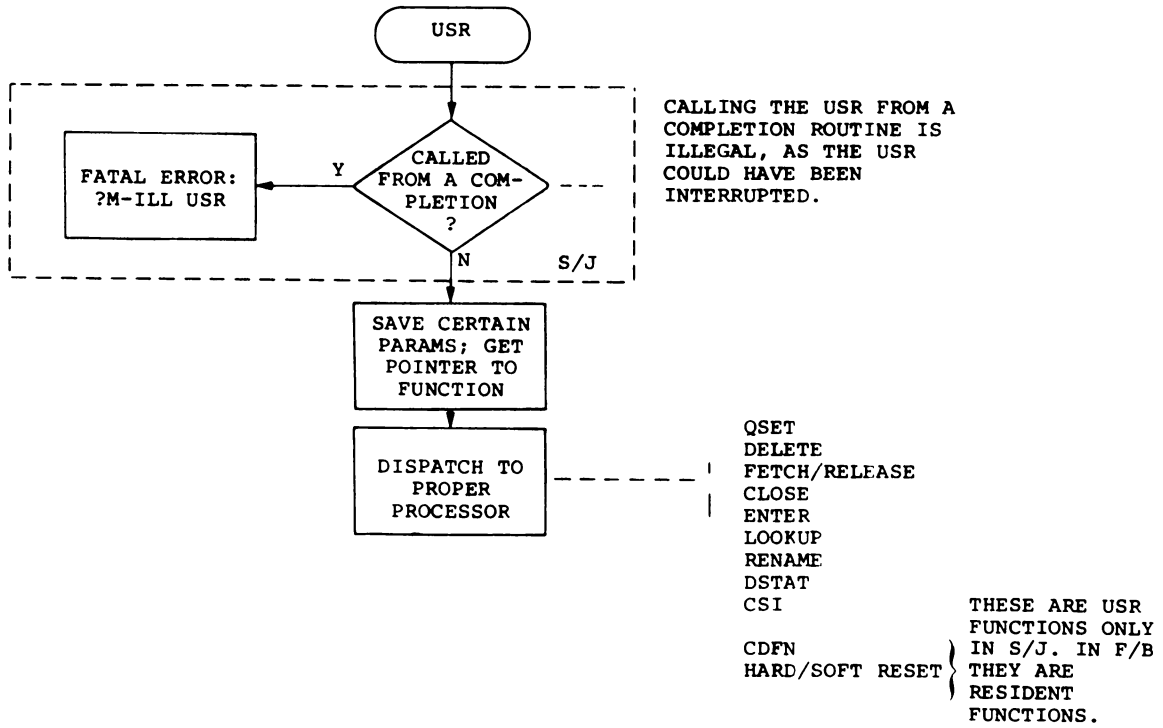
USRBUF/FATAL/CDFN (CONT.)

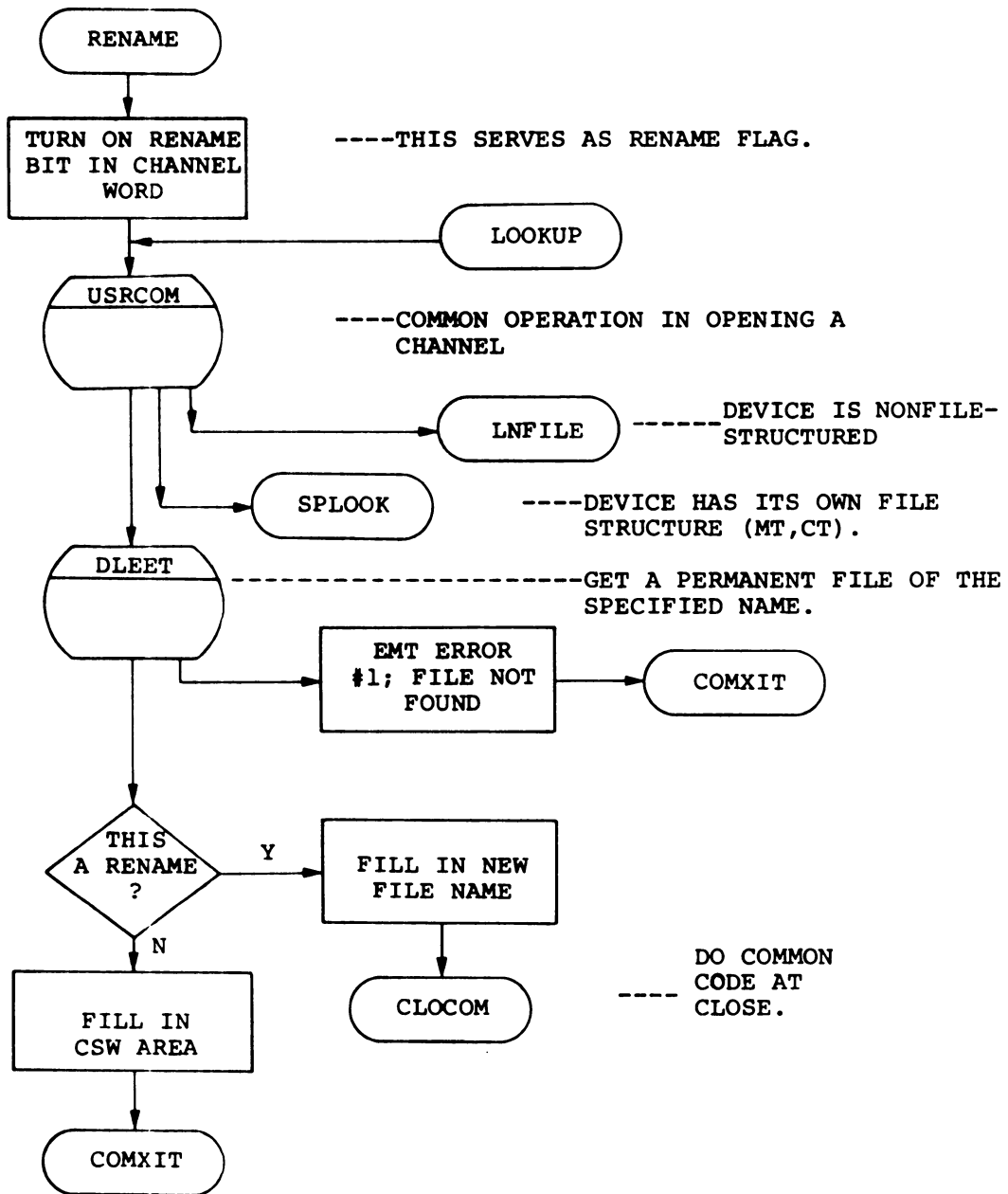
CDFN - A resident function in the F/B system.



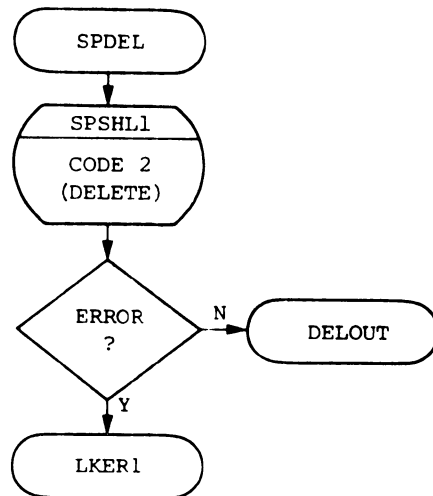
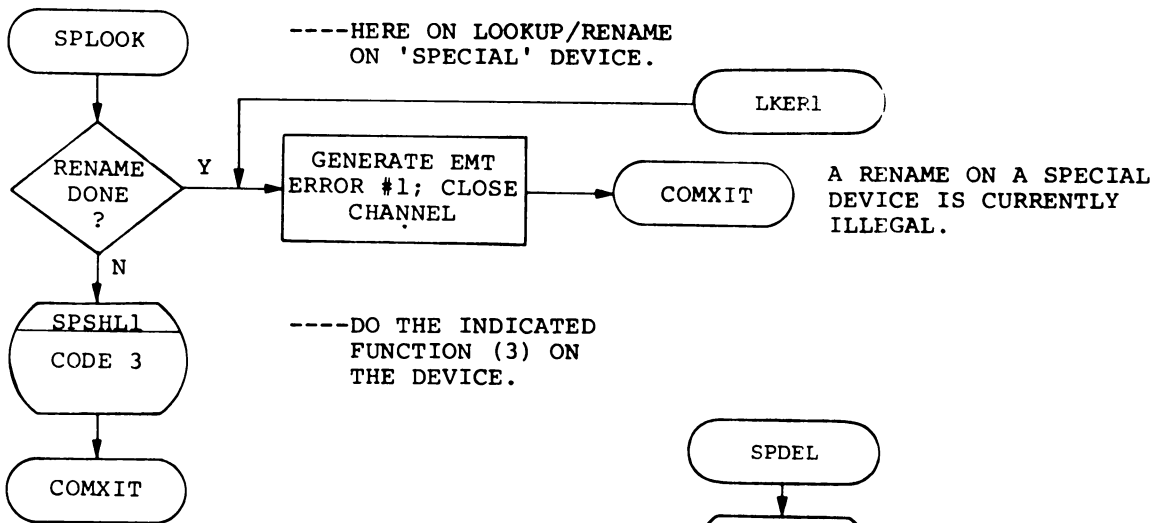
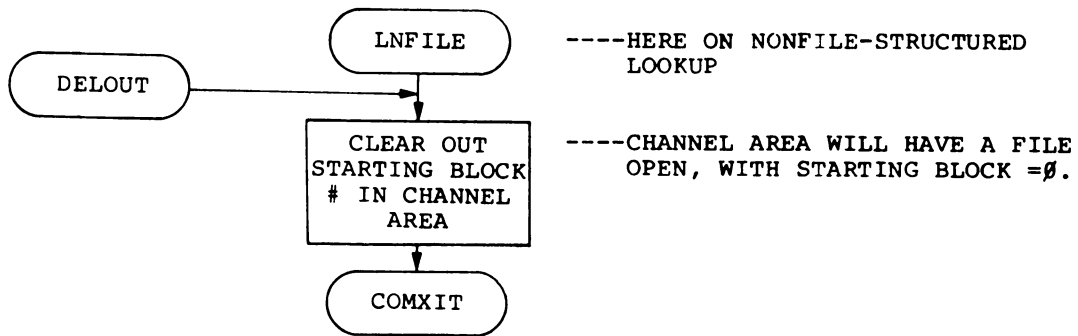
The following flowcharts detail the code contained in the main body of the USR. On entry to the USR, R2 contains an index representing the function to be performed. This is used to dispatch control to the proper processor.

USR CODE



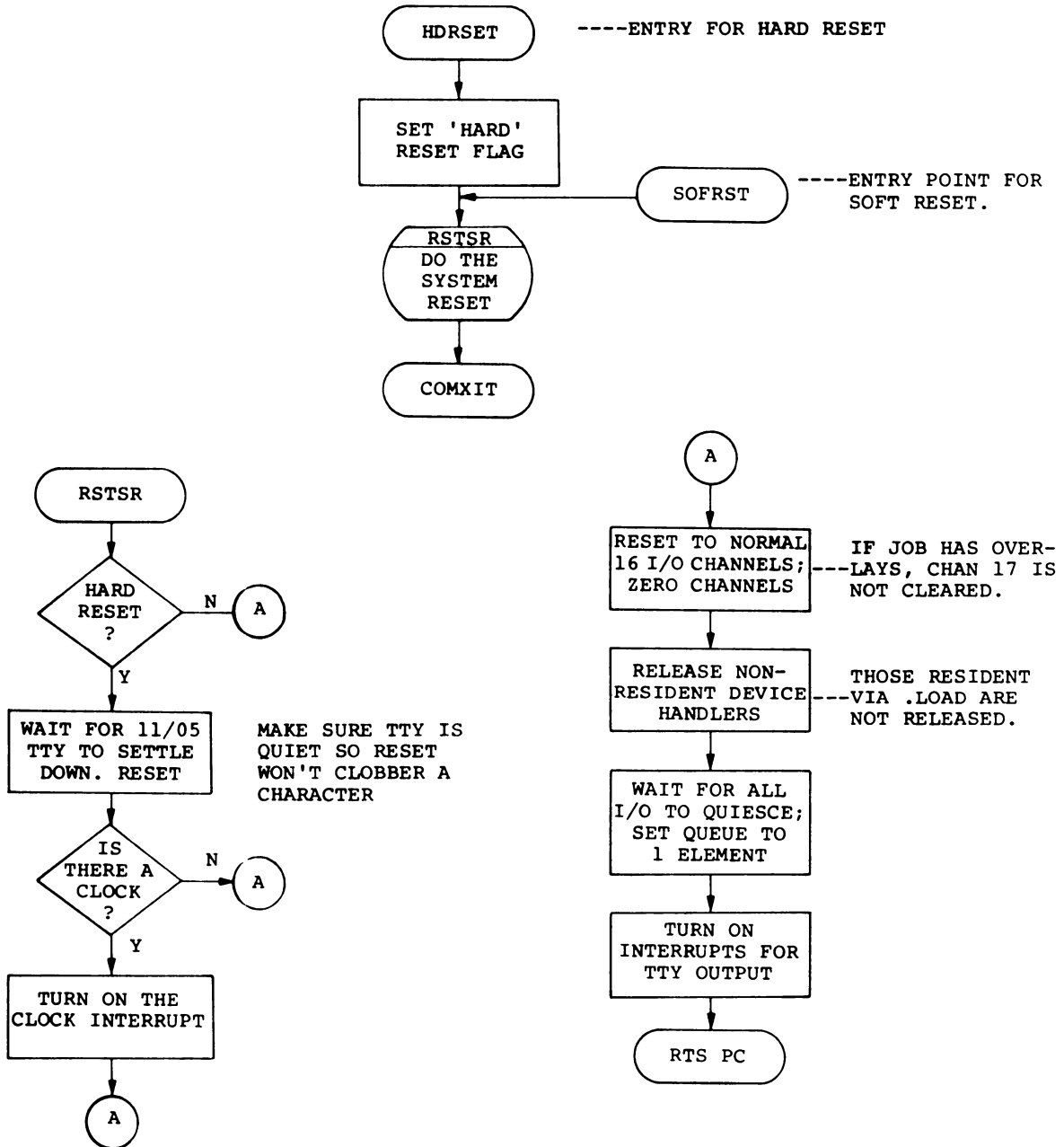


LOOKUP/RENAME (CONT.)

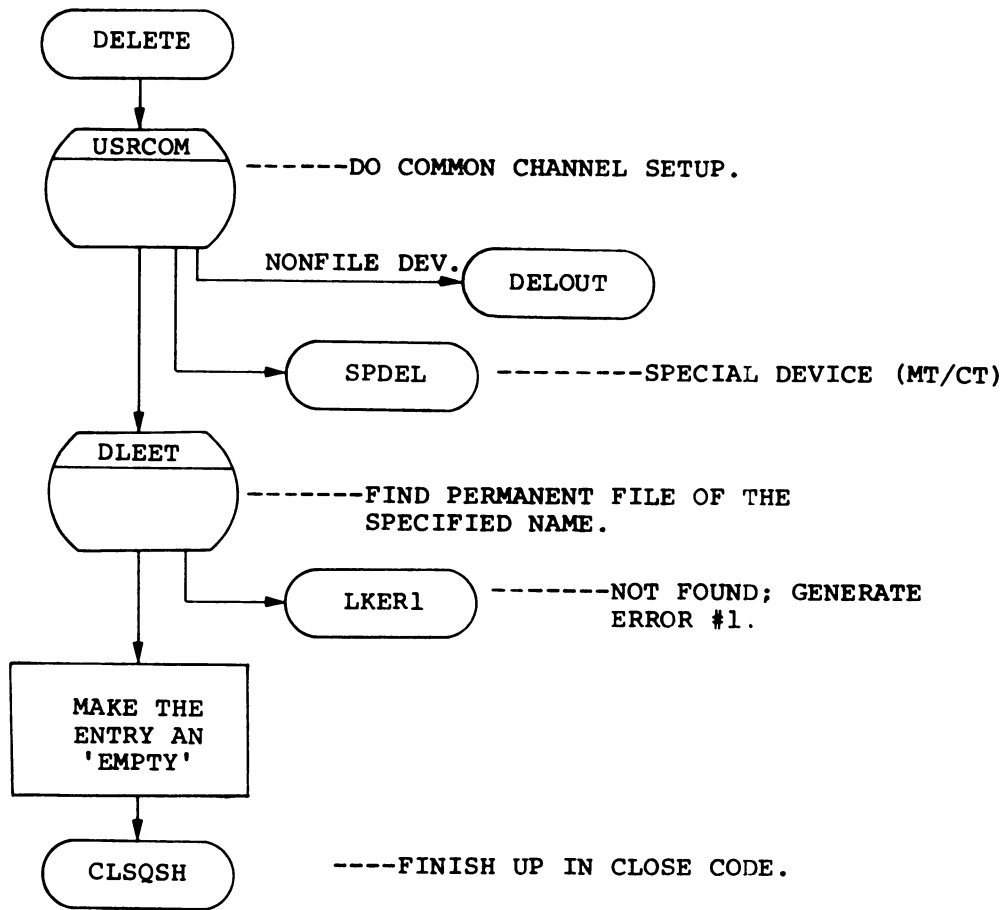


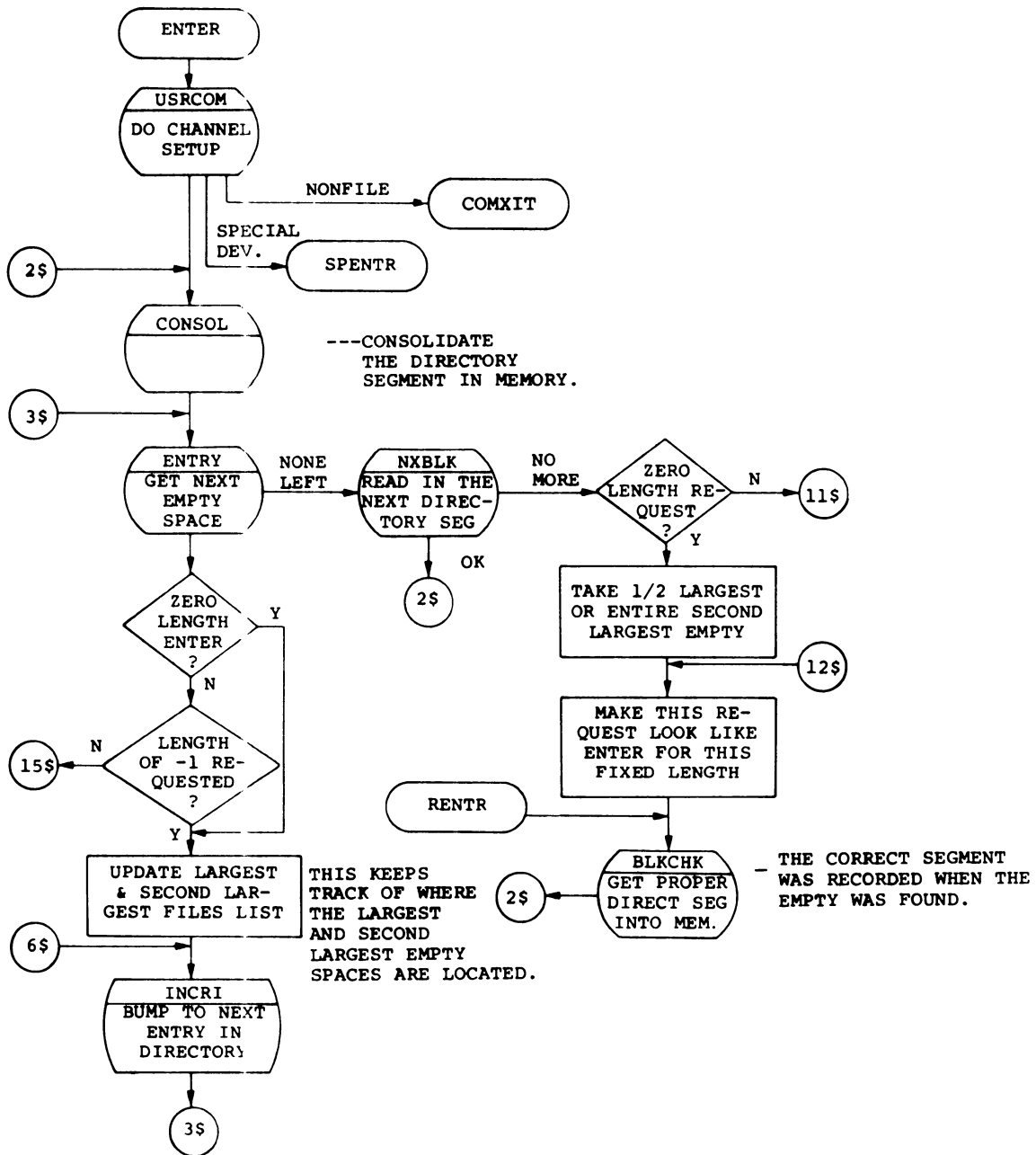
HARD/SOFT RESET

These are resident functions in F/B; USR functions in S/J.

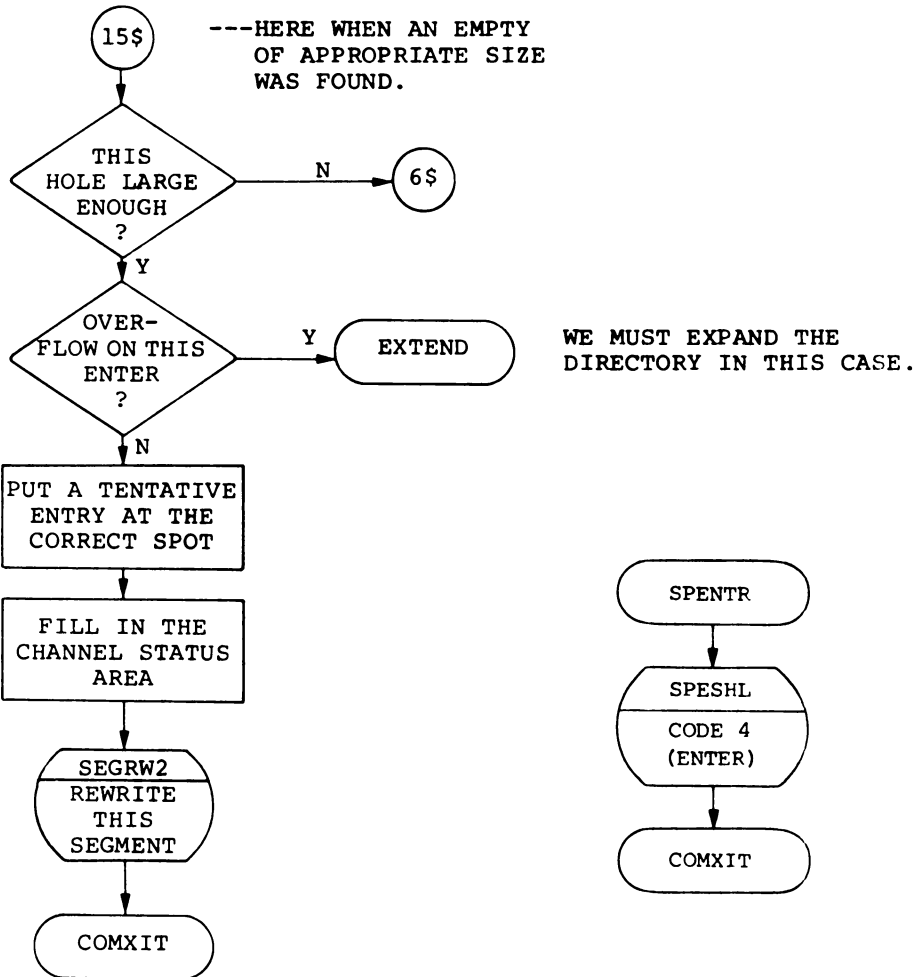
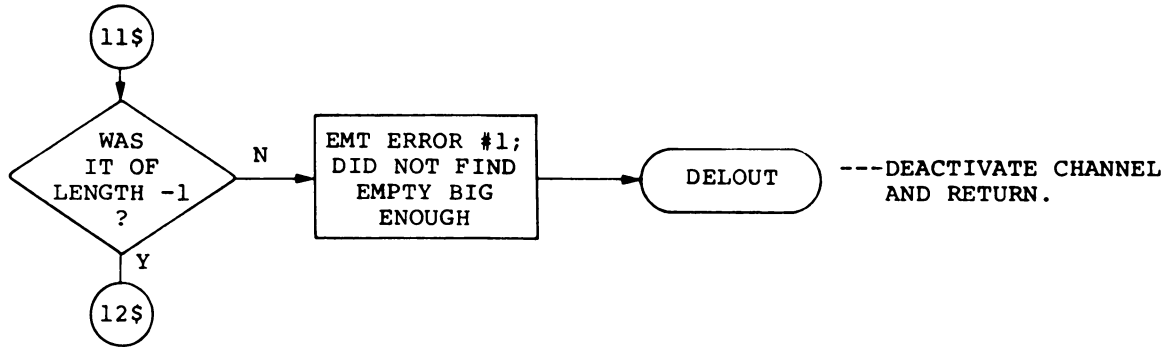


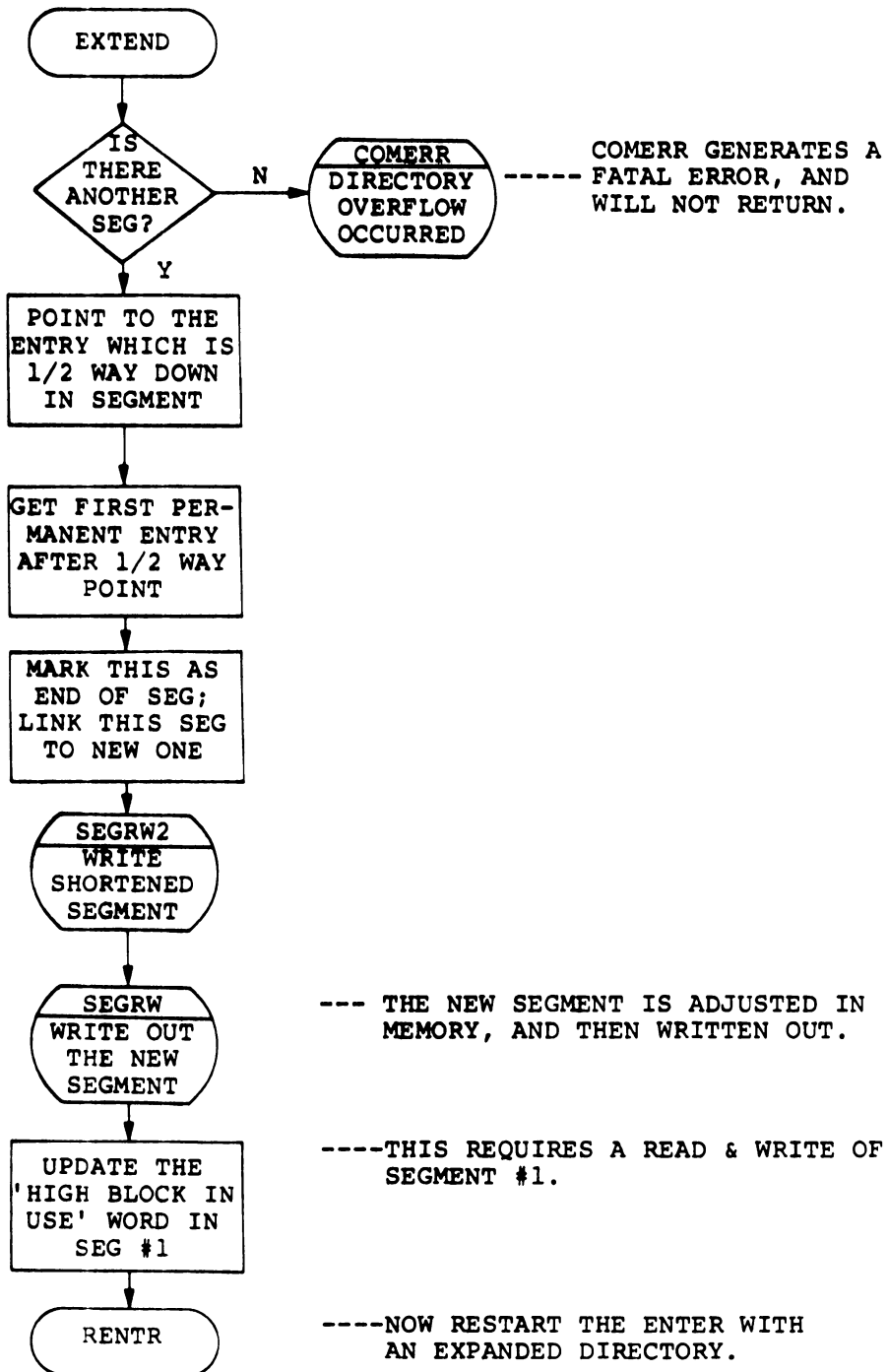
DELETE





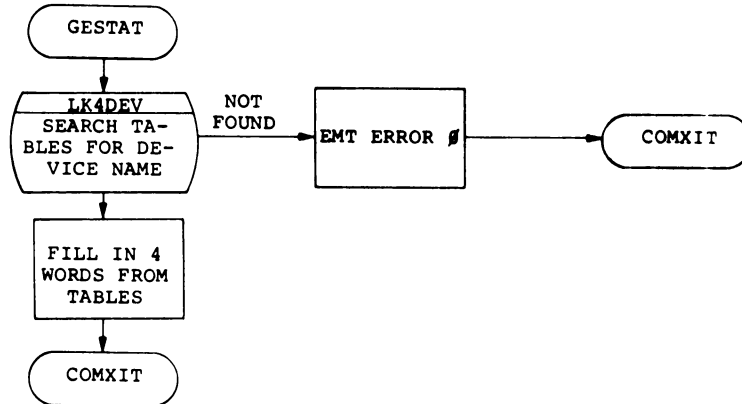
ENTER (CONT.)



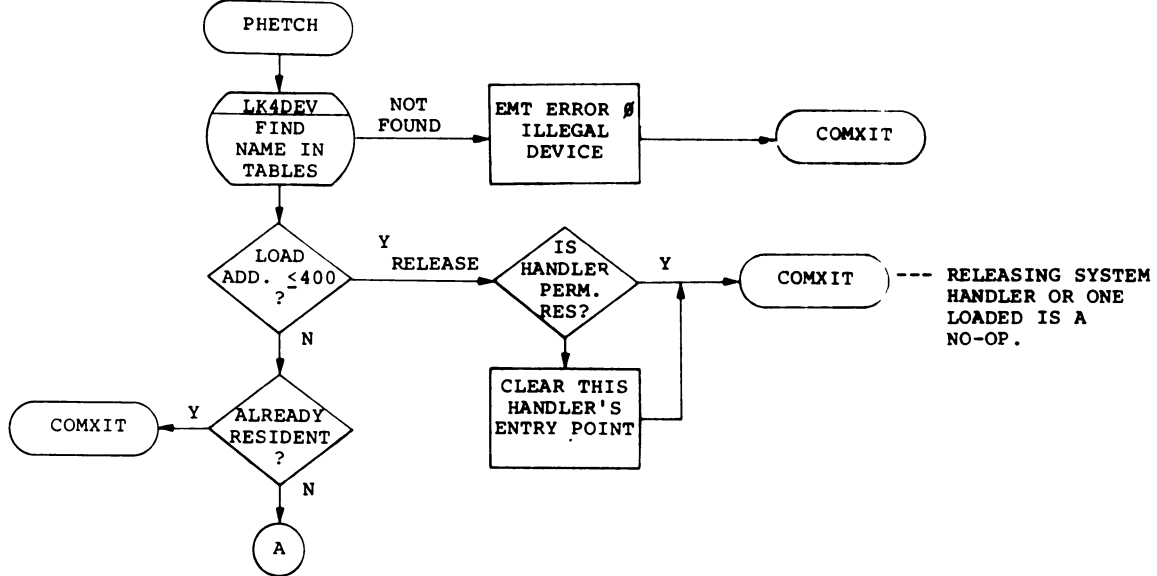


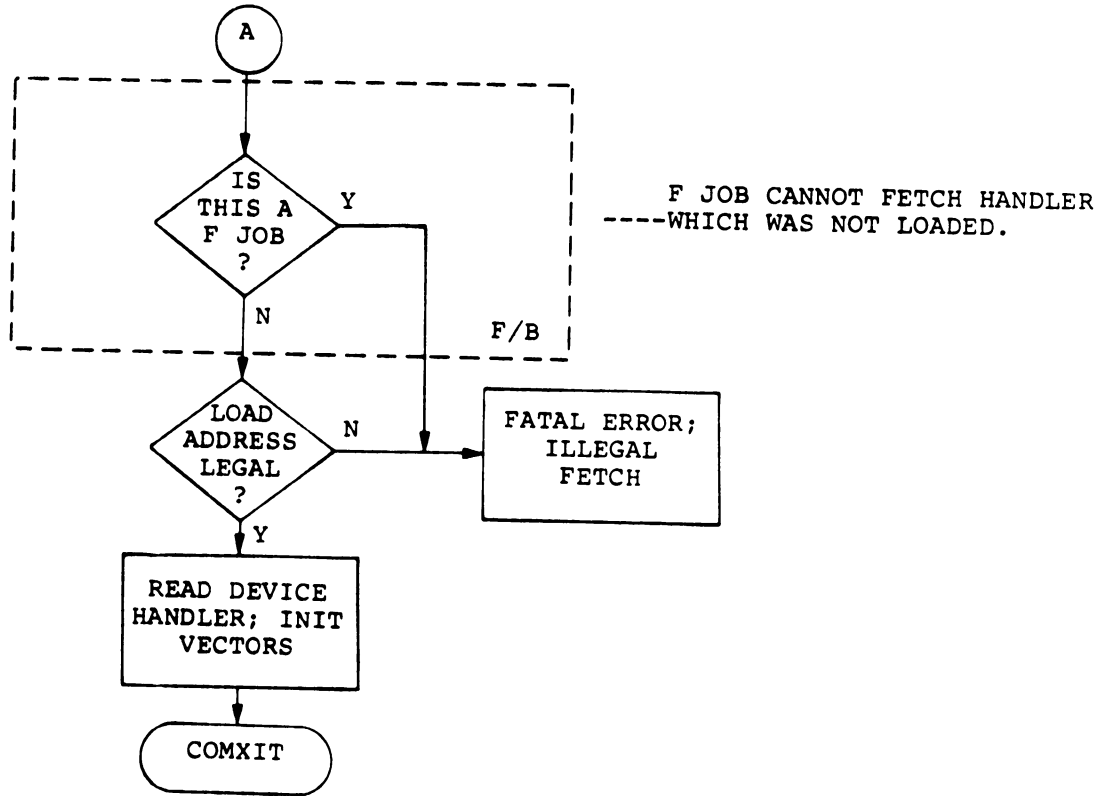
DSTAT/FETCH/RELEASE

DSTAT- GET DEVICE STATUS

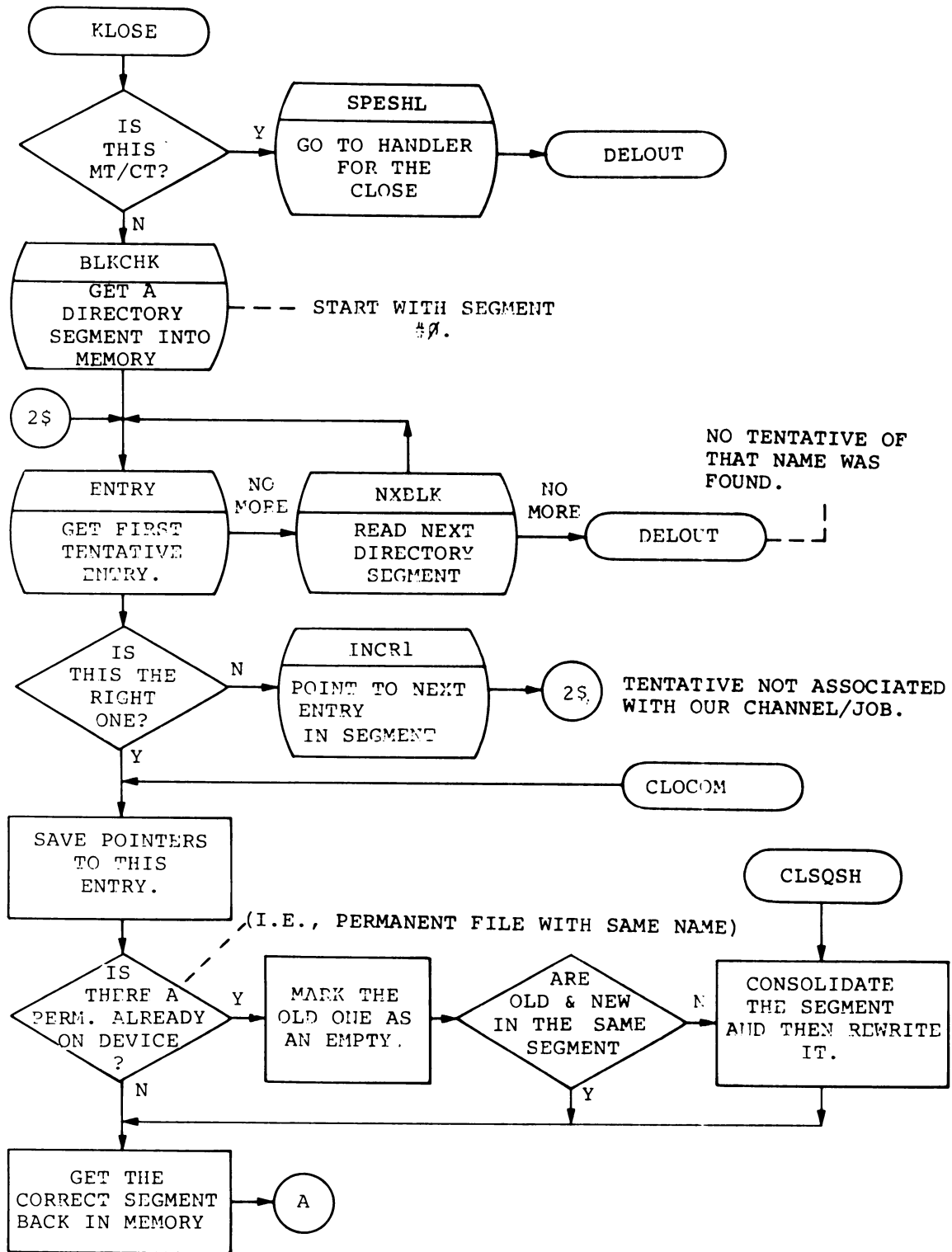


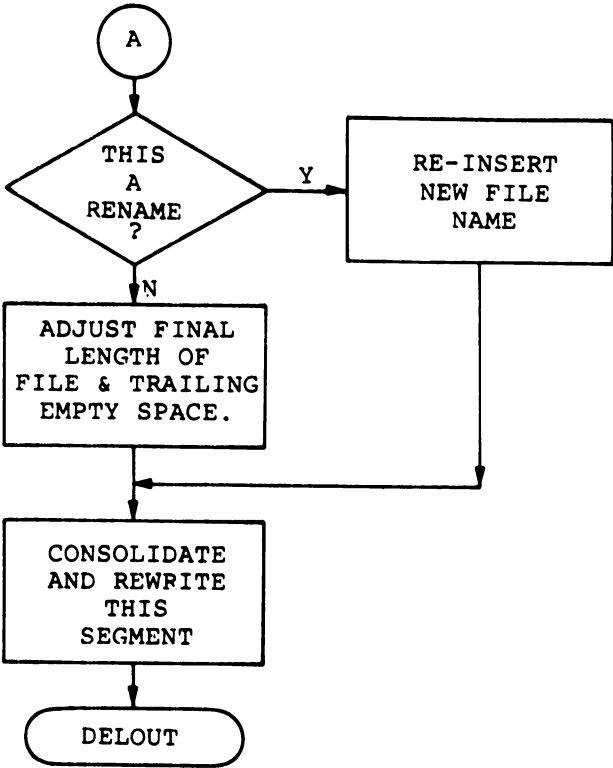
FETCH/RELEASE



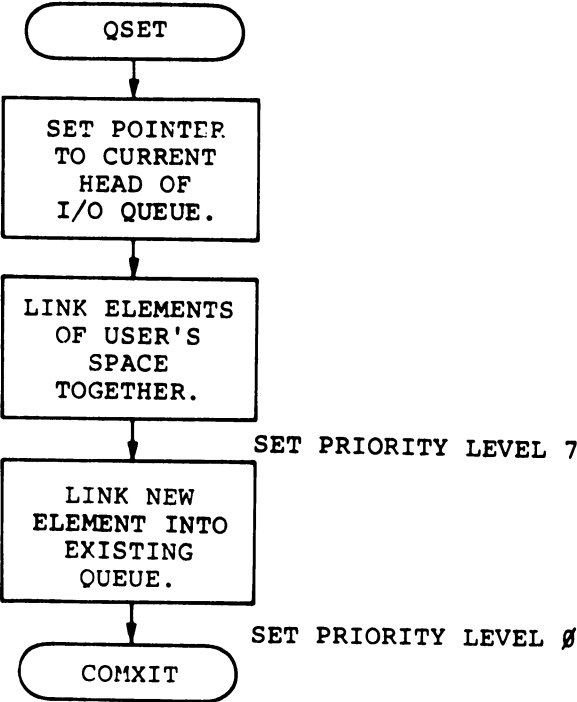


CLOSE



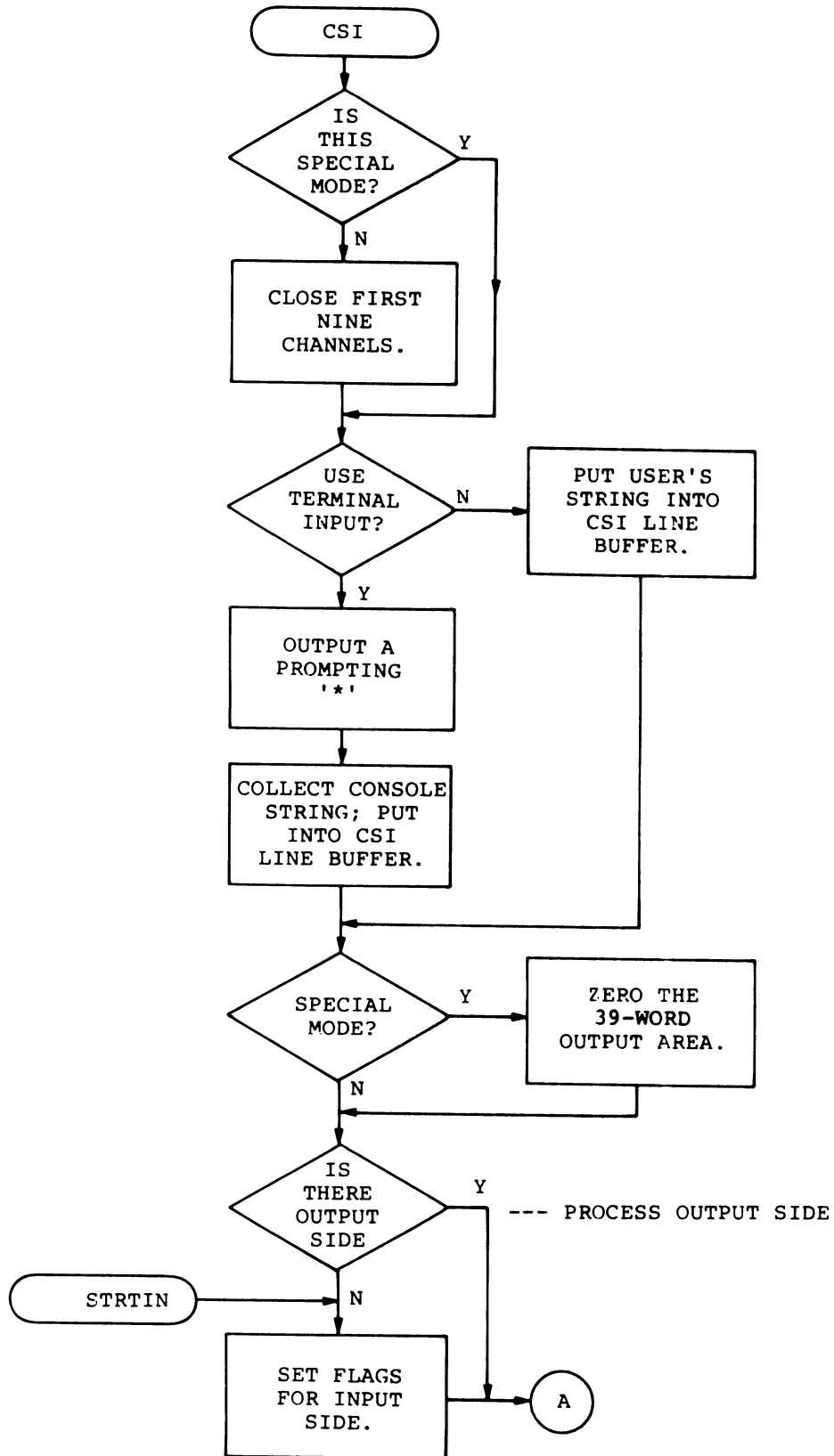


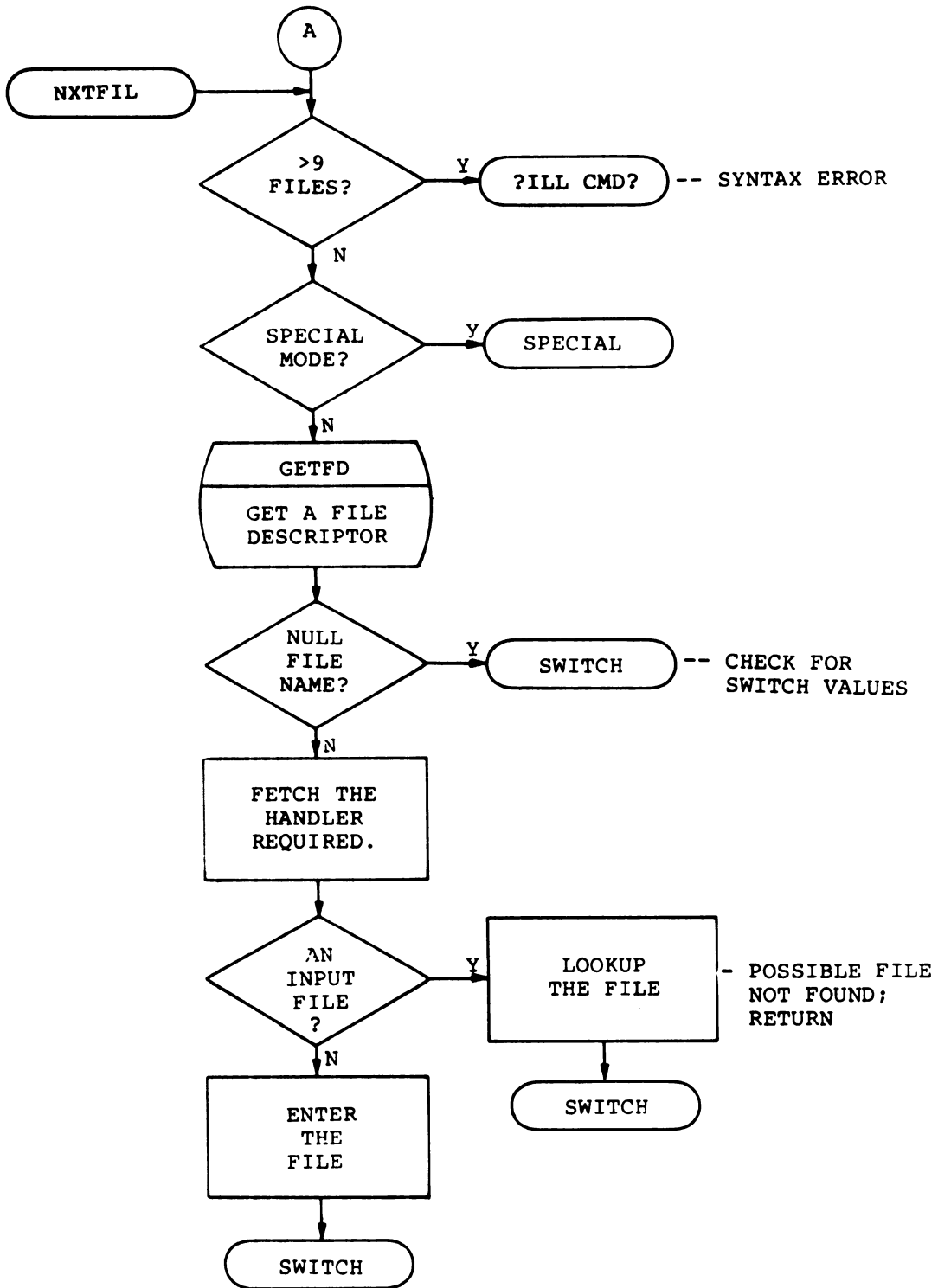
QUEUE EXTEND (QSET)



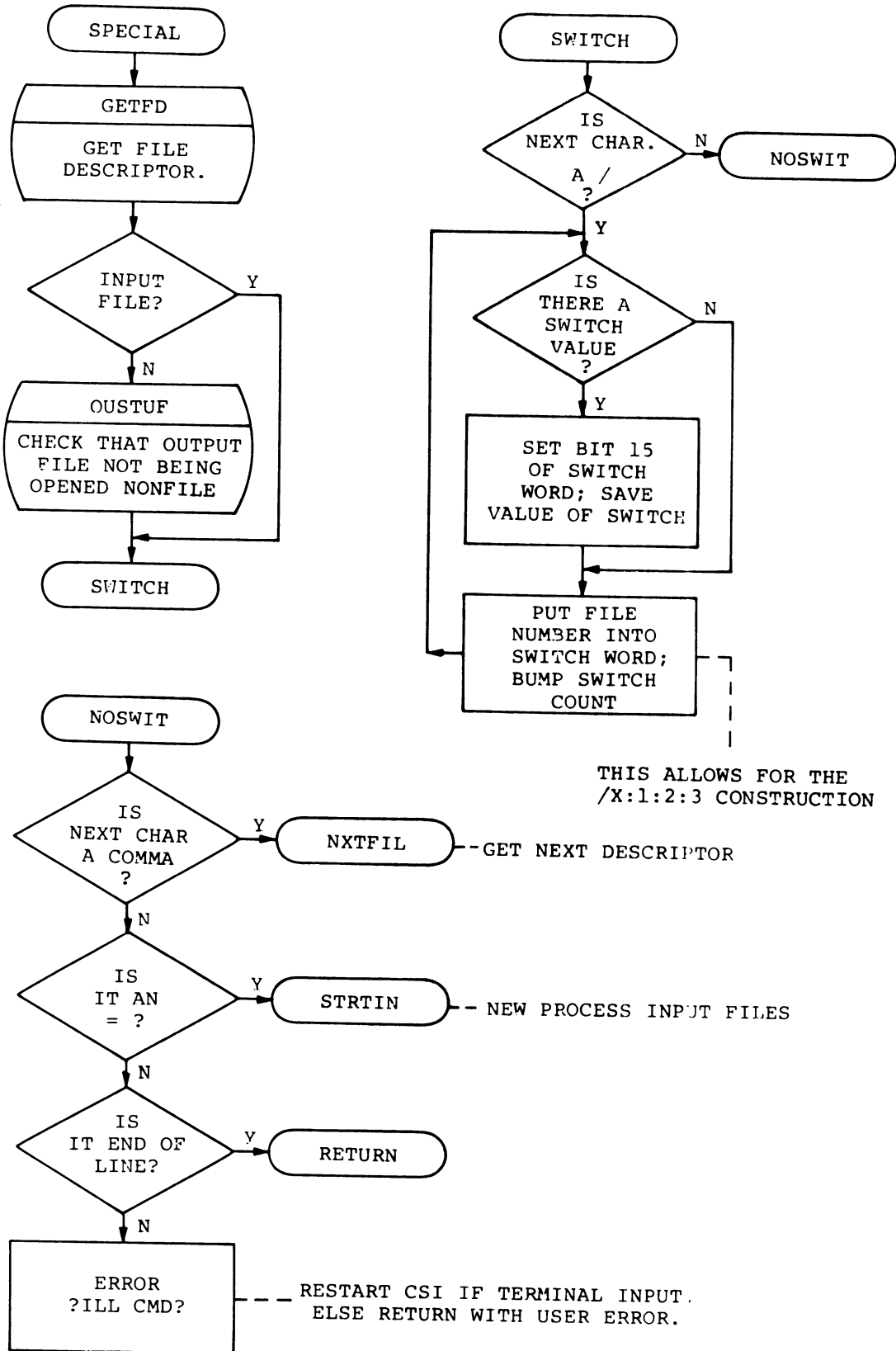
E.3 CSI (COMMAND STRING INTERPRETER) FLOWCHARTS

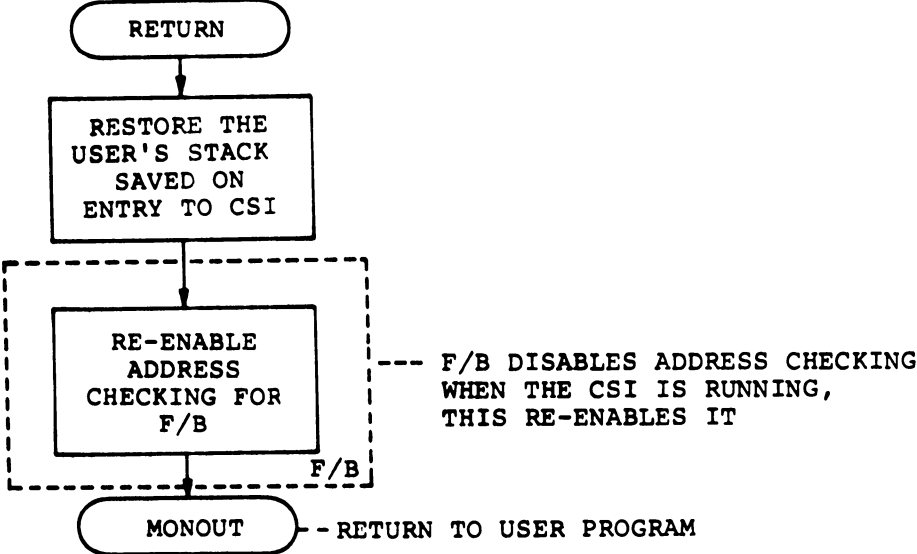
CSI CODE





CSI CODE (CONT.)



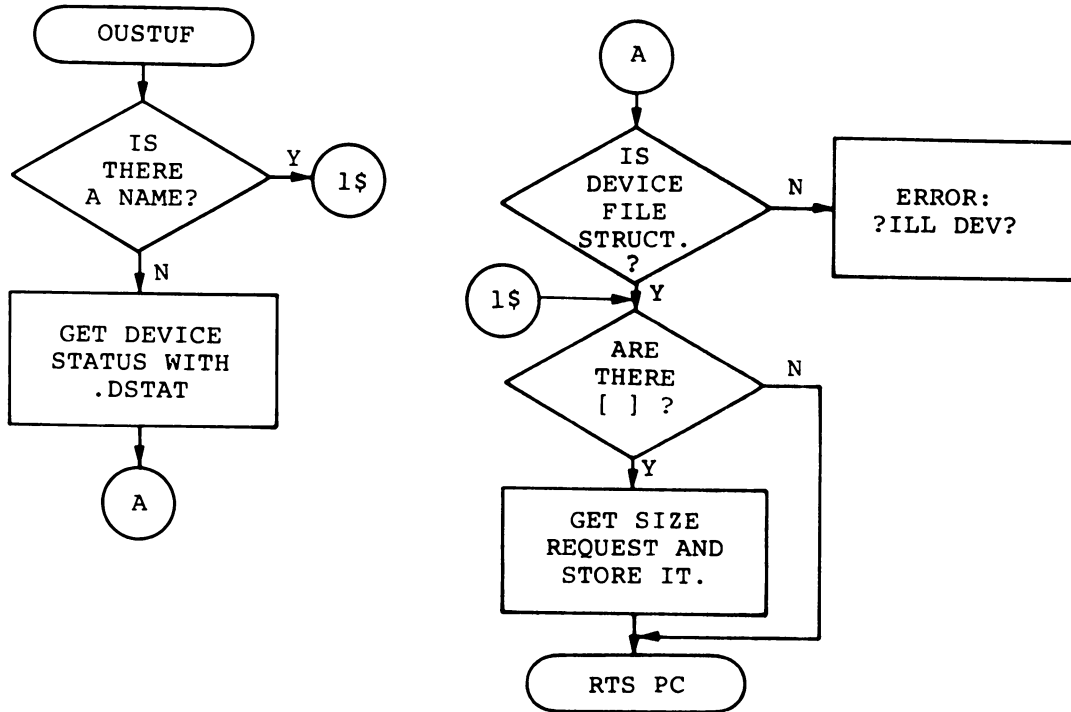


E.3.1 CSI Subroutines

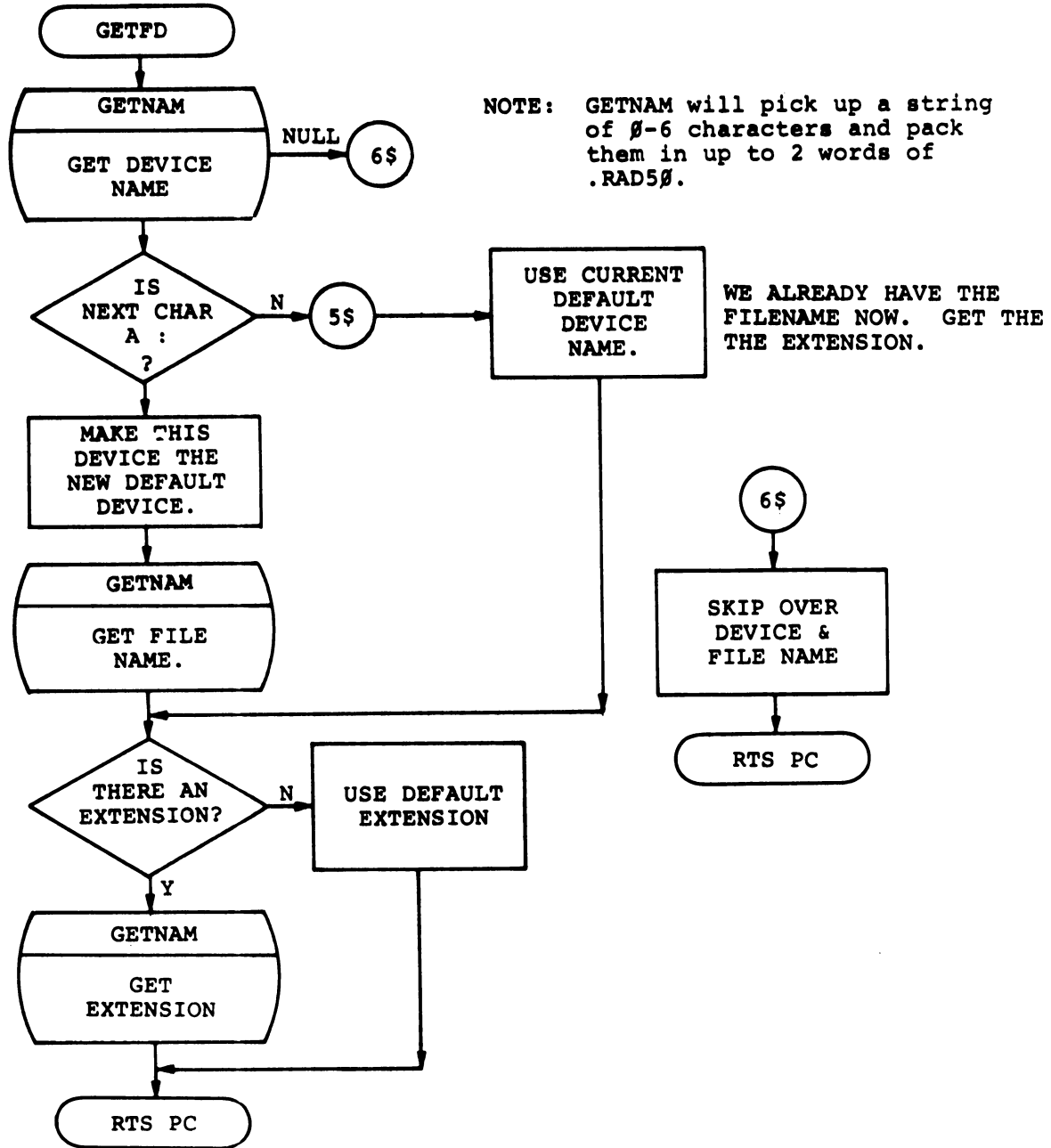
These subroutines are used by the CSI, and, in certain cases by the KMON.

OUSTUF

OUSTUF - This routine verifies that an output descriptor has a file name. If not, a syntax error is generated. It also will scan off the size in [] if it was specified.



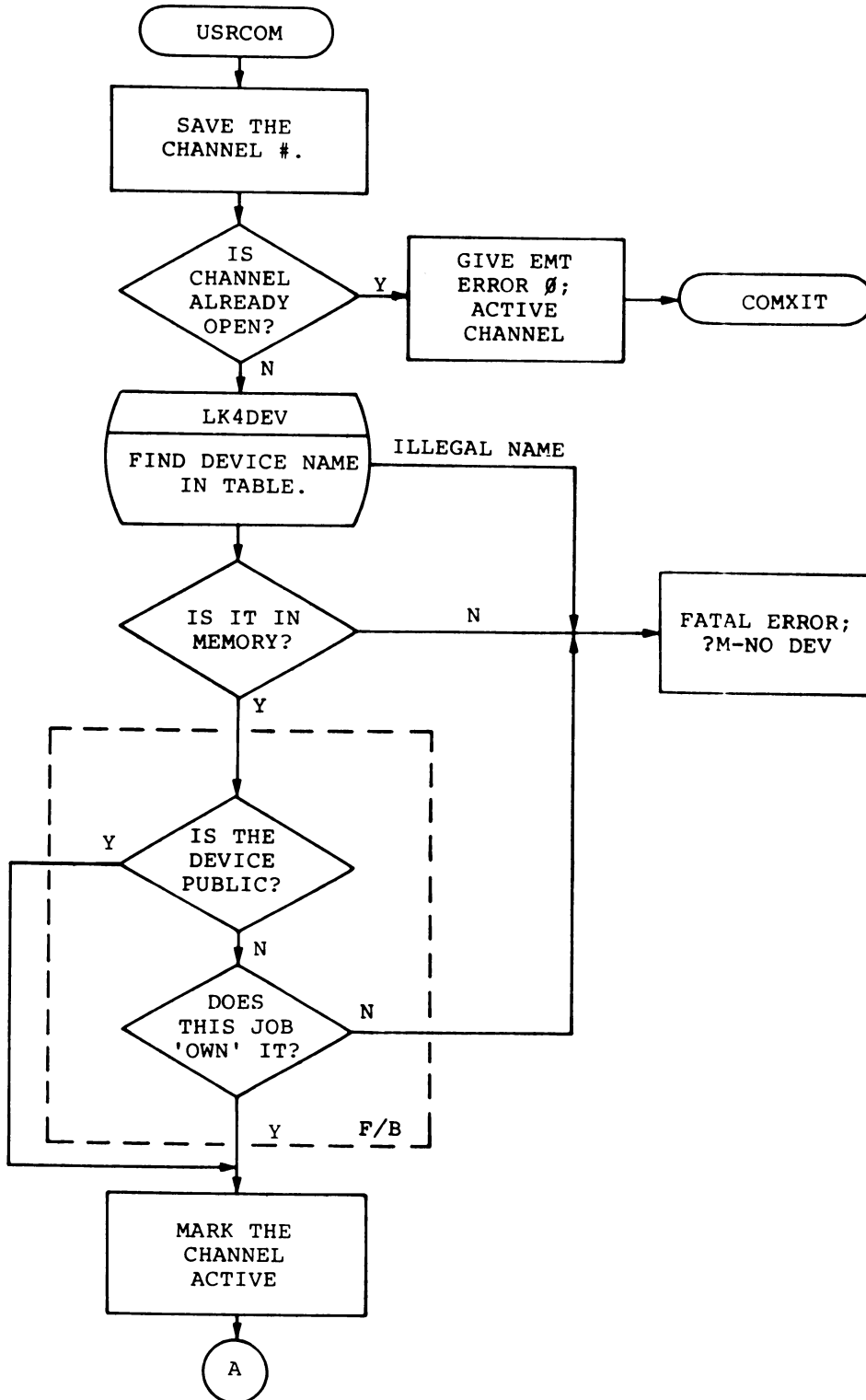
GETFD - Picks up a file descriptor (DEV:FILE.EXT) from an input string and packs it in 4 words of .RAD5Ø.

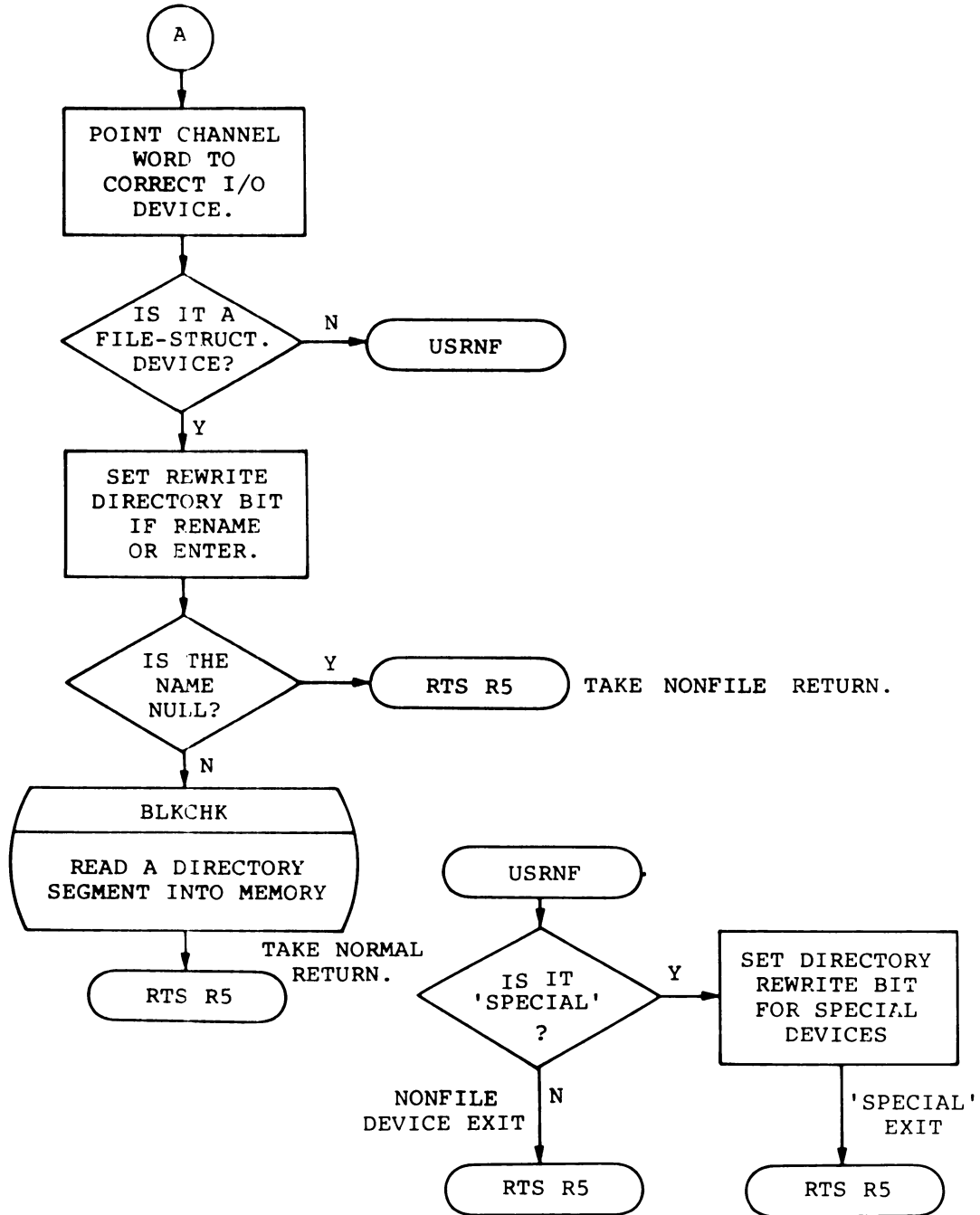


GETNAM - Converts a string of Ø-6 alphanumeric characters to a 2-word RAD5Ø group. The two words are zero filled when necessary. See code at GETNAM in the source listing if greater detail is necessary.

USRCOM

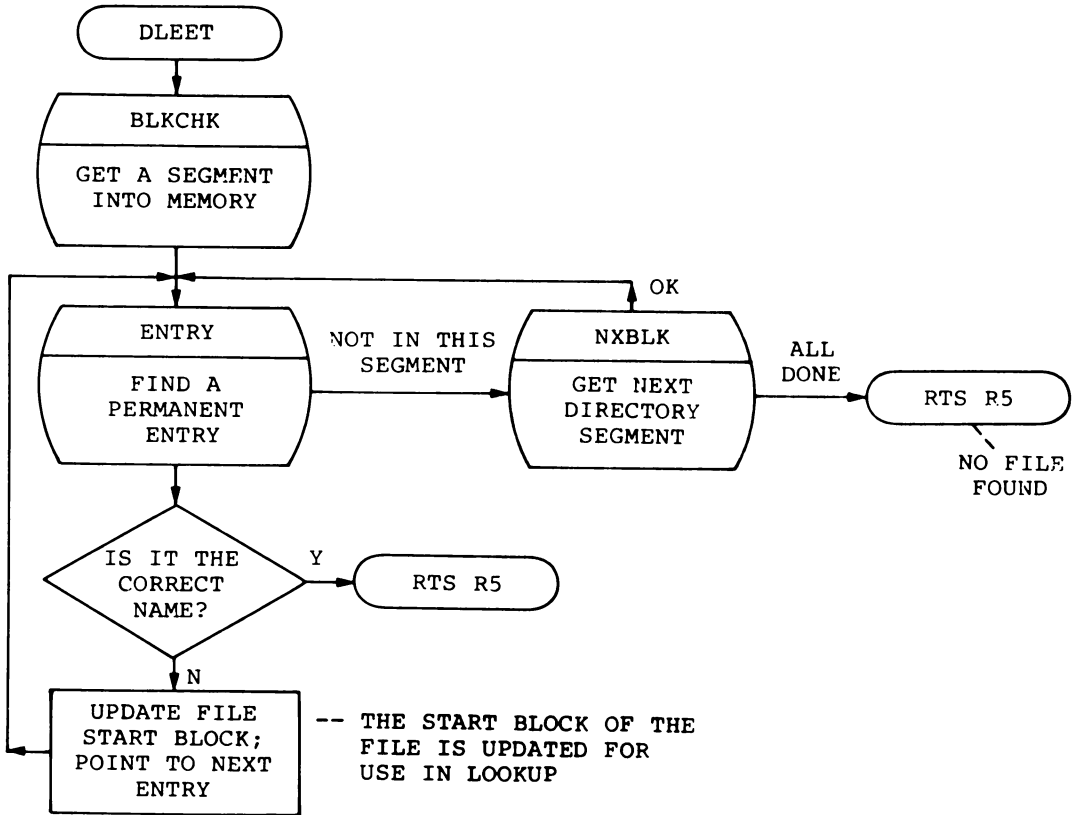
USRCOM - This routine is used to prepare a channel for I/O operations.



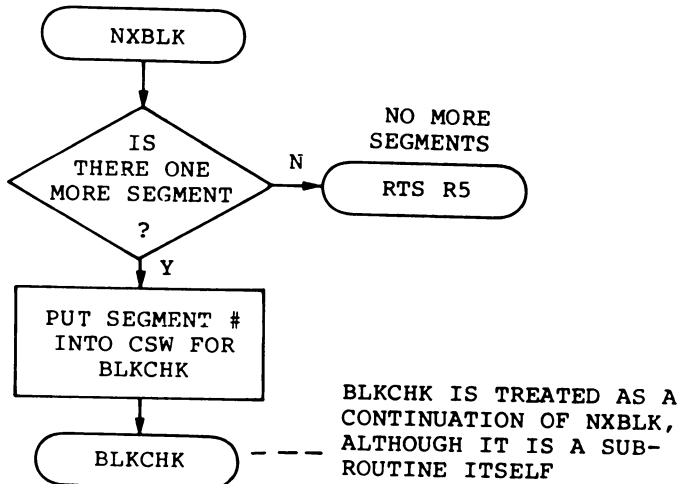


DLEET/NXBLK

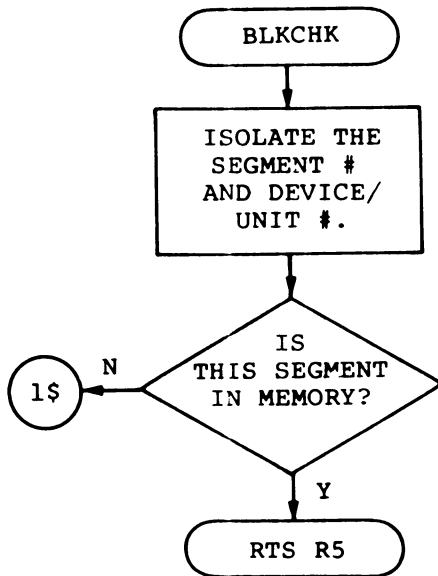
DLEET - This routine scans a device directory to find a file of a specified name.



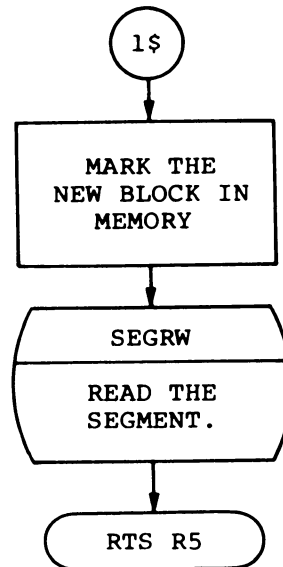
NXBLK - Gets the next in the series of directory segments, if one exists.



BLKCHK - This routine isolates the segment number contained in bits 8-12 of the CSW, and checks to see if that segment is in memory at the current time. If not, it is read in.



Note that not only must the segment numbers agree, but also the device and unit numbers must be the same.



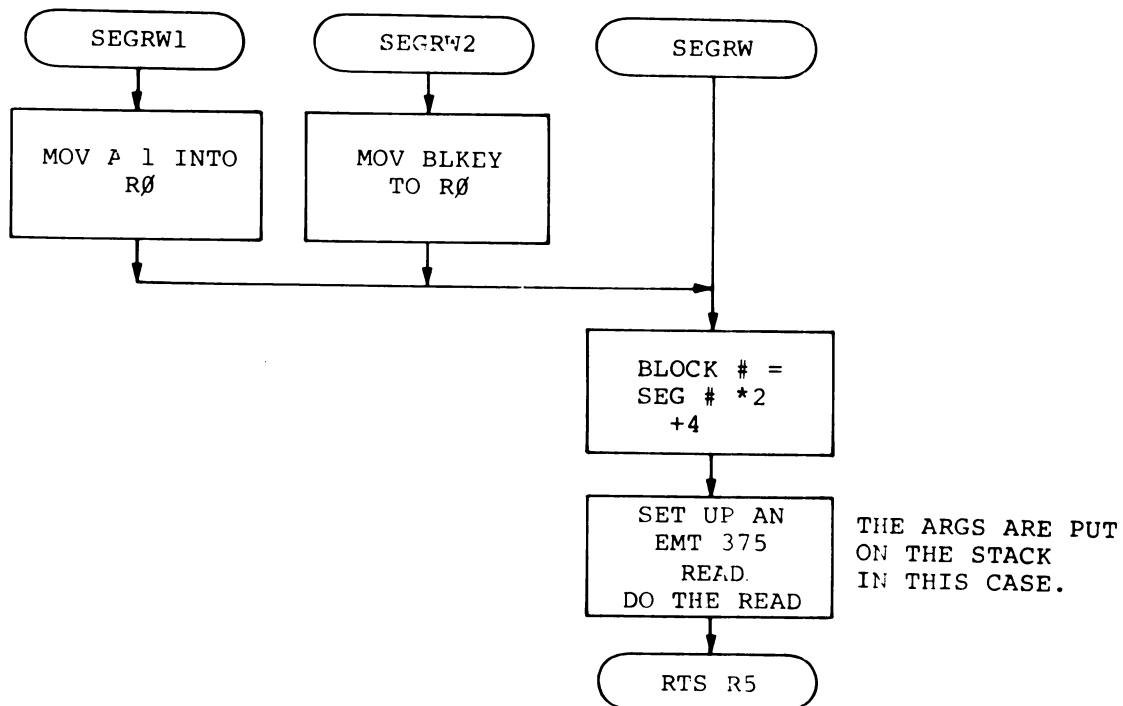
SEGRW

SEGRW - Segment Read/Write. This routine read/writes selected directory segments. There are three entry points:

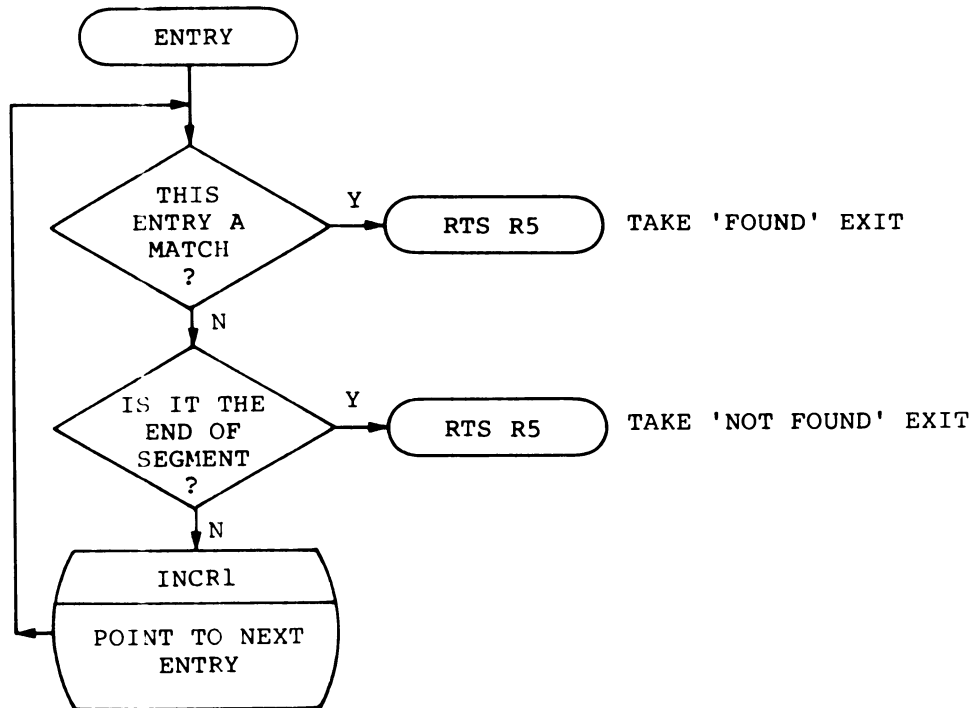
SEGRW1: Use segment #1

SEGRW2: Use the segment currently in memory (BLKEY)

SEGRW: Use the number in R0 as the segment #.



ENTRY - This routine uses R1 as a pointer into a directory segment to find a specified file type (Permanent, Tentative, Empty) or the end of segment mark.



INCR1 - This routine bumps R1 to the next entry in a directory segment.

COMERR - This routine generates a fatal error from the USR. The call is:

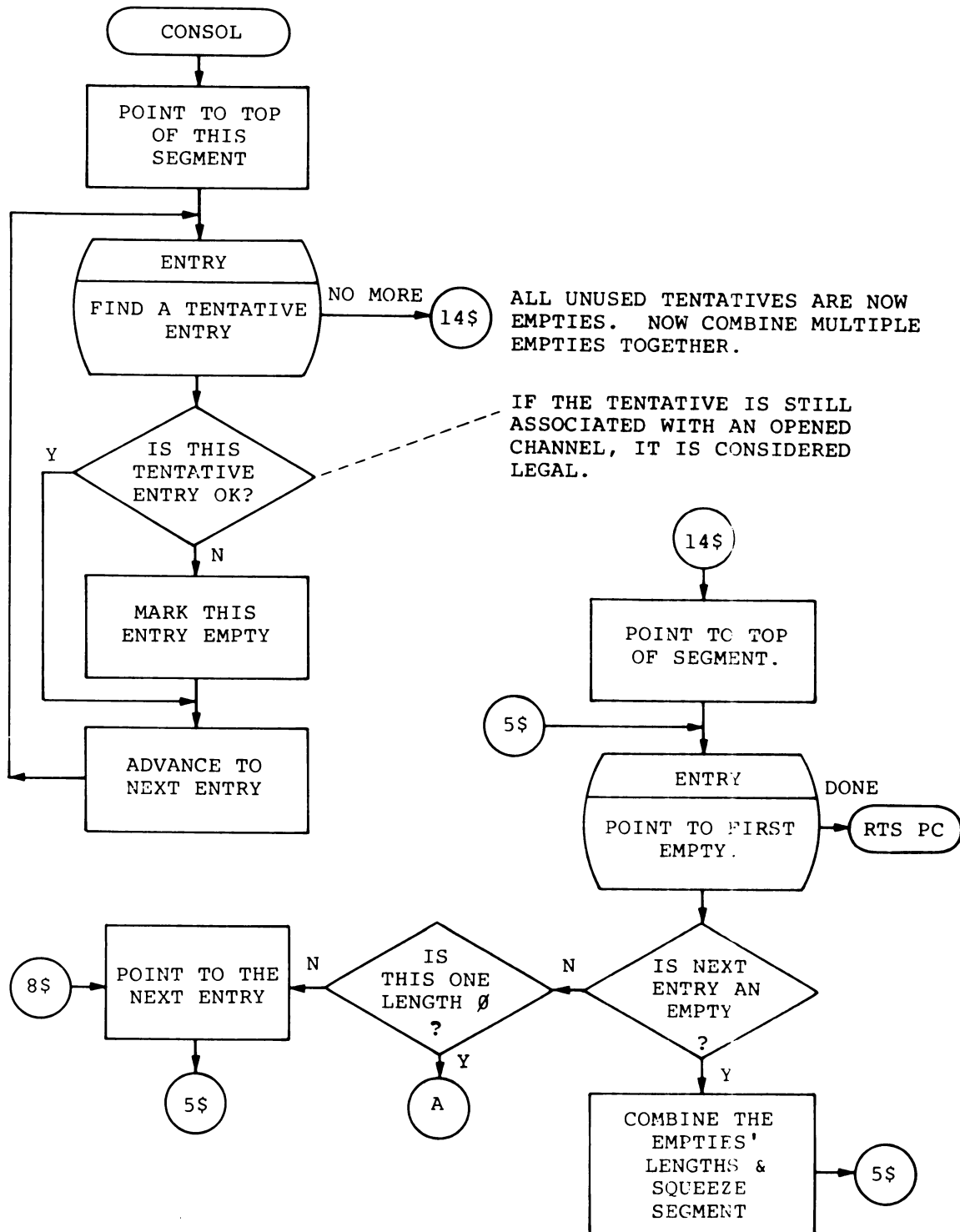
```
JSR R5,COMERR
code
```

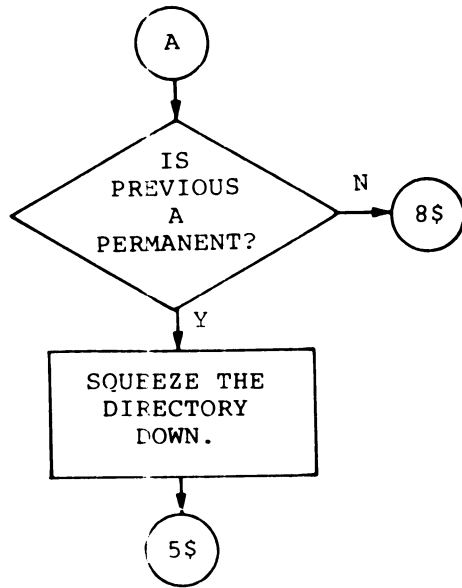
Code is used to indicate which error is to be generated. If .SERR is in effect, control passes to COMXIT, which returns to RMON.

SPESHL - This routine is used to effect file operations on MT/CT. This is done by passing a READ request to the Q manager. The even byte of the completion function will contain a 377. The queue manager detects this, and modifies the I/O queue element to indicate that the handler should perform a USR function.

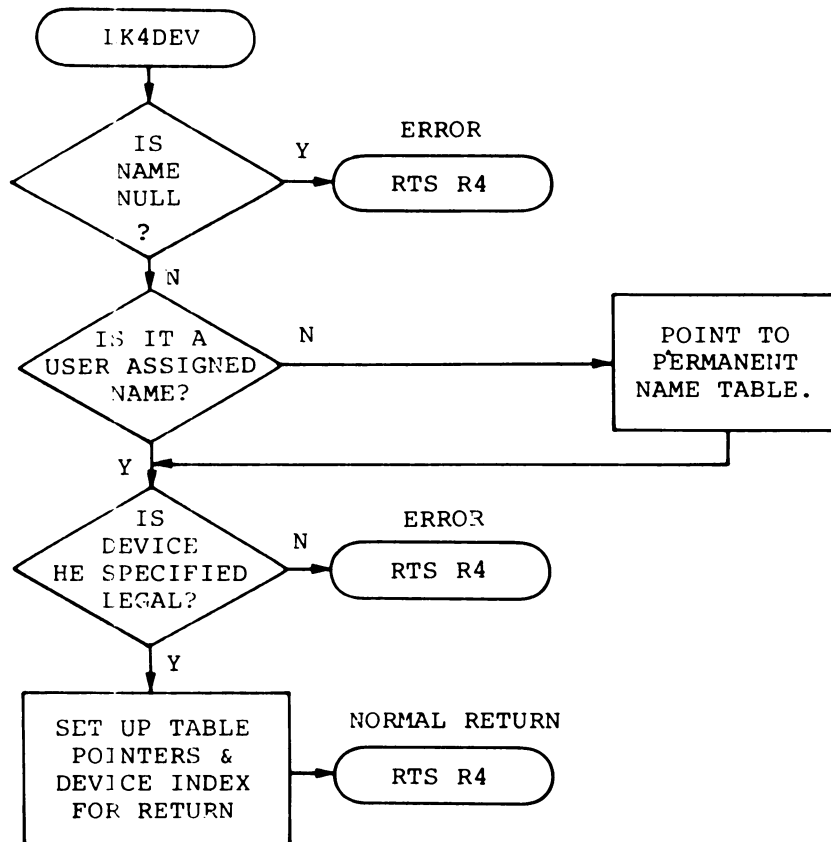
CONSOL

CONSOL - This routine is used to compact a directory segment. It combines consecutive empties into one, and makes empties out of tentative files which are not associated with an active channel.





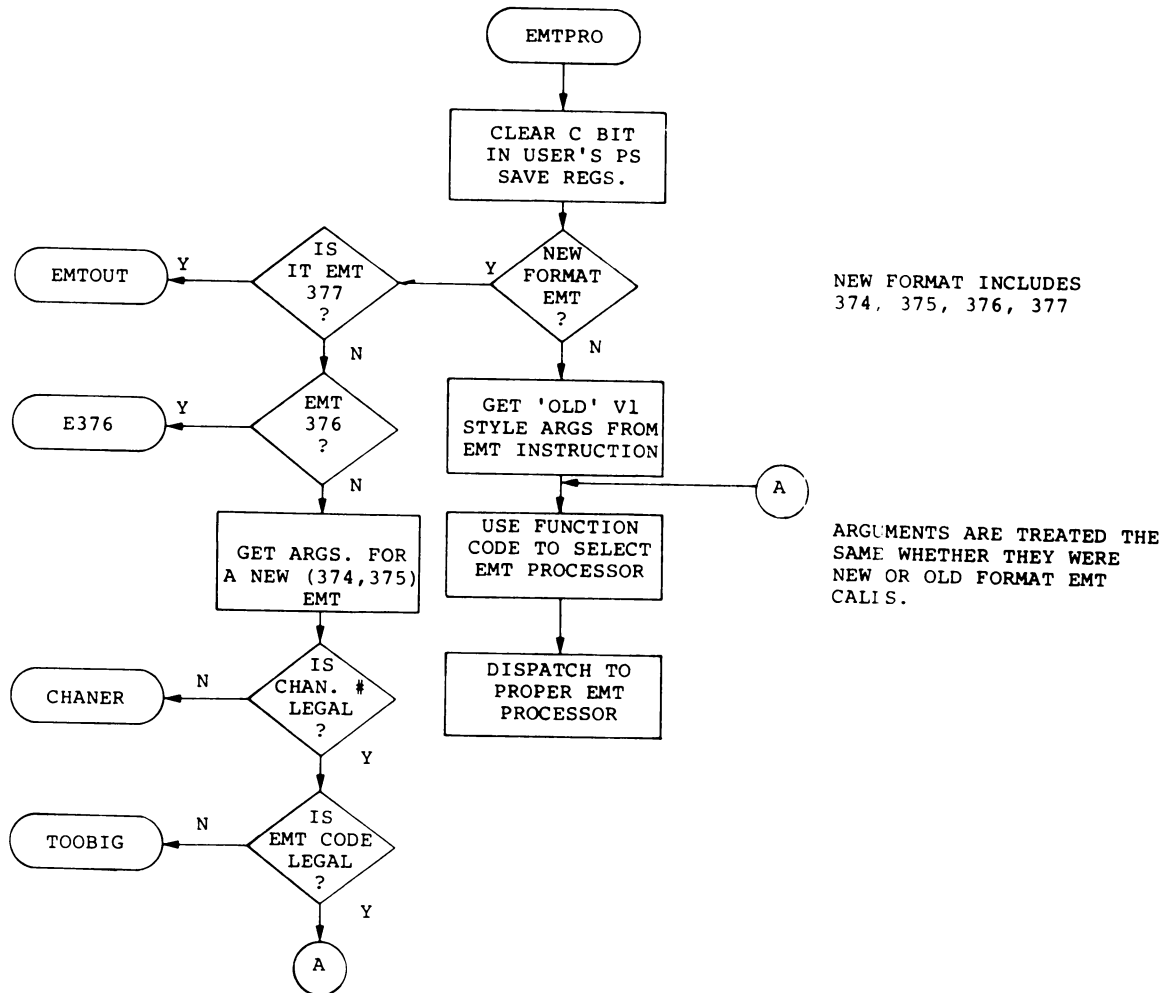
LK4DEV - This routine looks up a specified device name in the system tables. It first attempts to find the name in the user assigned name table; failing that, the permanent name table is searched.



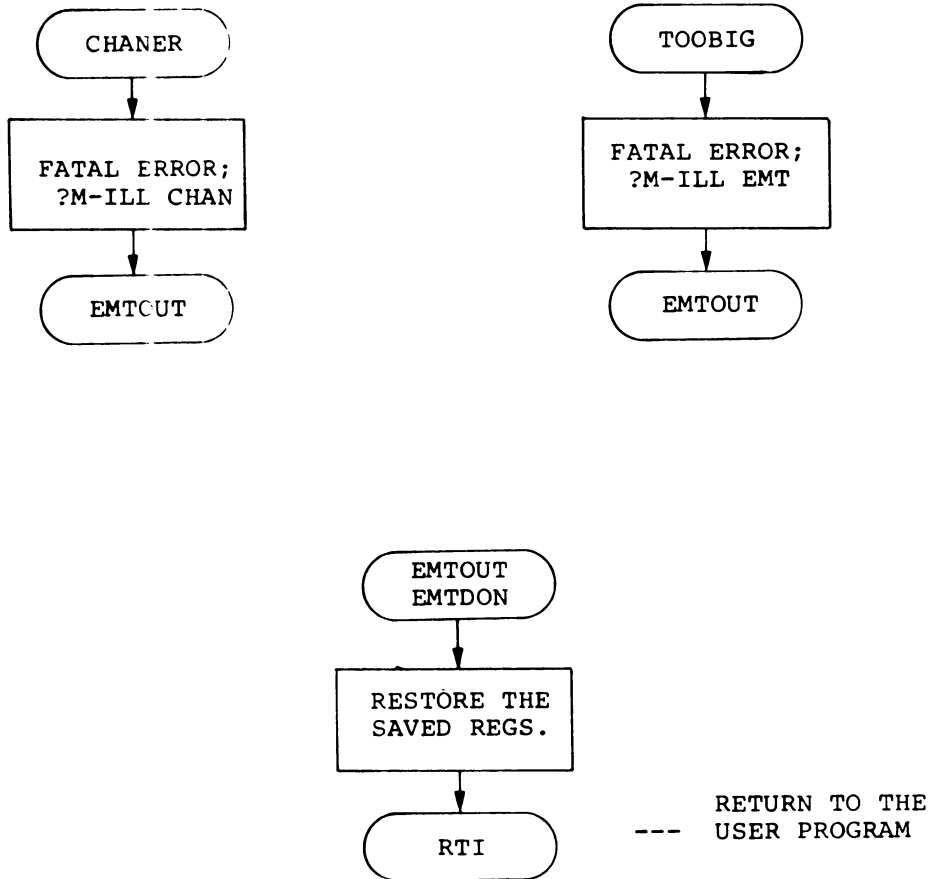
E.4 RMON (RESIDENT MONITOR) FLOWCHARTS FOR SINGLE-JOB MONITOR

EMT DISPATCHER

The code of the EMT dispatcher is entered when an EMT instruction is executed. The EMT instruction is decoded and control passes to the appropriate code for processing.



EMT DISPATCHER (CONT.)



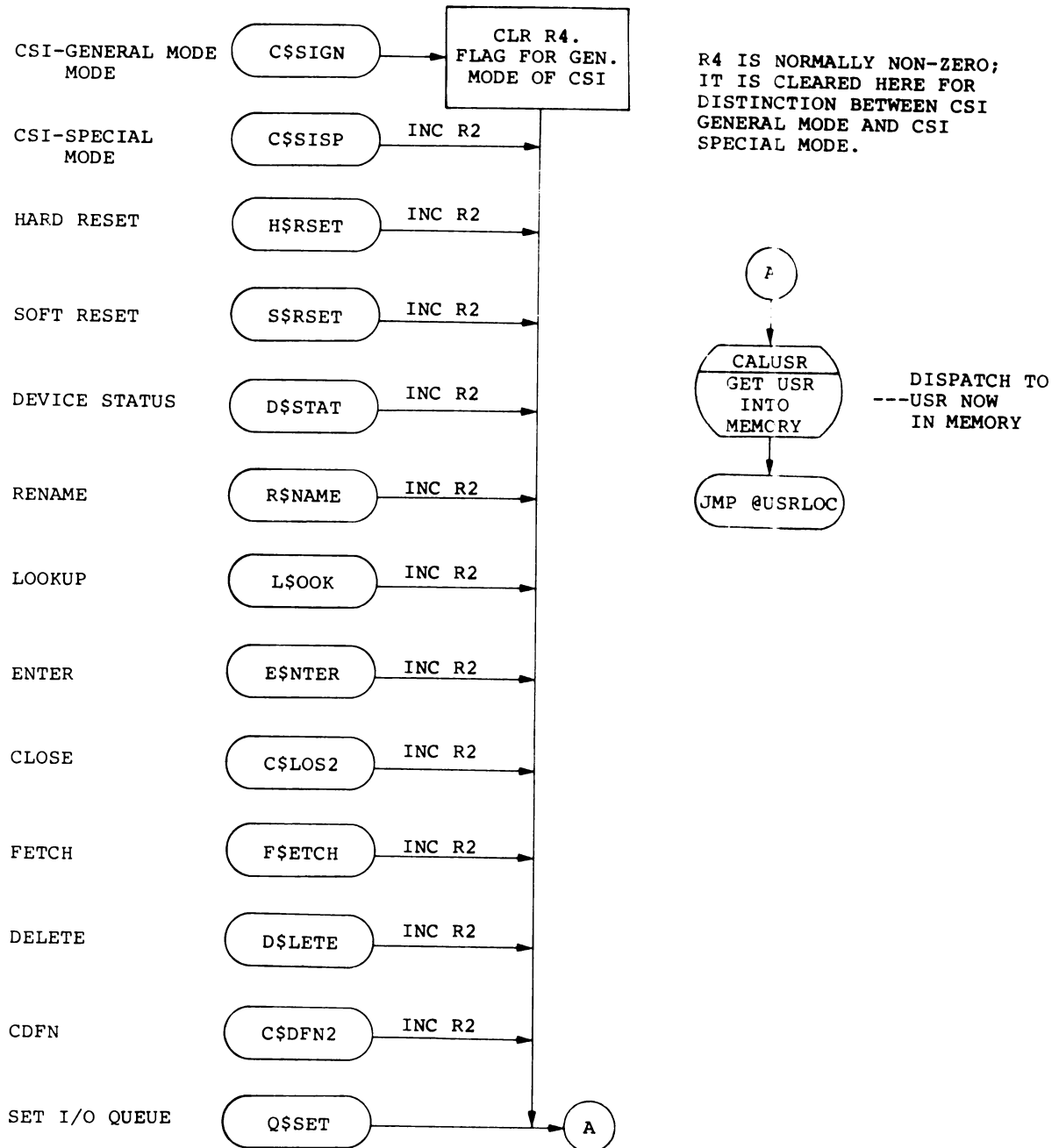
The following EMT requests are no-ops in the S/J Monitor:

Mark Time	.MRKT
Cancel Mark Time	.CMKT
Timed Wait	.TWAIT
Send Data	.SDAT
Receive Data	.RCVD
Channel Status	.CSTAT
Protect Vectors	.PROTECT
Channel Copy	.CHCOPY
Special Device	.DEVICE

Executing these requests in S/J will cause an immediate successful returns with no action taken.

USR DISPATCHER TABLE FOR EMT'S 340-357

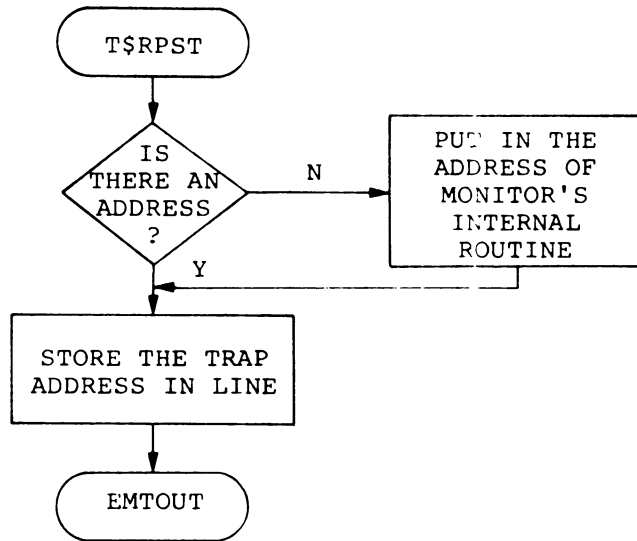
The USR Dispatch code handles dispatching those EMT's which require the USR. At each entry point, an INC R2 is performed. Thus, R2 acts as a function identifier once the USR is entered.



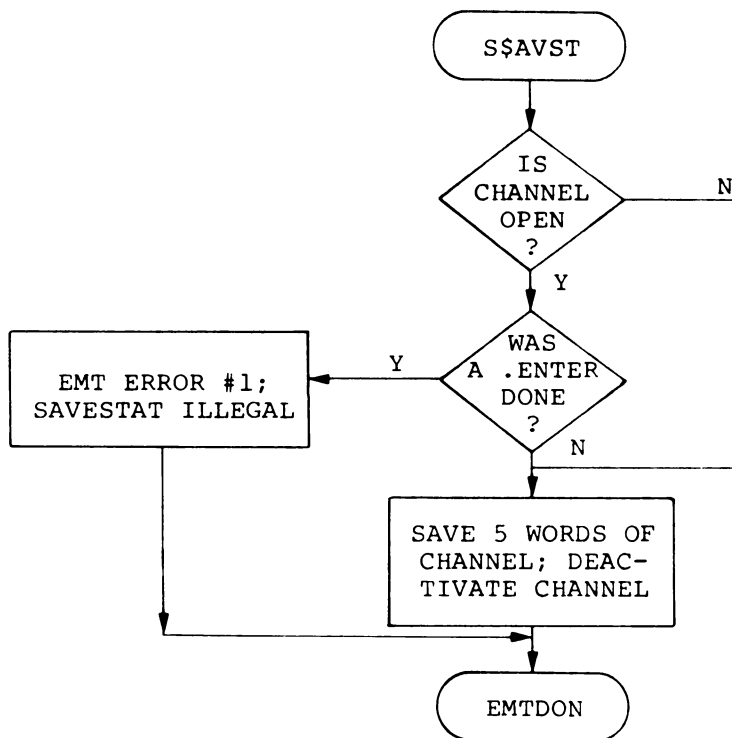
E.4.1 EMT Processors

SET TRAP/SAVE STATUS

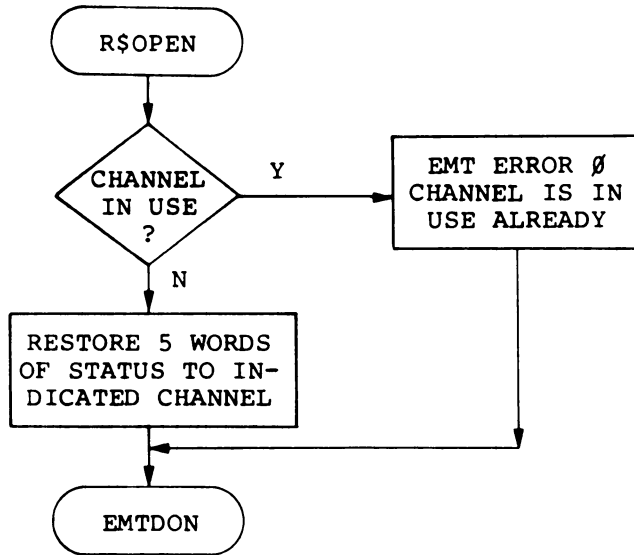
SET TRAP ADDRESS



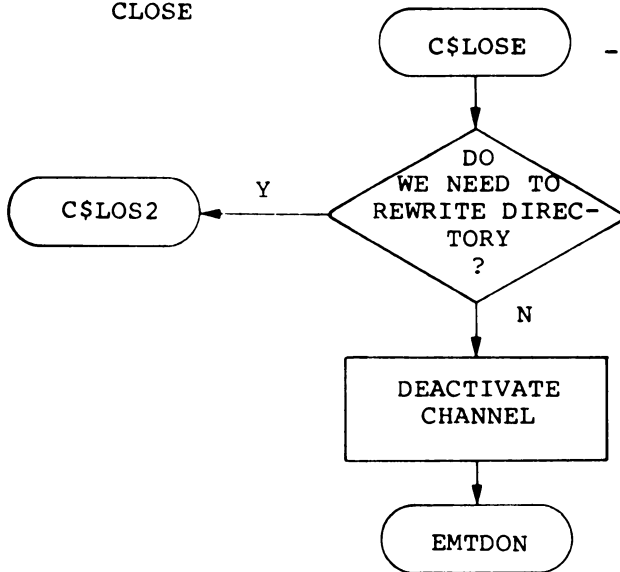
SAVESTATUS



REOPEN

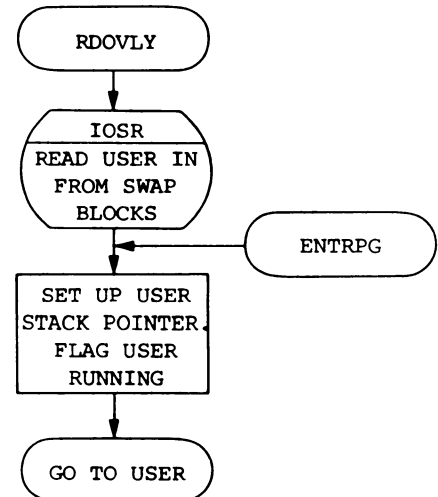


CLOSE

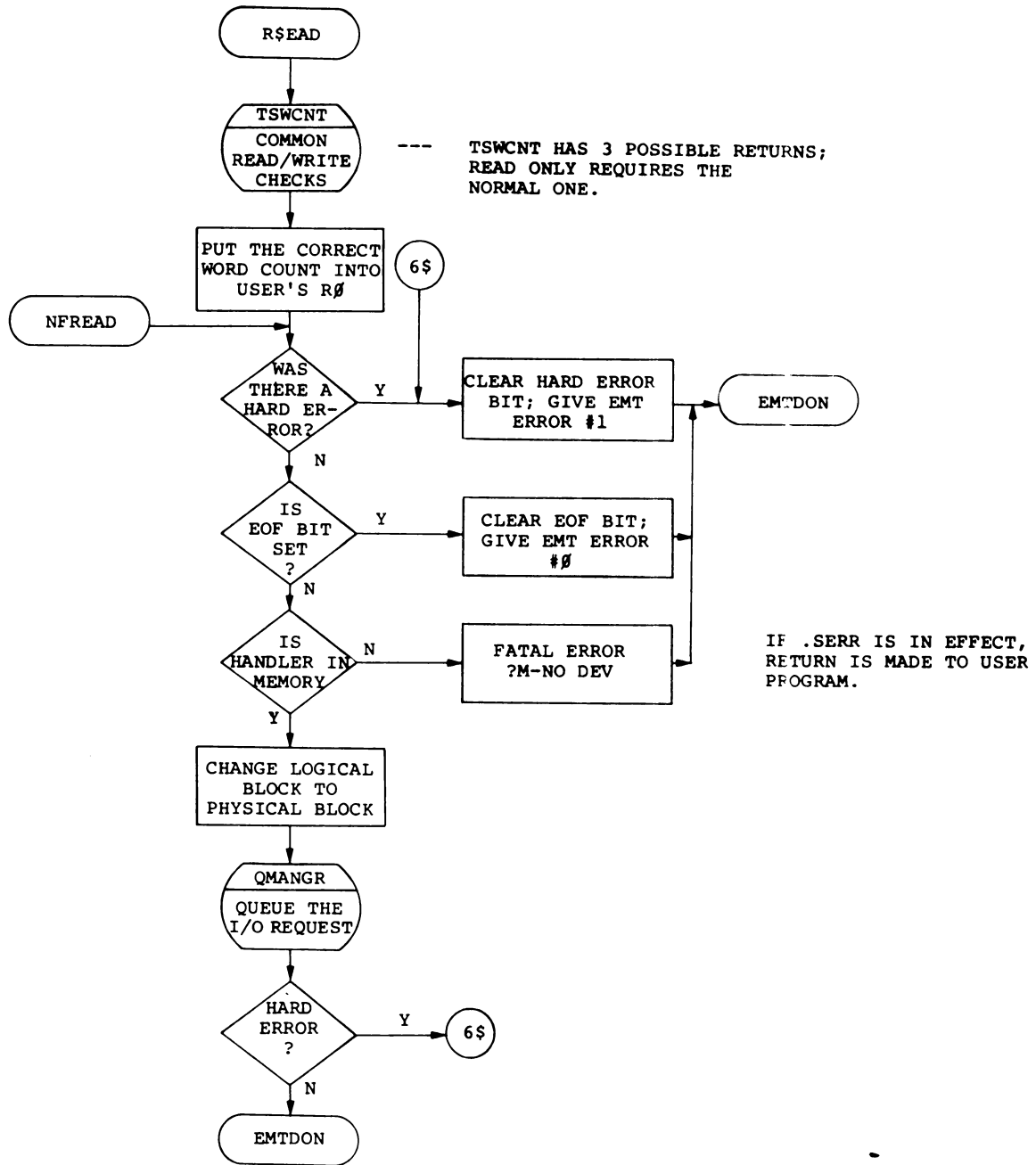


--- IF A LOOKUP WERE DONE, THE USR IS NOT REQUIRED.

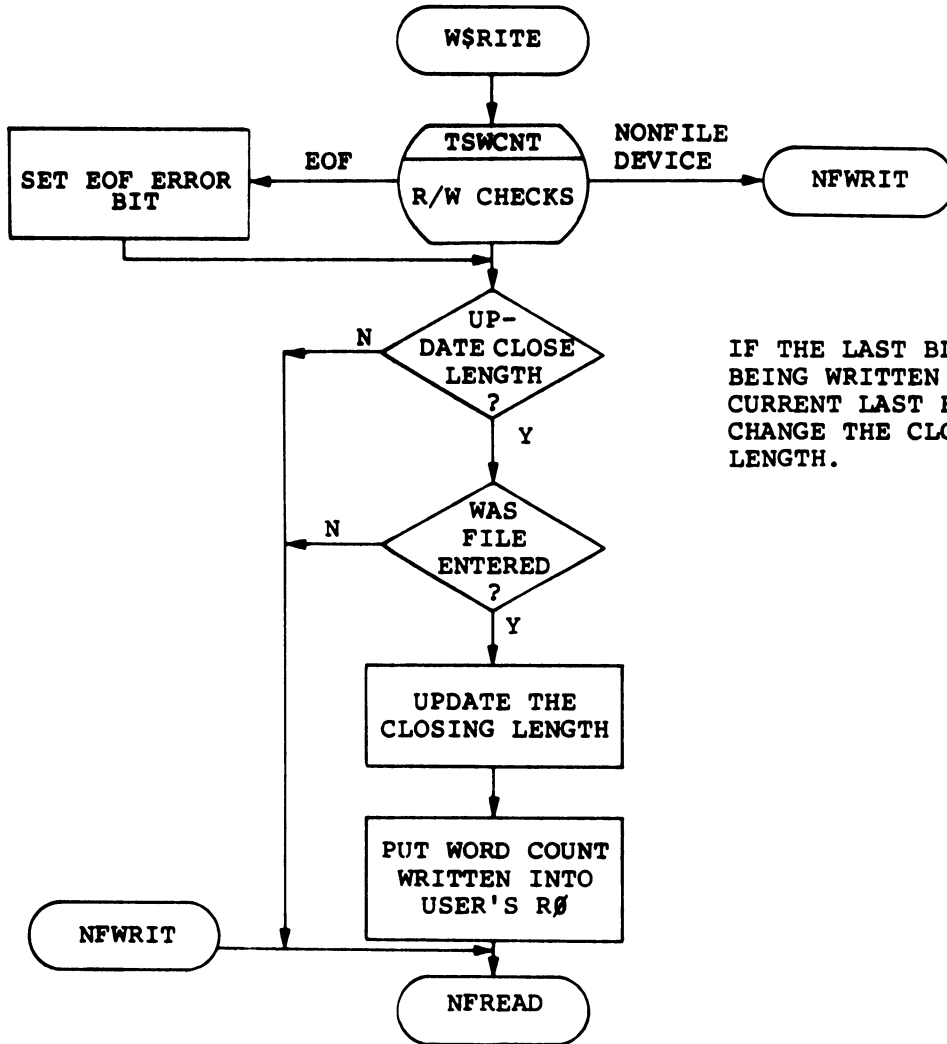
RDOVLY



READ

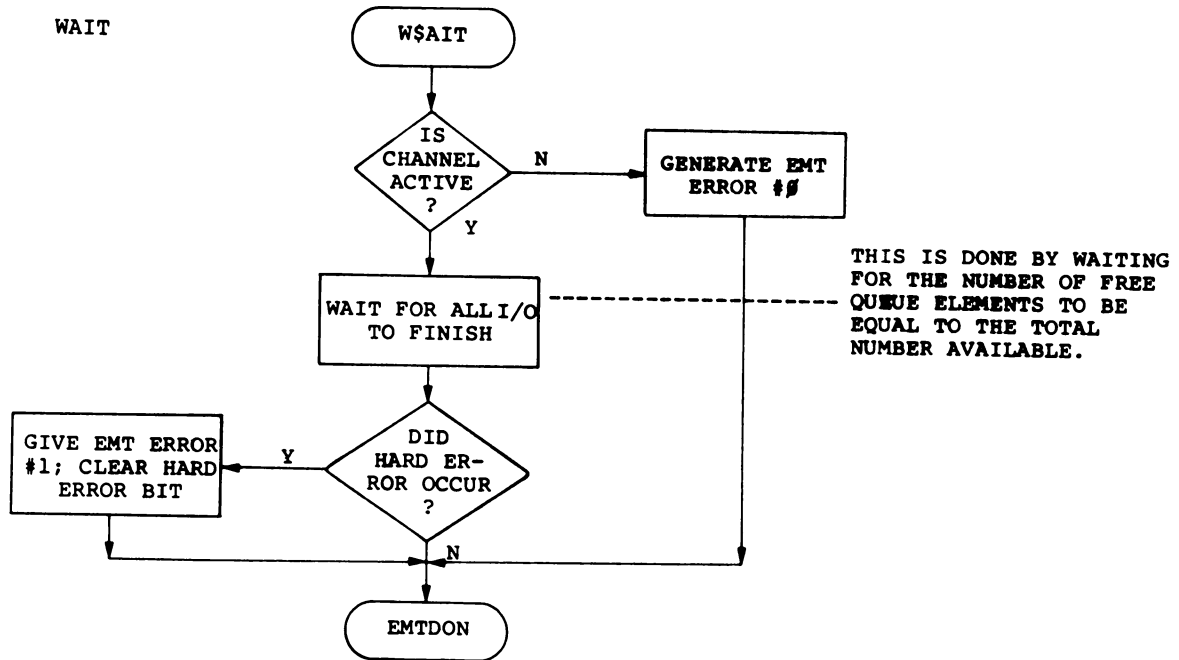


WRITE



IF THE LAST BLOCK BEING WRITTEN IS > THE CURRENT LAST BLOCK WRITTEN, CHANGE THE CLOSING FILE LENGTH.

WAIT/CDFN



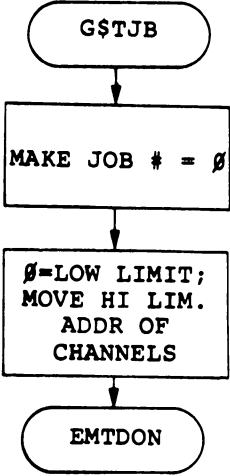
CDFN

Channel Define - the resident portion of CDFN causes a fresh copy of the USR to be read in, then enters the USR.

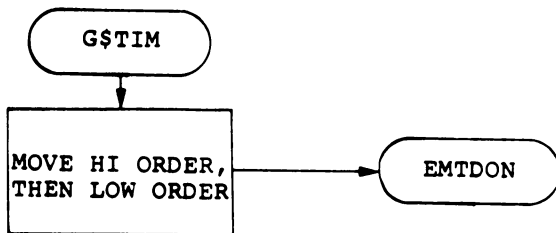


GET JOB PARAMETERS
GET TIME OF DAY
SET FPP EXCEPTION

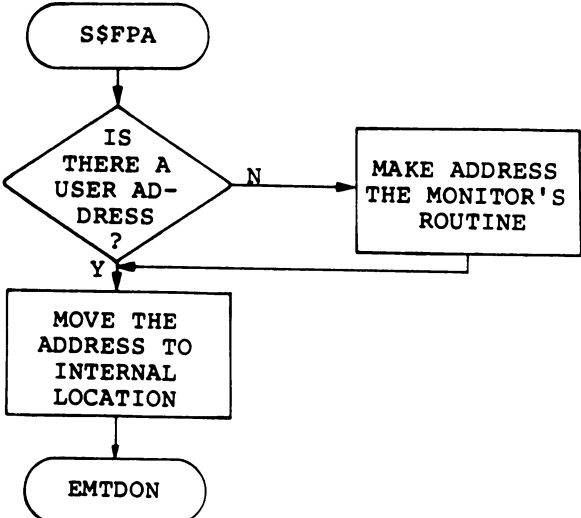
GET JOB PARAMETERS



GET TIME OF DAY

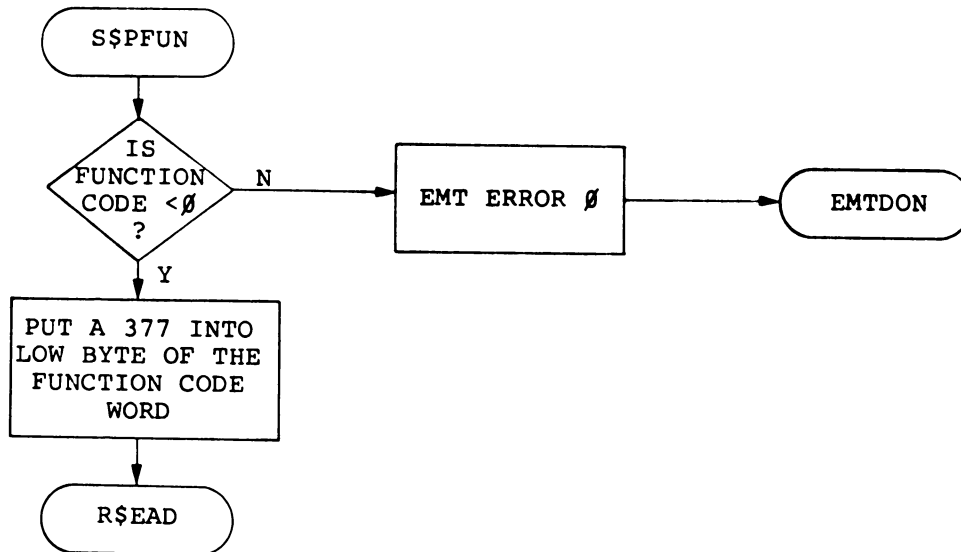


SET FPP EXCEPTION



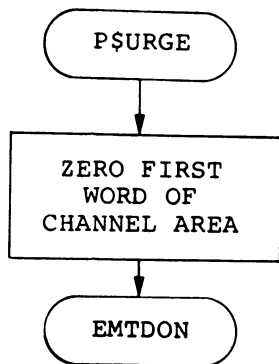
SPECIAL FUNCTIONS/PURGE
SOFT/HARD ERRORS

SPECIAL FUNCTIONS (MAGTAPE/CASSETTE)

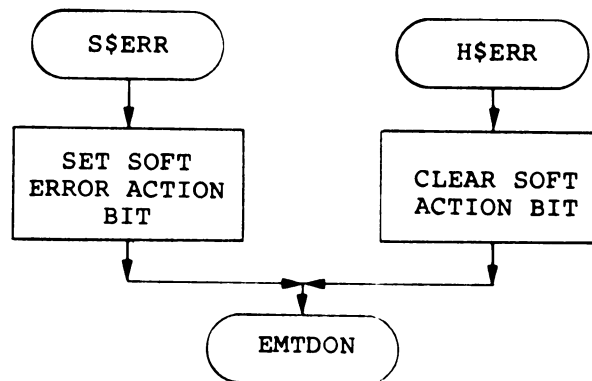


SPECIAL FUNCTIONS/PURGE
SOFT/HARD ERRORS

PURGE

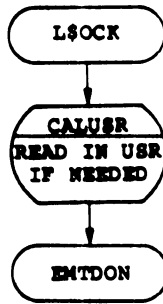


SOFT/HARD ERRORS



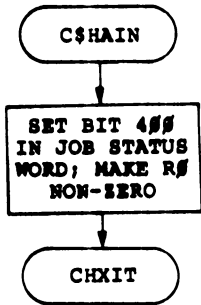
LOCK USR/CHAIN/UNLOCK USR

LOCK USR



----- THIS BUMPS A COUNTER WHICH IS DECREMENTED DURING A .UNLOCK. THE USR IS REALLY UNLOCKED ONLY WHEN THIS COUNT IS 0.

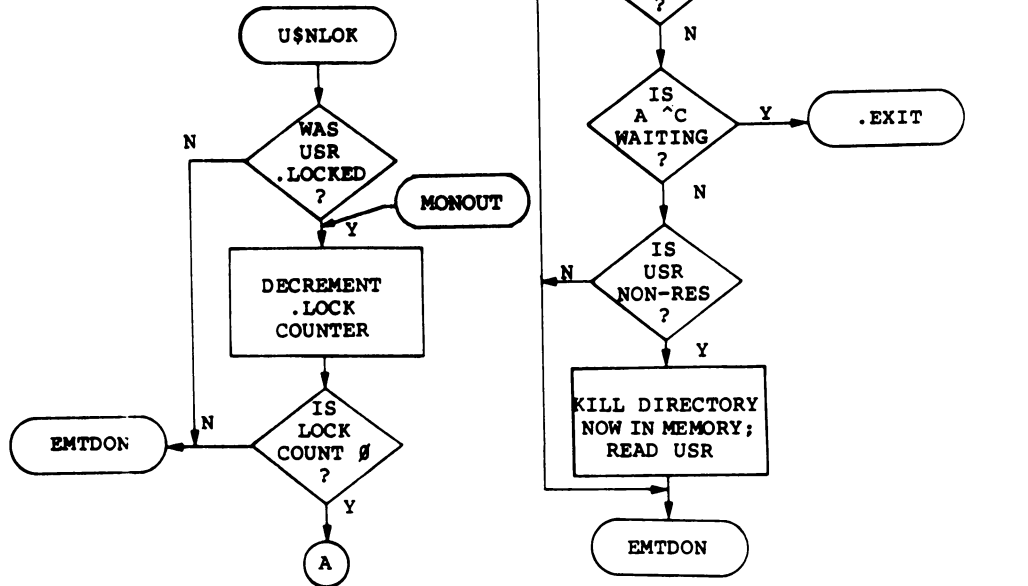
CHAIN



R\$ NOT ZERO FORCES A SOFT CONDITION AT EXIT

JOIN EXIT CODE AT CHXIT

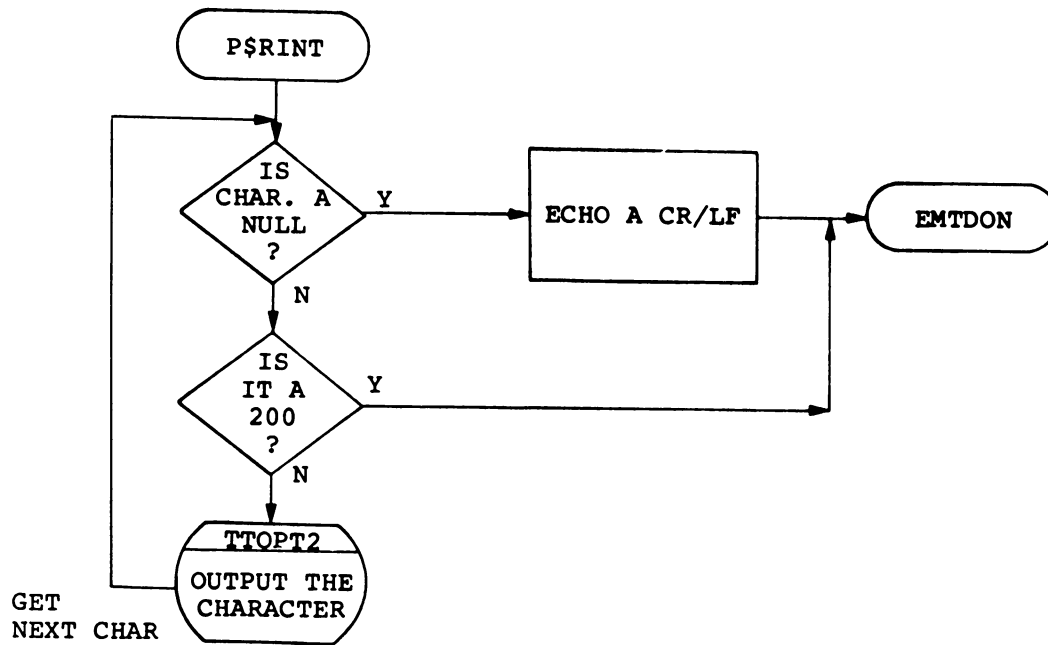
UNLOCK USR

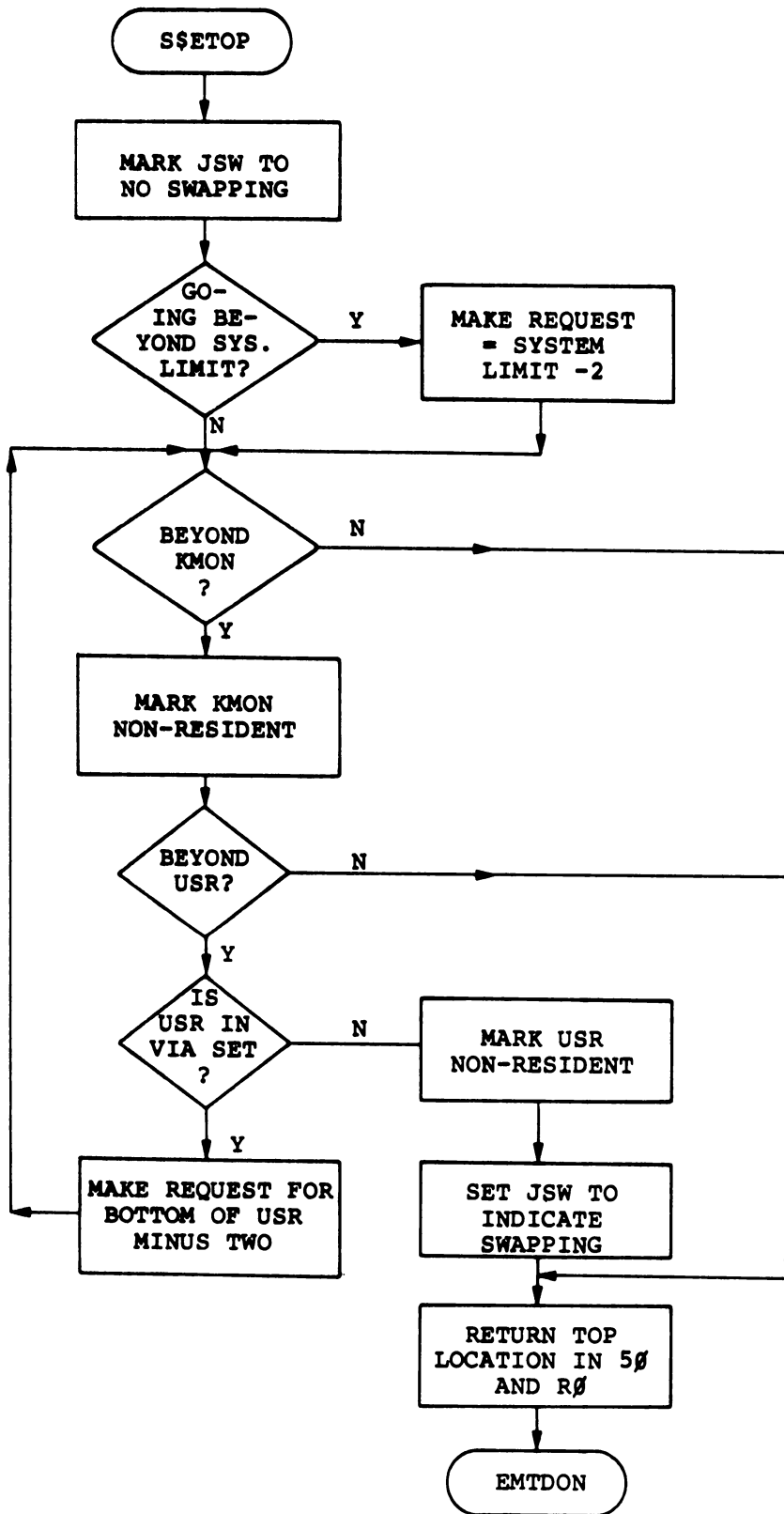


IF KMON IS IN, SO IS USR.

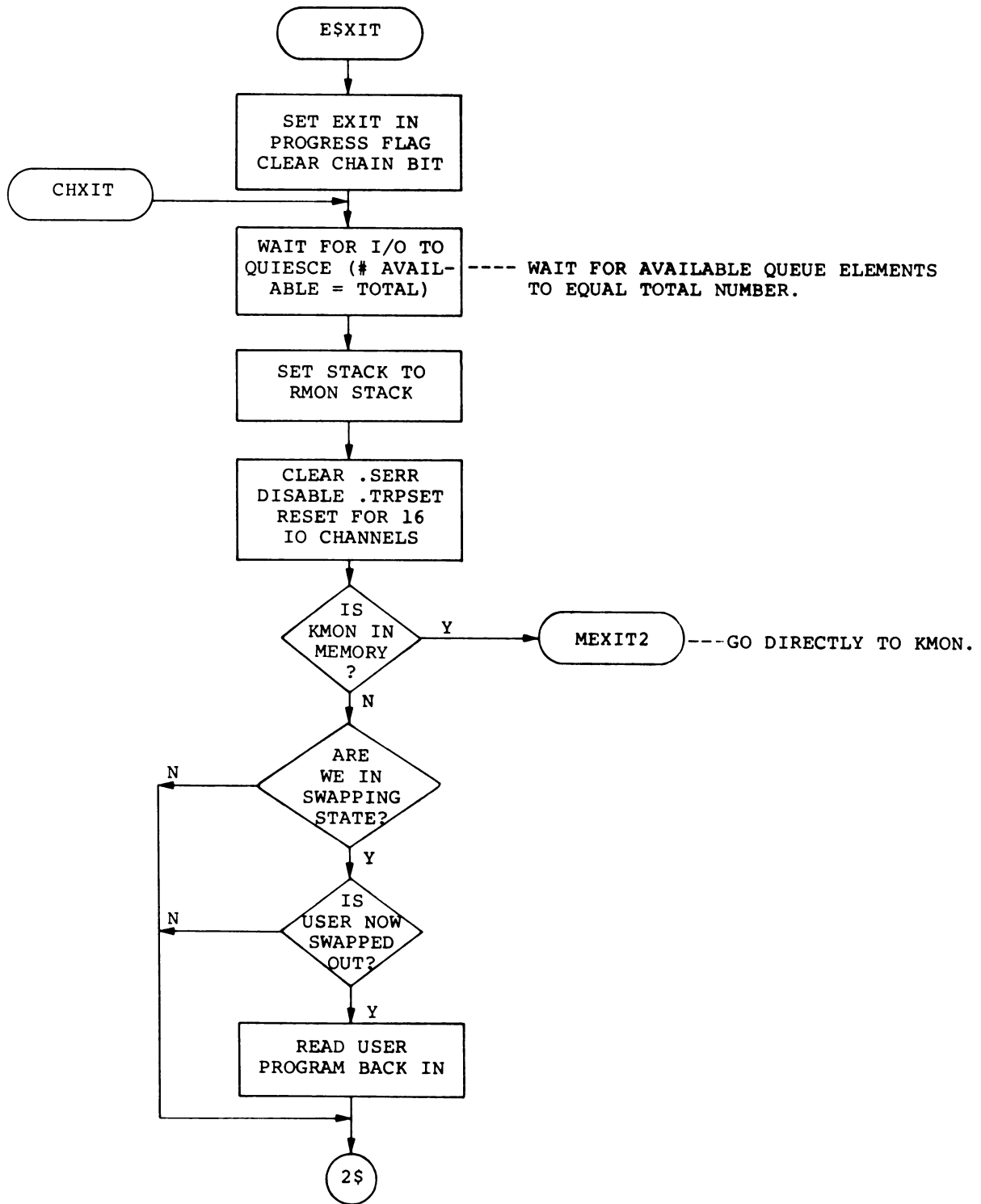
PRINT

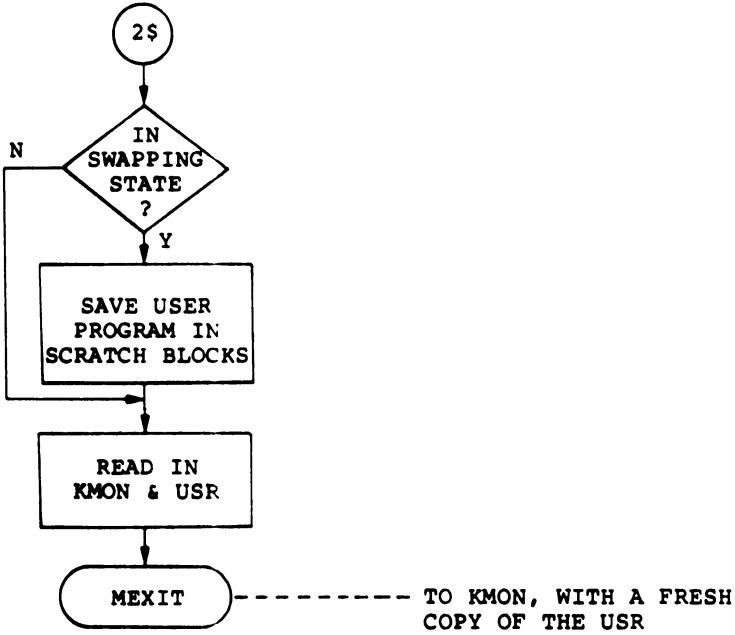
PRINT - Causes a line to be output to the console terminal.



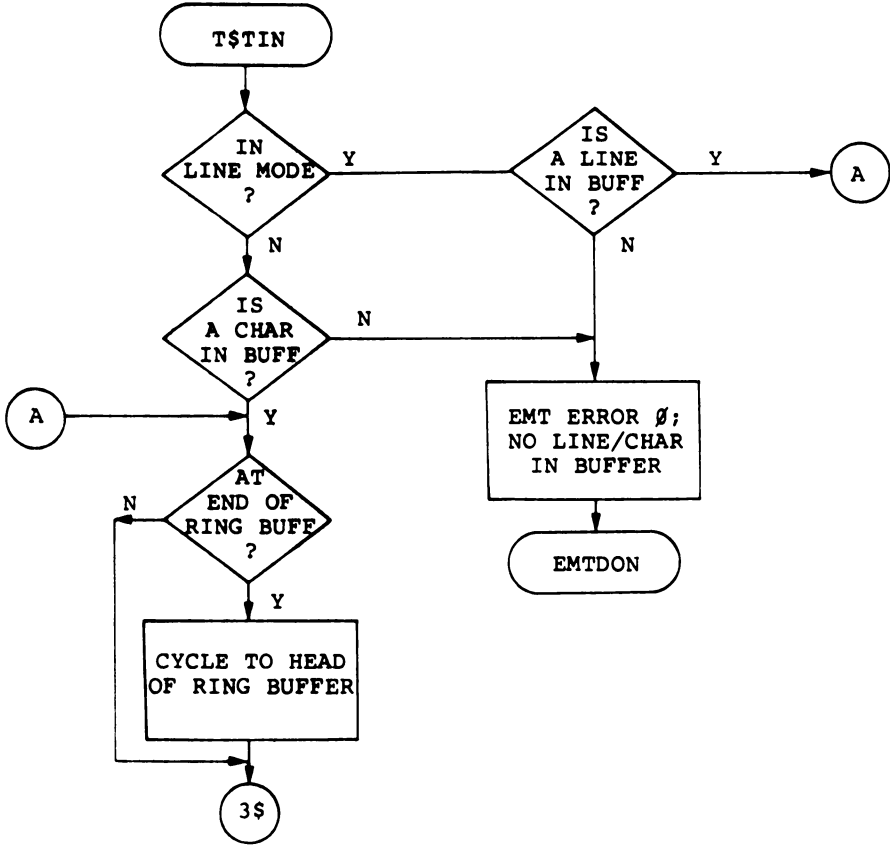


EXIT

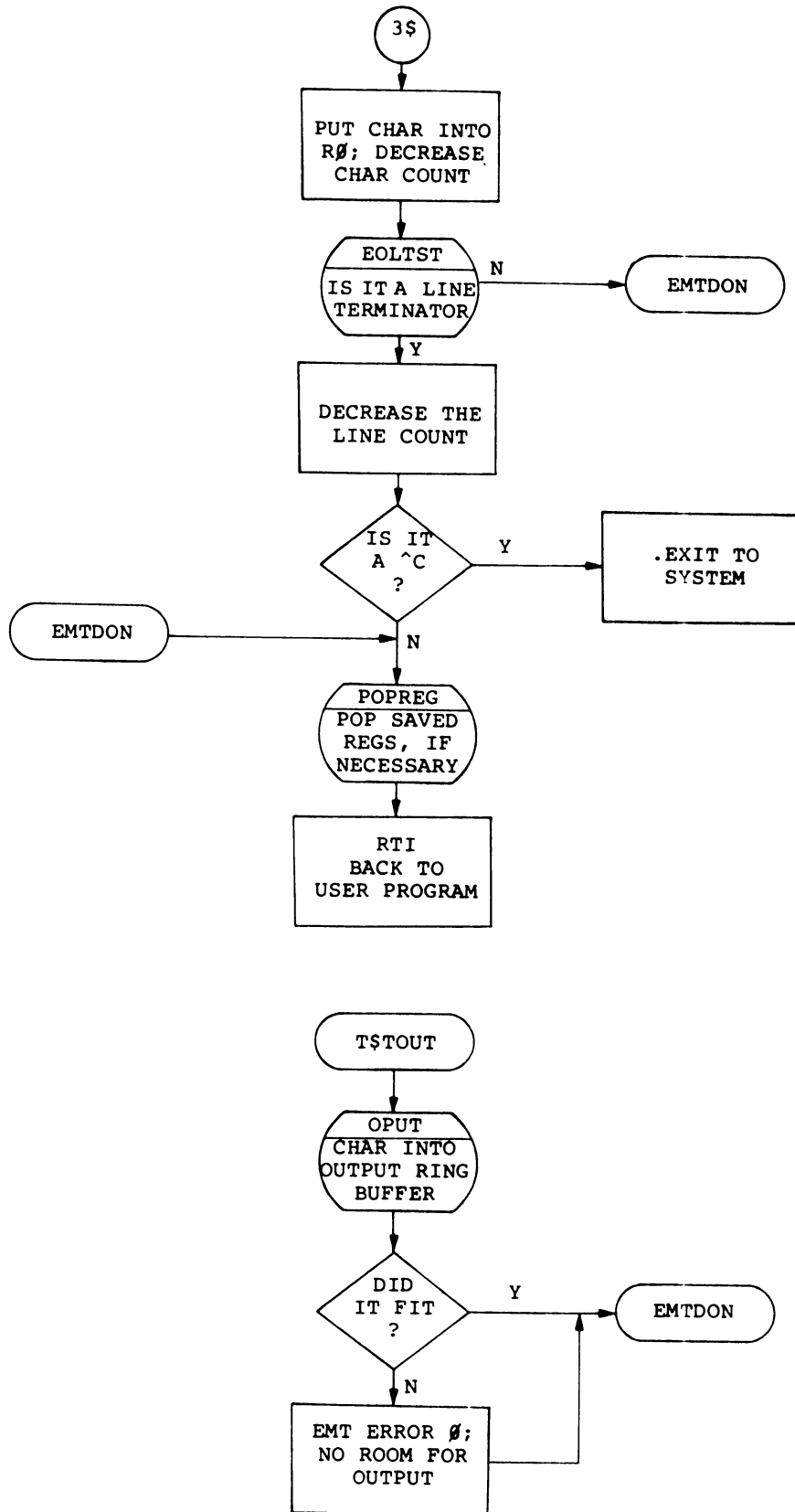




TTYIN



TTYIN (CONT.)/TTYOUT

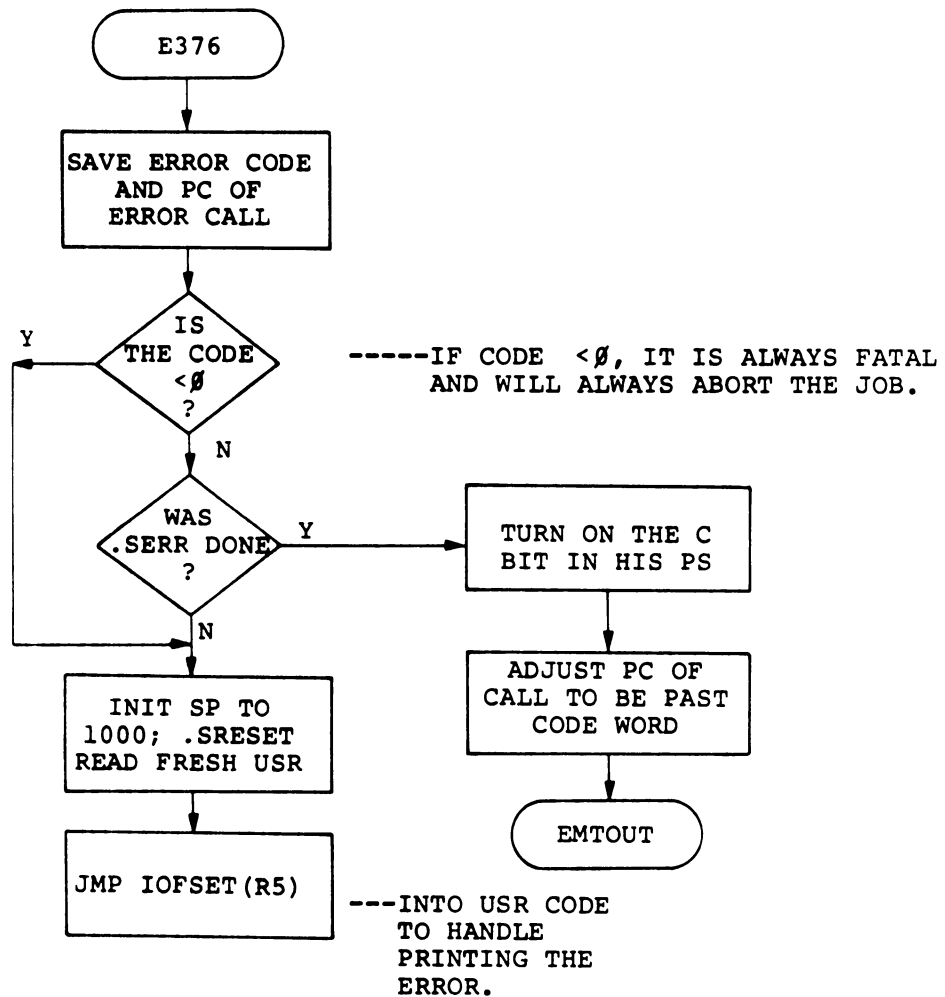


FATAL ERROR PROCESSING

EMT 376 is reserved for reporting fatal monitor errors. When a fatal error condition is encountered, a call of the form:

EMT 376
code

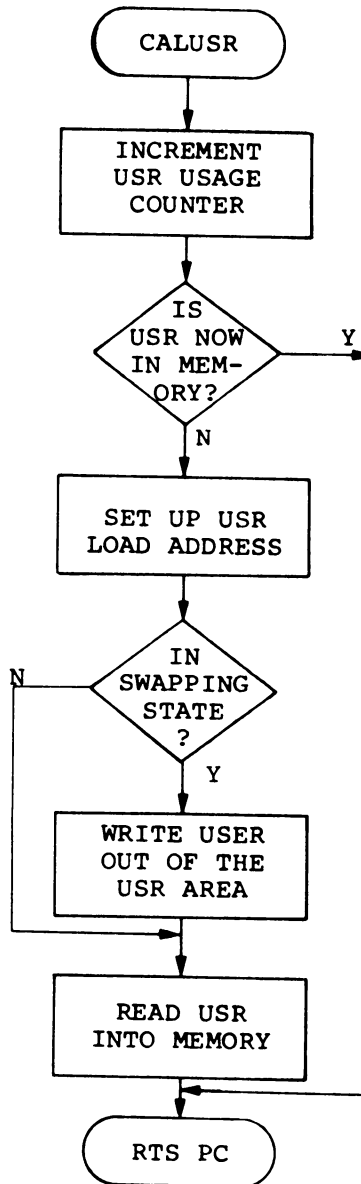
is executed. This indicates to RT-11 that a fatal error has occurred. The normal response is to print a ?M-error message and then abort the job. However, if a .SERR request has been done, no message will occur and control will pass to the user's program. The error bit (C bit) will be set and byte 52 will contain the negative of the error code.



CALUSR

CALUSR is used to ensure that the USR is in memory for a USR type request. It will handle the situation where the user program must be written to scratch blocks before the USR is read in. Entry is made at CLUSR2 when an error has occurred and the error processing code in the USR buffer is required.

EITHER 'NORMAL' VALUE
OR WHAT THE USER HAS
PUT IN LOCATION 46.



E.4.2 Clock Interrupt Service

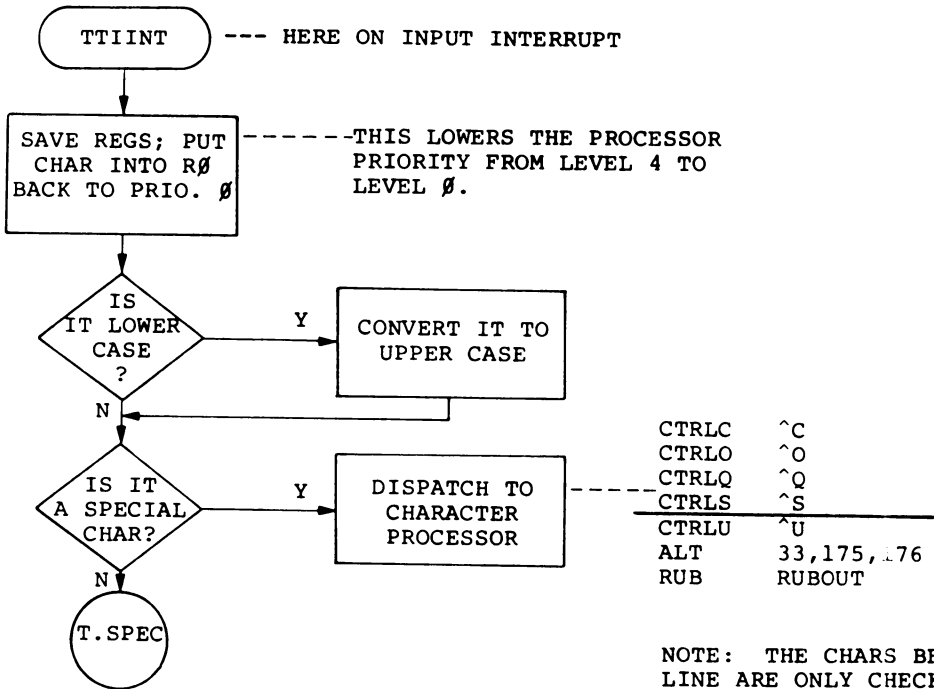
The interrupt service for the clock is primitive. The clock vector is set up such that the interrupt routine is always entered with the C bit = 1. At the interrupt routine, the code is:

```
ADC $TIME+2
ADC $TIME
RTI
```

Since the C bit is 1, \$TIME+2 is incremented by the ADC. When the low order word goes from 177777 to 0, the C bit remains on and \$TIME is then incremented. No 24 hour wrap around is provided.

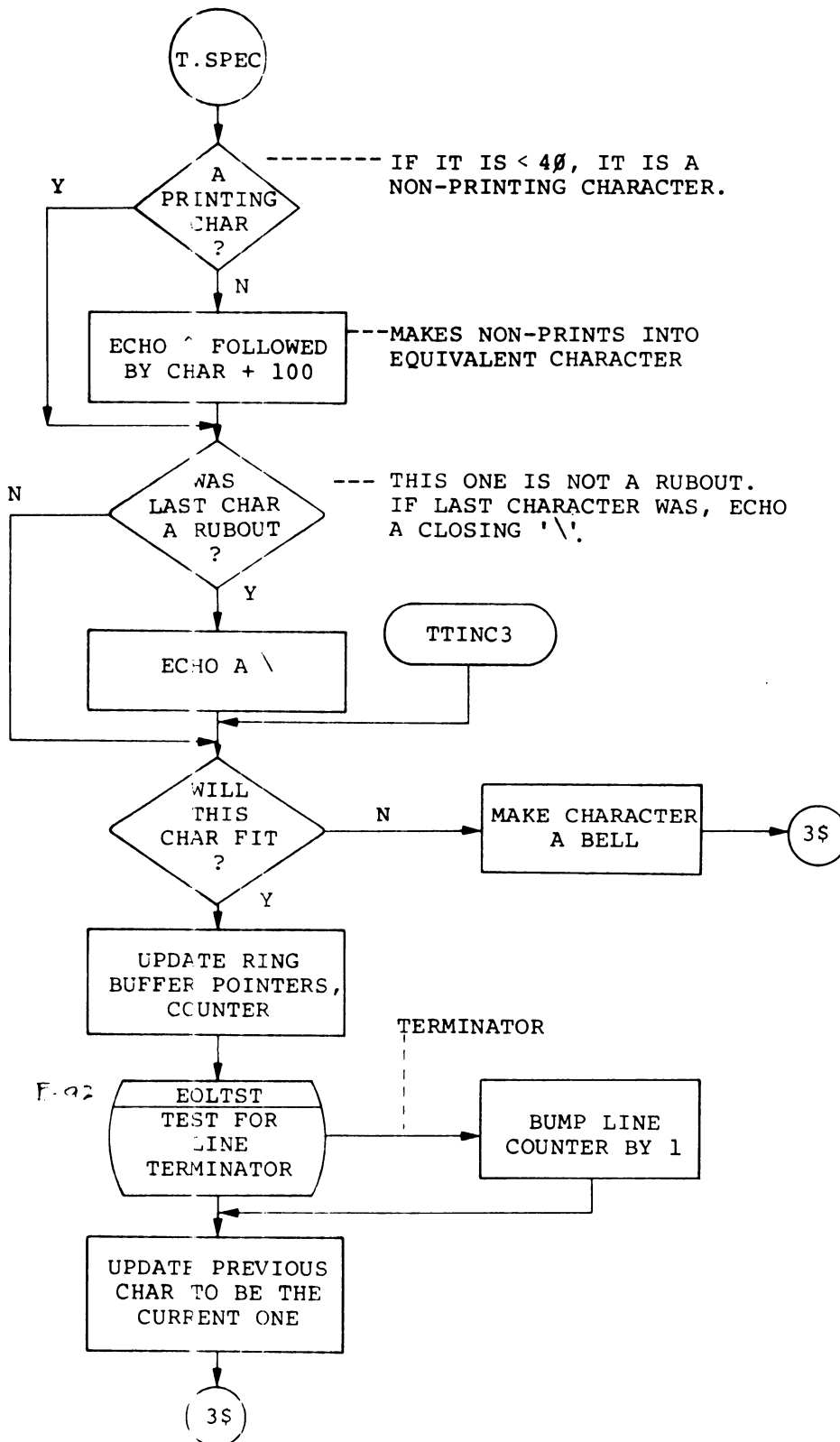
E.4.3 Console Terminal Interrupt Service

TT INPUT INTERRUPT SERVICE

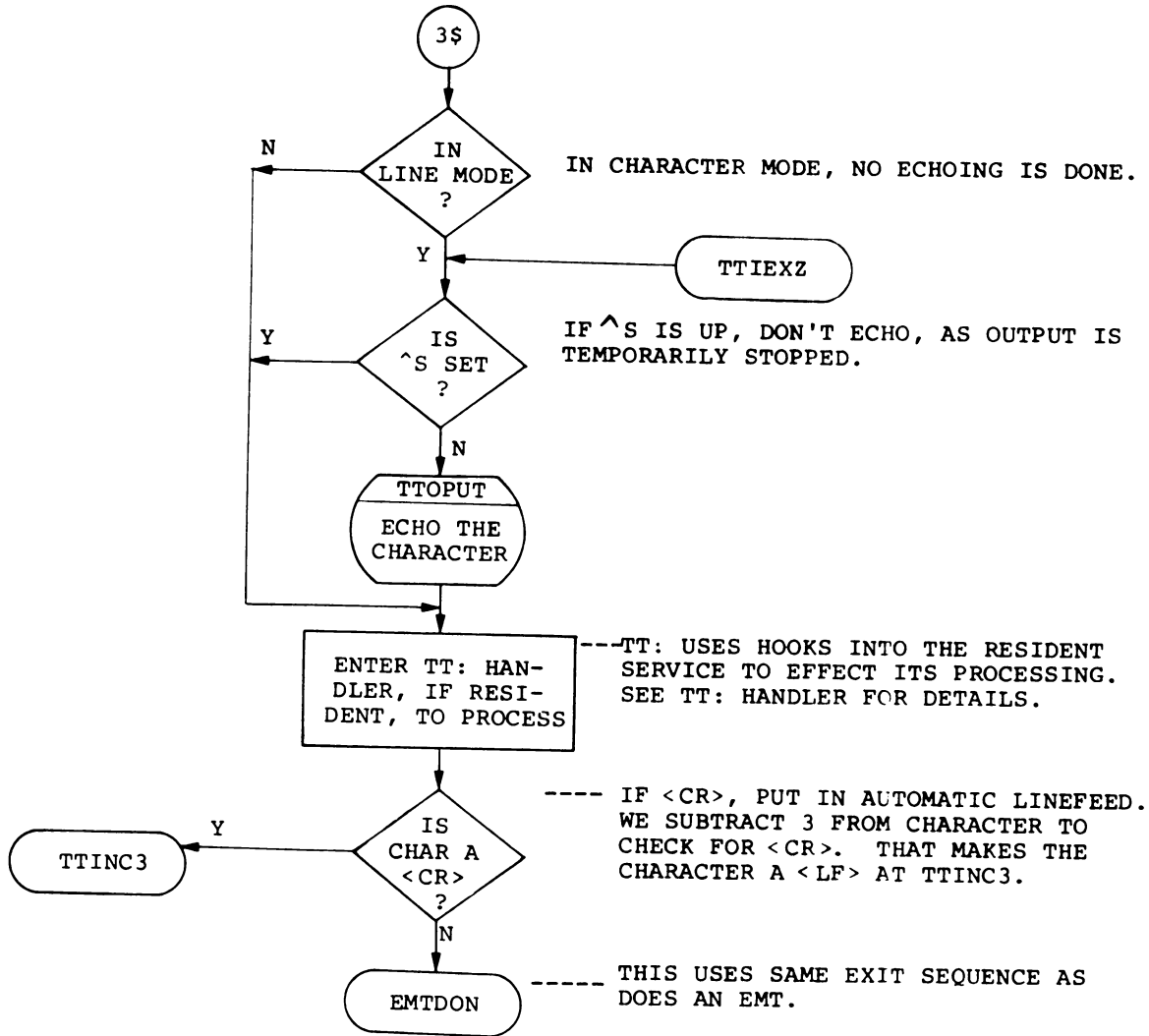


NOTE: THE CHARS BELOW THE LINE ARE ONLY CHECKED WHEN TT IS IN LINE MODE. IN CHARACTER MODE THEY ARE NOT ACTED UPON.

TT INPUT INTERRUPT SERVICE (CONT.)



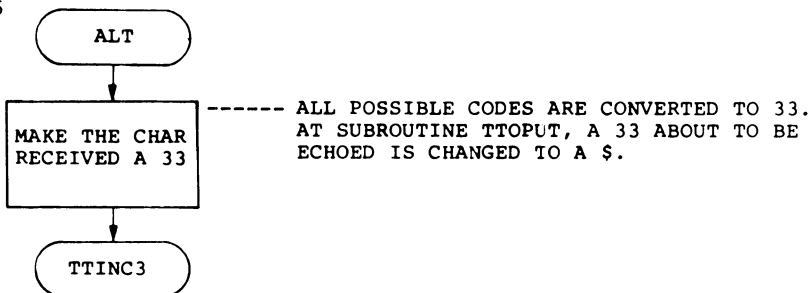
TT INPUT INTERRUPT SERVICE (CONT.)



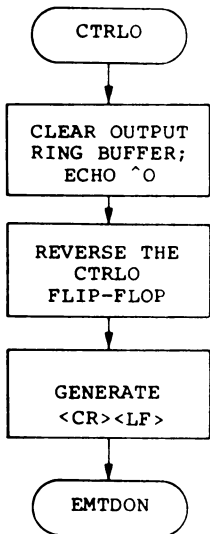
ALTMODE/CONTROL O, S, Q

These routines are entered when any of the corresponding special characters are struck.

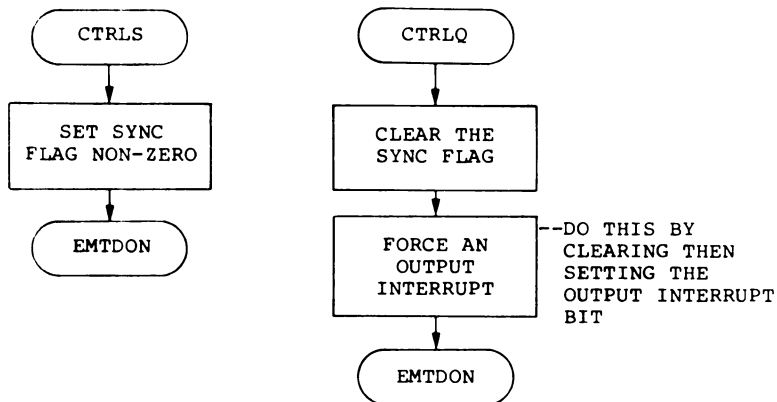
ALTMODE - 33,175,176



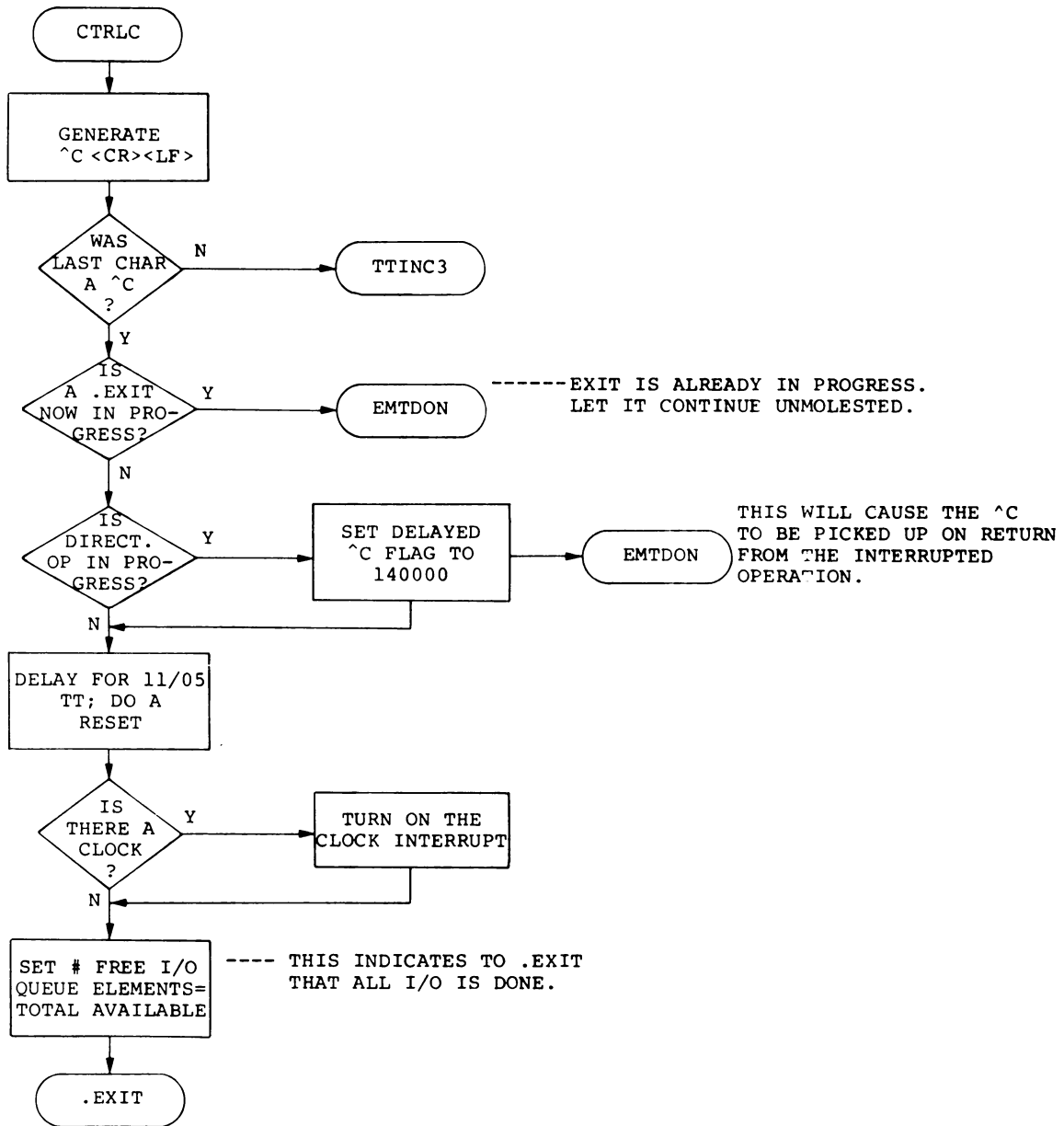
CONTROL O



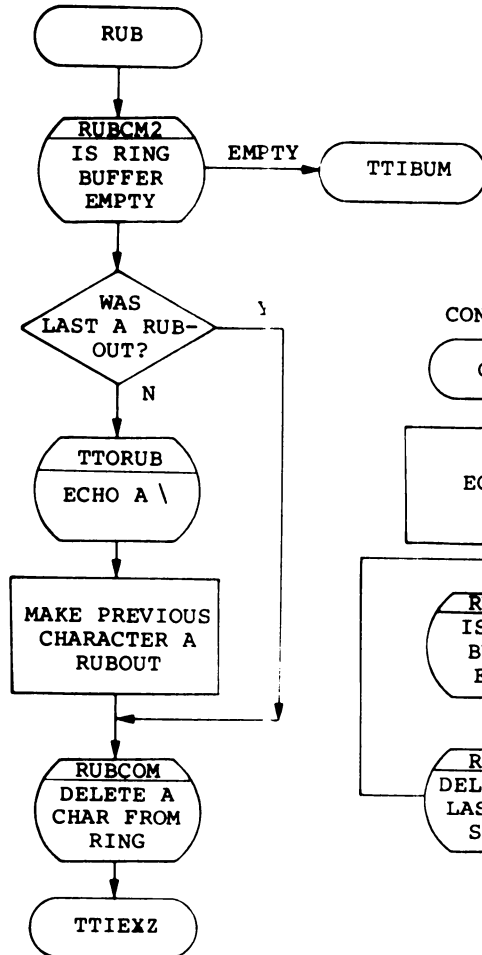
CONTROL S, CONTROL Q



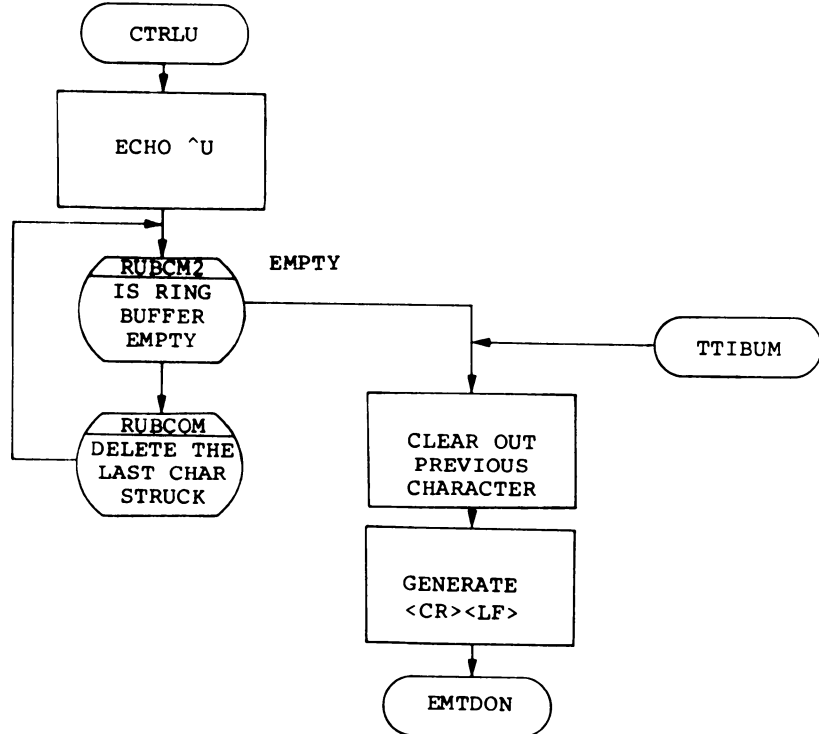
CONTROL C



RUBOUT

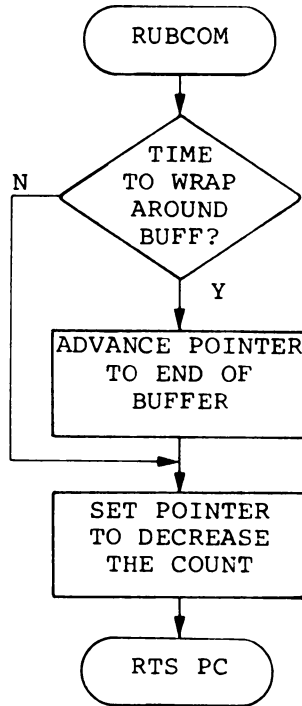


CONTROL U

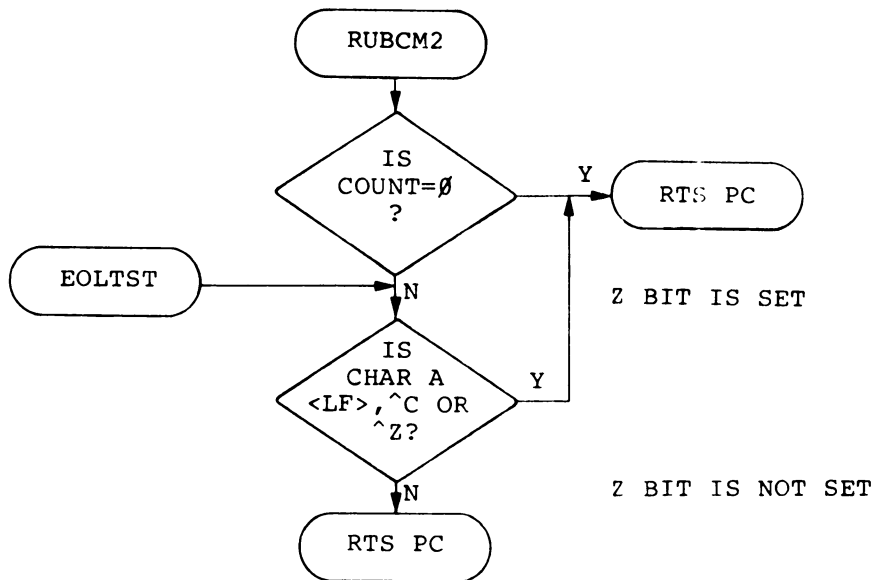


RUBCON/RUBCM2

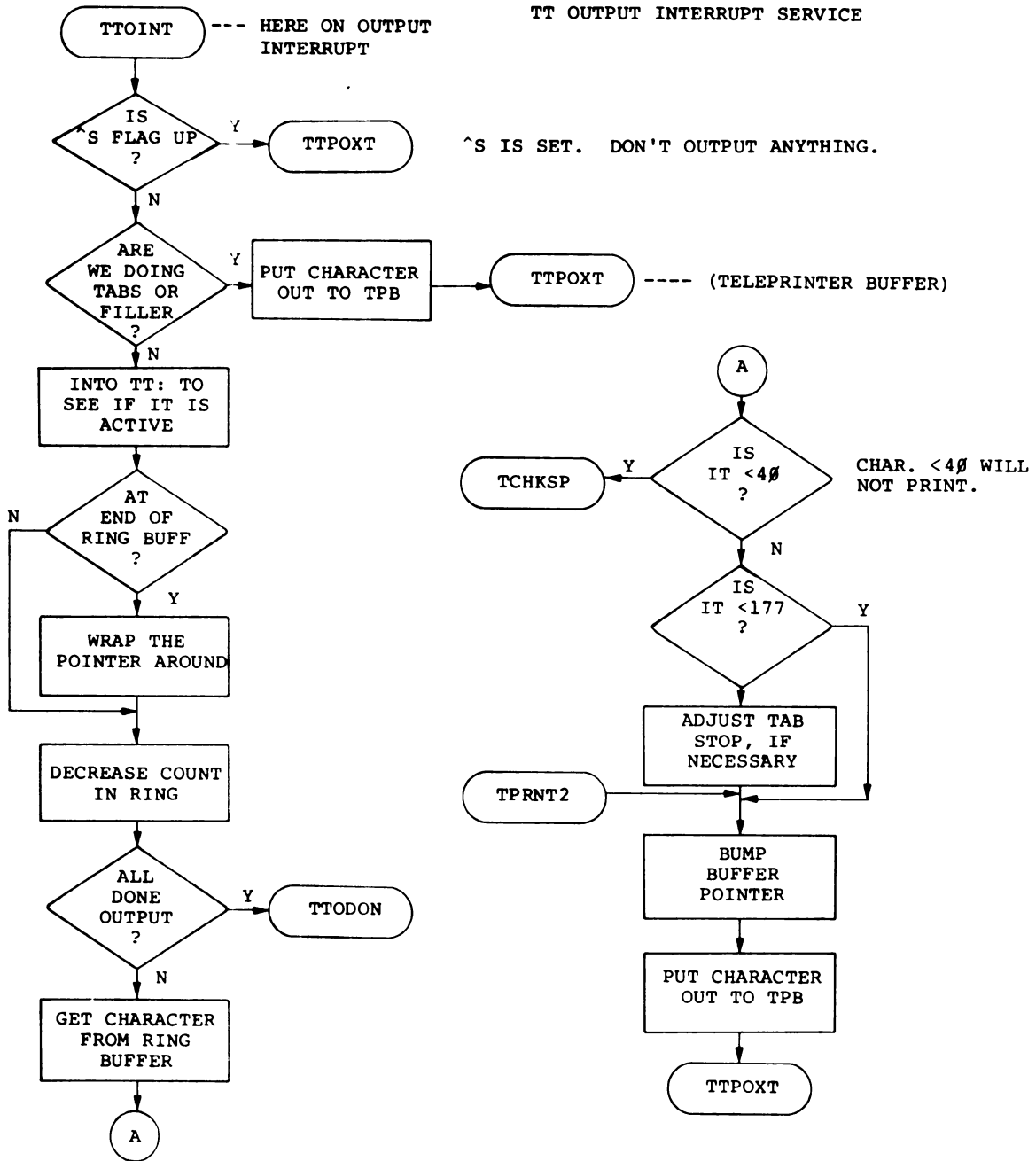
RUBCOM will update the input ring buffer pointers when a character is to be deleted.



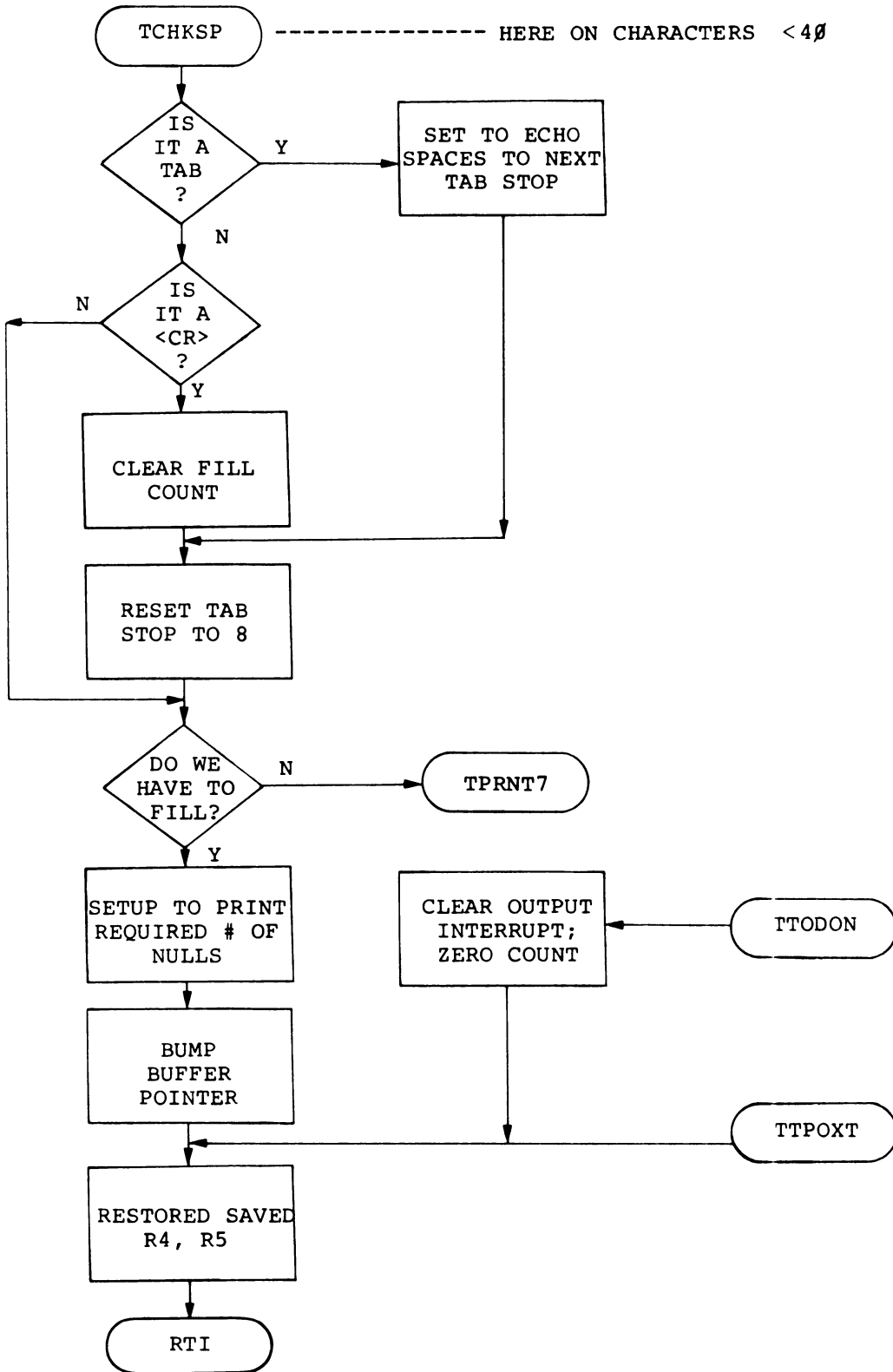
RUBCM2 checks to see if the ring buffer is empty. The buffer is empty if either the count = 0 or if the character to be deleted is a line terminator. This routine falls into routine EOLTST. The zero condition is returned if the buffer is empty.



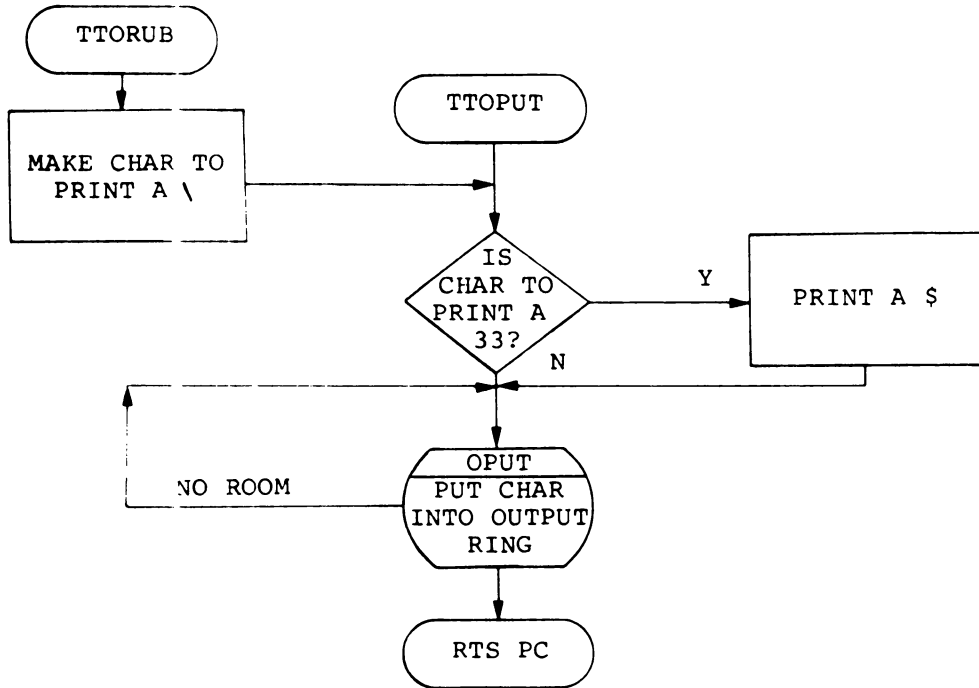
TT OUTPUT INTERRUPT SERVICE



TT OUTPUT INTERRUPT SERVICE (CONT.)

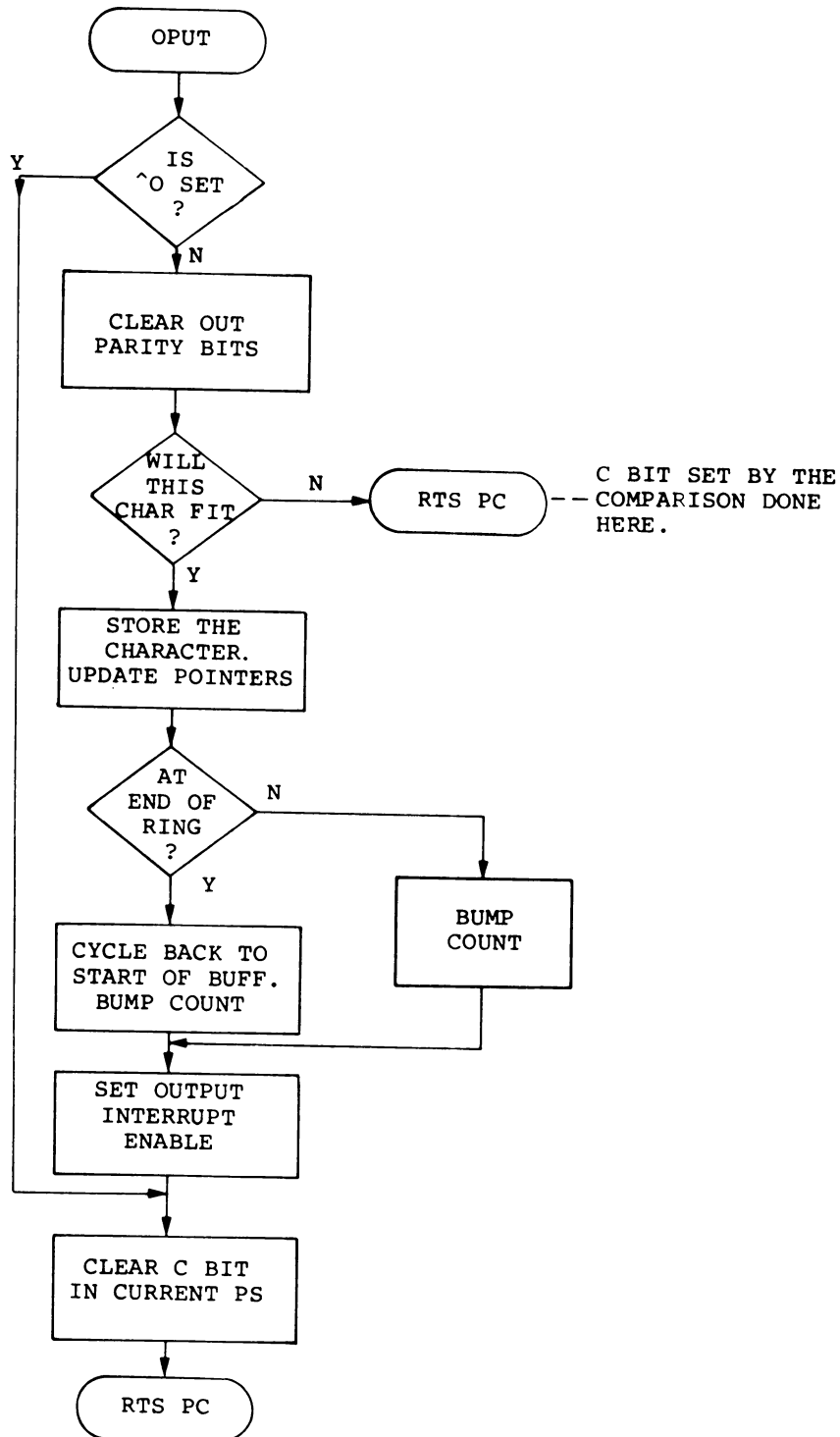


TTORUB and TTOPUT handle the printing of ALTMODE and RUBOUT. They print a \$ for ALTMODE and \ for RUBOUT.



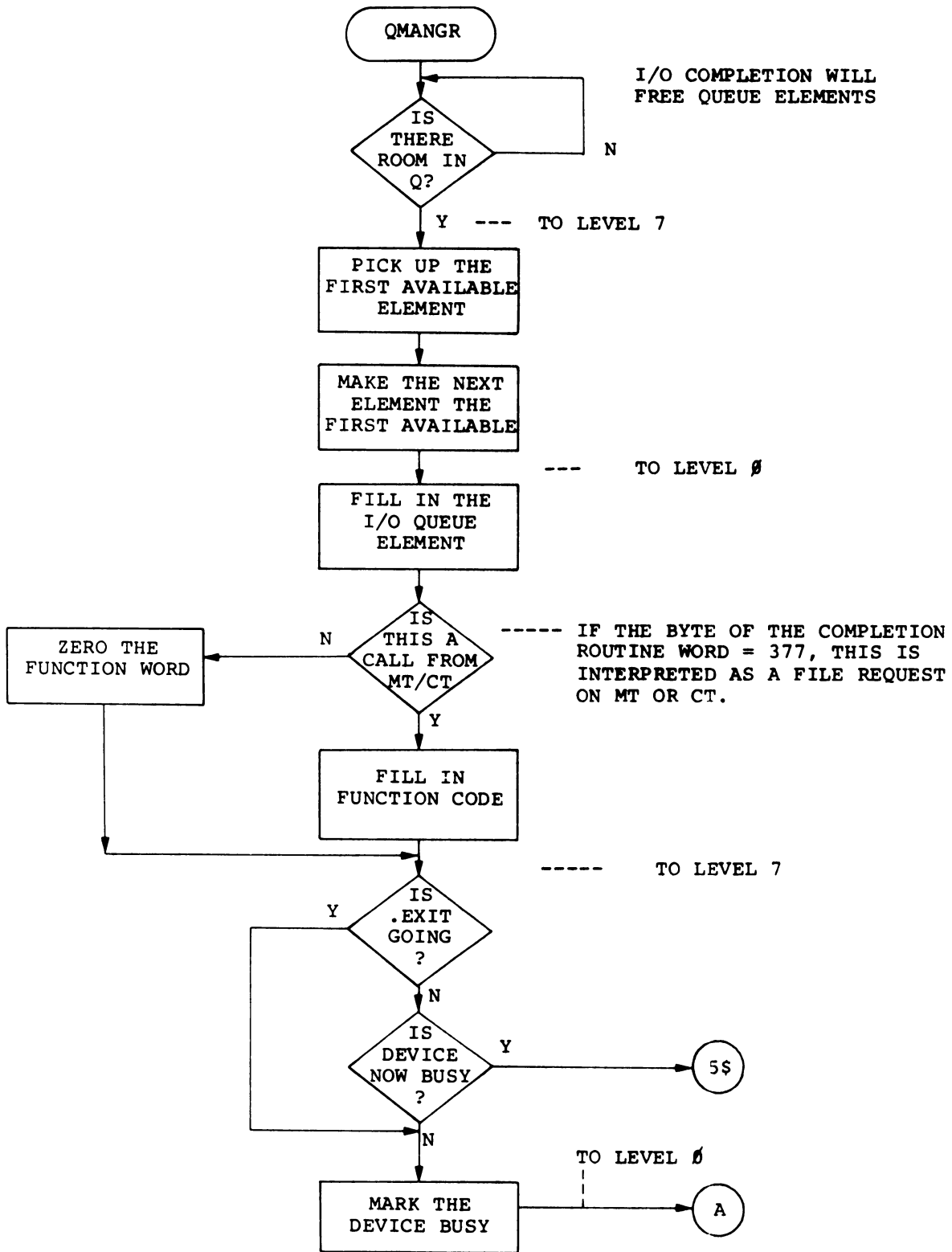
OPUT

OPUT actually puts the output character into the ring buffer. It updates the ring pointers and sets the interrupt enable bit. If the buffer is full, it returns with the C bit set.

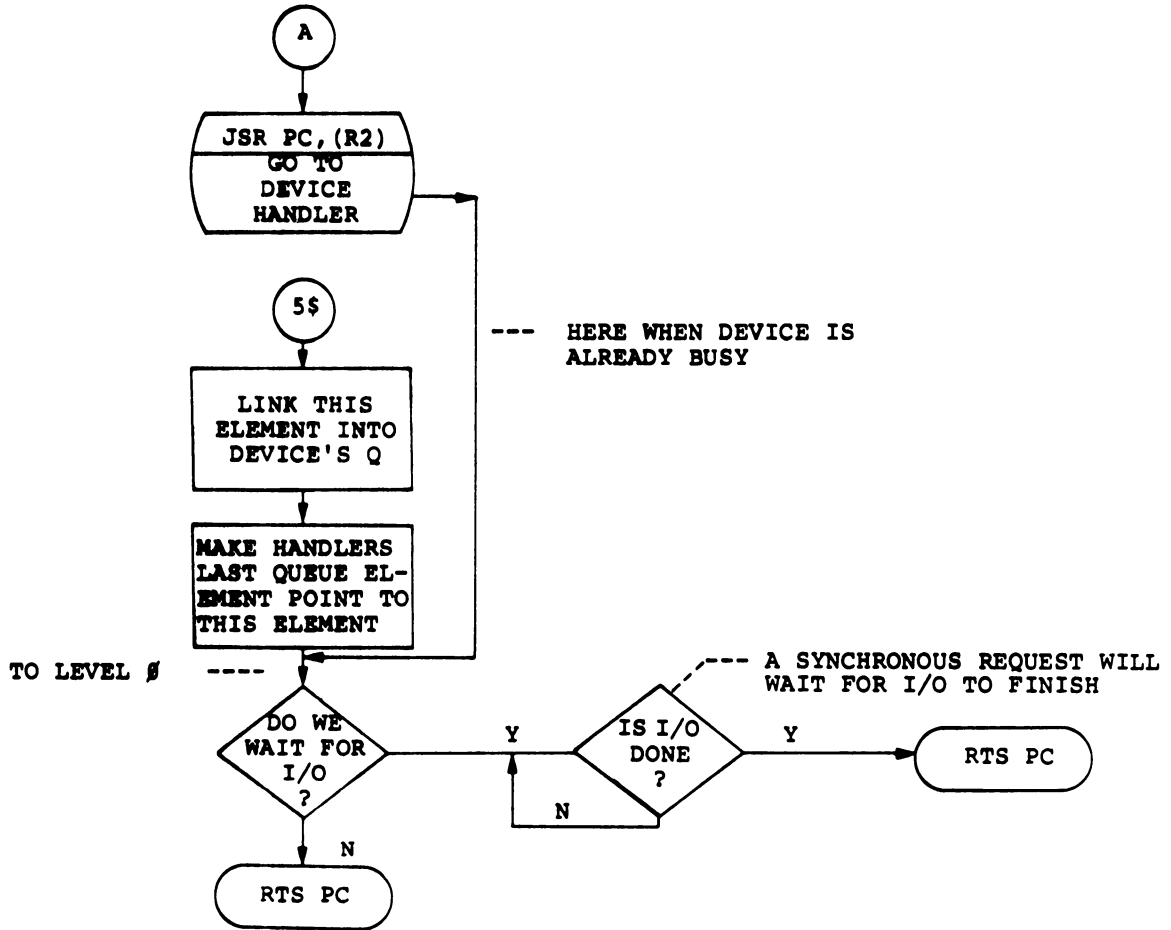


E.4.4 I/O Routines

I/O QUEUE MANAGEMENT ROUTINES

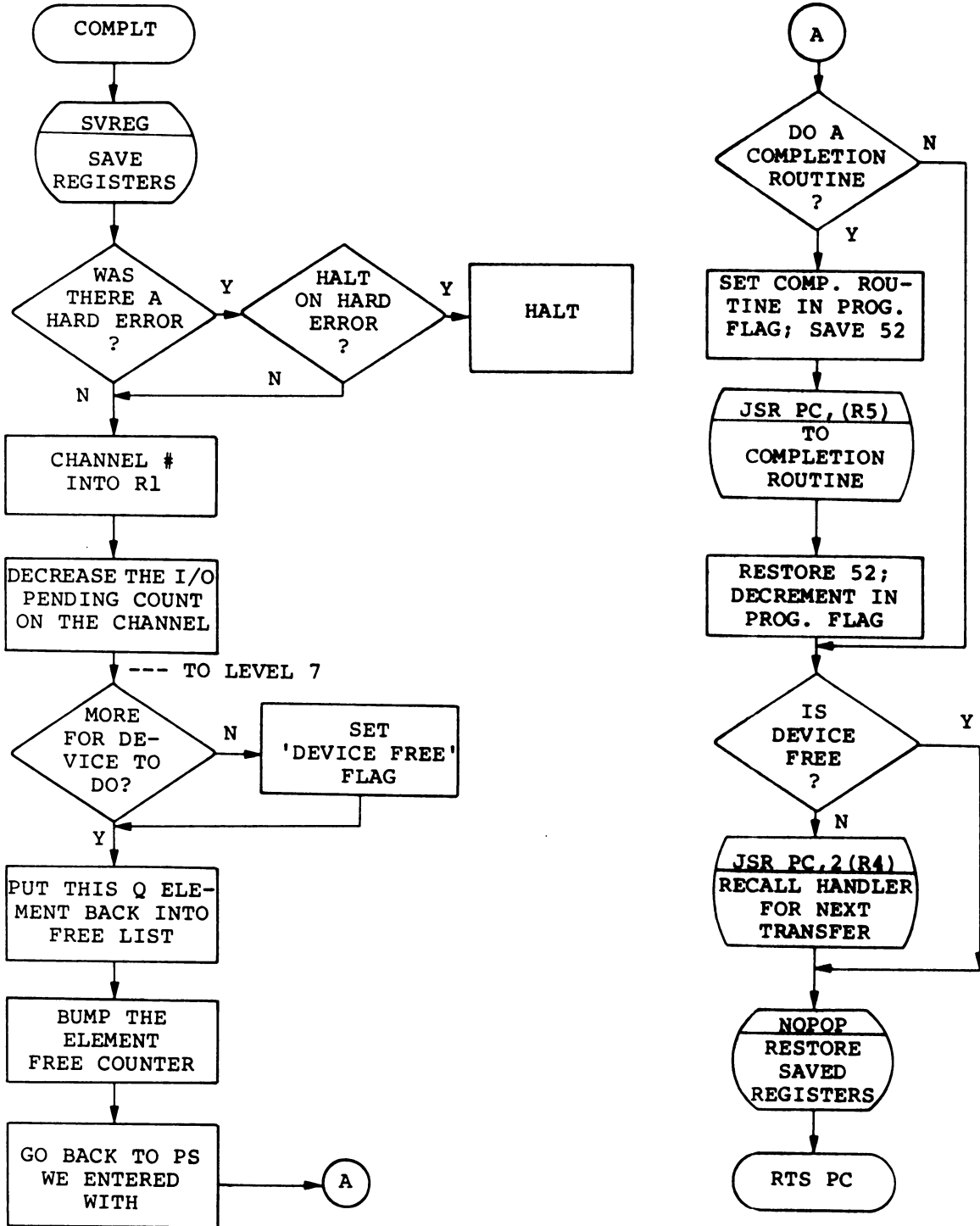


I/O QUEUE MANAGEMENT ROUTINES (CONT.)



I/O QUEUE COMPLETION

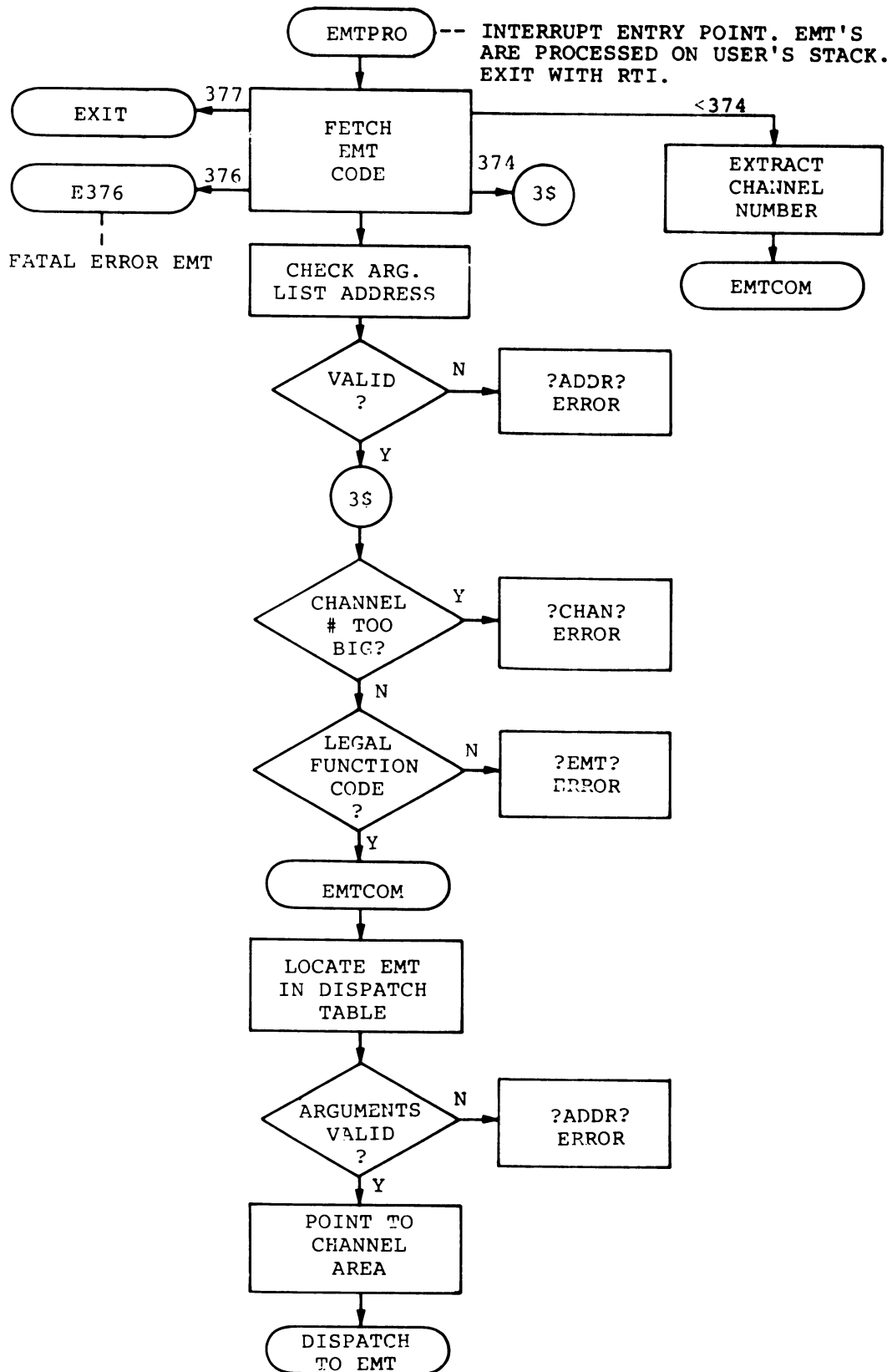
COMPLT is entered when an I/O transfer finishes.

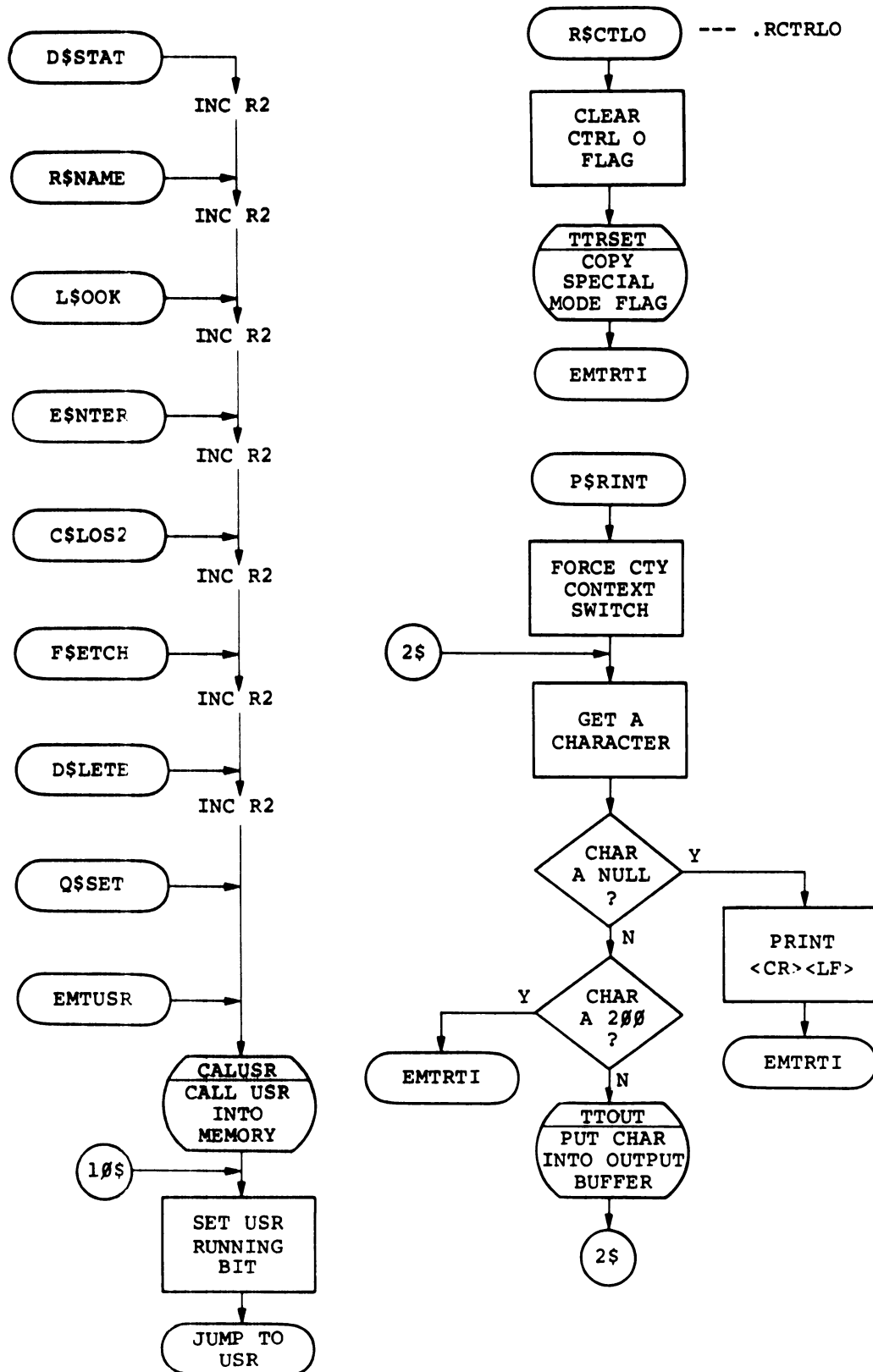


**E.5 RMON (RESIDENT MONITOR) FLOWCHARTS FOR
BACKGROUND/BACKGROUND MONITOR**

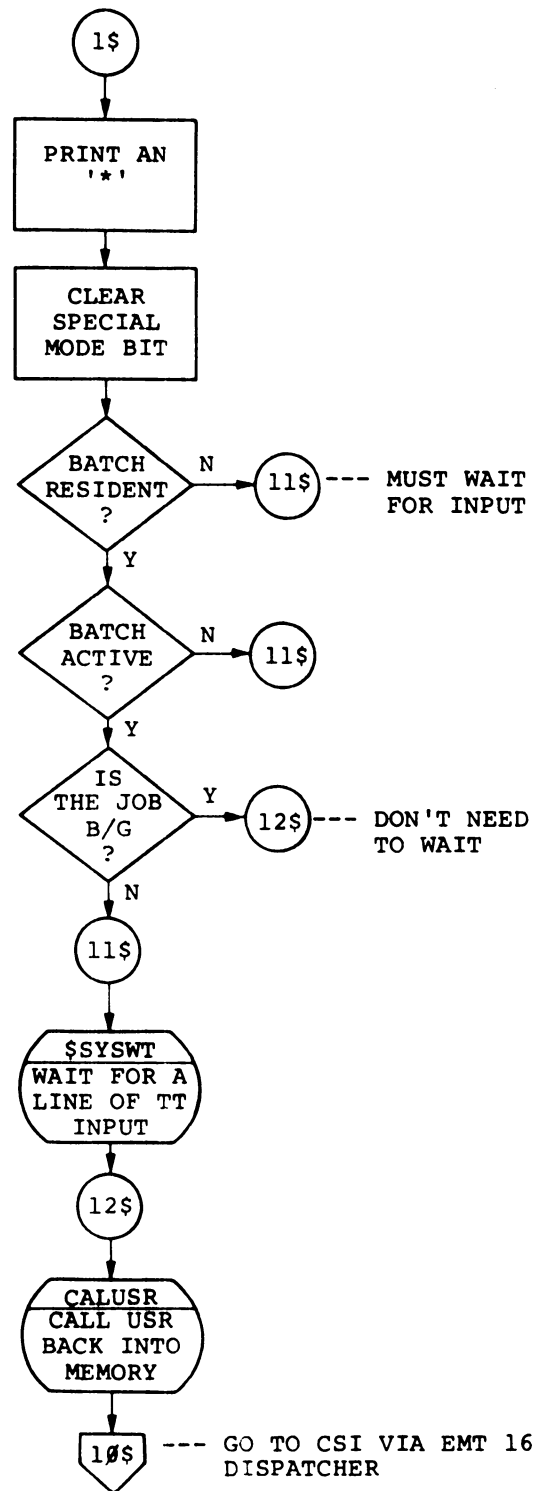
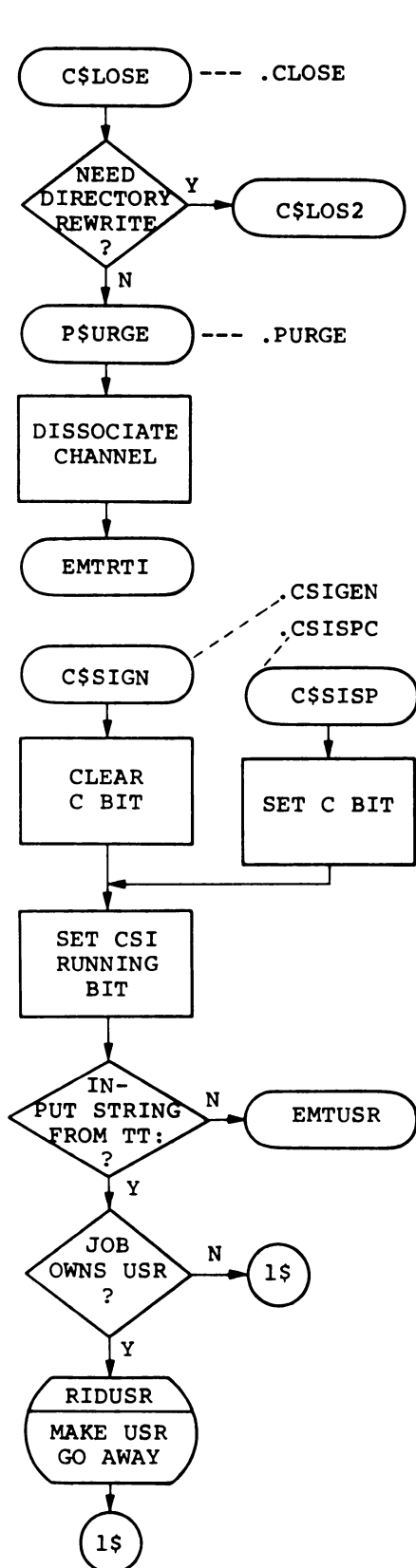
E.5.1 EMT Processors

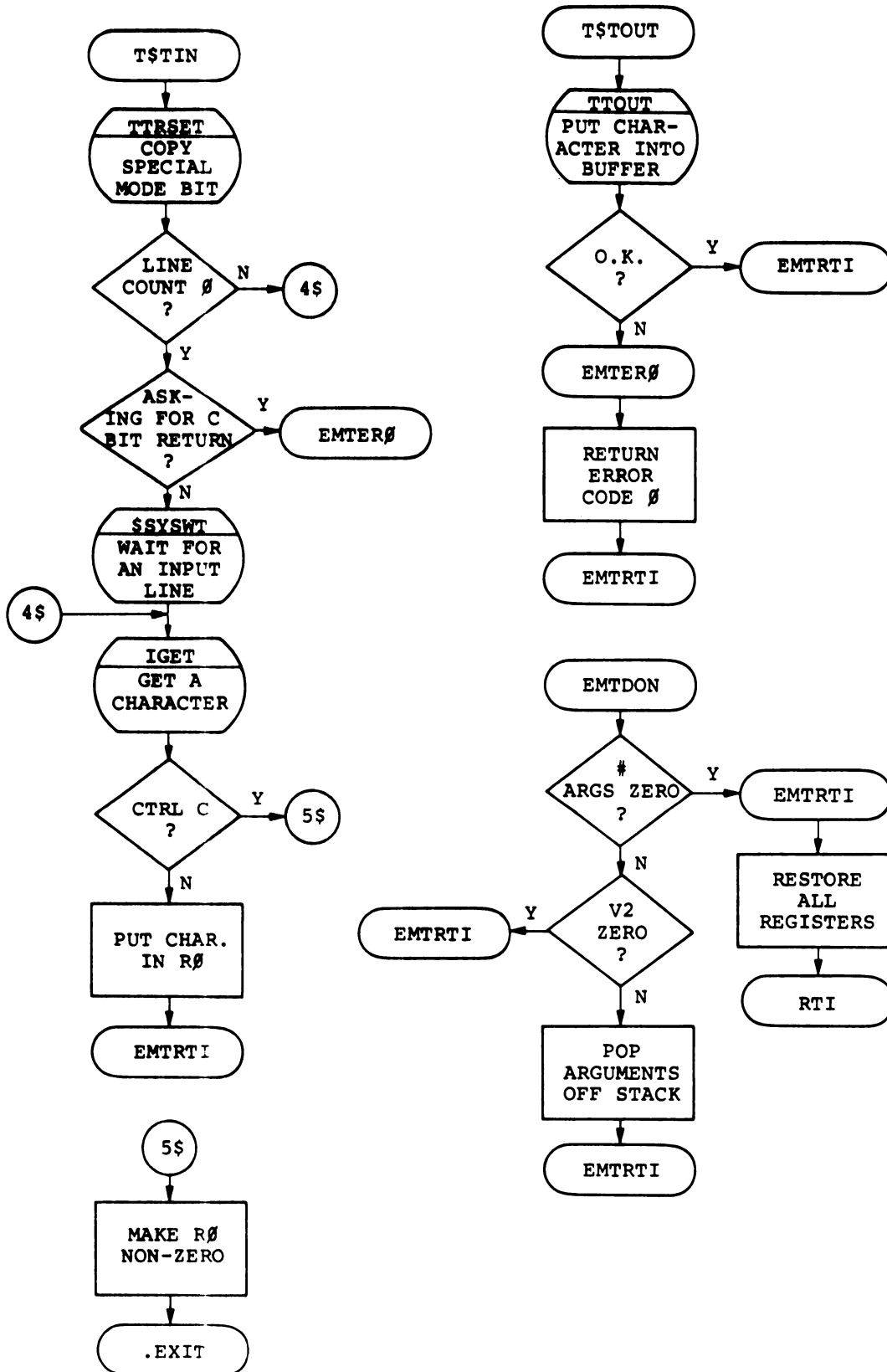
EMT DISPATCHER



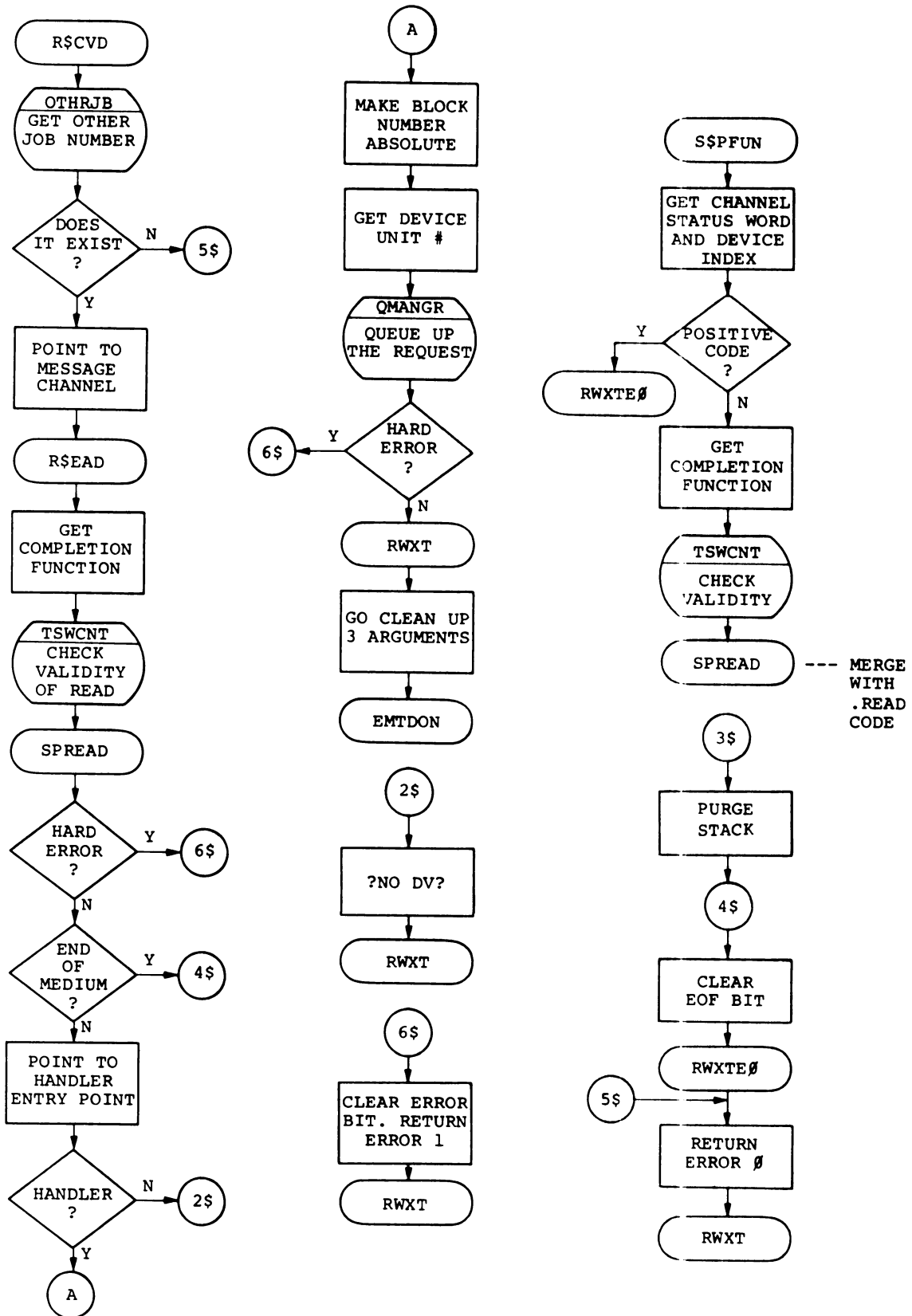


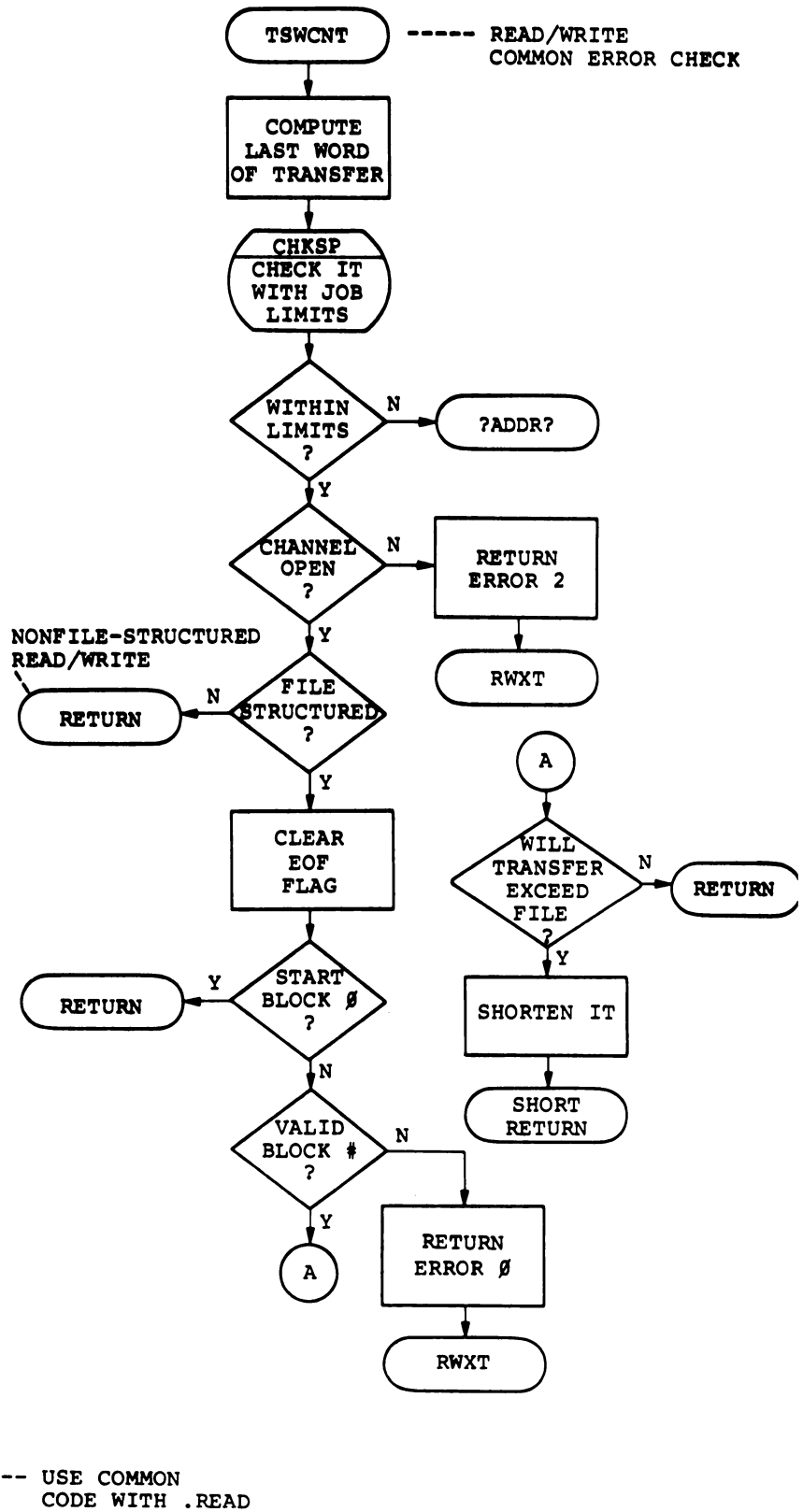
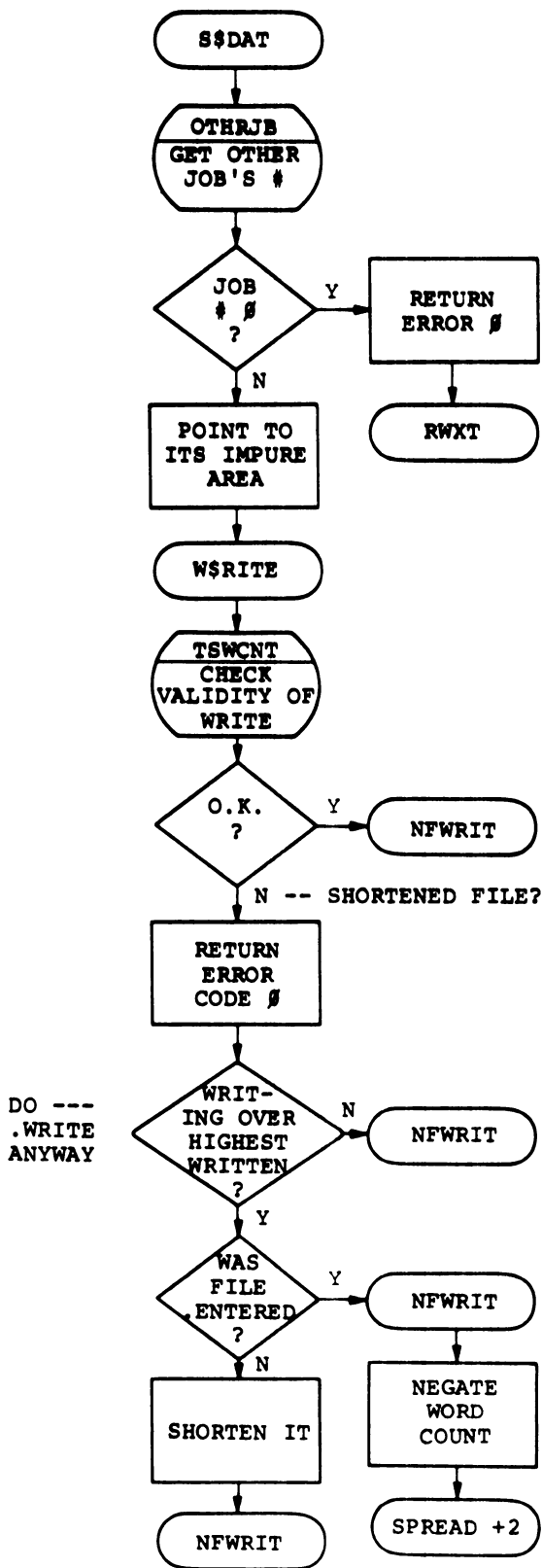
.CLOSE, .PURGE, .CSISPC, .CSIGEN



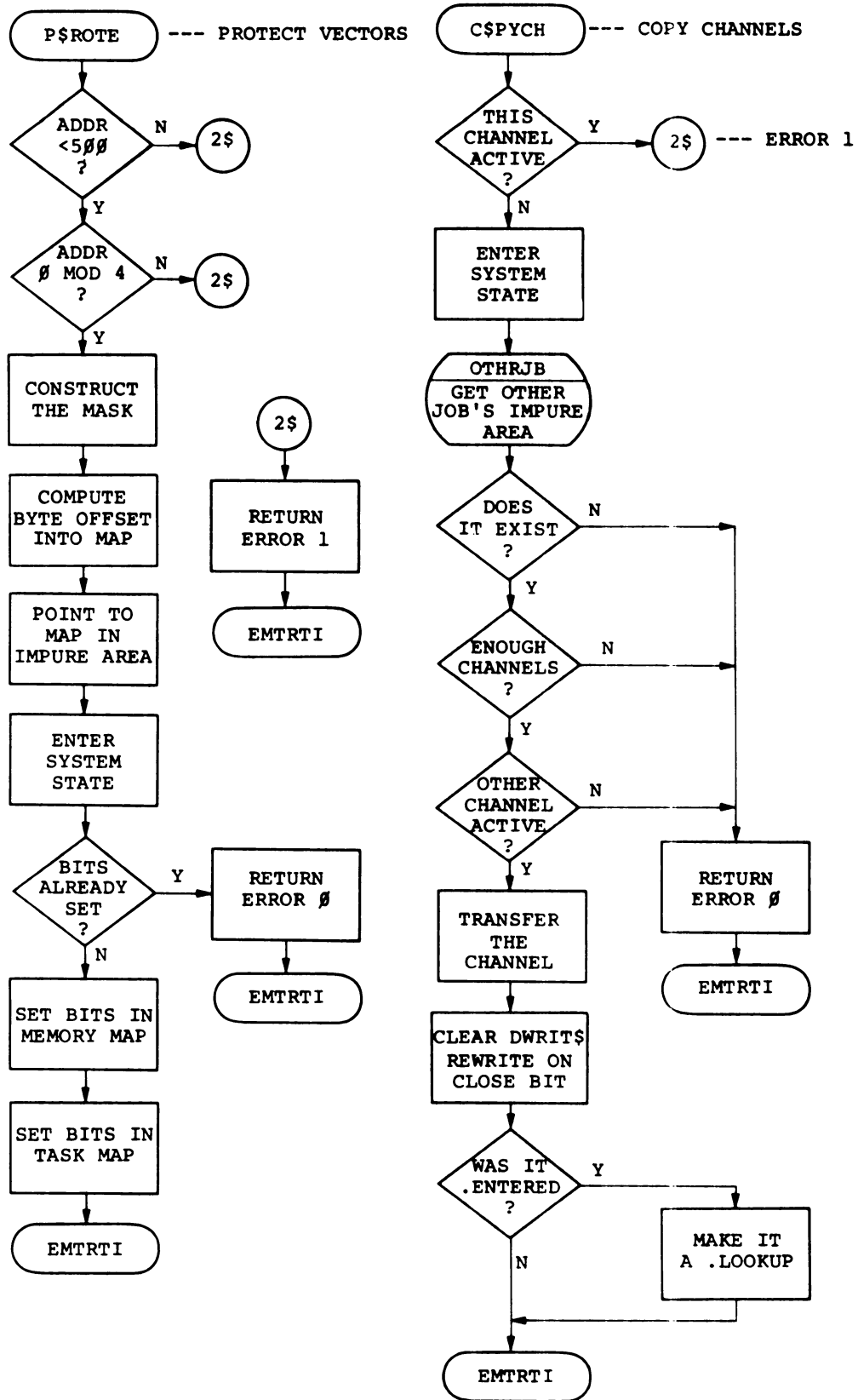


.READ, .RCVD, .SPFUN

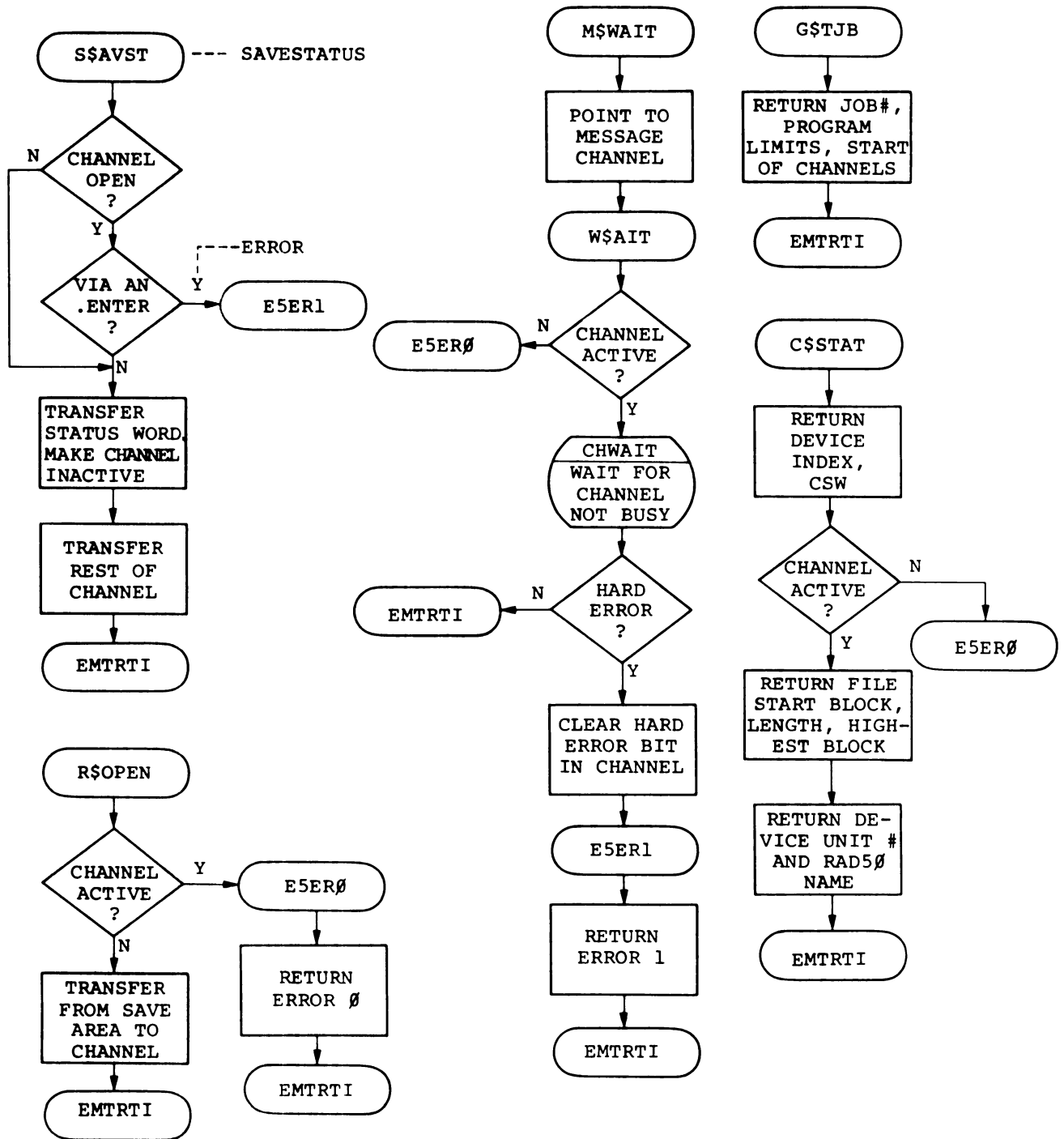




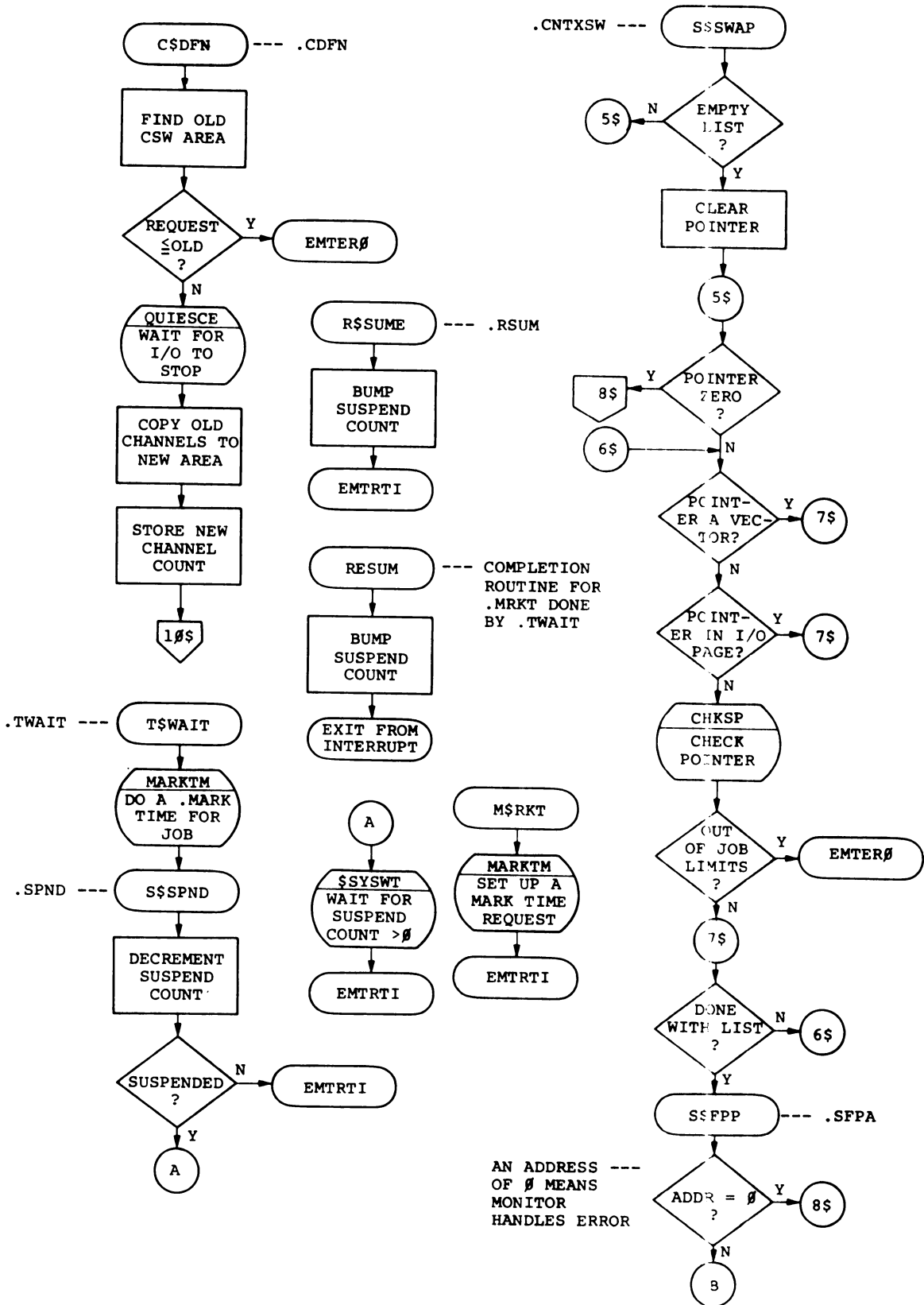
.PROTECT, .CPYCH

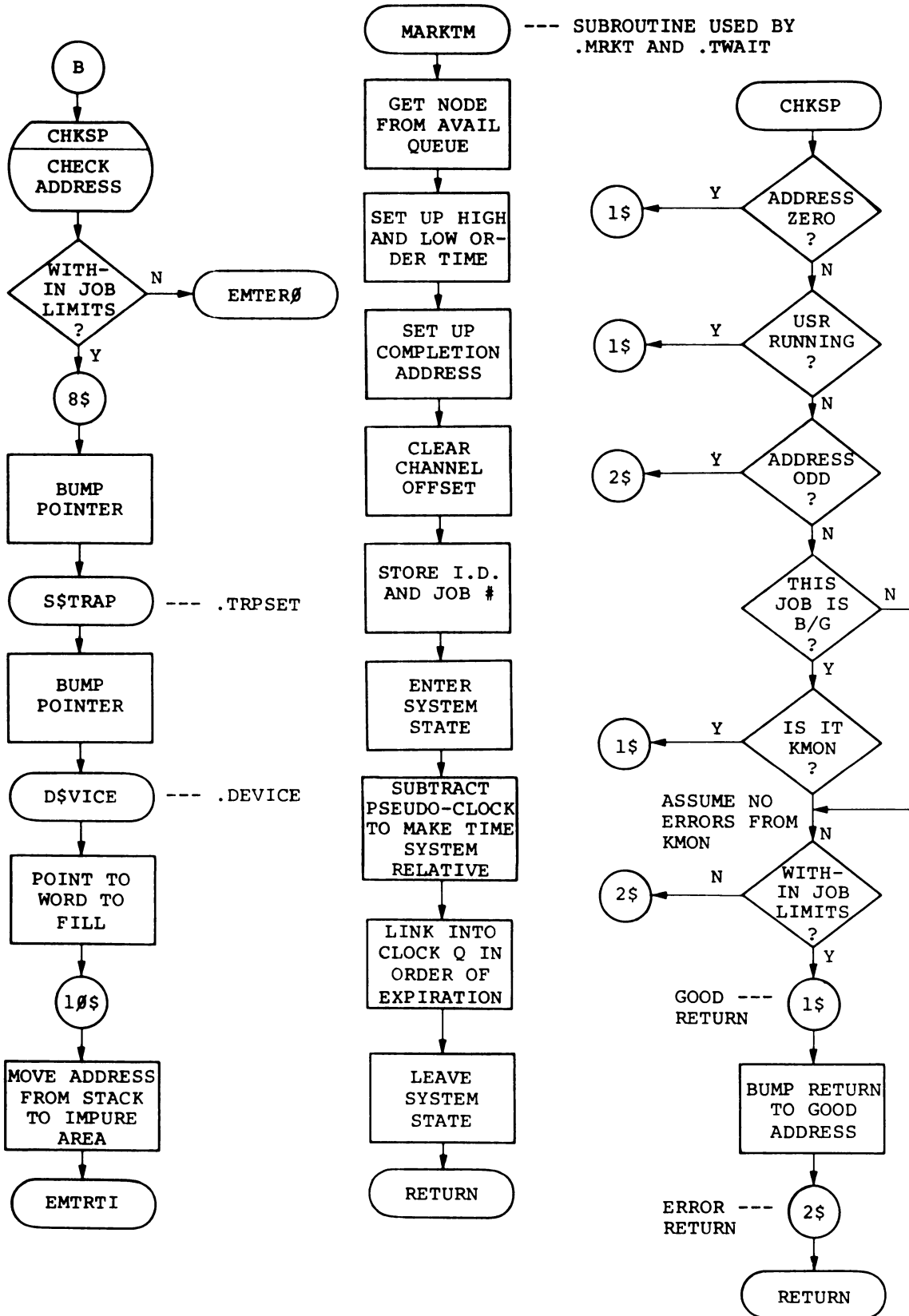


.SAVESTATUS, .REOPEN,
 .M\$WAIT, .WAIT,
 .GTJB, .CSTATUS

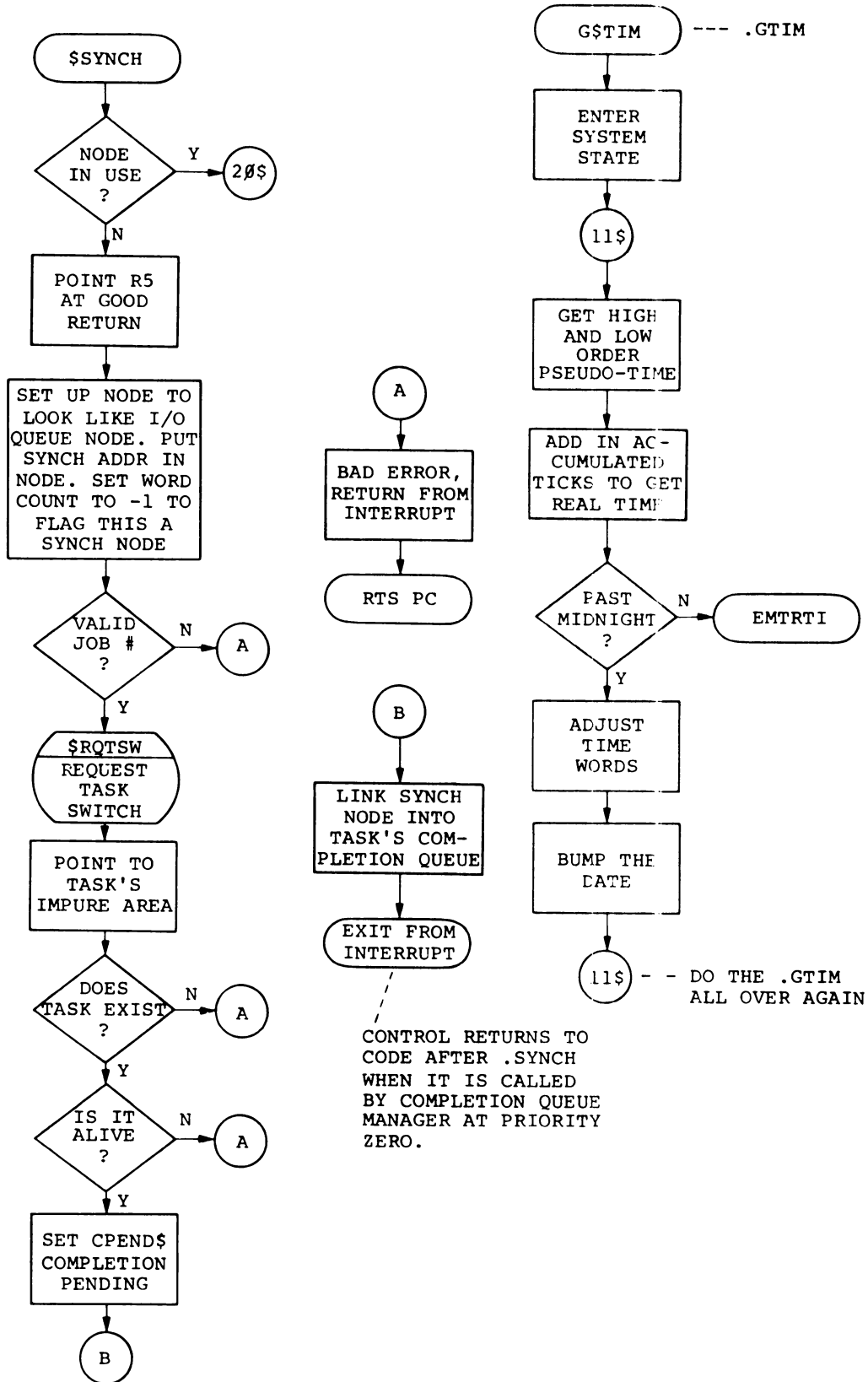


.CDFN, .TWAIT, .SPND, .RSUM,
.CNTXSW, .SFPA, .TRPSET, .DEVICE

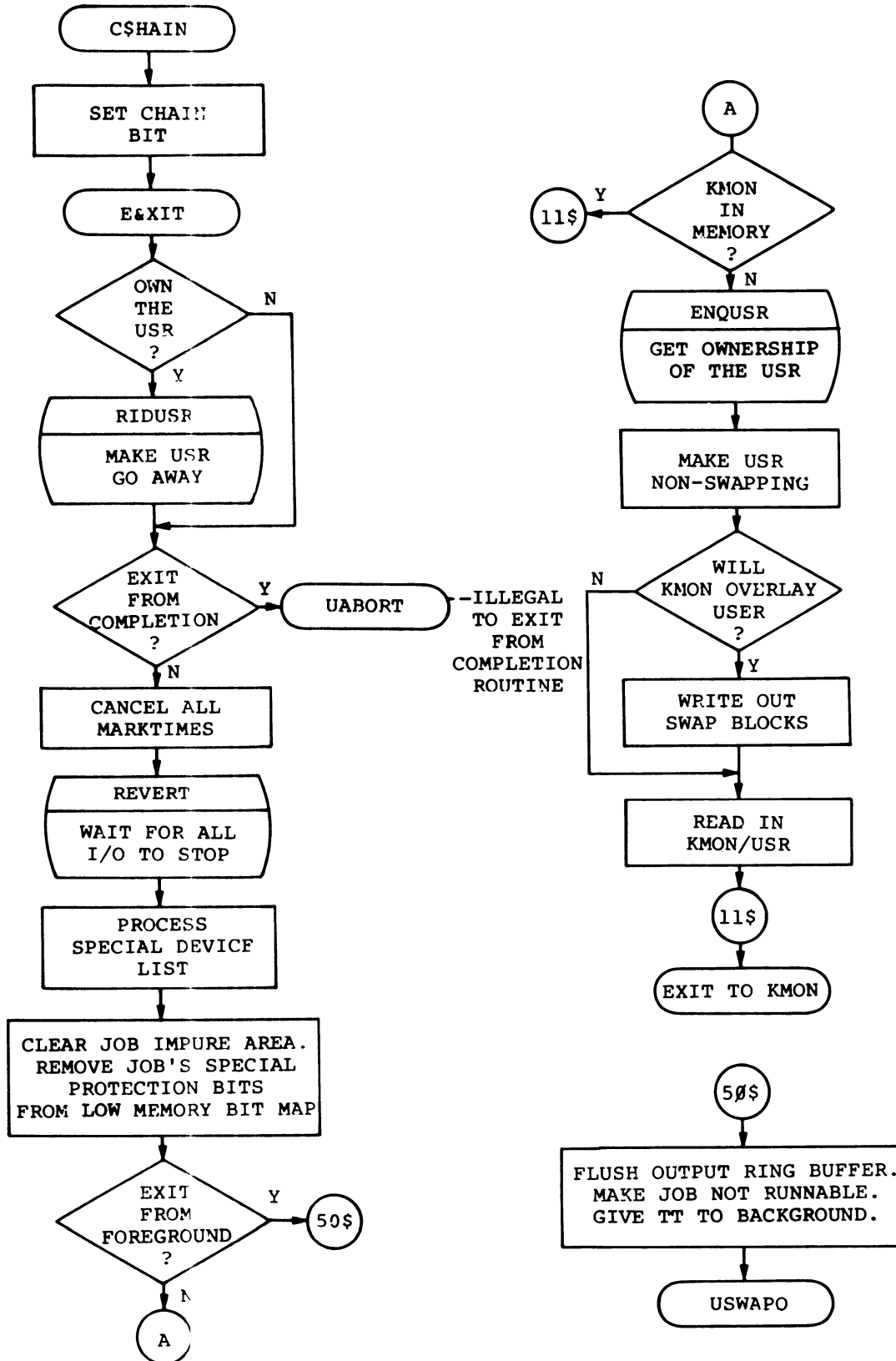




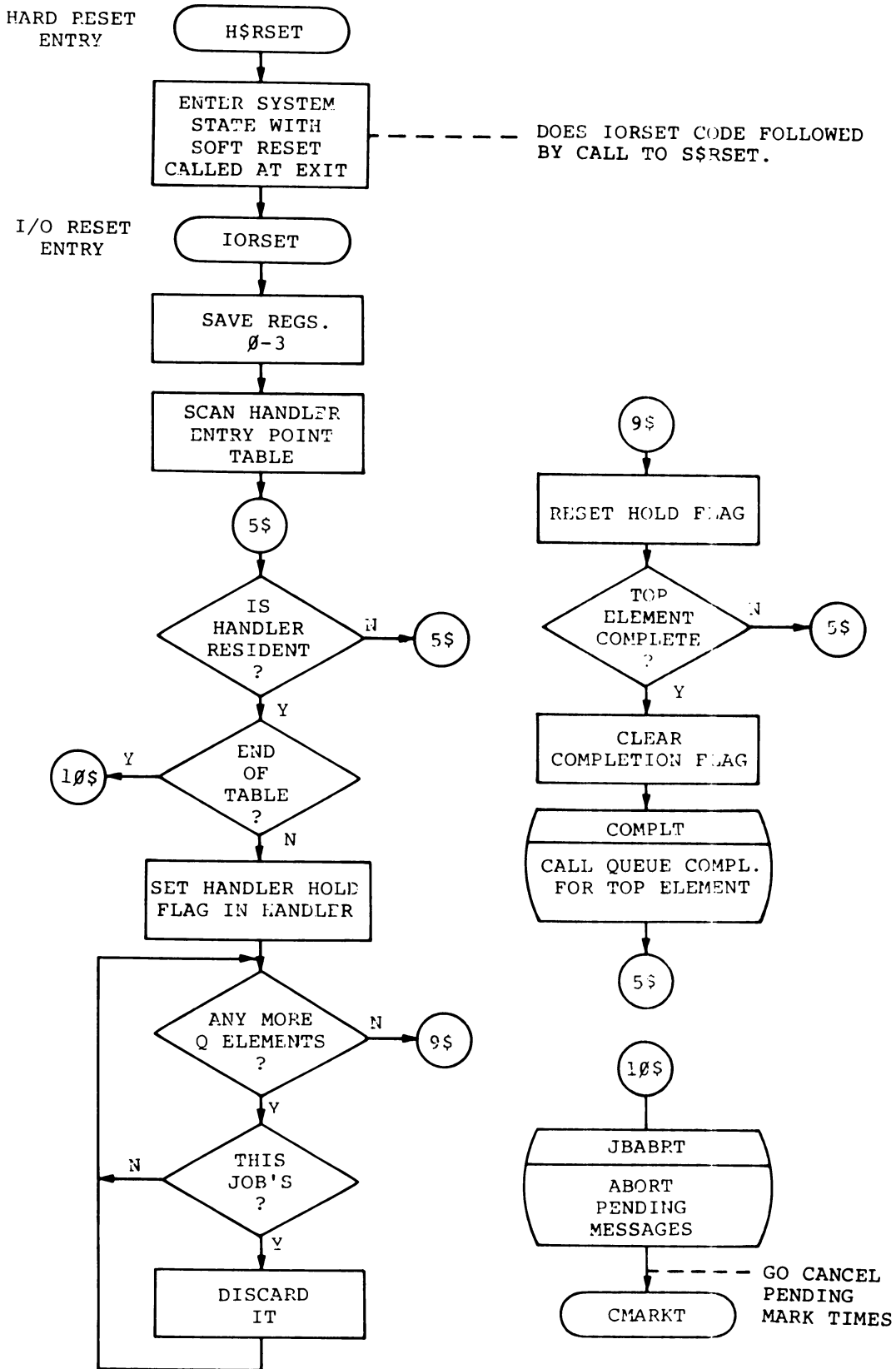
.SYNCH, .GTIM



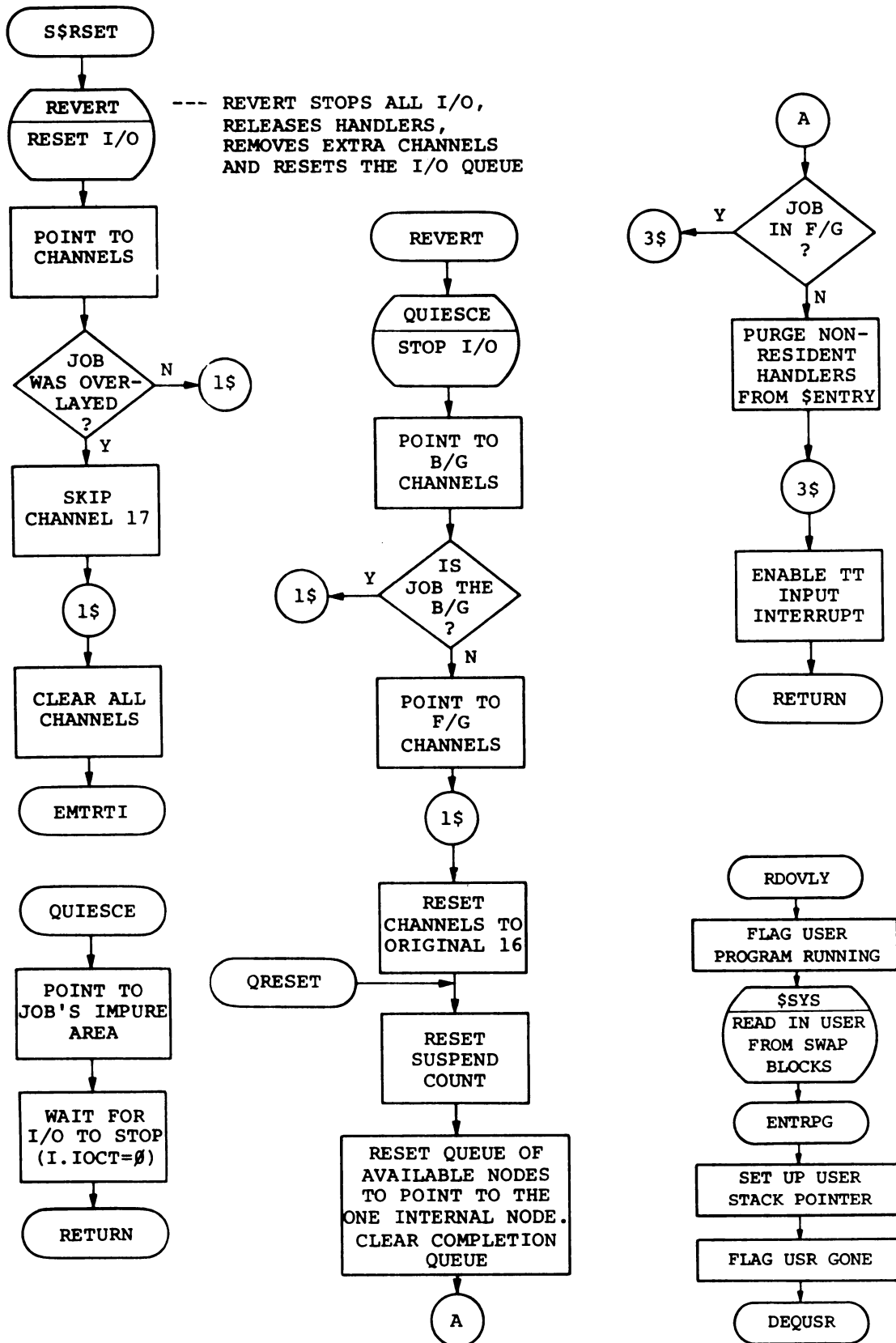
.EXIT
.CHAIN



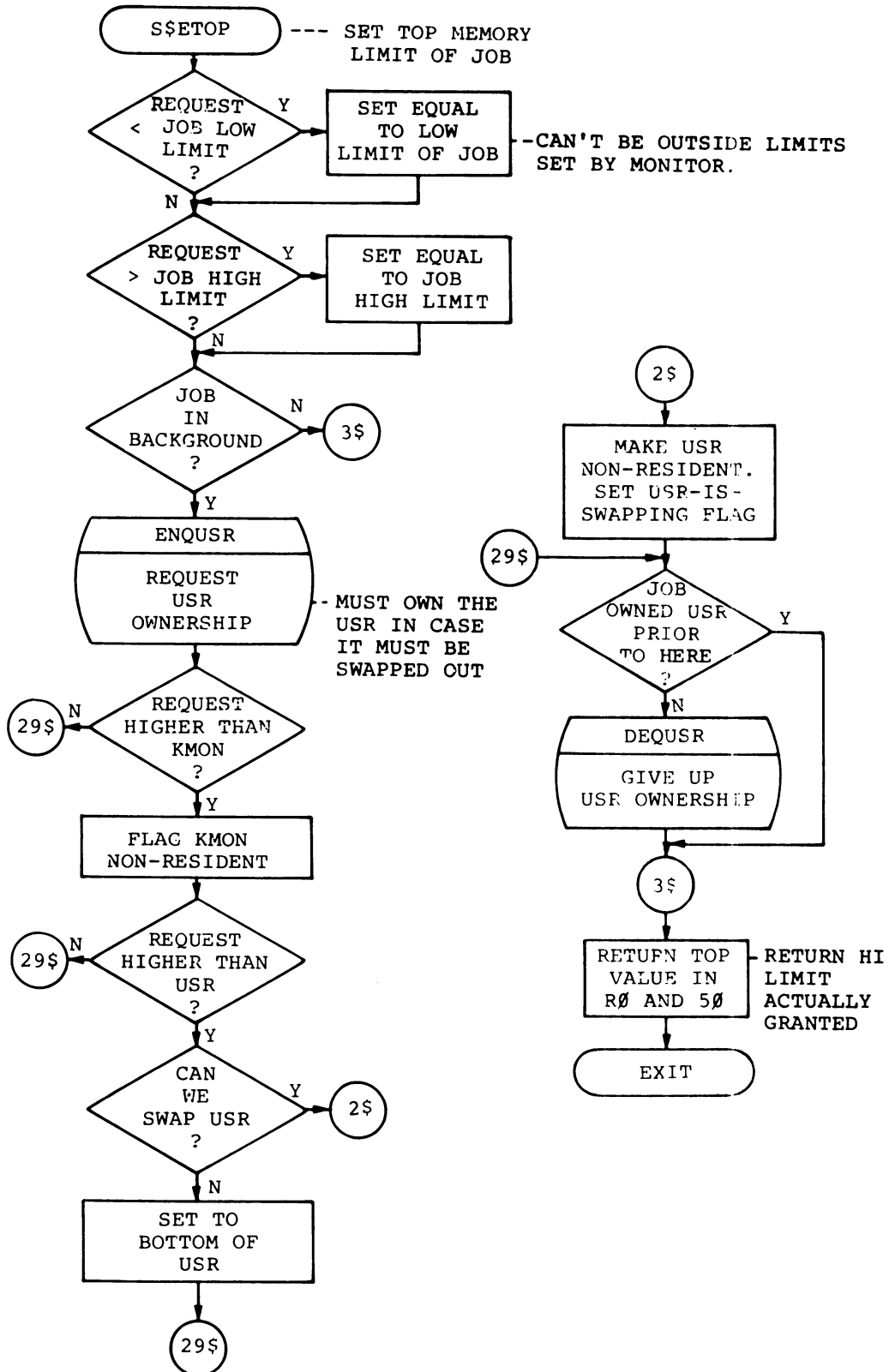
HARD AND SOFT RESET



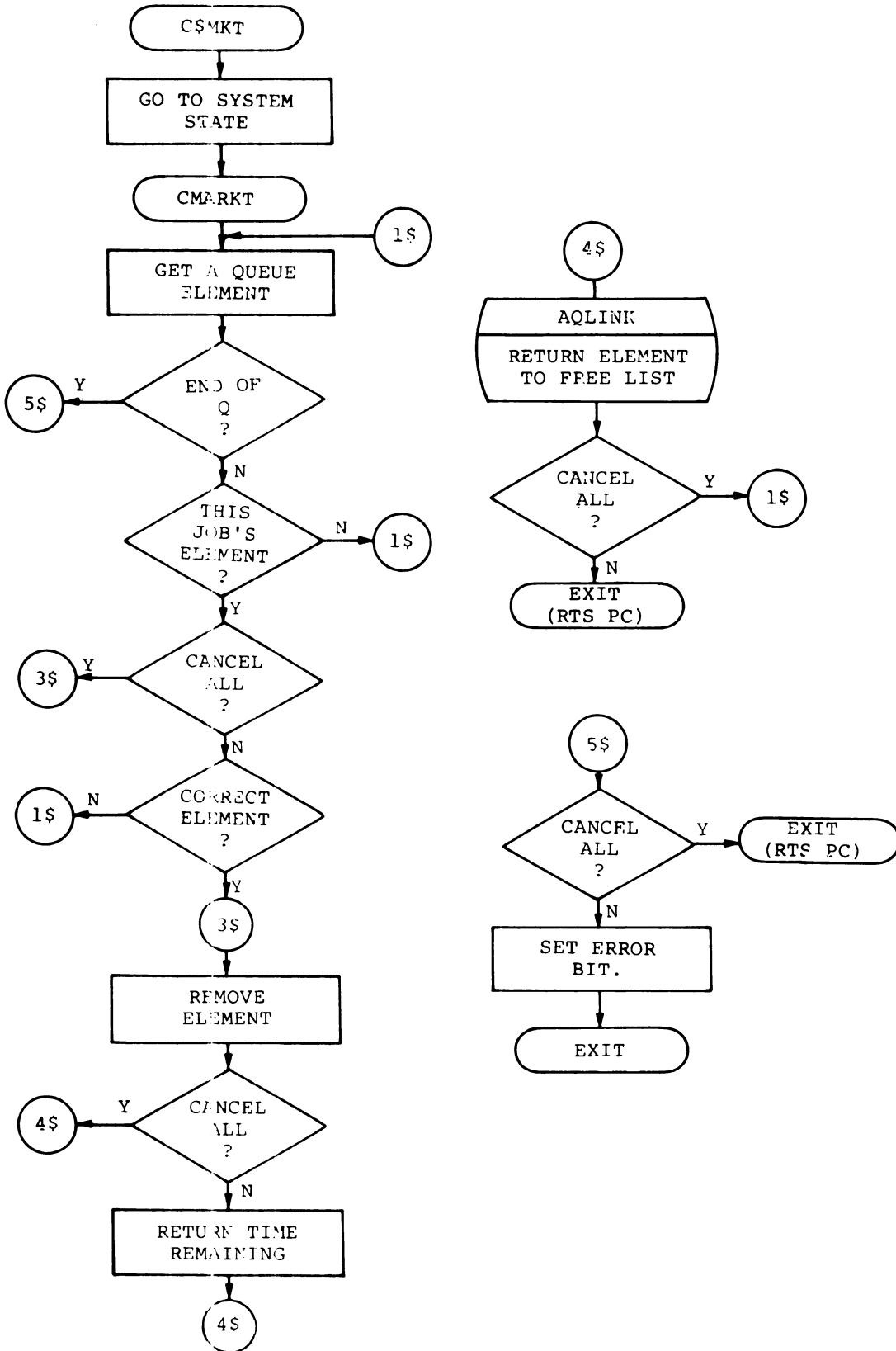
HARD AND SOFT RESET (CONT.)/RDOVLY



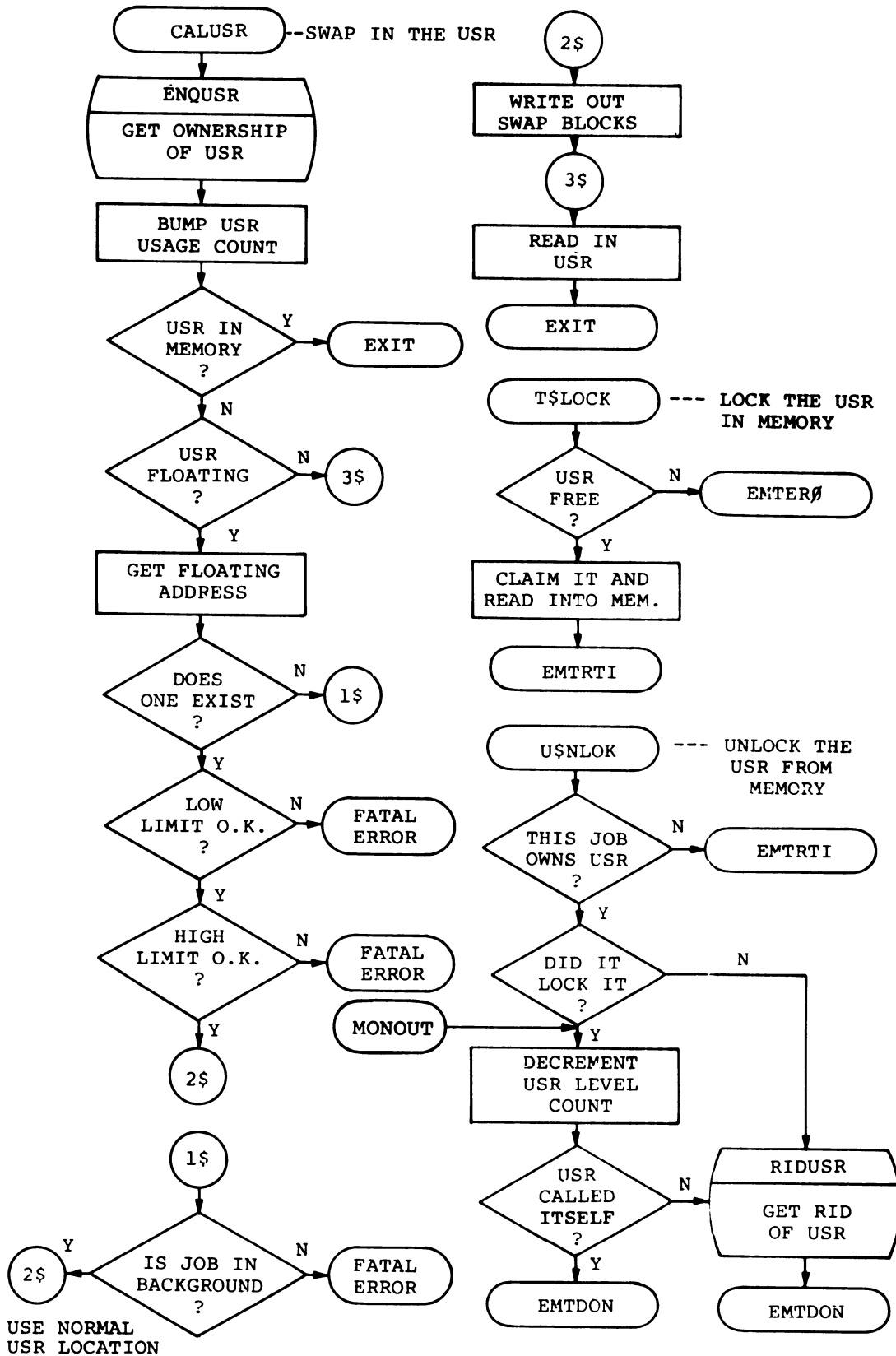
.SETTOP



CANCEL MARK TIME

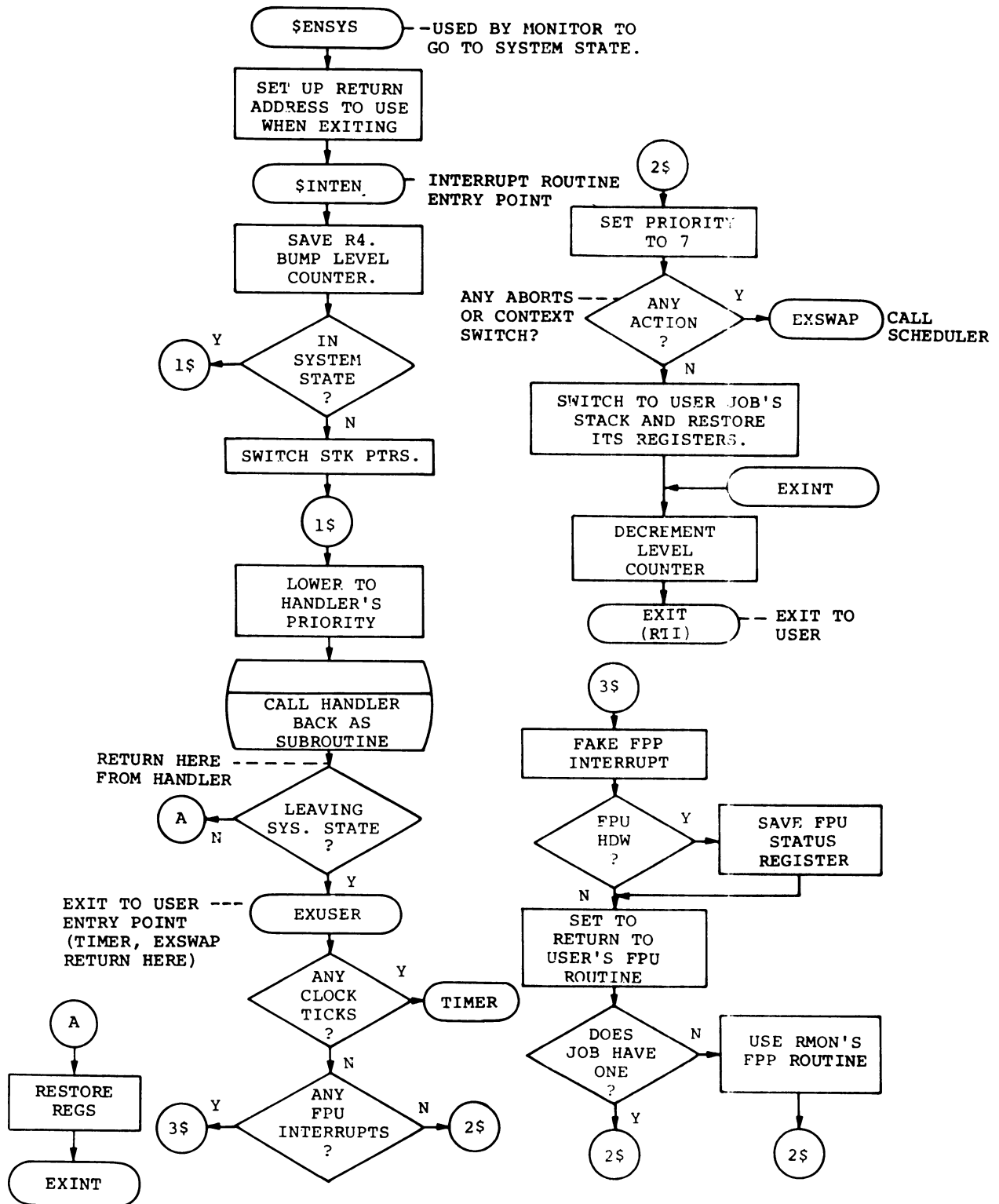


SWAP IN USR, LOCK/UNLOCK USR

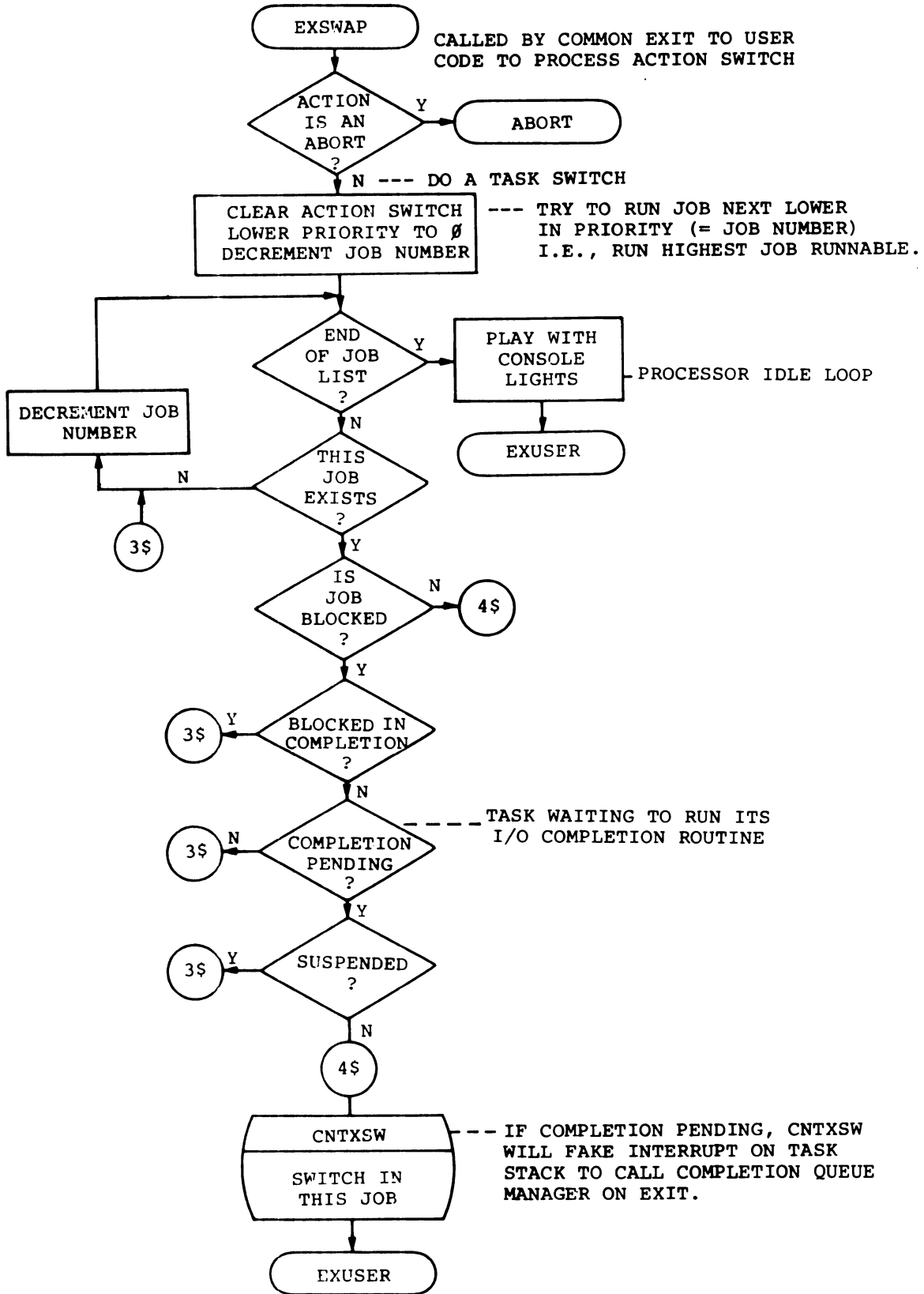


E.5.2 Job Arbitration, Error Processing

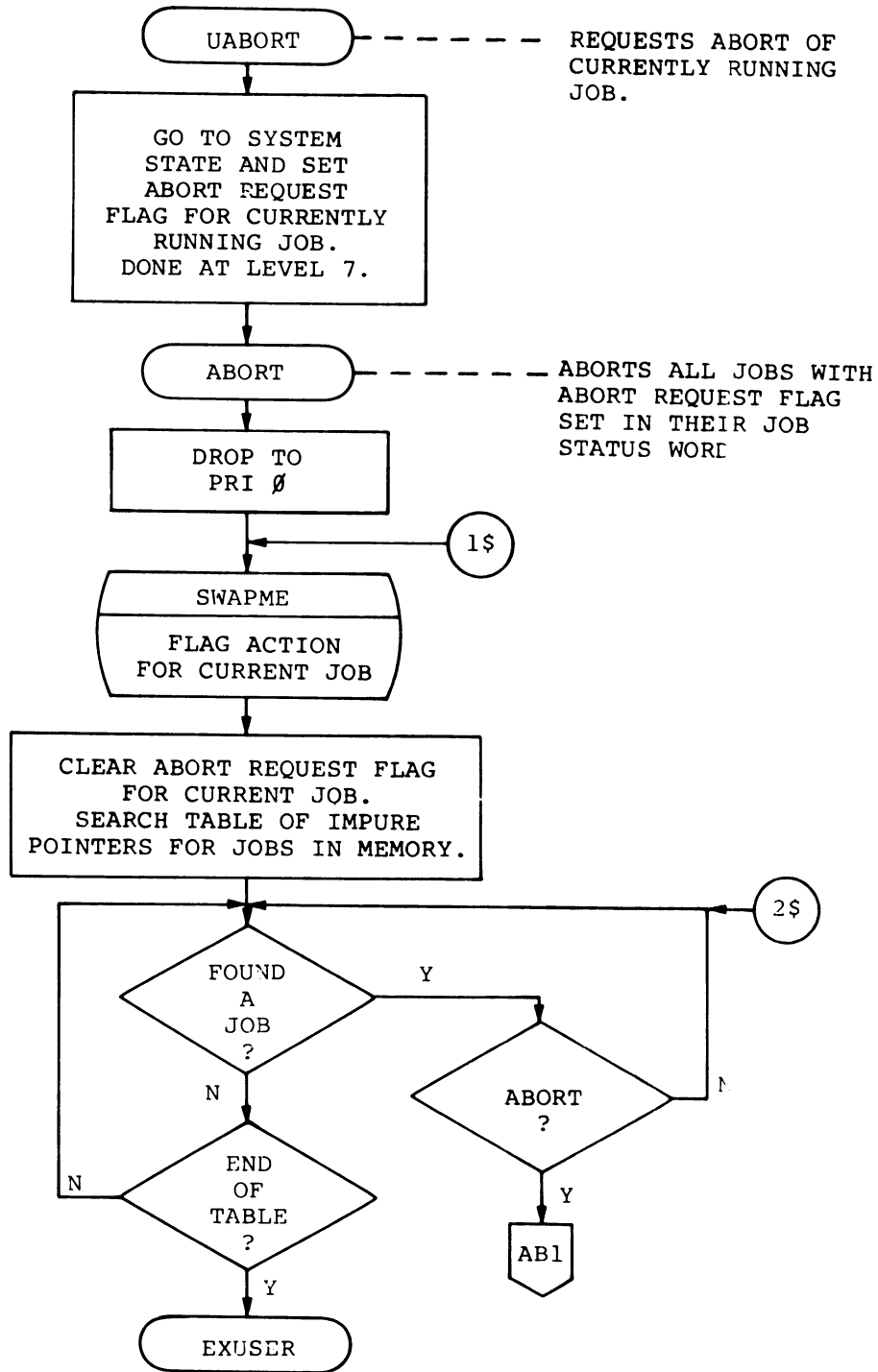
COMMON INTERRUPT
ENTRY AND EXIT



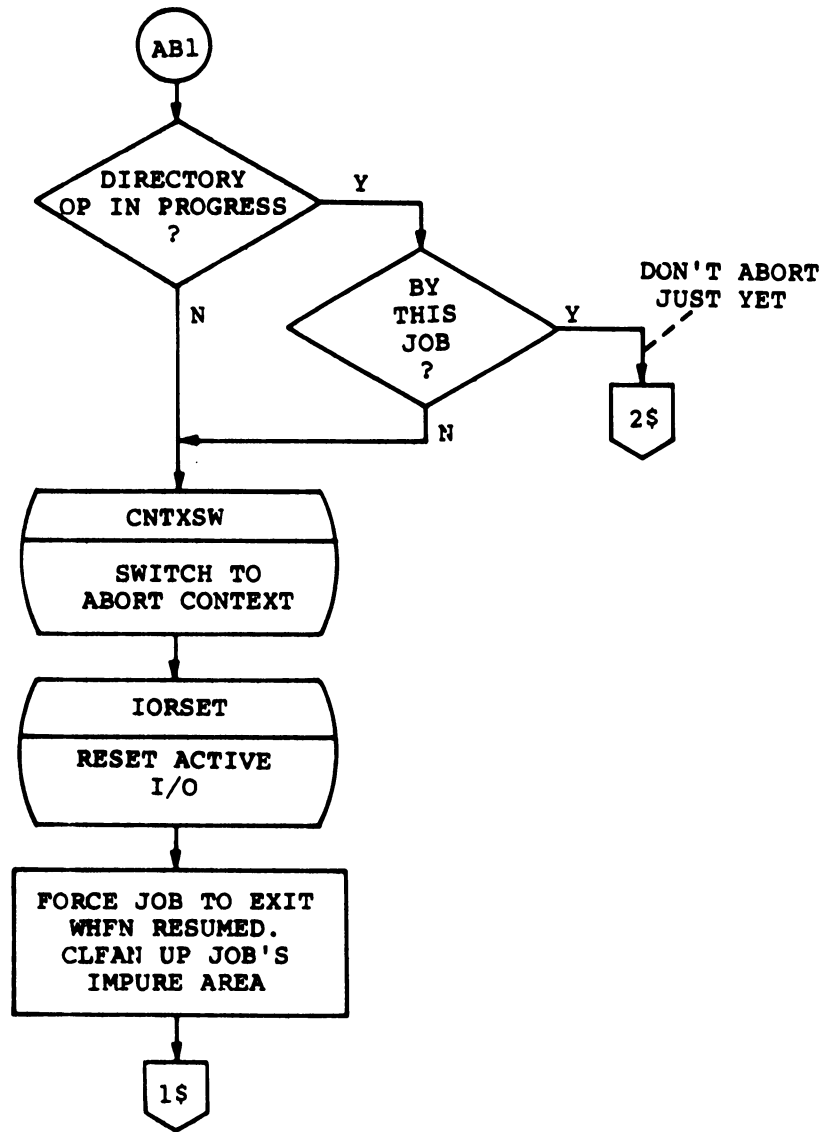
SCHEDULER



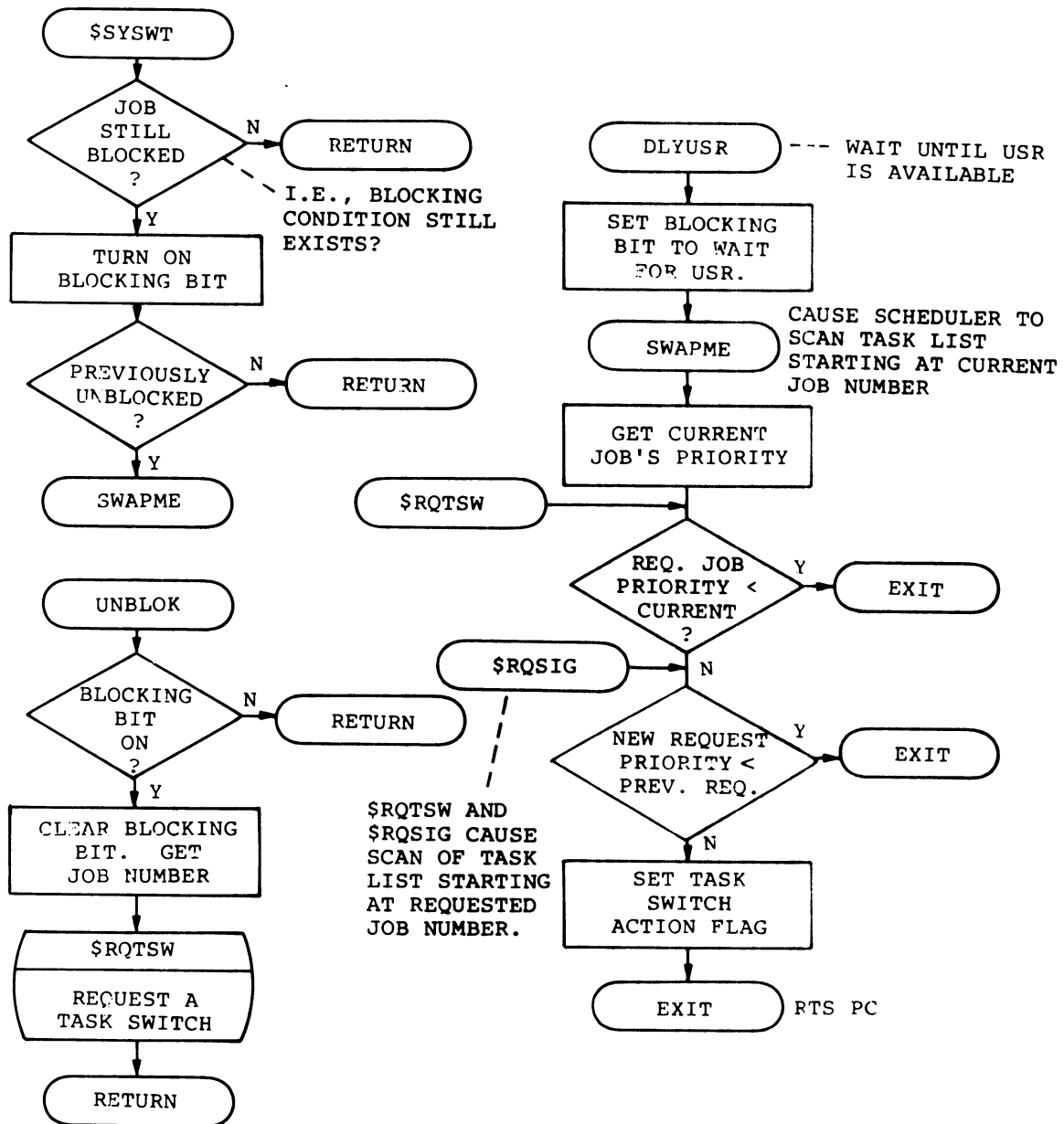
JOB ABORT



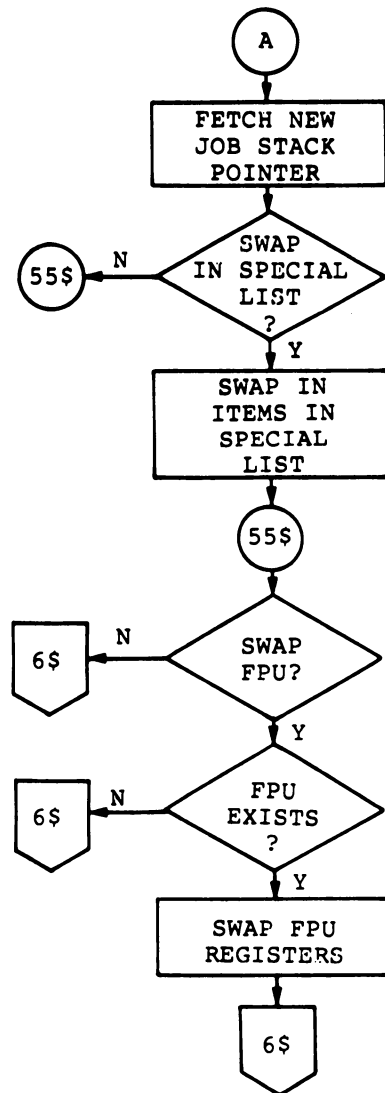
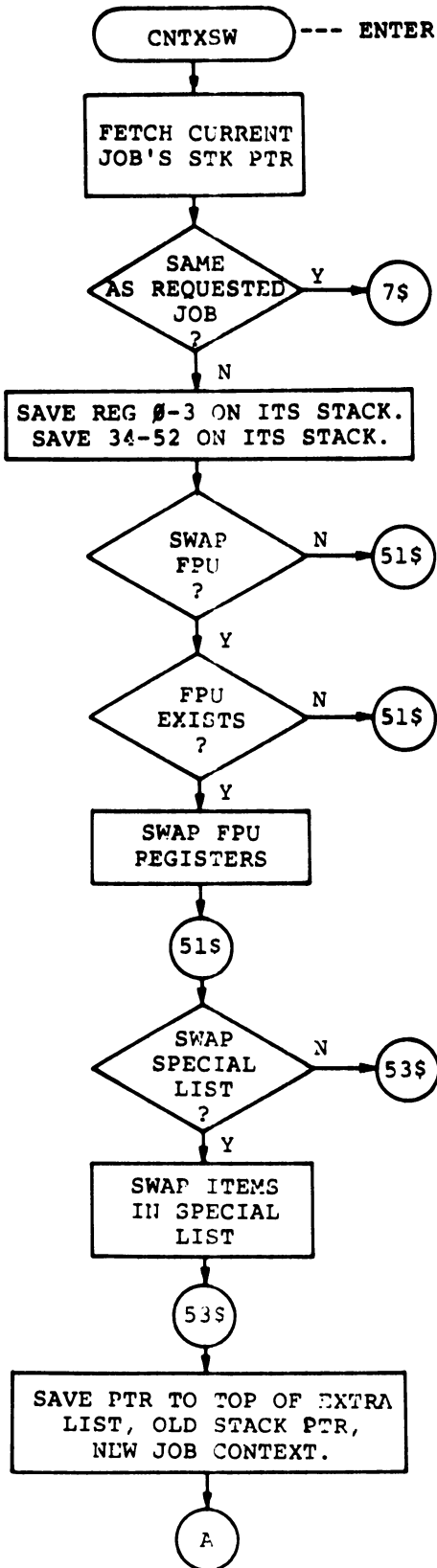
JOB ABORT (CONT.)



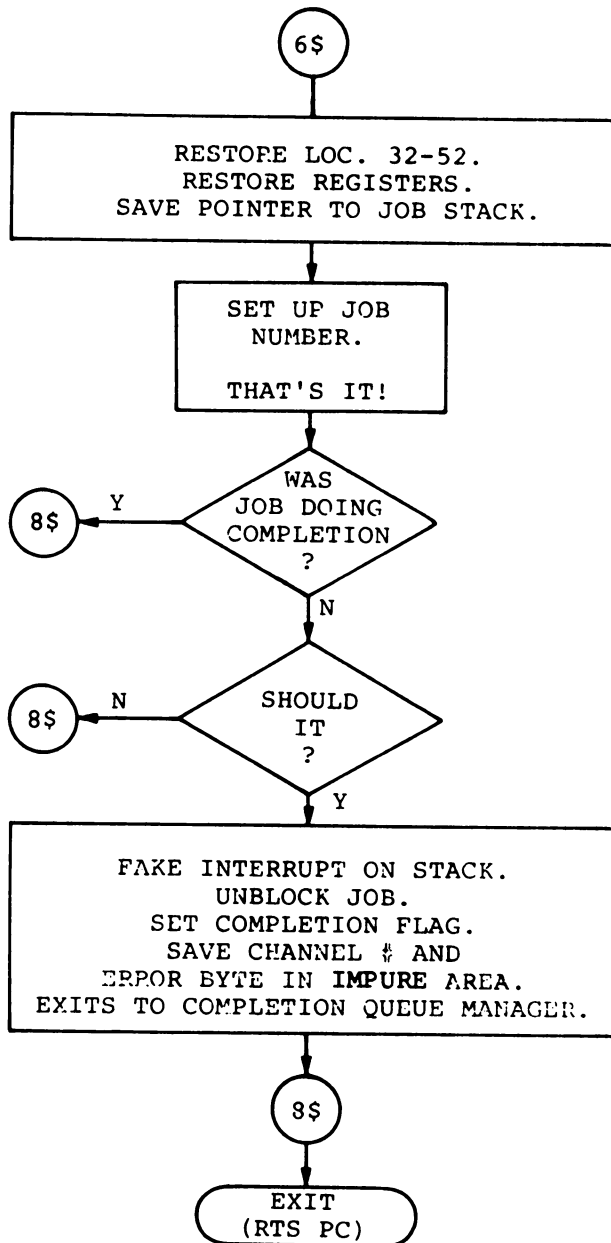
BLOCK A TASK/UNBLOCK A TASK
REQUEST TASK SWITCH



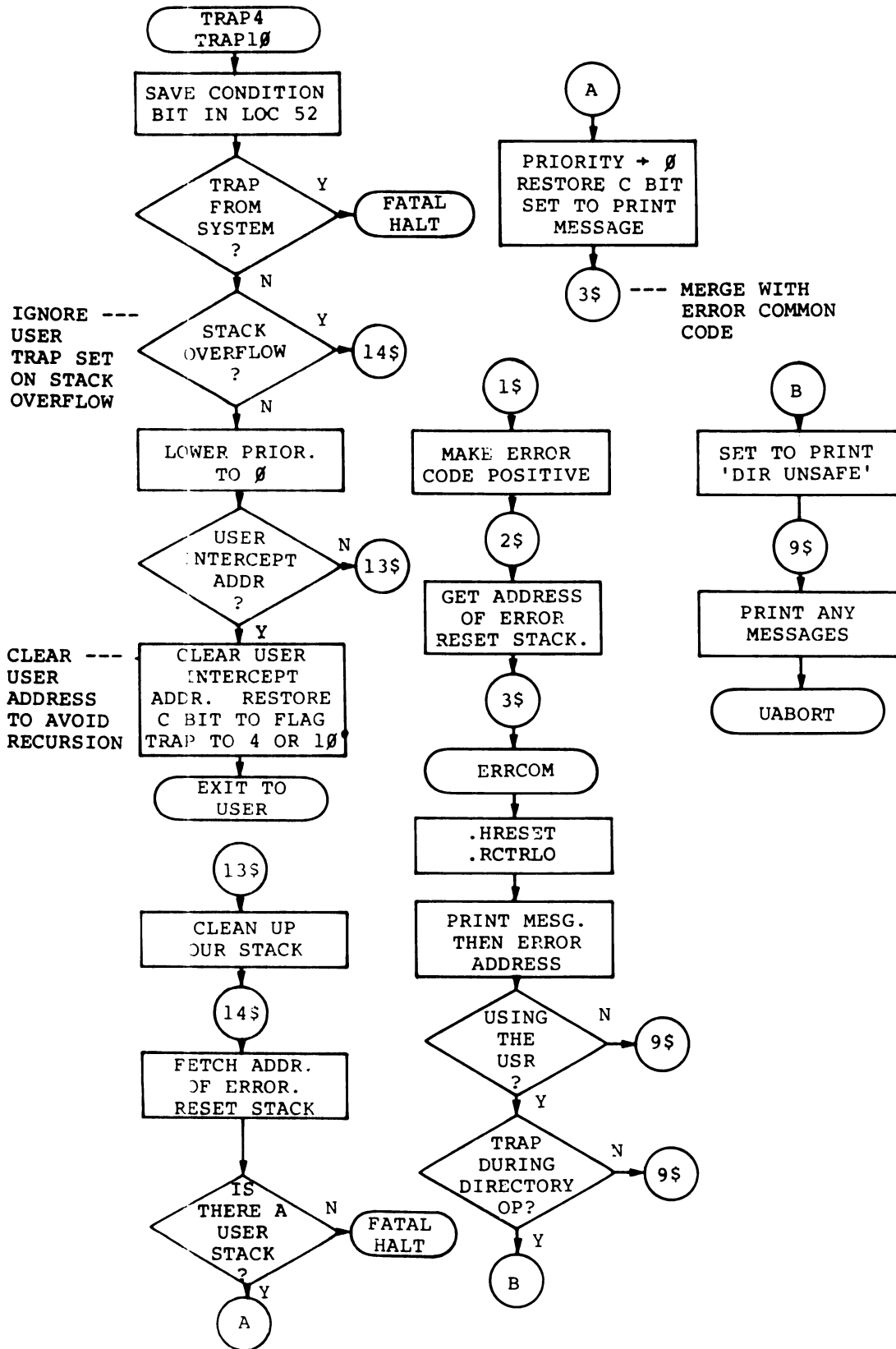
CHANGE CURRENT CONTEXT



CHANGE CURRENT CONTEXT (CONT.)

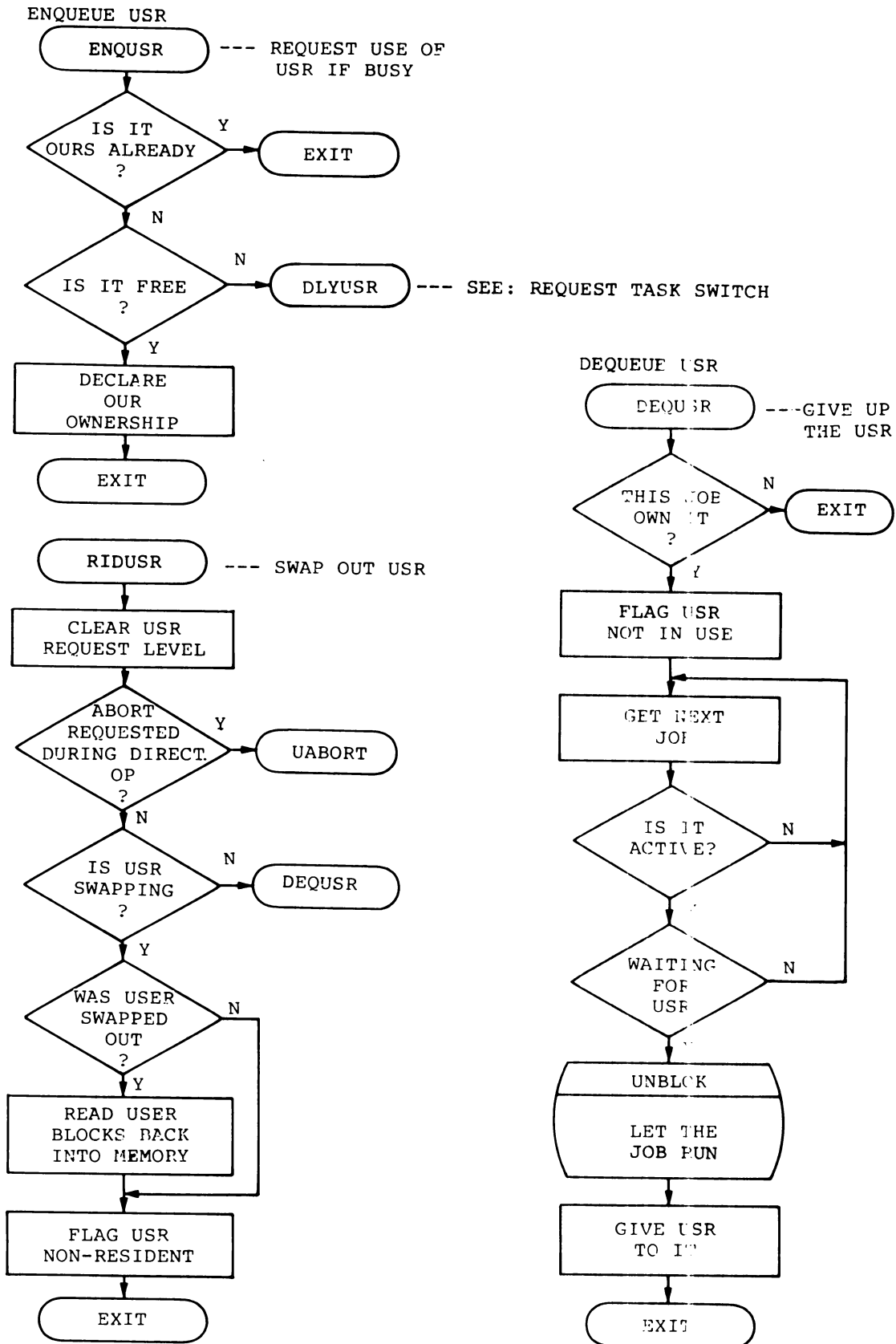


ERROR PROCESSING

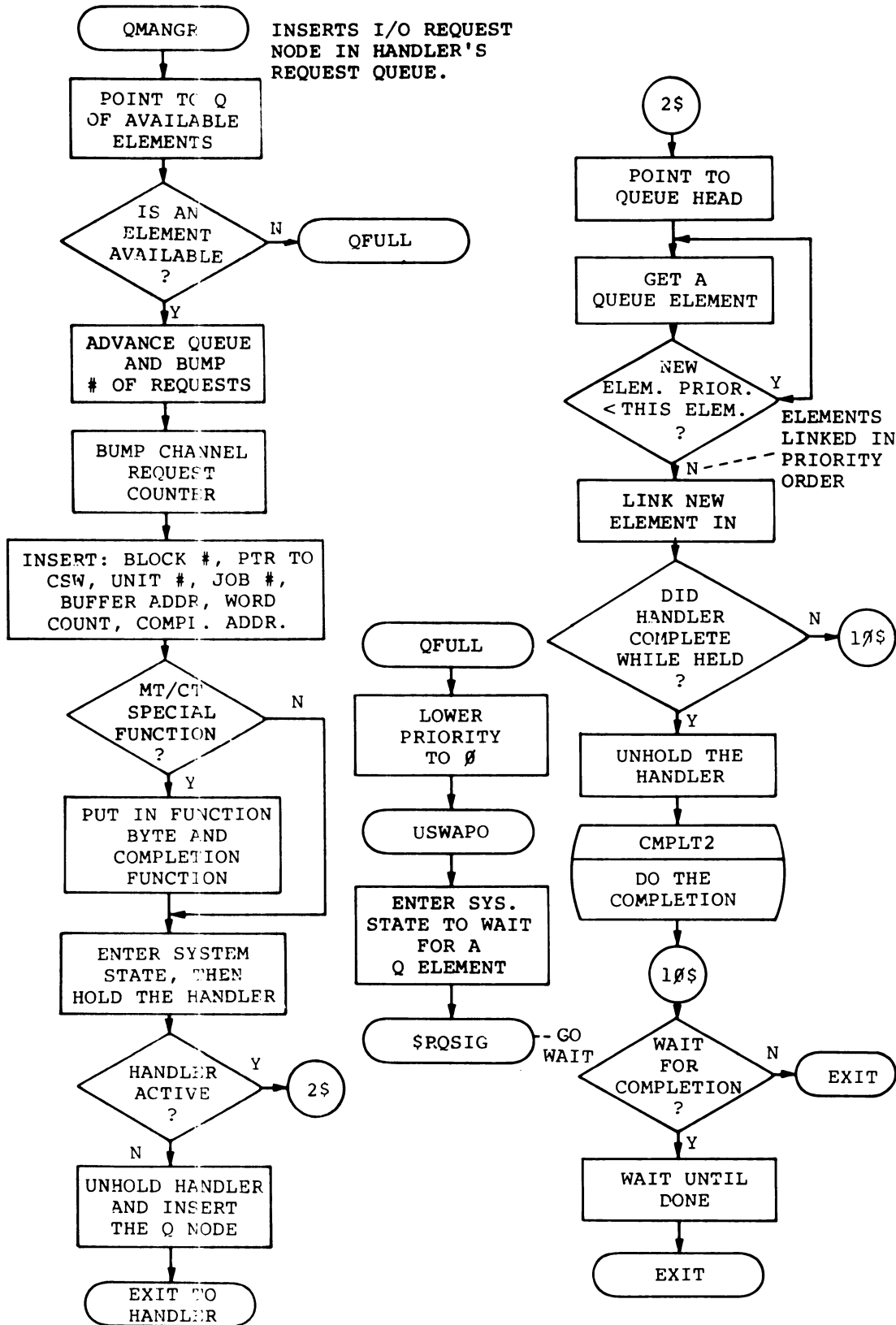


E.5.3 Queue Managers (I/O, USR, Completion)

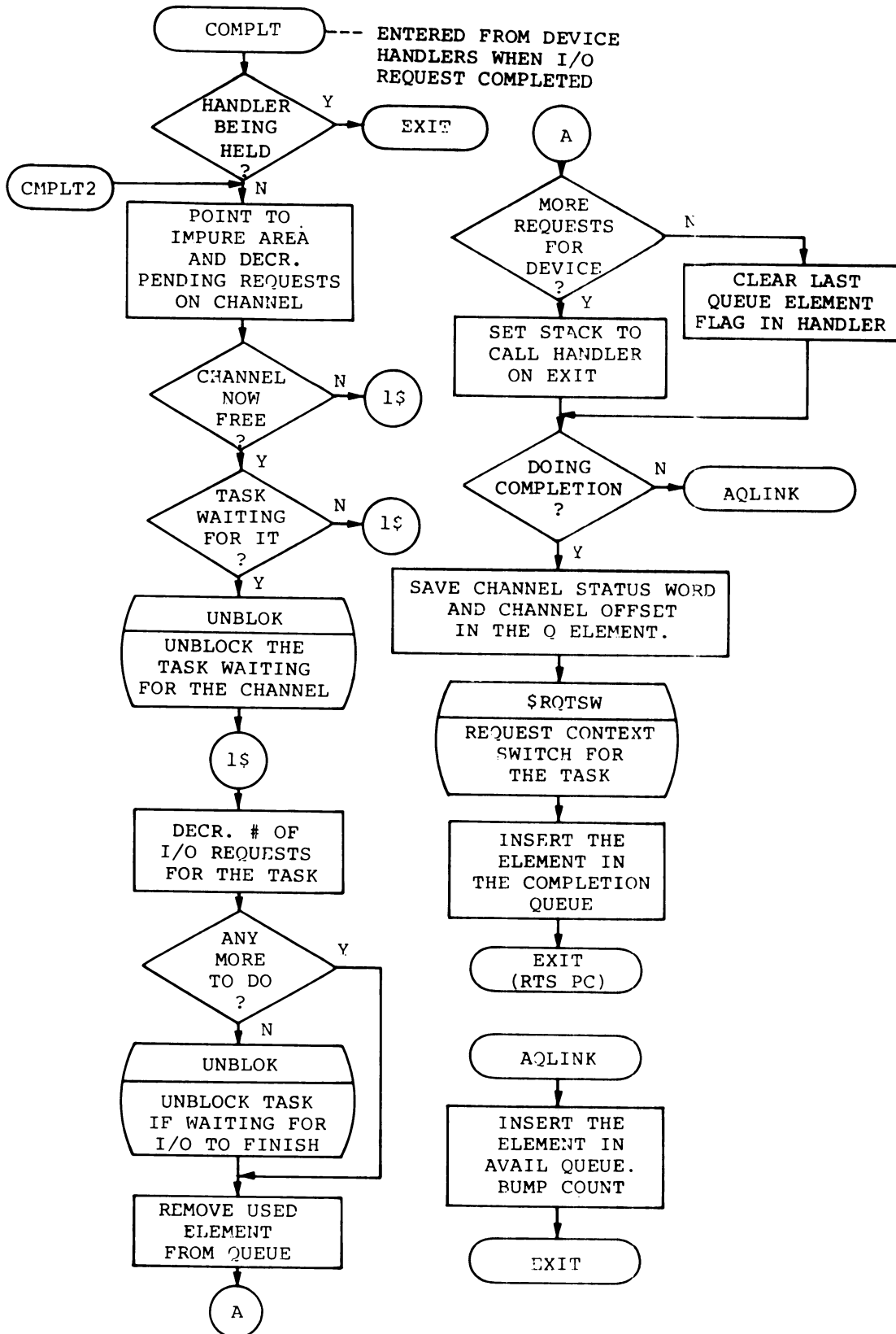
ENQUEUE/DEQUEUE USR



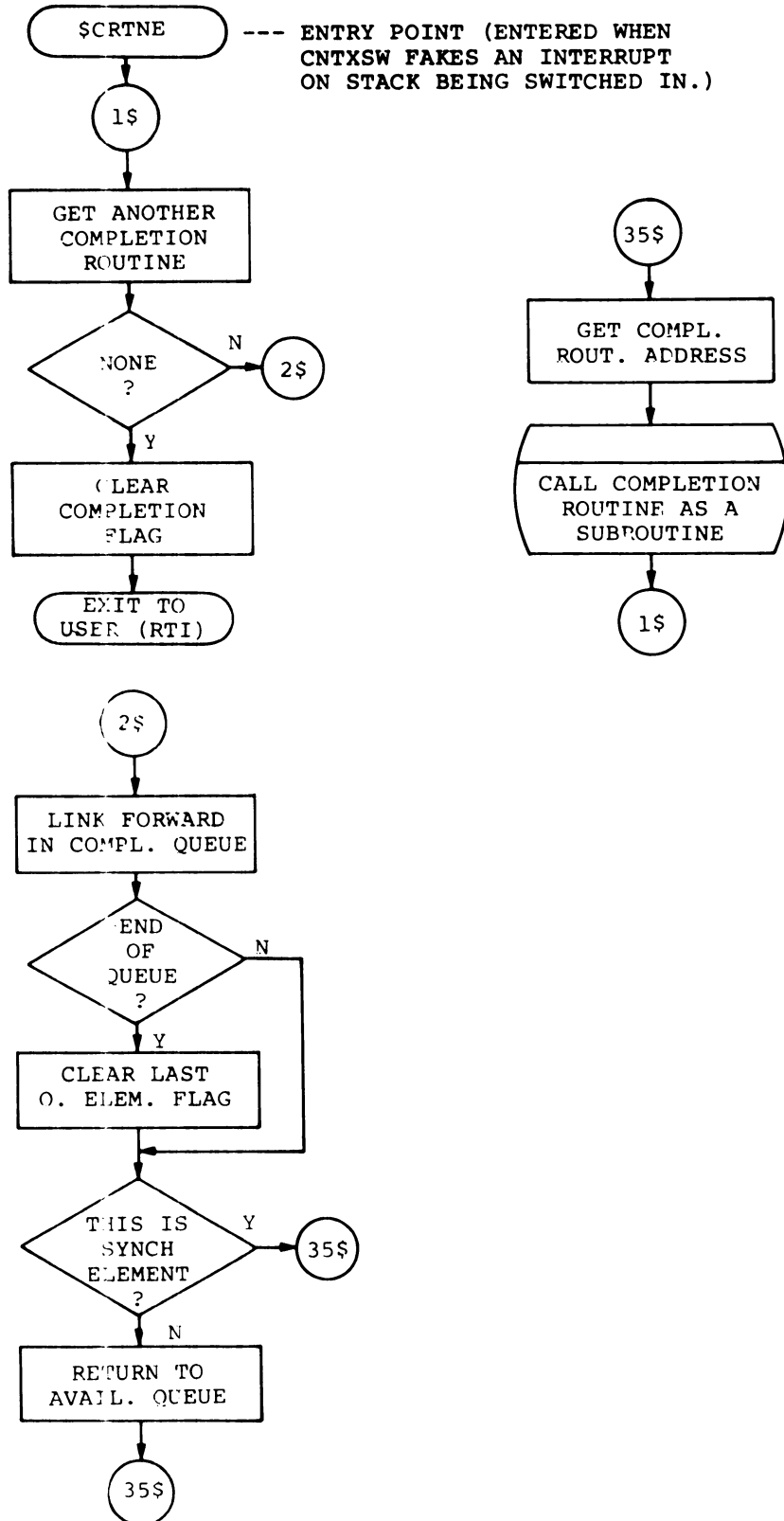
I/O QUEUE MANAGER



QUEUE COMPLETION

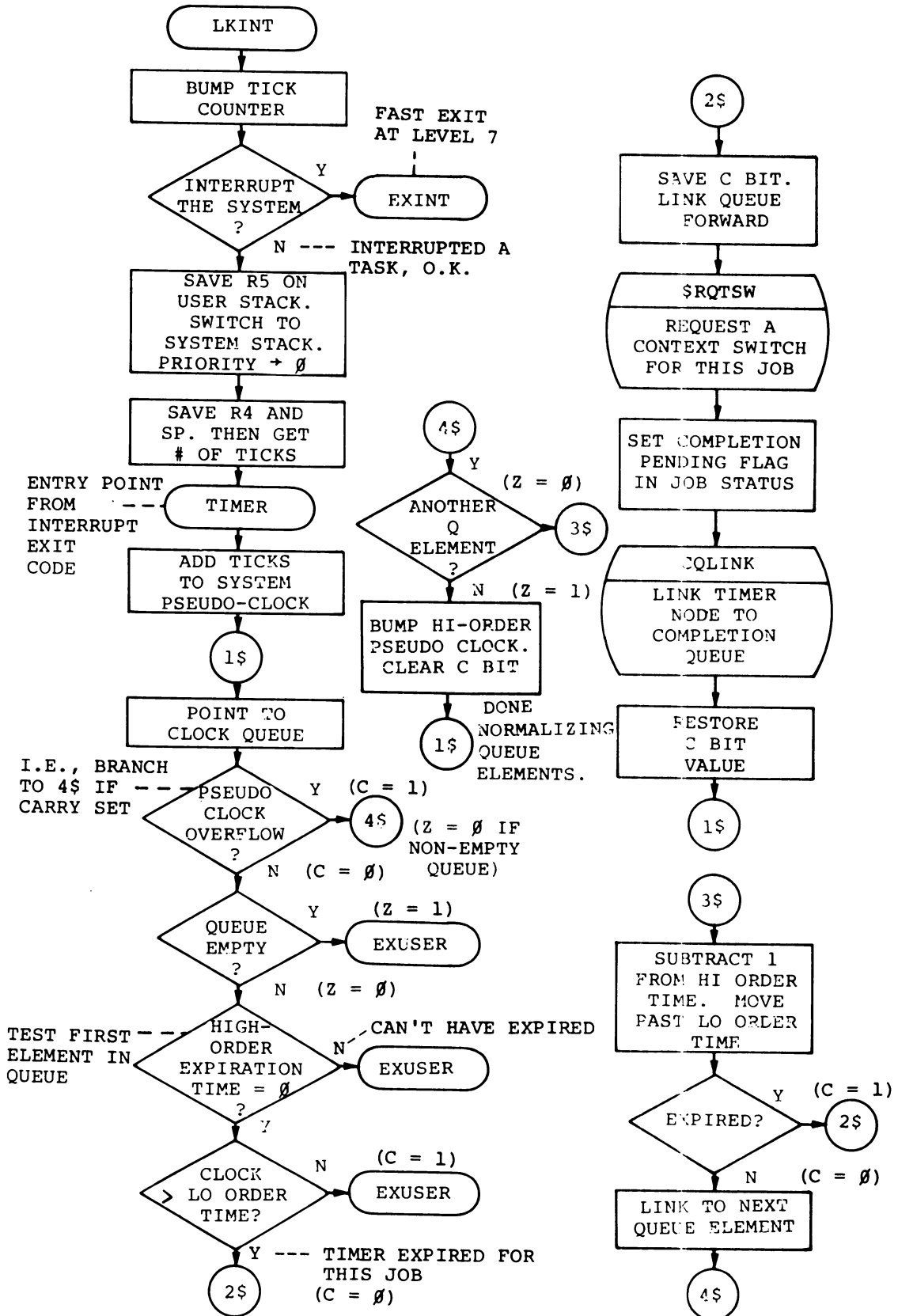


COMPLETION QUEUE MANAGER



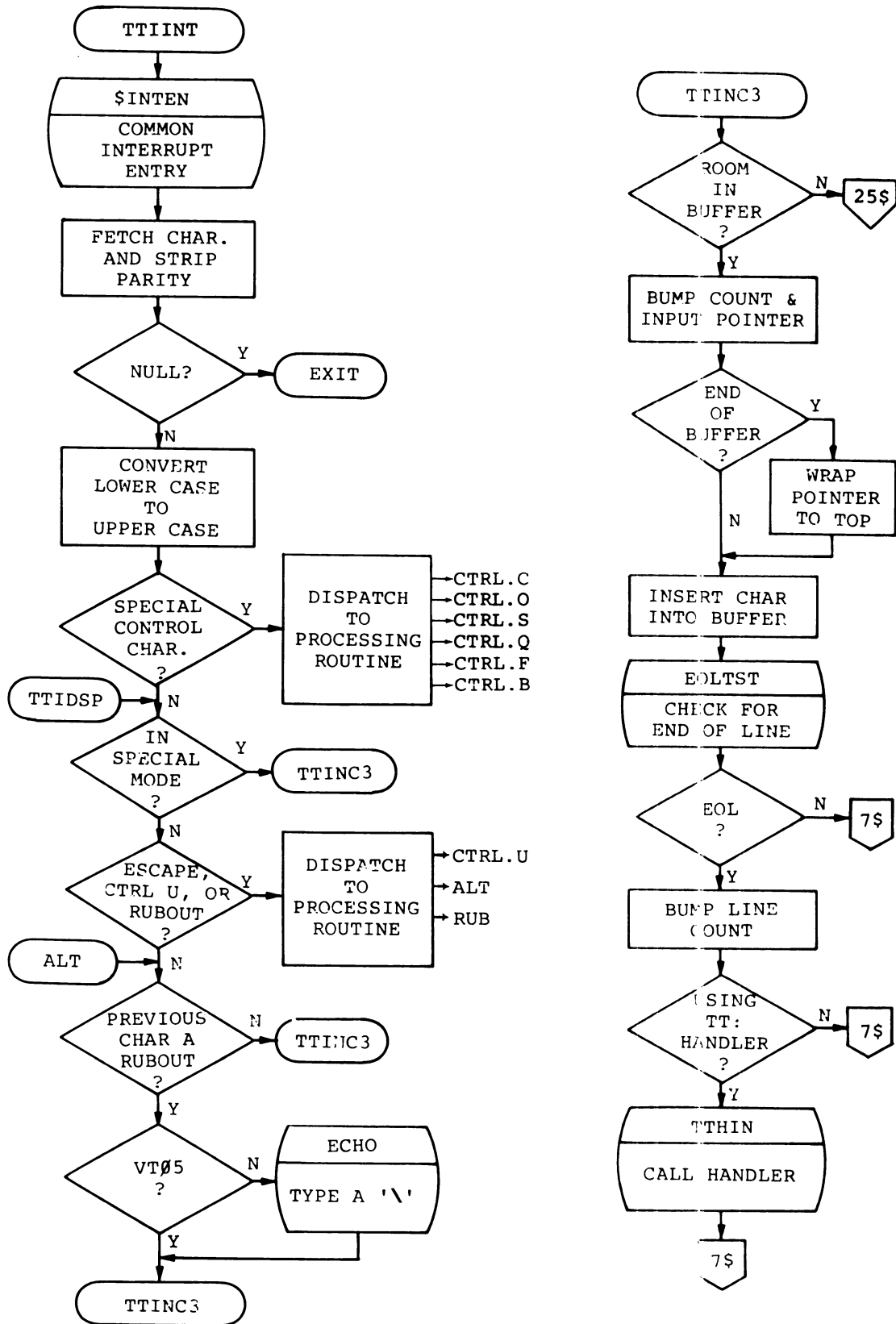
E.5.4 Clock Interrupt Service

CLOCK INTERRUPT HANDLER

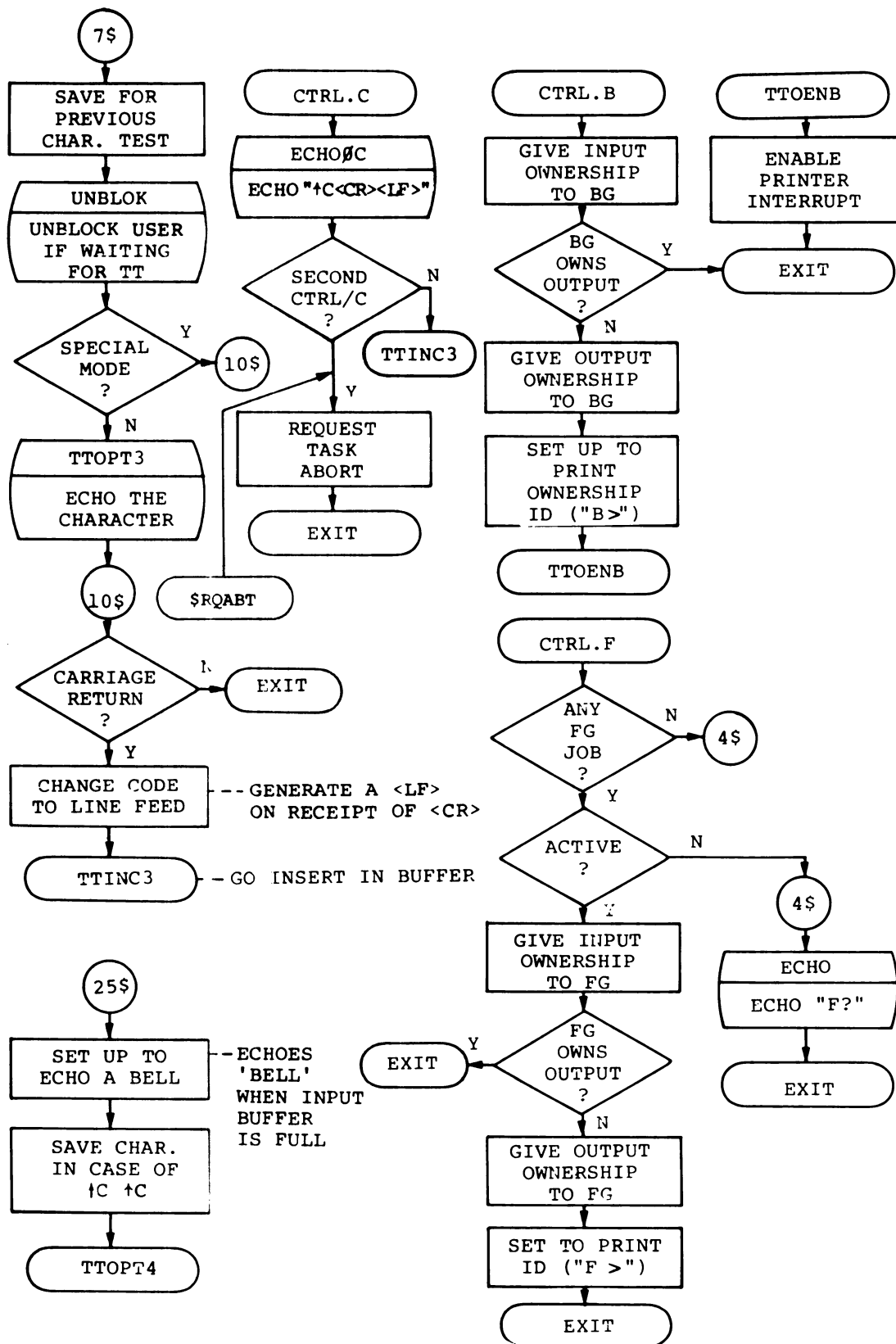


E.5.5 Console Terminal Interrupt Service

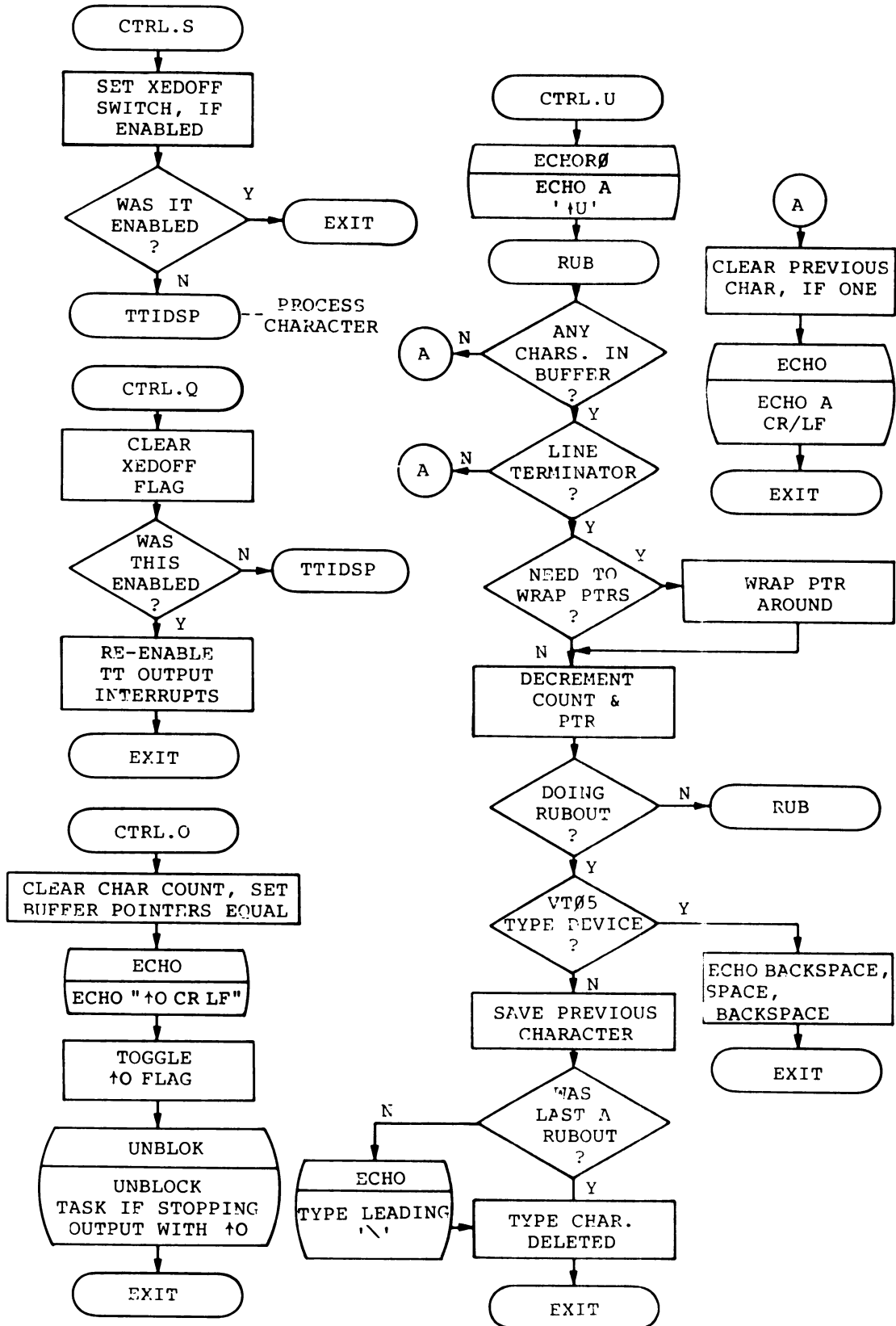
TT INPUT INTERRUPT ROUTINE

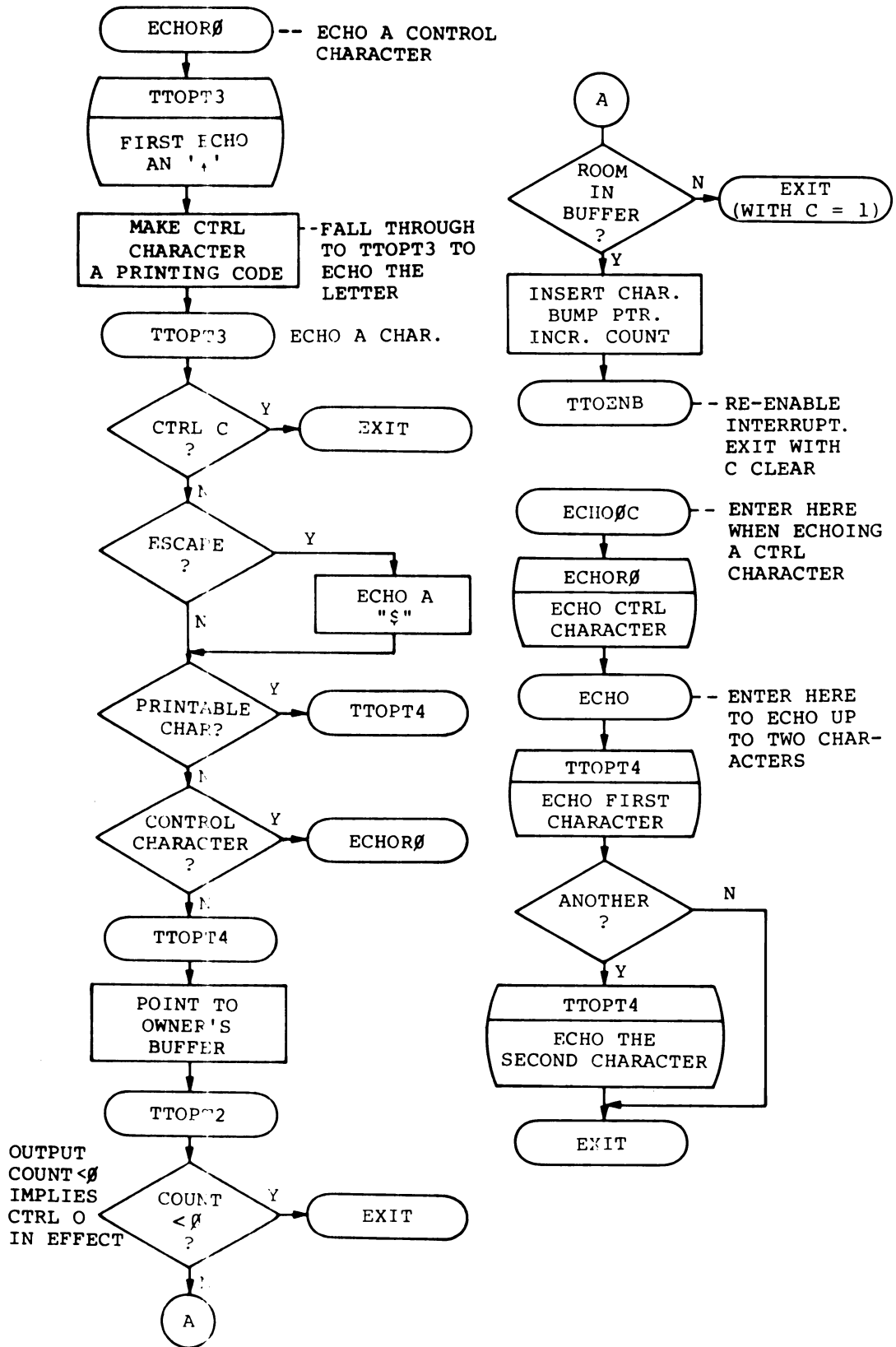


TT INPUT INTERRUPT ROUTINE (CONT.)

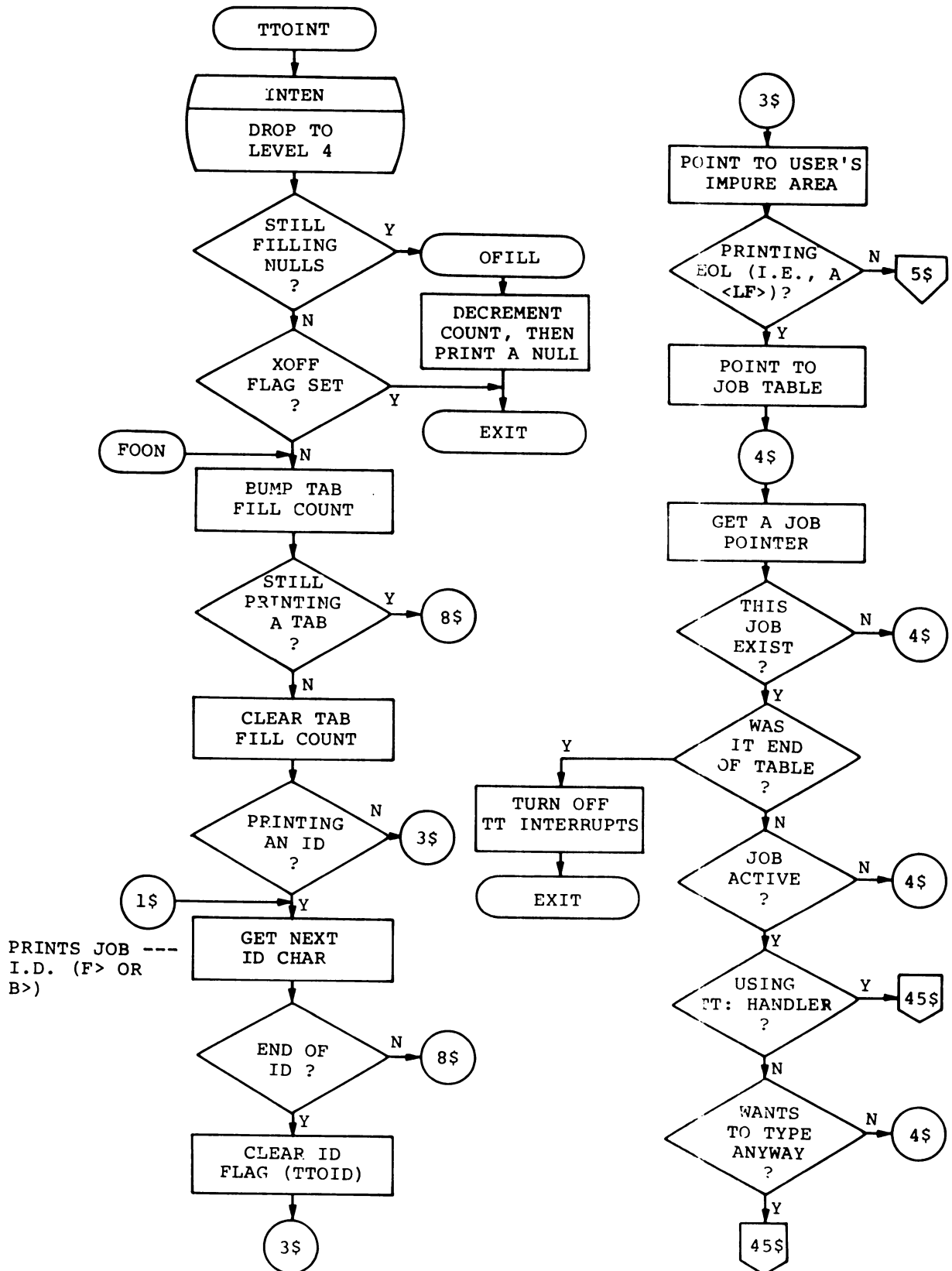


TT INPUT INTERRUPT ROUTINE (CONT.)

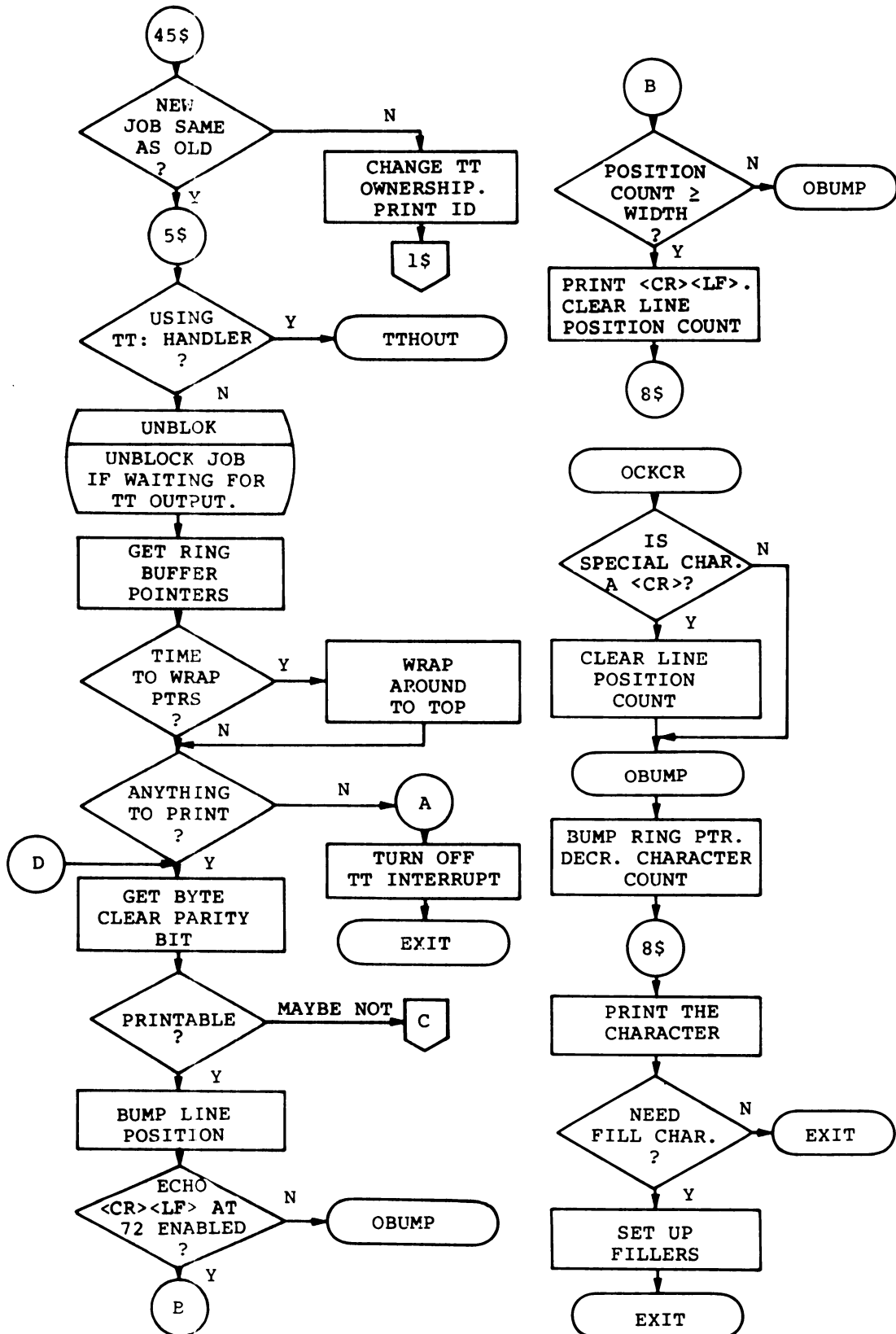




TT OUTPUT INTERRUPT ROUTINE

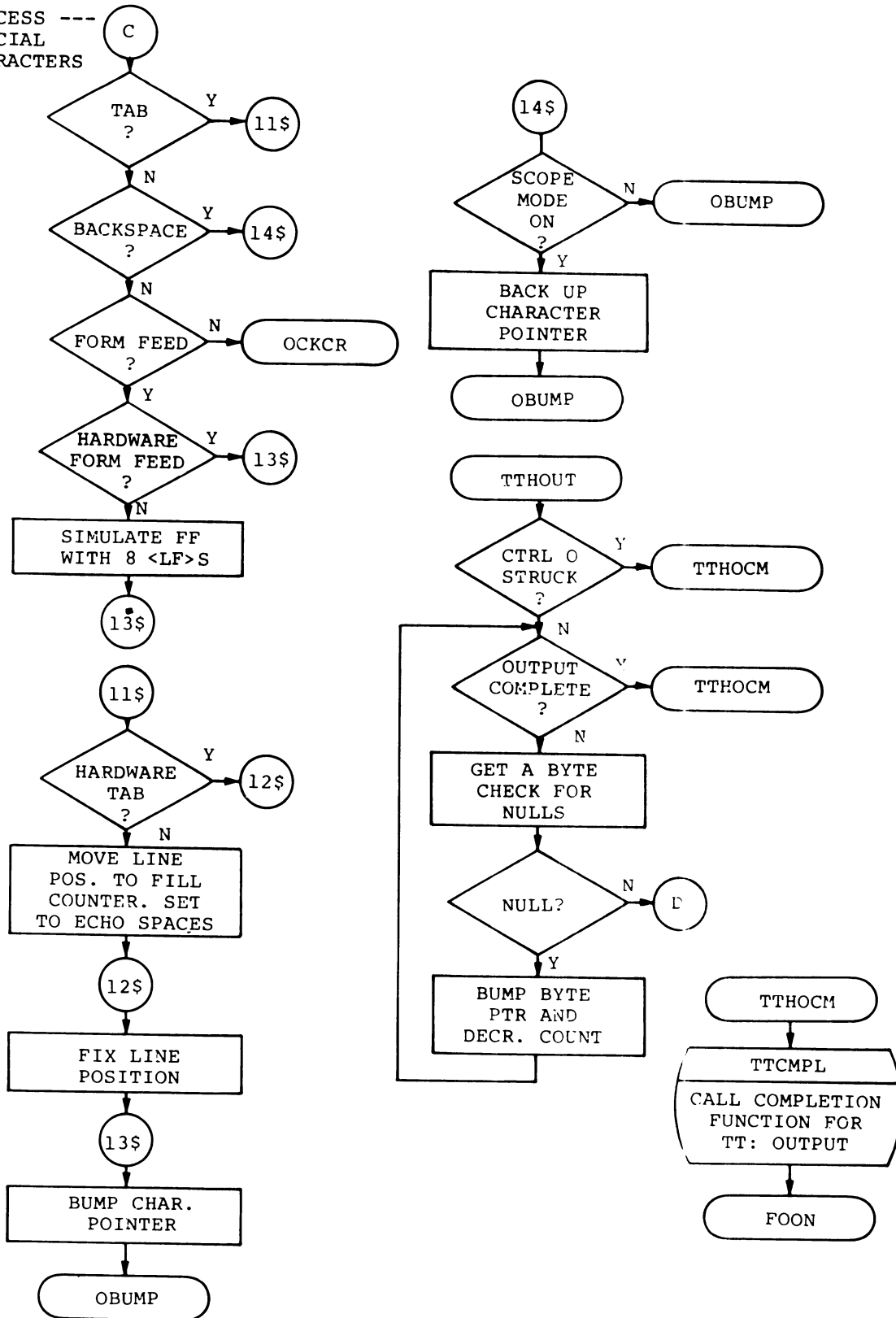


TT OUTPUT INTERRUPT ROUTINE (CONT.)



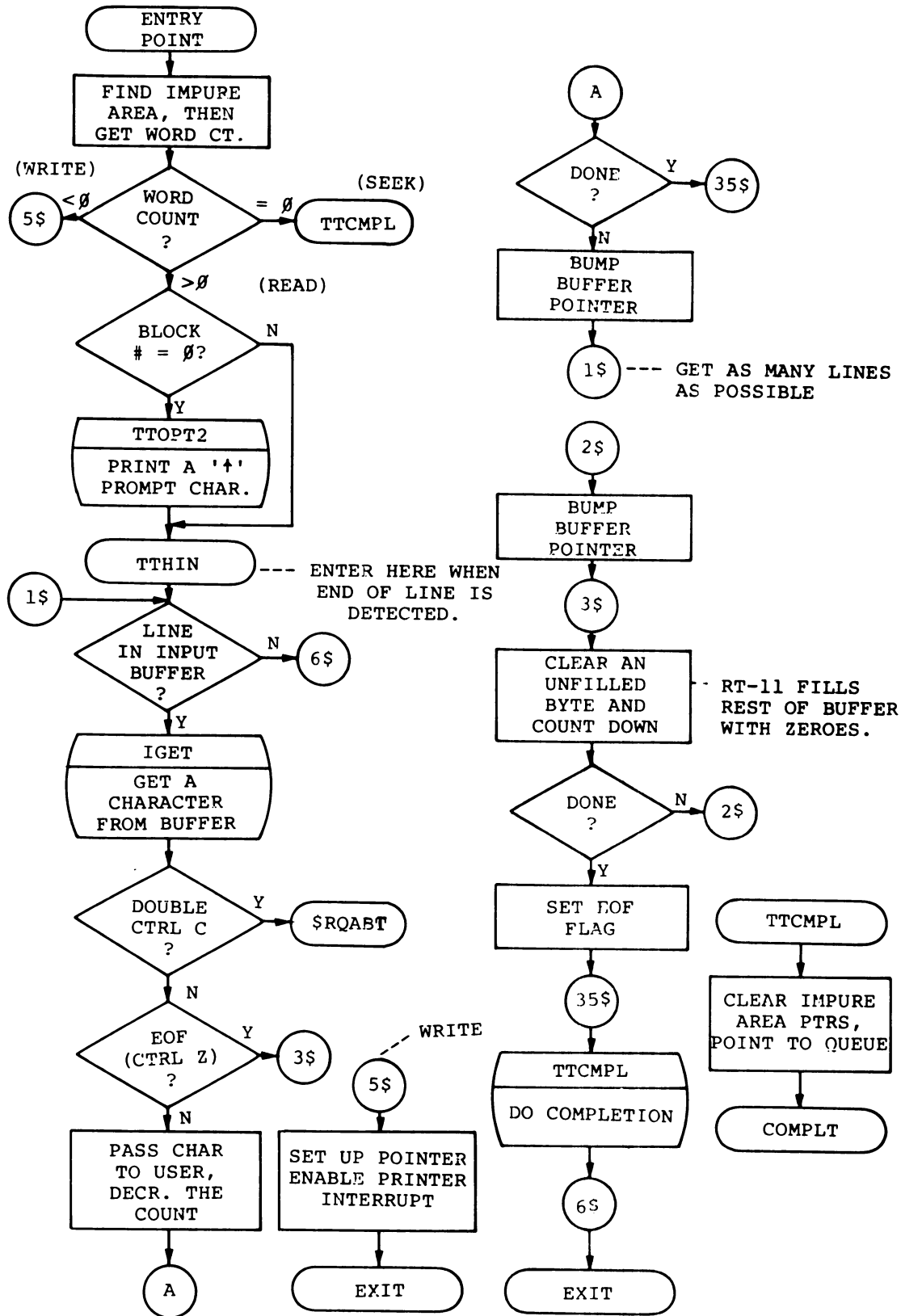
TT OUTPUT INTERRUPT ROUTINE (CONT.)

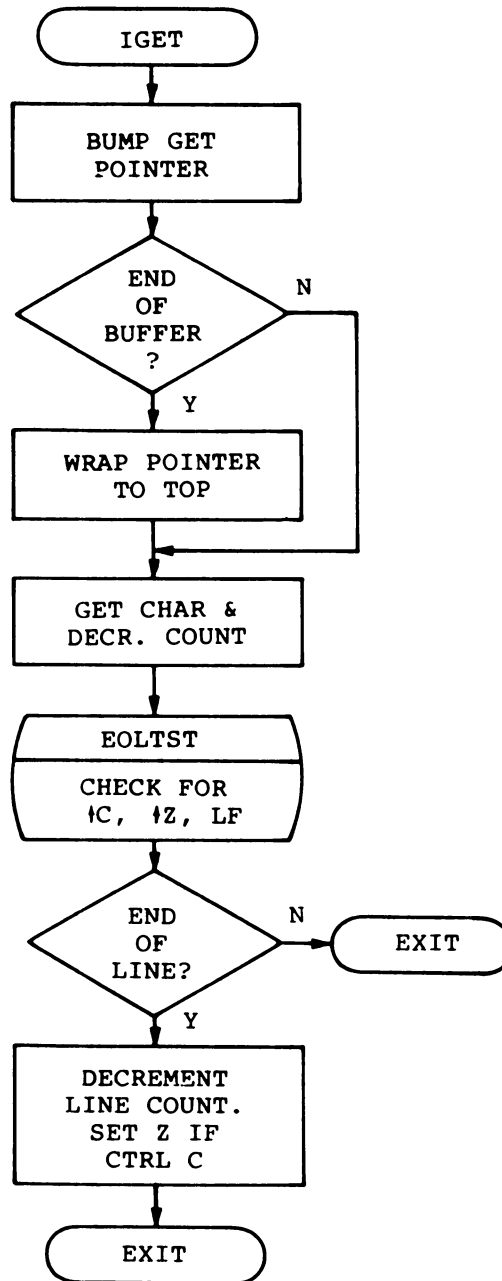
PROCESS ---
SPECIAL
CHARACTERS



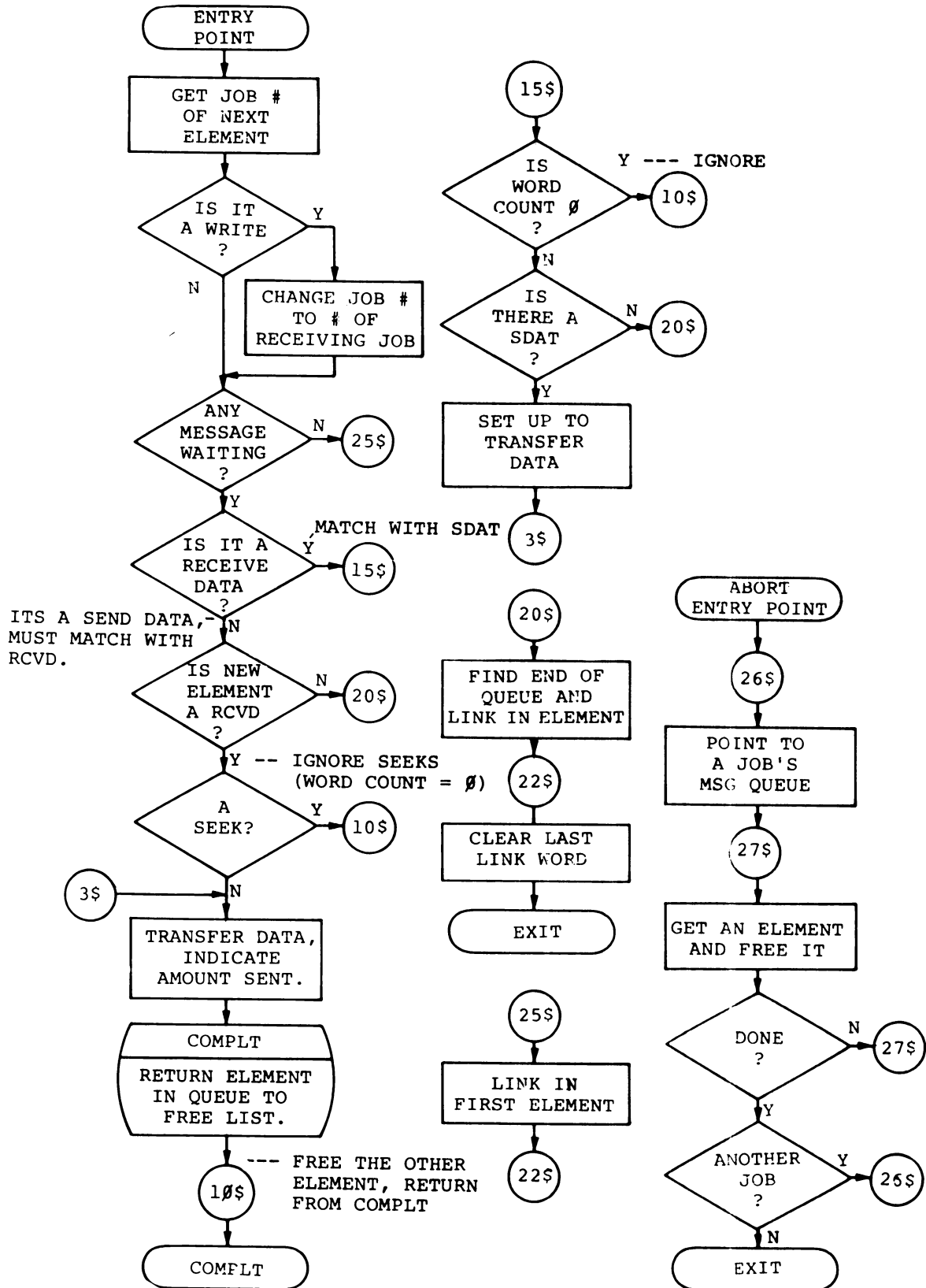
E.5.6 Resident Device Handlers (TT, Message)

TT: RESIDENT HANDLER





MESSAGE HANDLER



ENTRY POINT INDEX

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C\$PYCH, E-110*	EMTCOM, E-104*
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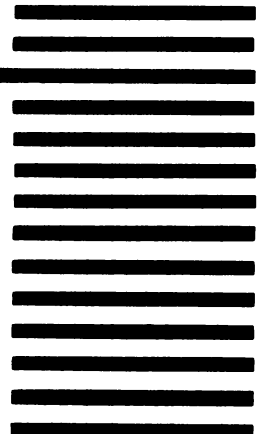
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