

```

%
Page Numbers: Yes First Page: 1
Heading:
MODEL 1: preamble.mc September 19, 1979 10:42 AM %
%
September 19, 1979 10:42 AM
    Try undoing zeroHold fix -- found a different bug that probably accounts for the behavior.
September 18, 1979 5:55 PM
    Try to make zeroHold more reliable: apparently 3 in a row is not enough.
June 17, 1979 5:11 PM
    Change holdValueLoc to accommodate ifu entry points
April 17, 1979 10:45 PM
    Add sim.holdMask, sim.taskMask, sim.holdShift, sim.taskShift.
January 25, 1979 11:05 AM
    Change simTaskLevel to 12B.
January 18, 1979 5:10 PM
    Add currentTaskNum.
%

```

```

TITLE[PREAMBLE];
%

```

NAMING CONVENTIONS:

LABELS BEGIN W/ THE OPERATION BEING TESTED:

```

    Aplus1
    cntFcn
    aluLTO
lsh => Left shift
rsh => Right shift
lcy => Left Cycle
rcy => Right Cycle

```

Register Read/write tests are suffixed with RW:

```

    cntRW
    shcRW

```

LOOP LABELS ARE SUFFIXED AS INNER (IL) AND OUTER(OL) LOOPS (L).

```

    cntFcnIL          * INNER LOOP
    cntFcnOL          * OUTER LOOP
    cntFcnXITIL       * LABEL FOR EXITING INNER LOOP
    cntFcnXITOL       * LABEL FOR EXITING OUTER LOOP
Aplus1L              * ONLY LOOP

```

```

%

```

RMREGION[DEFAULTREGION]:

```

rv[r0,0];          rv[r1,1];          rv[rm1,177777];
rv[r01,52525];    rv[r10,125252]; rv[rhigh1,100000];
rv[rscr,0];       rv[rscr2,0];       rv[rscr3,0];
rv[rscr4,0];     rv[stackPAddr, 0];   rv[stackPTopBits, 0];
rv[klink, 0];    rv[hack0,0];       rv[hack1, 0];
rv[hack2,0];

```

```

rvrel[rmx0, 0]; rvrel[rmx1, 1]; rvrel[rmx2, 2];
rvrel[rmx3, 3]; rvrel[rmx4, 4]; rvrel[rmx5, 5];
rvrel[rmx6, 6]; rvrel[rmx7, 7]; rvrel[rmx10, 10];

```

* Constants from FF

```

nsp[PNB0,100000];
nsp[PNB1,40000];          nsp[PNB2,20000];          nsp[PNB3,10000];
nsp[PNB4,4000]; nsp[PNB5,2000]; nsp[PNB6,1000];
nsp[PNB7,400]; nsp[PNB8,200]; nsp[PNB9,100];
nsp[PNB10,40]; nsp[PNB11,20]; nsp[PNB12,10];
nsp[PNB13,4]; nsp[PNB14,2]; nsp[PNB15,1];

```

```

mc[B0,100000];
mc[B1,40000]; mc[B2,20000]; mc[B3,10000];
mc[B4,4000]; mc[B5,2000]; mc[B6,1000];
mc[B7,400]; mc[B8,200]; mc[B9,100];
mc[B10,40]; mc[B11,20]; mc[B12,10];
mc[B13,4]; mc[B14,2]; mc[B15,1];

```

```

mc[NB0,PNB0];
mc[NB1,PNB1]; mc[NB2,PNB2]; mc[NB3,PNB3];
mc[NB4,PNB4]; mc[NB5,PNB5]; mc[NB6,PNB6];
mc[NB7,PNB7]; mc[NB8,PNB8]; mc[NB9,PNB9];

```

```
mc[NB10,PNB10]; mc[NB11,PNB11]; mc[NB12,PNB12];  
mc[NB13,PNB13]; mc[NB14,PNB14]; mc[NB15,PNB15];
```

```
mc[CM1,177777];mc[C77400,77400]; mc[C377,377];mc[CM2,-2];  
mc[getIMmask, 377]; * isolate IM data after getIm[!]
```

```
m[noop, BRANCH[.+1]];  
m[skip, BRANCH[.+2]];  
m[error, ILC[(BRANCH[ERR])]];  
m[skiperr, ILC[(BRANCH[.+2])] ILC[(BRANCH[ERR])]];  
m[skpif, BRGO[tsd] BAT[.+2,#1,#2]];  
m[skpUnless, BRGO[tsd] DBAT[.+1,+.2,#1,#2]];  
m[loopChk, BRGO[tsd] DBAT[#1,+.1,#2,#3]];  
m[loopUntil, BRGO[tsd] DBAT[.+1,#2,#1,#3]];* if #1 then goto .+1 else goto #2  
m[loopWhile, BRGO[tsd] DBAT[#2,+.1,#1,#3]];* if #1 then goto #2 else goto .+1
```

* April 17, 1979 10:47 PM

```

rmRegion[rm2ForKernelRtn];
knowRbase[rm2ForKernelRtn];

rv[randV,0]; * current value from random number generator
rv[randX,0]; * current index into random number jump table
rv[oldRandV,0]; * saved value
rv[oldRandX,0]; * saved value
rv[chkSimulatingRtn, 0];
rv[fixSimRtn, 0];
rv[chkRunSimRtn, 0];
rv[currentTaskNum, 0];
rmRegion[randomRM];
knowRbase[randomRM];

rv[rndm0, 134134]; rv[rndm1, 054206];
rv[rndm2, 036711]; rv[rndm3, 103625];
rv[rndm4, 117253]; rv[rndm5,154737];
rv[rndm6, 041344]; rv[rndm7, 006712];

knowRbase[defaultRegion];

mp[flags.conditiona1P, 200]; * bit that indicates conditional simulating
mp[flags.conditionOKp, 100]; * bit that indicates conditional simulating is ok

set[holdValueLoc, 3400]; mc[holdValueLocC, holdValueLoc];
mc[flags.taskSim, b15]; * NOTE: The "flags" manipulation code
mc[flags.holdSim, b14]; * works only so long as there are no more
mc[flags.simulating, flags.taskSim, flags.holdSim];
set[simTaskLevel, 12]; mc[simTaskLevelC, simTaskLevel];

mc[sim.holdMask, 377]; set[sim.holdShift, 0];
mc[sim.taskMask, 177400]; set[sim.taskShift, 10];

m[lh, byt0[ and[rshift[#1,10], 377] ] byt1[ and[#1, 377]]
]; * assemble data for left half of IM
m[rh, byt2[ and[rshift[#1,10], 377] ] byt3[ and[#1, 377]]
]; * assemble data for right half of IM

m[zeroHold, ilc[#1 ← A0]]
ilc[(hold&tasksim ← #1)]
ilc[(hold&tasksim ← #1)]
ilc[(hold&tasksim ← #1)]
];

```

```

* December 11, 1978 3:20 PM
%
    subroutine entry/exit macros
%
m[saveReturn, ilc[(t ← link)]
    top level[]
    ilc[(#1 ← t)]
    ];

m[saveReturnAndT, ilc[(#2 ← t)]
    ilc[(t ← link)]
    top level[]
    ilc[(#1 ← t)]
    ];

m[returnUsing, subroutine[]
    ilc[(RBASE ← rbase[#1])]
    ilc[(link ← #1)]
    ilc[(return, RBASE ← rbase[defaultRegion])]
    ];

m[returnAndBranch, subroutine[]
    ilc[(RBASE ← rbase[#1])]
    ilc[(link ← #1)]
    ilc[(RBASE ← rbase[defaultRegion])]
    ilc[(return, PD ← #2)]
    ];

m[pushReturn, subroutine[]
    ilc[(stkp+1)]
    top level[]
    ilc[(stack ← link)]
    ];
    * notice that this macro doesn't clobber T !!!
m[pushReturnAndT, subroutine[]
    ilc[(stkp+1)]
    ilc[(stack&+1 ← link)]
    top level[]
    ilc[(stack ← t)]
    ];

m[returnP, subroutine[]
    ilc[(link ← (stack&-1))]
    ilc[(return)]
    ];

m[pReturnP,
    ilc[(stkp-1)]
    subroutine[]
    ilc[(link ← (stack&-1))]
    ilc[(return)]
    ];

m[returnPAndBranch, subroutine[]
    ilc[(link ← (stack&-1))]
    ilc[(return, PD ← #1)]
    ];

m[pReturnPAndBranch, subroutine[]
    ilc[(stkp-1)]
    ilc[(link ← (stack&-1))]
    ilc[(return, PD ← #1)]
    ];

m[getRandom, ilc[(RBASE ← rbase[randX])]
    ilc[(randX ← (randX)+1, Bdispatch ← randX)]
    ilc[(call[random], RBASE ← rbase[rndm0])]
    ilc[(RBASE ← rbase[defaultRegion])]
    ];
    * Returns random number in T, leaves RBASE= defaultRegion

knowRbase[defaultRegion];

```

%
Page Numbers: Yes First Page: 1
Heading:
kernelAlu.mc October 20, 1978 10:38 AM %
TITLE[KernelALU-model1.8/30/78];

%
October 19, 1978 1:13 PM
PD← for LU←

This file defines ALU operations for the program. Definitions for all possible ALU operations are given here with one line/definition. The lines for the 20 operations enabled should have the "n" replaced by consecutive values from 0 to 17. 16 should be "NOT A" for the shifter and 17 is normally used as a variable by BitBlt (i.e., not defined here), but can be used otherwise, if desired. Other lines should be commented out.

Note that the letter "E" added as an argument where noted below converts the operation to emulator-only. This is intended for locations used as variables by BitBlt, which must restore the smashed ALUFM instructions to the proper value before exit.

The "A" and "B" operations must both be defined because many macros optionally route using either A or B for RB, T, MD, and Q (B path is preferred). Since ALUF[0] is the default for microinstructions, it should contain a non-arithmetic operation to preserve CARRY and OVERFLOW branch conditions.

Also, when running Midas the "B" operation is stored in ALUFM 0 and the NOT A operation in ALUFM 16, so these should be loaded the same way by the microprogram to avoid confusion, unless there is some good reason to do otherwise.

Note that an important choice must be made between the arithmetic and logical versions of ALU←A. The arithmetic version allows XORCARRY with ALU←A but smashes OVERFLOW and CARRY branch conditions on ALU←A.
%

```

ABOPB[PD<,,0,25];      *ALU←B (use n=0)
ABOPA[PD<,,1,0];      *ALU←A
                        *(arithmetic--so XORCARRY ok. Requires no more
                        *time than boolean "A" in absence of carryin)
*      ABOPA[PD<,,n,37];      *ALU←A (logical, so XORCARRY illegal but OVERFLOW
                        *and CARRY branch conditions not smashed).

```

```

*"NOT B" and "NOT A" must be both defined also.
ABOPB[NOT,,15,13];    *NOT B
ABOPA[NOT,,16,1];     *NOT A (use n=16 for shifter)

```

```

*Operations of no args (4th arg = letter E makes emulator-only)
ZALUOP[A0,2,31];      *A0 [= all zeroes]
ZALUOP[A1,3,7];       *A1 [= all ones]

```

```

*Arithmetic operations of two args [do not accept "slow" (external)
*B sources] (5th arg = letter E makes emulator-only)
SALUOP[+,4,,14];      *A+B
*      SALUOP[+,+1,n,,214];    *A+B+1
SALUOP[-,,5,,222];    *A-B
*      SALUOP[-,-1,n,,22];     *A-B-1

```

```

*Boolean operations of two args (5th arg = letter E makes emulator-only)
**Note how synonyms are defined for # (XOR) and = (XNOR, EQV)
ALUOP[AND,,6,,35];    *A AND B
ALUOP[OR,,7,,27];     *A OR B
ALUOP[#,10,,23];      XALUOP[XOR,,n];      *A # B (A XOR B)
*      ALUOP[=,,n,,15];        XALUOP[XNOR,,n];      XALUOP[EQV,,n];
ALUOP[AND NOT,,11,,33]; *A AND NOT B [= A AND (NOT B)]
ALUOP[OR NOT,,12,,17]; *A OR NOT B [= A OR (NOT B)]
*      ALUOP[NOT AND,,n,,21];   *NOT A AND B [= (NOT A) AND B]
*      ALUOP[NOT OR,,n,,5];     *NOT A OR B [= (NOT A) OR B]
*      ALUOP[NOT AND NOT,,n,,11]; *NOT A AND NOT B [= (NOT A) AND (NOT B)]
*      ALUOP[NOT OR NOT,,n,,3]; *NOT A OR NOT B [= (NOT A) OR (NOT B)]

```

```

*Operations of one arg (5th arg = letter E makes emulator-only)
*      AOP[2,,n,6];           *2A
*      AOP[2,+1,n,206];       *2A+1
AOP[+1,13,200];       *A+1
AOP[-1,14,36];        *A-1

```

```

%
Page Numbers: Yes First Page: 1
Heading:
kernel.mc      February 1, 1980  9:20 PM      %
TITLE[KERNEL];
IM[ILC,0];
TOP LEVEL;
* February 1, 1980  9:20 PM
restartDiagnostic:
BEGIN:
    goto[im0];
afterKernel1:
    goto[beginKernel2];
afterKernel2:
    goto[beginKernel3];
afterKernel3:
    goto[beginKernel4];
afterKernel4:

    T←R0;      * R0 SHOULD HAVE ZERO IN IT
    BRANCH[.+2,ALU=0];
    ERROR;

    T←(R1)-1;      * R1 SHOULD HAVE ONE IN IT.
    BRANCH[.+2,ALU=0];
    ERROR;

    T←R1;
    T←T+(RM1);      * RM1 SHOULD HAVE -1 IN IT;
    BRANCH[.+2,ALU=0];
    ERROR;

    T←100000C;
    T←T#(RHIGH1);      * RHIGH1 SHOULD HAVE 100000B
    BRANCH[.+2, ALU=0];
    ERROR;

    T←R10;      * R10 SHOULD HAVE 125252B
    BRANCH[.+2, ALU<0];
    ERROR;

    T←R01;
    DBLBRANCH[.+1, .+2, ALU<0];
    ERROR; * R01 SHOULD HAVE 52525B IN IT

    T←NOT(R01);
    T←T#(R10);      * R01 SHOULD EQUAL NOT(R10);
    BRANCH[.+2,ALU=0];
    ERROR; * NOTE THIS IS NOT A COMPLETELY ACCURATE
* TEST FOR CONTENTS OF R10, R01!
    goto[done];
* CODE for midas debugging
    top level;
set[dbgTbls,100];
11:  branch[11],      at[dbgTbls,0];
12:  noop,      at[dbgTbls,1];
     branch[12],      at[dbgTbls,2];
13:  noop,      at[dbgTbls,3];
     noop,      at[dbgTbls,4];
     branch[13],      at[dbgTbls,5];
14:  noop,      at[dbgTbls,6];
     noop,      at[dbgTbls,7];
     noop,      at[dbgTbls,10];
     branch[14],      at[dbgTbls,11];
15:  noop,      at[dbgTbls,12];
     noop,      at[dbgTbls,13];
     noop,      at[dbgTbls,14];
     noop,      at[dbgTbls,15];
     branch[15],      at[dbgTbls,16];
16:  noop,      at[dbgTbls,17];
     noop,      at[dbgTbls,20];
     noop,      at[dbgTbls,21];
     noop,      at[dbgTbls,22];
     noop,      at[dbgTbls,23];
     branch[16],      at[dbgTbls,24];

```

END;


```

%
Page Numbers: Yes First Page: 1
Heading:
kernel1.mc      May 8, 1979  11:40 AM  %
* INSERT[D1ALU.MC];
* TITLE[PROG1];
* INSERT[PREAMBLE.MC];
%
May 8, 1979  11:40 AM
      Add RoddByPass tests at enbd of xorBypass
March 26, 1979  10:58 AM
      Add overflow test
March 10, 1979  6:43 PM
      Add test of branch conditions when reschedule is ON.
January 18, 1979  5:23 PM
      Remove checkTaskNum, a temporary kludge that caused reschedTest to fail during task circulation
**,
January 9, 1979  10:44 AM
      add reschedTest
%

```

CONTENTS

| TEST | DESCRIPTION |
|--------------------|---|
| (singlestep) | Chec RM to T, T to RM movement |
| aluEQ0 | check the fast branch code |
| aluLT0 | check the fast branch code |
| rEven | check the fast branch code |
| rGEO | check the fast branch code |
| reschedTest | check the reschedule/noreschedule fast branches |
| xorNoBypass | test XOR alu op |
| bypass | test bypass decision logic |
| xorBypass | test XOR alu op, ALLOW BYPASS; R odd bypass test here, too. |
| (alu ops) | Test various alu operations (A+1, A+B, A-1,A-B) |
| Carry | Test carry fast branch |
| (resched+branches) | test effect of resched upon fast branches. |
| freezeBCtest | Test Freeze BC function (emulator only) |
| overflowTest | Test the overflow fast branch function |
| % | |

* September 15, 1978 10:18 AM

%

SINGLE STEP THIS CODE: A AND B MULTIPLEXORS

The point is to determine if it is possible to move data values thru
the alu into different registers.

%

top level;

kernel1:

IM0: T←RM1; *TEST ALL ONES, ALL ZEROS, ALTER. 01, 10
NOOP; * USE NOOP TO AVOID BYPASS LOGIC

IM2: RSCR←T;
NOOP;

IM4: T←R0; * TEST 0
NOOP;

IM6: RSCR ← T;
NOOP;

* NOW MOVE IT THRU A MUX

IM14: T←A←RM1; * TEST ALL ONES
NOOP;

IM16: RSCR ←A← T;
NOOP;

IM20: T←A←R0; * TEST ALL ZEROS
NOOP;

IM22: RSCR← A←T;

* CHECK B MUX THRU FF FIELD: SINGLE STEP THIS CODE

IM23: T←B0; * check that FF,0 works

IM24: T←77400C;

IM25: T←B15; * check that 0,FF works

IM26: T←376C;

* September 15, 1978 10:19 AM

%

END SINGLE STEPPING !!!

GIVEN SIMPLE A AND B PATHS, VALIDATE:

RESULT=0

RESULT<0

R>=0

R EVEN

CNT=0&+1

%

% TEST ALU=0 BY CHECKING EVERY BIT IN THE WORD: GET CONSTANTS FROM
FF AND CHECK THEM FOR =0. USE BYPASS LOGIC!!

These tests assume that there is no difference between amux source and bmux source for fast branches. A
**CTUALLY, the initial set of tests will check amux
sources too!

T contains the value received.

%

aluEQ0FF:

t←B0;

skpUnless[ALU=0],rscr←(A+t); * check it thru Amux

error; * Thinks bit0 is zero

skpUnless[ALU=0];

error; * Thinks bit0 is zero

t←B1;

skpUnless[ALU=0],rscr←(A+t); * check it thru Amux

aluEq0FFB1:

error; * Thinks bit1 is zero

skpUnless[ALU=0];

error; * Thinks bit1 is zero

noop; *here for placement.

t←B2;

skpUnless[ALU=0],rscr←(A+t); * check it thru Amux

aluEq0FFB2:

error; * Thinks bit2 is zero

skpUnless[ALU=0];

error; * Thinks bit2 is zero

t←B3;

skpUnless[ALU=0],rscr←(A+t); * check it thru Amux

error; * Thinks bit3 is zero

skpUnless[ALU=0];

error; * Thinks bit3 is zero

t←B4;

skpUnless[ALU=0], rscr←(A+t); * check it thru Amux

aluEq0FFB4:

error; * Thinks bit4 is zero

skpUnless[ALU=0];

error; * Thinks bit4 is zero

t←B5;

skpUnless[ALU=0], rscr←(A+t); * check it thru Amux

error; * Thinks bit5 is zero

skpUnless[ALU=0];

error; * Thinks bit5 is zero

noop; * here for placement.

t←B6;

skpUnless[ALU=0], rscr←(A+t); * check it thru Amux

aluEq0FFB6:

error; * Thinks bit6 is zero

skpUnless[ALU=0];

error; * Thinks bit6 is zero

t←B7;

skpUnless[ALU=0], rscr←(A+t); * check it thru Amux

error; * Thinks bit7 is zero

skpUnless[ALU=0];

```

        error; * Thinks bit7 is zero

        t←B8;
        skipUnless[ALU=0], rscr←(A←t); * check it thru Amux
aluEq0FFB8:
        error; * Thinks bit8 is zero
        skipUnless[ALU=0];
        error; * Thinks bit8 is zero

        t←B9;
        skipUnless[ALU=0], rscr←(A←t); * check it thru Amux
        error; * Thinks bit9 is zero
        skipUnless[ALU=0];
        error; * Thinks bit9 is zero

        noop; * here for placement.
        t←B10;
        skipUnless[ALU=0], rscr←(A←t); * check it thru Amux
aluEq0FFB10:
        error; * Thinks bit10 is zero
        skipUnless[ALU=0];
        error; * Thinks bit10 is zero

        t←B11;
        skipUnless[ALU=0], rscr←(A←t); * check it thru Amux
        error; * Thinks bit11 is zero
        skipUnless[ALU=0];
        error; * Thinks bit11 is zero

        t←B12;
        skipUnless[ALU=0], rscr←(A←t); * check it thru Amux
aluEq0FFB12:
        error; * Thinks bit12 is zero
        skipUnless[ALU=0];
        error; * Thinks bit12 is zero

        t←B13;
        skipUnless[ALU=0], rscr←(A←t); * check it thru Amux
        error; * Thinks bit13 is zero
        skipUnless[ALU=0];
        error; * Thinks bit13 is zero

        noop; * here for placement.
        t←B14;
        skipUnless[ALU=0], rscr←(A←t); * check it thru Amux
aluEq0FFB14:
        error; * Thinks bit14 is zero
        rscr←(A←t); * check it thru Amux
        skipUnless[ALU=0];
        error; * Thinks bit14 is zero

        t←B15;
        skipUnless[ALU=0], rscr←(A←t); * check it thru Amux
        error; * Thinks bit15 is zero
        skipUnless[ALU=0];
        error; * Thinks bit15 is zero

```

```

%
TEST ALU=0 BY PASSAGE THRU RM AND PASSAGE THRU T

```

For all the alu=0 tests, an error implies the wrong branch was taken.
The known values in RM are used to test the branch

AVOID BYPASS LOGIC!

```

%

```

```

aluEq0RT:
        t←r0;
        skipif[alu=0];
        error; * Thinks r0 is zero
        rscr←t;
        skipif[alu=0],t←r1;
        error;

        skipUnless[alu=0];

```

```
error; * Thinks r1 is zero
rscr←t;
skpUnless[alu=0],t←rm1;
error;
```

```
aluEq0RTM1:
skpUnless[alu=0];
error; * Thinks rm1 is zero
rscr←t;
skpUnless[alu=0],t←r1;
error;
```

```
skpUnless[alu=0];
error; * Thinks r1 is zero
rscr←t;
skpUnless[alu=0],t←r01;
error;
```

```
aluEq0RT01:
skpUnless[alu=0];
error; * Thinks r01 is zero
rscr←t;
skpUnless[alu=0],t←r10;
error;

skpUnless[alu=0];
error; * Thinks r10 is zero
rscr←t;
skpUnless[alu=0],t←rhigh1;
error;

skpUnless[alu=0];
error; * Thinks rhigh1 is zero
rscr←t;
skpUnless[alu=0];
error;
```

```

%
TEST RESULT <0

    For all the alu<0 tests, an error implies the wrong branch was taken.
    The known values in RM are used to test the branch

    AVOID BYPASS LOGIC
%
aluLTORT:
    t←rhigh1;
    skipif[alu<0];
    error; * Thinks rhigh1 >=0
    rscr←t;
    skipif[alu<0];
    error; * Thinks T (=RIGH1) >=0

    t←r10;
    skipif[alu<0];
aluLTORT10:
    error; * Thinks r10 >= 0
    rscr←t;
    skipif[alu<0];
    error; * Thinks T (=r10) >=0

    t←r1;
    skipUnless[alu<0];
aluLTORT1:
    error; * Thinks r1<0
    rscr←t;
    skipUnless[alu<0];
    error; * Thinks T (=r1) >=0

    t←r01;
    skipUnless[alu<0];
aluLTORT01:
    error; * Thinks r10 >= 0
    rscr←t;
    skipUnless[alu<0];
    error; * Thinks T (=r10) >=0

* TEST FOR RESULT EVEN

rEven:
    skipif[r even], t←r0;
    error; * thinks r0 odd

    skipUnless[r even], t←r1;
    error; * Thinks r1 EVEN

    skipif[r even], t←rhigh1;
    error; * Thinks rhigh1 ODD

    skipUnless[r even], t←r01;
    error; * Thinks r01 EVEN

    skipif[r even], t←r10;
    error; * Thinks r10 ODD

rGEO:
    skipif[r >=0],t←r1;
    error; * Thinks r1 <0

    skipif[r>=0],t←r01;
    error; * Thinks r01 <0

    skipif[r>=0],t←r0;
    error; * Thinks r0 <0

    skipUnless[r>=0],t←rm1;
    error; * Thinks RM1>=0

    skipUnless[r>=0],t←rhigh1; * Thinks rhigh1 >=0

```

error;



* January 9, 1979 10:42 AM

```
%
    reschedTest
Set and clear resched; see if we can branch on its value.
%
reschedTest:
    call[checkTaskNum], t←t-t;
    skipif[ALU=0];
    branch[afterResched];

    noreschedule[];
    skipif[reschedule'];
reschedErr1:    * we just cleared resched, yet
error;    * branch condition thinks it is set.

    reschedule[];
    skipif[reschedule];
reschedErr2:    * we just set resched, yet the
error;    * branch condition doesn't realize it.
    noreschedule[];
afterResched:
```


%

September 15, 1978 10:56 AM
TEST XOR USING ALU=0. USE NOOP TO AVOID BYPASS.

Generally, $T \leftarrow RSCR \oplus \text{someFFconstant}$;
 $T \leftarrow T \#(RSCR)$
IF T is non zero, there was an error: one bits in T indicate
the problem.

%

xorNoBypass:

$t \leftarrow (rscr) \oplus B0$;
noop; $t \leftarrow t \#(rscr)$;
skpif[alu=0];
error; * ($T \leftarrow B0 \text{ xor } (RSCR)$) NE 0

$t \leftarrow (rscr) \oplus B1$;
noop; $t \leftarrow t \#(rscr)$;
skpif[alu=0];
error; * ($T \leftarrow B1 \text{ xor } (RSCR)$) NE 0

xorNoBypassB2:

$t \leftarrow (rscr) \oplus B2$;
noop; $t \leftarrow t \#(rscr)$;
skpif[alu=0];
error; * ($T \leftarrow B2 \text{ xor } (RSCR)$) NE 0

$t \leftarrow (rscr) \oplus B3$;
noop; $t \leftarrow t \#(rscr)$;
skpif[alu=0];
error; * ($T \leftarrow B3 \text{ xor } (RSCR)$) NE 0

$t \leftarrow (rscr) \oplus B4$;
noop; $t \leftarrow t \#(rscr)$;
skpif[alu=0];

xorNoBypassB4:

error; * ($T \leftarrow B4 \text{ xor } (RSCR)$) NE 0

$t \leftarrow (rscr) \oplus B5$;
noop; $t \leftarrow t \#(rscr)$;
skpif[alu=0];
error; * ($T \leftarrow B5 \text{ xor } (RSCR)$) NE 0

xorNoBypassB6:

$t \leftarrow (rscr) \oplus B6$;
noop; $t \leftarrow t \#(rscr)$;
skpif[alu=0];
error; * ($T \leftarrow B6 \text{ xor } (RSCR)$) NE 0

$t \leftarrow (rscr) \oplus B7$;
noop; $t \leftarrow t \#(rscr)$;
skpif[alu=0];
error; * ($T \leftarrow B7 \text{ xor } (RSCR)$) NE 0

xorNoBypassB8:

$t \leftarrow (rscr) \oplus B8$;
noop; $t \leftarrow t \#(rscr)$;
skpif[alu=0];
error; * ($T \leftarrow B8 \text{ xor } (RSCR)$) NE 0

$t \leftarrow (rscr) \oplus B9$;
noop; $t \leftarrow t \#(rscr)$;
skpif[alu=0];
error; * ($T \leftarrow B9 \text{ xor } (RSCR)$) NE 0

xorNoBypassB10:

$t \leftarrow (rscr) \oplus B10$;
noop; $t \leftarrow t \#(rscr)$;
skpif[alu=0];
error; * ($T \leftarrow B8 \text{ xor } (RSCR)$) NE 0

$t \leftarrow (rscr) \oplus B11$;
noop; $t \leftarrow t \#(rscr)$;
skpif[alu=0];

```
error; * (T ← B10 xor (RSCR) ) NE 0

t←(rscr)←B12;
noop;t←t#(rscr);
skpif[alu=0];
xorNoBypassB12:
error; * (T ← B12 xor (RSCR) ) NE 0

t←(rscr)←B13;
noop;t←t#(rscr);
skpif[alu=0];
error; * (T ← B13 xor (RSCR) ) NE 0

xorNoBypassB14:
t←(rscr)←B14;
noop;t←t#(rscr);
skpif[alu=0];
error; * (T ← B14 xor (RSCR) ) NE 0

t←(rscr)←B15;
noop;t←t#(rscr);
skpif[alu=0];
error; * (T ← B15 xor (RSCR) ) NE 0
```

* October 25, 1978 4:06 PM

% bypass

This code checks the decision portion of the bypass circuitry. There are at least two different
** issues associated with bypass: 1) should a bypass be done, and 2) do the bypass data paths work. Thi
**s test addresses point 1.

%

rvrel[rmx10, 10];

bypass:

```
RBASE ← 0s;
t←rmx0←cm1;    * this is the old, stable version
rmx0←t-t;      * this is the new version
t←rmx0; * should use bypassed version of rmx0
skpif[alu=0];
```

bypassErr0: * bypass associated w/ rm addr 0 doesn't
error; * seem to work

```
t←rmx1←cm1;    * this is the old, stable version
rmx1←t-t;      * this is the new version
t←rmx1; * should use bypassed version of rmx1
skpif[alu=0];
```

bypassErr1: * bypass associated w/ rm addr 1 doesn't
error; * seem to work

```
t←rmx2←cm1;    * this is the old, stable version
rmx2←t-t;      * this is the new version
t←rmx2; * should use bypassed version of rmx2
skpif[alu=0];
```

bypassErr2: * bypass associated w/ rm addr 2 doesn't
error; * seem to work

```
t←rmx4←cm1;    * this is the old, stable version
rmx4←t-t;      * this is the new version
t←rmx4; * should use bypassed version of rmx4
skpif[alu=0];
```

bypassErr4: * bypass associated w/ rm addr 4 doesn't
error; * seem to work

```
t←rmx10←cm1;   * this is the old, stable version
rmx10←t-t;     * this is the new version
t←rmx10; * should use bypassed version of rmx10
skpif[alu=0];
```

bypassErr10: * bypass associated w/ rm addr 10 doesn't
error; * seem to work

%

This section of the test works by changing Rbase.

%

```
RBASE ← 2s;
t←rmx0←cm1;    * this is the old, stable version
rmx0←t-t;      * this is the new version
t←rmx0; * should use bypassed version of rmx0
skpif[alu=0];
```

bypassErr20: * bypass associated w/ rm addr 20 doesn't
error; * seem to work

```
RBASE ← 4s;
t←rmx0←cm1;    * this is the old, stable version
rmx0←t-t;      * this is the new version
t←rmx0; * should use bypassed version of rmx0
skpif[alu=0];
```

bypassErr40: * bypass associated w/ rm addr 40 doesn't
error; * seem to work

```
RBASE ← 10s;
t←rmx0←cm1;    * this is the old, stable version
rmx0←t-t;      * this is the new version
t←rmx0; * should use bypassed version of rmx0
skpif[alu=0];
```

bypassErr100: * bypass associated w/ rm addr 100 doesn't
error; * seem to work

RBASE ← rbase[defaultRegion];

```
%
August 30, 1977 6:29 PM
TEST XOR USING ALU=0.

    Generally, T ← RSCR ← someFFconstant;
    T ← T#(RSCR)
    IF T is non zero, there was an error: one bits in T indicate
    the problem.
%
* TEST XOR USING ALU=0. CHECK BYPASS.

xorBypass:
    t←(rscr)←B0;
    t←t#(rscr);
    skipif[ALU=0];
    error; * ( T ← B0 xor (RSCR) ) NE 0

    t←(rscr)←B1;
    t←t#(rscr);
    skipif[ALU=0];
    error; * ( T ← B1 xor (RSCR) ) NE 0

    t←(rscr)←B2;
    t←t#(rscr);
    skipif[ALU=0];
xorBypass2:
    error; * ( T ← B2 xor (RSCR) ) NE 0

    t←(rscr)←B3;
    t←t#(rscr);
    skipif[ALU=0];
    error; * ( T ← B3 xor (RSCR) ) NE 0

    t←(rscr)←B4;
    t←t#(rscr);
    skipif[ALU=0];
xorBypass4:
    error; * ( T ← B4 xor (RSCR) ) NE 0

    t←(rscr)←B5;
    t←t#(rscr);
    skipif[ALU=0];
    error; * ( T ← B5 xor (RSCR) ) NE 0

    t←(rscr)←B6;
    t←t#(rscr);
    skipif[ALU=0];
xorBypassB6:
    error; * ( T ← B6 xor (RSCR) ) NE 0

    t←(rscr)←B7;
    t←t#(rscr);
    skipif[ALU=0];
    error; * ( T ← B7 xor (RSCR) ) NE 0

    t←(rscr)←B8;
    t←t#(rscr);
    skipif[ALU=0];
xorBypassB8:
    error; * ( T ← B8 xor (RSCR) ) NE 0

    t←(rscr)←B9;
    t←t#(rscr);
    skipif[ALU=0];
    error; * ( T ← B9 xor (RSCR) ) NE 0

    t←(rscr)←B10;
    t←t#(rscr);
    skipif[ALU=0];
xorBypassB10:
    error; * ( T ← B10 xor (RSCR) ) NE 0

    t←(rscr)←B11;
```

```
t←t#(rscr);
skpif[ALU=0];
error; * (T ← B11 xor (RSCR) ) NE 0

t←(rscr)←B12;
t←t#(rscr);
skpif[ALU=0];
xorBypassB12:
error; * (T ← B12 xor (RSCR) ) NE 0

t←(rscr)←B13;
t←t#(rscr);
skpif[ALU=0];
error; * (T ← B13 xor (RSCR) ) NE 0

t←(rscr)←B14;
t←t#(rscr);
skpif[ALU=0];
xorBypassB14:
error; * (T ← B14 xor (RSCR) ) NE 0

t←(rscr)←B15;
t←t#(rscr);
skpif[ALU=0];
error; * (T ← B15 xor (RSCR) ) NE 0

rscr ← t-t;
rscr ← 1c;
skpif[R ODD], rscr;
RoddByPassErr0: * fast branch r odd bypass doesn't work
error; * Rscr has 1 in it

rscr ← t-t;
skpUnless[R ODD], rscr;
RoddByPassErr1: * fast branch r odd bypass doesn't work.
error; * rscr has zero in it.
```

* August 30, 1977 6:29 PM

%
TEST ALU ADDITION AND SUBTRACTION:
A+1 A+B A-1 A-B

%

Aplus1:

t←(r0)+1; * 1=t←r0+1
t←(r1)#t;
skpif[alu=0];
error;

Aplus1b:

t←(rm1)+1; * 0=t←rm1+1
skpif[alu=0];
error;

Aplus1c:

rscr←CM2;
t←(rscr)+1;
t←t#(rm1); * (T=-1)=-2+1
skpif[alu=0];
error;

Aplus1d:

rscr←5C;
t←(rscr)+1;
rscr←6C;
t←t#(rscr); * (T=6)=5+1
skpif[alu=0];
error;

AplusB:

t←r0;
t←t+(r0); * 0=0+0
skpif[alu=0];
error;

AplusBb:

t←r1;
t←t+(r0);
t←t#(r1); * 1=1+0
skpif[alu=0];
error;

AplusBc:

t←r0;
t←t+(r1);
t←t#(r1); * 1=0+1
skpif[alu=0];
error;

AplusBd:

t←rm1;
t←t+(r0);
t←t#(rm1); * -1=-1+0
skpif[alu=0];
error;

AplusBe:

t←rm1;
t←t+(rm1);
rscr←17776C;
t←t#(rscr); * -2=-1+-1
skpif[alu=0];
error;

AplusBf:

t←rm1;
t←t+(r1); * 0=-1+1
skpif[alu=0];
error;

AplusBg:

t←r1;

```
t←t+(r1);      * 0=1+-1
skpif[alu=0];
error;

AplusBh:
t←r01;
t←t+(r10);
t←t#(r1);      * -1=52525+125252
skpif[alu=0];
error;

AplusBi:
t←r10;
t←t+(r01);
t←t#(r1);      * -1=125252+52525
skpif[alu=0];
error;

AplusBj:
t←rhigh1;
t←t+(rhigh1); *0=100000+100000
skpif[alu=0];
error;

AplusBk:
t←rhigh1;
t←t+(r1);
rscr←77777C;
t←t#(rscr);    * 77777=100000+177777
skpif[alu=0];
error;

AplusBl:
t←r1;
t←t+(rhigh1);
rscr←77777C;
t←t#(rscr);    * 77777=177777+100000
skpif[alu=0];
error;

* August 9, 1977 12:30 PM
%
TEST A-1
%

Aminus1:
rscr←r0;
Aminus1L:      * CHECK A-1 IN A LOOP FOR ALL 16 BIT VALUES.
t←(rscr)-1;
t←t+1;
t←t#(rscr);    * t← (rscr-1+1) xor rscr
skpif[alu=0];
error;
rscr←(rscr)+1; * rscr IS LOOP CTRL
db1Branch[. +1,Aminus1L,ALU=0];
```

```
*
%
%      TEST A-B
%
aMinusB:
    t←r0;
    t←t-(r1);      * T ← 0 -1
    t←t#(r1);
    skipif[alu=0];
    error; * T SHOULD HAVE BEEN -1

aMinusBb:
    t←r0;
    t←t-(r1);      * t←0 - (-1)
    t←t#(r1);
    skipif[alu=0];
    error; * T SHOULD HAVE BEEN 1

aMinusBc:
    t←r0;
    t←t-(rhigh1); * t← 0 - (100000)
    t←t#(rhigh1);
    skipif[alu=0];
    error; * T SHOULD HAVE BEEN 100000

aMinusBd:
    t←100C;
    t←t-(r1);      * t ← 100 -1
    rscr←77C;
    t←t#(rscr);
    skipif[alu=0];
    error; * T SHOULD HAVE BEEN 77

aMinusBe:
    t←rscr←17C;
    t←t-(rscr);    * t ← 17 - 17
    t←t#(r0);
    skipif[alu=0];
    error; * T SHOULD HAVE BEEN 0

aMinusBf:
    t←rscr←177C;
    t←t-(rscr);    * t ← 177 - 177
    t←t#(r0);
    skipif[alu=0];
    error; * T SHOULD HAVE BEEN 0

aMinusBg:
    t←rscr←377C;
    t←t-(rscr);    * t ← 377 - 377
    t←t#(r0);
    skipif[alu=0];
    error; * T SHOULD HAVE BEEN 0

aMinusBh:
    t←rscr←400C;
    t←t-(rscr);    * t ← 400 - 400
    t←t#(r0);
    skipif[alu=0];
    error; * T SHOULD HAVE BEEN 0

aMinusBi:
    t←rscr←777C;
    t←t-(rscr);    * t ← 777 - 777
    t←t#(r0);
    skipif[alu=0];
    error; * T SHOULD HAVE BEEN 0

aMinusBj:
    t←rscr←1777C;
    t←t-(rscr);    * t ← 1777 - 1777
    t←t#(r0);
    skipif[alu=0];
    error; * T SHOULD HAVE BEEN 0
```




... ..

* January 18, 1979 5:24 PM

%

TEST FAST BRANCH CONDITION: CARRY

FIRST TEST WHEN WE KNOW THERE IS NO CARRY, THEN TRY TO
GENERATE A CARRY AND BRANCH ON IT. NOTE: THIS CODE DEPENDS UPON
THE ALU FUNCTIONS PLUS AND MINUS WORKING.

%

carryNo:

t←(r0)+(r0);
skpUnless[carry];
error; * r0 + r0 SHOULD NOT CAUSE CARRY

carryNob:

t←rm1;
t←t+(r0);
skpUnless[carry];
error; * rm1 + r0 SHOULD NOT CAUSE CARRY

carryNoc:

t←r10;
t←t+(r01);
skpUnless[carry];
error; * r10 + r01 SHOULD NOT CAUSE CARRY

carryNod:

t←77777C;
t←t+(r0);
skpUnless[carry];
error; * 77777C + r0 SHOULD NOT CAUSE CARRY

carryNoe:

t←r0;
t←t+(r1);
skpUnless[carry];
error; * r0 + r1 SHOULD NOT CAUSE CARRY

carryNof:

t←r01;
t←t-(r10);
skpUnless[carry];
error; * r01 - r10 SHOULD NOT CAUSE CARRY

carryNog:

t←r0;
t←t-(r10);
skpUnless[carry];
error; * r0 + r10 SHOULD NOT CAUSE CARRY

* NOW TRY SOMETHINGS THAT SHOULD GENERATE A CARRY

carryYes:

t←rm1;
t←t-(r1);
skpif[carry];
error; * -1 -(+1) SHOULD CAUSE CARRY

carryYesb:

t←rm1;
t←t+(rhigh1);
skpif[carry];
error; * -1 + 100000 SHOULD CAUSE CARRY

carryYesc:

t←rm1;
t←t+(r1);
skpif[carry];
error; * -1 + 1 SHOULD CAUSE CARRY

carryYesd:

t←rhigh1;
t←t-(r01);
skpif[carry];
error; * 100000 - r01 SHOULD CAUSE CARRY

carryYese:

```
t←rhigh1;
t←t+(rhigh1);
skipif[carry];
error; * 100000 + 100000 SHOULD CAUSE CARRY
```

* NOW COMPLICATE THINGS INTERLEAVING ALU OPS W/ TESTS

carryOps:

```
t←r0;
t←t-(r1); * t←0-1
skipUnless[carry],t←t+(r1); * t←-1+-1
error; * 0-1 SHOULD NOT CAUSE CARRY
```

carryOpsb:

```
skipif[carry], t←t+(rhigh1); * T ← -2+100000
error; * -1+-1 SHOULD CAUSE CARRY
```

carryOpsc:

```
skipif[carry],t←t+(rhigh1); * t← 77776 + 100000
error; * -2 + 100000 SHOULD CAUSE CARRY
```

carryOpsd:

```
skipUnless[carry];
error; * 77776 + 100000 SHOULD NOT CAUSE CARRY
```

carryOpse:

```
t←r1;
t←t-(r1); * t← -1-(+1)
skipif[carry],t←t-(r01); * t← -2 -r01
error; * -1-1 SHOULD CAUSE CARRY
```

carryOpsf:

```
skipif[carry];
error; * 177776 - 52525 SHOULD CAUSE CARRY
reschedule;
```

```
* March 10, 1979 6:42 PM
* test the branch conditions when reschedule is ON
  t←r0;
  t←t-(r1);      * t←0-1
  skipUnless[carry],t←t+(r1);      * t←-1+-1
  error; * 0-1 SHOULD NOT CAUSE CARRY
carryOpsRb:
  skipif[carry], t←t+(rhigh1);      * T ← -2+100000
  error; * -1+-1 SHOULD CAUSE CARRY

carryOpsRc:
  skipif[carry],t←t+(rhigh1);      * t← 77776 + 100000
  error; * -2 + 100000 SHOULD CAUSE CARRY

carryOpsRd:
  skipUnless[carry];
  error; * 77776 + 100000 SHOULD NOT CAUSE CARRY

carryOpsRe:
  t←r1;
  t←t-(r1);      * t← -1-(+1)
  skipif[carry],t←t-(r01); * t← -2 -r01
  error; * -1-1 SHOULD CAUSE CARRY

carryOpsRf:
  skipif[carry];
  error; * 177776 - 52525 SHOULD CAUSE CARRY

  t←r0;
  skipif[ALU=0];
rescheq0br:
  error;
  t←r1;
  skipif[ALU#0];
reschne0br:
  error;
  skipif[r even], B←r0;
reschevenbr:
  error;
  skipif[r odd], B←r1;
reschoddbr:
  error;
  t←rhigh1;
  skipif[alu<0];
reschl1t0br:
  error;
  t←r0;
  skipif[alu>=0];
reschge0br:
  error;

  noreschedule;
```

* September 15, 1978 11:38 AM

%

TEST FREEZEBC FUNCTION

Generate the two different branch conditions and freeze them. Force the carry to be explicitly different, see if the frozen branch is still there. Unfreeze and make sure the expected results happen.

%

freezeBCtest:

t←rm1;
t←t+(r1); * t← 0 ← -1+1 (SHOULD CAUSE carry)
skpif[carry],t←t+(r1),freezeBC; * FREEZE[carry=1]
error;

* carry WAS FROZEN. CONTINUE THAT WAY (carry←1, RESULT←0)

freezeBC1a:

skpif[alu=0],freezeBC; * (result was ZERO)
error;

t←(rm1)+(rm1),freezeBC; * Would normally CAUSE RESULT <0
skpif[alu>=0],freezeBC; * (result was ZERO)
error;

freezeBC1b:

t←(r1)+(r1),freezeBC; * t← 0+1 (carry Would NORMALLY BE ZERO)
skpif[carry],freezeBC;
error;

t←tAND(rm1),freezeBC; * t←1 and -1 (TEST IT A FEW MORE TIMES)
skpif[carry],freezeBC;
error; * carry SHOULD HAVE BEEN 1

freezeBC1c:

t←t+(r0),freezeBC; * t←1+0
skpif[carry],freezeBC;
error; * carry SHOULD HAVE BEEN 1

* ALLOW A NEW alu RESULT CONDITION, KEEP carry THE SAME(carry=1, RESULT=77777)

freezeBC2a:

t←rm1,freezeBC;
t←t+(rhigh1); * carry←1, RESULT←77777
t←(r0)+(r0),freezeBC;
skpif[alu#0],freezeBC; * result was 77777
error;

freezeBC2b:

t←t+(r0),freezeBC; * Would normally ZERO carry
skpif[carry],freezeBC; * carry SHOULD BE ONE
error;

freezeBC2c:

t←rm1,freezeBC;
t←t+(r0),freezeBC;
skpif[alu>=0],freezeBC; * result was 77777
error;

* FORCE carry←0, RESULT←0

t←r0,freezeBC;
t←t+(r0); * t←0+0 (SHOULD CAUSE carry←0, RESULT←0)
t←(rhigh1)+(rhigh1),freezeBC; * Would NORMALLY CAUSE carry←1

freezeBC3a:

skpUnless[carry],freezeBC; * FREEZE IT AT ZERO
error; * EXPECTED 0 carry GOT 1 carry

t←(r1)+(r1),freezeBC;
skpif[alu=0],freezeBC; * test it again just to see
error;

freezeBC3b:

t←(rm1)+(rm1),freezeBC;
skpif[alu>=0],freezeBC; * test it again just to see
error;

```
* FORCE carry←0, RESULT ← -1
  t←(r1)+(r1);
  t←t+(r1); * -2+1 ==> carry←0, RESULT←-1

freezeBC4a:
  t←t+(r1),freezeBC; * -1+1 Would normally CAUSE carry←1
  skpUnless[carry],freezeBC;
  error;

freezeBC4b:
  t←(r1)+(r1),freezeBC; *Would normally CAUSE alu>=0
  skpUnless[alu>=0],freezeBC;
  error;

freezeBC4c:
  skpUnless[alu=0];
  error;
```

* March 26, 1979 11:04 AM

%*****

overflowTest

Perform an exhaustive test of the overflow condition. Even though we expect the arithmetic result of R**M+T to be identical to T+RM, we test all possible combinations since the arithmetic gets implemented**inside a rather complicated chip.

The tables below show the aluA and aluB inputs, and the carry out values for b0, b1. Notice the conten**ts of the table are not the sum of a,b, but the carry out values. The subtraction table shows the or**iginal input for B and then its converted value after the number gets converted to a twos complement**value (the chip converts it to the twos complement form, then adds).

For Addition

| | | | | | |
|----------|----|----|----|----|--|
| B input= | 00 | 01 | 10 | 11 | |
| A input | | | | | |
| 00 | 00 | 00 | 00 | 00 | |
| 01 | 00 | 01 | 00 | 11 | |

** out values

| | | | | | |
|-----------------|----|----|----|----|--|
| 10 | 00 | 00 | 10 | 10 | |
| **e carry-in to | | | | | |
| 11 | 00 | 11 | 10 | 11 | |

Notice these values represent the carry

for b0,b1 during addition. They presum

b1 is zero

%*****

mc[x01, 40000];

mc[x10, 100000];

mc[x11, 140000];

overflowTest:

t ← t-t;

t ← t + t;

skipif[overflow'];

overflErr0: * 0+0 should not cause overflow

error;

t←t-t;

t←t+(x01);

skipif[overflow'];

overflErr1: * see 0 + 01 entry

error;

t ← t-t;

t ← t + (x10);

skipif[overflow'];

overflErr2: * see 0 + 10 entry

error;

t←t-t;

t←t+(x11);

skipif[overflow'];

overflerr3: * see 0 + 11 entry

error;

t←rscr+x01; * keep x01 in rscr for a while

t ← t + (0c);

skipif[overflow'];

overflErr4: * see 01 + 0 entry

error;

t←rscr;

t ← t + (x01);

skipif[overflow']; * FIRST TRY FOR OVERFLOW

overflErr5: * see 01 + 01 entry

error;

t←rscr;

t ← t + (x10);

skipif[overflow'];

overFLerr6:

error; * see 01 + 10 entry

t←rscr;

t←t+(x11);

skipif[overflow'];

overflErr7: * see 01 + 11 entry

error;

t←rscr+x10; * keep x10 in rscr for a while

t ← t + (0c);

skipif[overflow'];

overflErr10: * see 10 + 0 entry

error;

t←rscr;

t ← t + (x01);

skipif[overflow'];

overflErr11: * see 10 + 01 entry

```
    error;
    t←rscr;
    t ← t + (x10);
    skipif[overflow];
overfLerr12:
    error; * see 10 + 10 entry
    t←rscr;
    t←t+(x11);
    skipif[overflow];
overf1Err13: * see 10 + 11 entry
    error;

    t←rscr←x11; * keep x11 in rscr for a while
    t ← t + (0c);
    skipif[overflow'];
overf1Err14: * see 11 + 0 entry
    error;
    t←rscr;
    t ← t + (x01);
    skipif[overflow'];
overf1Err15: * see 11 + 01 entry
    error;
    t←rscr;
    t ← t + (x10);
    skipif[overflow];
overfLerr16:
    error; * see 11 + 10 entry
    t←rscr;
    t←t+(x11);
    skipif[overflow'];
overf1Err17: * see 11 + 11 entry
    error;
goto[afterKernel11];
```



```
%
Page Numbers: Yes First Page: 1
Heading:
kernel2.mc      January 18, 1979  2:07 PM      %
* INSERT[D1ALU.MC];
* TITLE[KERNEL2];
* INSERT[PREAMBLE.MC];
top level;
beginKernel2:
%
      January 20, 1978  3:13 PM
%

%
TEST  CONTENTS
cntRW  read and write CNT
cntFFrw read and write CNT, load from FF
cntFcn test CNT=0&+1 fast branch
NotAtest      test alu op, NOT A
NotBtest      test alu op, NOT B
AandBtest     test alu op, A AND B
AorBtest      test alu op, A OR B
LINKRW read and write LINK
callTest     global and local subroutine calls
QtestRW read and write Q, q lsh 1, q rsh 1
tioaTest     load and read tioa from FF and from bmux
STKPttestRW  read and write STKP, perform TIOA&STKP
rstkTest0    write different RM address from one read
%
%
January 18, 1979  2:07 PM
      Add tioaTest
%
```

* October 19, 1978 5:22 PM

```
%  
    TEST ALL THE BITS IN CNT: REMEMBER THAT CNT CAN BE LOADED FROM  
    BOTH B AND FF.  
%
```

```
cntRW:  
    t ← cnt ← r0; * test loading cnt w/ 0  
    rscr ← cnt;  
    t ← t # (rscr); * t ← bits read from cnt # expected bits  
    skipif[ALU=0];  
cntErr1: * t = bad bits, rscr = expected  
    error; * value of cnt  
  
    t ← cnt ← rml; * test loading cnt w/ -1  
    rscr ← cnt;  
    t ← t # (rscr); * t ← bits read from cnt # expected bits  
    skipif[ALU=0];  
cntErr2: * t = bad bits, rscr = expected  
    error; * value of cnt  
  
    t ← cnt ← r01; * test loading cnt w/ alternating 01  
    rscr ← cnt;  
    t ← t # (rscr); * t ← bits read from cnt # expected bits  
    skipif[ALU=0];  
cntErr3: * t = bad bits, rscr = expected  
    error; * value of cnt  
  
    t ← cnt ← r10; * test loading cnt w/ alternating 10  
    rscr ← cnt;  
    t ← t # (rscr); * t ← bits read from cnt # expected bits  
    skipif[ALU=0];  
cntErr4: * t = bad bits, rscr = expected  
    error; * value of cnt
```

* October 19, 1978 8:33 PM

```
cntFFrw:  
    * TEST FF BITS FOR LOADING cnt  
    cnt←1s;  
    t←cnt;  
    t←t#(r1); * we set it to 1; check the val.  
    skipif[ALU=0];  
cntFFrw1: * t=bad bits, 1=expected value  
    error;  
  
    cnt ← 2s;  
    t←cnt;  
    t ← t # (2c);  
    skipif[ALU=0];  
cntFFrw2: * t = bad bits, 2 = expected value  
    error;  
  
    cnt ← 4s;  
    t ← cnt;  
    t ← t # (4c);  
    skipif[ALU=0];  
cntFFrw3: * t = bad bits, 4 = expected  
    error;  
  
    cnt ← 10s;  
    t ← cnt;  
    t ← t # (10c);  
    skipif[ALU=0];  
cntFFrw4: * t = bad bits, 10c = expected  
    error;
```

* October 30, 1978 1:54 PM

%

TEST cnt BY LOOPING FOR ALL VALUES OF cnt
AT POINTS TESTED, cnt AND rscr SHOULD BE EQUAL.

```
rscr←cnt←-1;    -- test cnt for maximum iterations
WHILE cnt NE 0 DO
    cnt←cnt-1;
    IF rscr=0 THEN ERROR;
    rscr←rscr-1;
    ENDLOOP;
IF rscr NE 0 THEN ERROR
%
cntFcn:
    t ← rscr←cm1;    * t ← rscr ← initial value into cnt
    cnt←t;    * cnt ← initial value

cntFcnIL:
    branch[cntFcnXitIL, cnt=0&-1], PD ← rscr;
    skipUnless[ALU=0], PD←rscr;
cntFcnErr1:    * value of rscr suggests we
    error;    * should have exited
    branch[cntFcnIL], rscr←(rscr)-1;
cntFcnXitIL:
    skipif[ALU=0];
cntFcnErr2:    * rscr#0. value of rscr suggests we
    error;    * should not have exited.

    cnt ← r0;    * test cnt for initial value = zero
    skipif[cnt=0&-1];
cntFcnErr3:    * didn't notice first value we loaded
    error;    * was zero
```

```
* August 31, 1977 12:48 PM
% Test not A, not B
%
```

```
NotAtest:
```

```
t←not(A←r0);
t←t#(rm1);
skipif[ALU=0];
error; * ~R0 # RM
```

```
NotAb:
```

```
t←not(A←rm1);
t←t#(r0);
skipif[ALU=0];
error; * ~rm1 # r0
```

```
NotAc:
```

```
t←not(A←r01);
t←t#(r10);
skipif[ALU=0];
error; * ~r01 # r10
```

```
NotAd:
```

```
t←not(A←r10);
t←t#(r01);
skipif[ALU=0];
error; * ~r10 # r01
```

```
NotAe:
```

```
rscr←177776C;
t←not(A←r1);
t←t#(rscr);
skipif[ALU=0];
error; * ~r1 # 177776
```

```
NotBtest:
```

```
t←not(B←r0);
t←t#(rm1);
skipif[ALU=0];
error; * ~R0 # RM
```

```
NotBTestb:
```

```
t←not(B←rm1);
t←t#(r0);
skipif[ALU=0];
error; * ~rm1 # r0
```

```
NotBTestc:
```

```
t←not(B←r01);
t←t#(r10);
skipif[ALU=0];
error; * ~r01 # r10
```

```
NotBTestd:
```

```
t←not(B←r10);
t←t#(r01);
skipif[ALU=0];
error; * ~r10 # r01
```

```
NotBTeste:
```

```
rscr←177776C;
t←not(B←r1);
t←t#(rscr);
skipif[ALU=0];
error; * ~r1 # 177776
```

```
%
```

```
Test A AND B
Assume a,b source dont matter. Ie.,
t←(b←t) and (a←r) =
t←(a←t) and (b←r)
```

```
%
```

```
AandBtest:
```

```
t←rm1;
t←tAND(rm1);
```

```
t←t#(rm1);
skpif[ALU=0];
error; * (rm1 AND rm1) #rm1
```

```
t←r01;
t←tAND(r10);
skpif[ALU=0];
error; * (r01 AND r10)
```

```
t←r0;
t←tAND(rm1);
skpif[ALU=0];
error; * r0 AND rm1
```

%

```
Test A orB.
Assume same as AandB test.
```

%

AorBtest:

```
t←rm1;
t←tOR(r0);
t←t#(rm1);
skpif[ALU=0];
error; * (rm1 OR r0) # rm1
```

AorBtestb:

```
t←r01;
t←tOR(r10);
t←t#(rm1);
skpif[ALU=0];
error; * (r01 OR r10) # rm1
```

AorBtestc:

```
t←rm1;
t←tOR(rm1);
t←t#(rm1);
skpif[ALU=0];
error; * (rm1 OR rm1) # rm1
```

AorBtestd:

```
t←r01;
t←tOR(r01);
t←t#(r01);
skpif[ALU=0];
error; * (r01 OR r01) # r01
```

AorBteste:

```
t←r10;
t←tOR(r10);
t←t#(r10);
skpif[ALU=0];
error; * (r10 OR r10) # r10
```

AorBtestf:

```
t←(r0)OR(r0);
t←t#(r0);
skpif[ALU=0];
error; * (r0 OR r0) # r0
```

* February 17, 1978 8:51 AM

%

LINK READ/WRITE TEST + MINOR TEST OF CALL

FOR I IN[0..7777B] DO

LINK←I;

CHECK←LINK;

CHECK ← BITAND[CHECK,7777B];

IF CHECK NE LINK THEN ERROR;

ENDLOOP;

minor test of LINK, call

%

linkRW:

rscr←7777C; * BEGIN W/ MAX LINK VALUE & COUNT DOWN

linkL:

link ← rscr;

t ← link;

t ← t and (77777C); * ISOLATE 15 BITS 'CAUSE OF DMUX DATA

t←t#(rscr);

skpif[alu=0];

linkErr1:

error; * LINK DOESN'T HAVE THE VALUE WE LOADED

rscr←(rscr)-1;

db1Branch[linkL,afterLink, alu<0];

afterLink:

```

* November 3, 1978 6:40 PM
%
    TEST Q: READ AND WRITE
    FOR I IN [0..177777B] DO
        Q←I;
        t←Q XOR I;
        IF T #0 THEN error;
    ENDLOOP;
then test q lsh 1, q rsh 1 w/ selected values
%
QtestRW:
    rscr←r0;

QRWL:
    Q←(rscr);
    t←(A←rscr)#(B←Q);
    skipif[ALU=0];

QrwErr:
    error;
    rscr←(rscr)+1;
    db1Branch[.+1,QRWL,ALU=0];

* now check rsh1, lsh1
    q ← r0;
    q lsh 1;          * q ← 0 lsh 1
    PD ← q;
    skipif[ALU=0];
qr0Lerr:          * r0 lsh 1 should be zero
    error;

    q ← r01;
    q lsh 1;          * q ← r01 lsh 1
    (q) # (r10);
    skipif[alu=0];
qr10Lerr:         * r10 lsh1 should be r01. (zero fill)
    error;

    q ← rm1;
    q lsh 1;
    t ← cm2;
    (q) # t;
    skipif[ALU=0];
qrm1Lerr:         * -1 lsh1 w/ zero fill should be -2
    error;

    q ← rhigh1;
    q lsh 1;          * q ← 100000B lsh 1
    PD ← q;
    skipif[ALU=0];
qrhigh1Lerr:     * rhigh1 (100000B) lsh1 w/ zero fill should
    error; * zero

    q ← r0;
    q rsh 1;          * zero rsh1 should be zero
    PD ← q;
    skipif[ALU=0];
qr0Rerr:
    error; * zero rsh1 should be zero

    q ← r10;
    q rsh 1;          * q ← r10 rsh 1
    (q) # (r01);
    skipif[ALU=0];
qr10Rerr:        * r10 rsh 1 w/ zero fill
    error; * should be r01

    q ← rm1;
    q rsh 1;          * q ← -1 rsh 1;
    t ← 77777c;
    (q) # t;
    skipif[ALU=0];
qrm1Rerr:        * -1 rsh 1 w/ zero fill should be 77777B

```

```
error;
q ← rhigh1;
q rsh 1;      * q ← 100000B rsh 1
t ← 40000C;
(q) # t;
skpif[ALU=0];
qrhigh1Rerr: * rhigh1 rsh1 should be 40000B
error;
```



```

* January 18, 1979  1:29 PM
%
tioaTest
Test the processor's ability to read and write TIOA.  Write TIOAk from both FF constants and from RM.
%
tioaTest:
    t ← 377c;
    cnt ← t;
    rscr2 ← t-t;
tioaL:
    tioa ← rscr2;  * RSCR2 = value we load into Tioa
    call[getTioa]; * rtn Tioa, still left justified, in t
    rscr ← (rscr2) # t;
    skipif[ALU=0];
tioaErr1:
    * We wrote tioa w/ contents of rscr2, got
    error; * back the value in t.  Bad bits in rscr.
    loopUntil[cnt=0&-1, tioaL], rscr2 ← (rscr2) + (b7);  * increment rscr2

* Here are device ddeclarations to keep micro happy. We use them to set Tioa directly from FF.
device[dvc5, b13!]; device[dvc6, b14!]; device[dvc7, b15!];
mc[tioa.0thru4C, b0,b1,b2,b3,b4];
mc[tioa.mask, 177400];

    tioa ← r0;  * zero all the bis of tioa
    tioa[dvc7];  * should set tioa[5:7] to 1
    call[getTioa];
    rscr ← (t) # (b7);  * only one bit should be set
    skipif[ALU=0];
tioaErr2:
    * tioa should be 1, (= 1 lshift 8 = 400)
    error; * t = value of tioa, rscr = bad bits.

    tioa[dvc6];  * should set tioa[5:7] to 2
    call[getTioa];
    rscr ← t # (b6);  * tioa should be 2, (= 1 lshift 9 = 1000)
    skipif[ALU=0];
tioaErr3:
    error; * rscr = bad bits, t = tioa left justified

    tioa[dvc5];  * tioa should be 4 (= 1 lshift 10 = 2000)
    call[getTioa];
    rscr ← t # (b5);
    skipif[ALU=0];
tioaErr4:
    error; * rscr = bad bits, t = tioa left justified

    tioa ← rm1;  * all ones into tioa
    tioa[dvc7];
    call[getTioa];
    rscr ← tioa.0thru4C;
    rscr ← (rscr) or (b7);  * only should have set tioa[5:7];
    rscr ← t # (q←rscr);  * q = expected value
    skipif[ALU=0];
tioaErr5:
    error; * t = tioa, left justified; rscr = bad bits, q = expected value

    tioa[dvc6];
    call[getTioa];  * set tioa[5:7] to 2
    rscr ← tioa.0thru4C;
    rscr ← (rscr) or (b6);  * only should have set tioa[5:7];
    rscr ← t # (q←rscr);  * q = expected value
    skipif[ALU=0];
tioaErr6:
    * q = expected value
    error; * t = tioa, left justified; rscr = bad bits

    tioa[dvc5];
    call[getTioa];  * set tioa[5:7] to 4
    rscr ← tioa.0thru4C;
    rscr ← (rscr) or (b5);  * only should have set tioa[5:7];
    rscr ← t # (q←rscr);  * q = expected value
    skipif[ALU=0];
tioaErr7:
    * q = expected value
    error; * t = tioa, left justified; rscr = bad bits

    branch[afterTioa];
getTioa: subroutine;

```

```
t ← TIOA&STKP;  
return, t ← t and (177400C); * isolate left byte  
top level;
```

afterTioa:

* October 19, 1978 8:54 PM

```
%
    TEST STKP: READ AND WRITE

    FOR I IN[0..377B] DO
        STKP←I;
        t←TIOA&STKP[]
        t←t and (stkpMask);
        t←t XOR I;
        IF T # 0 THEN error;
    ENDLOOP;
%

STKPttestRW:
    t←r0;
    rscr←t; * rscr = values loaded into stackp
    rscr2 ← t+377C; * MASK TO ISOLATE STACKP
    cnt←t; * mask just happens to be count, too

stkpL:
    STKP←rscr;      * LOAD STKP FROM rscr
    t ← (TIOA&STKP);
    t←t AND (rscr2);      * READ AND MASK THE VALUE
    t←t#(rscr);
    skpif[ALU=0];

stkpErr:
    error; * error: DIDN'T READ WHAT WE LOADED
    db1Branch[.+1,stkpL,CNT=0&-1],rscr←(rscr)+1;
```

* October 26, 1978 12:03 PM

%

rstkFF Test the FF operation that replaces rstk with a value from the FF field during rm Writi
 **ng. Test each bitpath only.

%

rstkFF:

```
q ← rmx0;      * save rmx0
rmx0 ← t-t;    * background test rm location w/ zero
t ← rmx7 ← cm1; * KEEP -1 IN RMX7, AND T
rmx0 ← rmx7;  * write into RM w/ rstk from FF field
t # (rmx0);   * compare target RM w/ expected value
skipif[alu=0];
```

rstkFF0Err: * can't write into rstk0 w/ ff
 error;

```
rmx0 ← q;      * restore old value
q ← rmx1;      * save rmx1
rmx1 ← t-t;    * background test rm location w/ zero
rmx1 ← rmx7;  * write into RM w/ rstk from FF field
t # (rmx1);   * compare target RM w/ expected value
skipif[ALU=0];
```

rstkFF1Err: * can't write into rstk1 w/ ff
 error;

```
rmx1 ← q;      * restore old value
q ← rmx2;      * save rmx2
rmx2 ← t-t;    * background test rm location w/ zero
rmx2 ← rmx7;  * write into RM w/ rstk from FF field
t # (rmx2);   * compare target RM w/ expected value
skipif[ALU=0];
```

rstkFF2Err: * can't write into rstk2 w/ ff
 error;

```
rmx2 ← q;      * restore old value
q ← rmx4;      * save rmx4
rmx4 ← t-t;    * background test rm location w/ zero
rmx4 ← rmx7;  * write into RM w/ rstk from FF field
t # (rmx7);   * compare target RM w/ expected value
skipif[ALU=0];
```

rstkFF4Err: * can't write rstk4 w/ ff
 error;

```
rmx4 ← q;      * restore old value
q ← rmx10;     * save rmx10
rmx10 ← t-t;   * background test rm location w/ zero
rmx10 ← rmx7; * write into RM w/ rstk from FF field
t # (rmx10);  * compare target RM w/ expected value
skipif[ALU=0];
```

rstkFF10Err: * can't write rstk10 w/ ff
 error;

```
rmx10 ← q;
```

* October 26, 1978 6:14 PM

%

rbaseFF test the facility that changes the value of rbase when rm storing occurs.

%

* sibling[FoosBrotherInRegion5, 5, foo] * declare FoosBrotherInRegion5 as an RM
 * location in rmRegion 5 with its rstk value the same as the one for foo. Eg., if foo is
 * located at rm addr 17,,12 (rbase = 17, rstk = 12) then FoosBrotherInRegion5 is located
 * at rm addr 5,,12

```
m[sibling,
  rm[#1, add[lshift[#2,4], and[17,ip[#3]]]]
];
```

```
sibling[rb0rm0, 0, rmx0];      sibling[rb1rm1, 1, rmx1];      sibling[rb2rm2, 2, rmx2];
sibling[rb4rm4, 4, rmx4]; sibling[rb10rm10, 10, rmx10];
```

rbaseFF:

```
rbase ← rbase[defaultRegion];
q ← rmx0;      * save current value for "source" rm
rb0rm0 ← t-t; * zero "destination" rm
rmx0 ← cm1;   * t ← "source rm" ← -1
rb0rm0 ← rmx0; * "destin" rm (different rbase)← source rm
rbase ← 0s;   * check the result. First fetch the value in
t ← rmx0, RBASE ← rbase[defaultRegion]; * the destination rm, then compare it to
t # (rmx0);   * the source rm. An error means we didn't
skpif[ALU=0];
```

```
rbaseFF0Err: * succeed in writing rm with rbase=0 from
error; * ff field. t = real val, rmx0=expected val.
```

```
rmx0 ← q;      * restore old value
q ← rmx1;      * save current value for "source" rm
t ← rmx1 ← cm1; * t ← "source rm" ← -1
rb1rm1 ← t-t; * zero "destination" rm
rb1rm1 ← rmx1; * "destin" rm (different rbase)← source rm
RBASE ← 1s;   * check the result. First fetch the value in
t ← rmx1, RBASE ← rbase[defaultRegion]; * the destination rm, then compare it to
t # (rmx1);   * the source rm. An error means we didn't
skpif[ALU=0];
```

```
rbaseFF1Err: * succeed in writing rm with rbase=0 from
error; * ff field. t = real val, rmx0=expected val.
```

```
rmx1 ← q;      * restore old value
q ← rmx2;      * save current value for "source" rm
t ← rmx2 ← cm1; * t ← "source rm" ← -1
rb2rm2 ← t-t; * zero "destination" rm
rb2rm2 ← rmx2; * "destin" rm (different rbase)← source rm
RBASE ← 2s;   * check the result. First fetch the value in
t ← rmx2, RBASE ← rbase[defaultRegion]; * the destination rm, then compare it to
t # (rmx2);   * the source rm. An error means we didn't
skpif[ALU=0];
```

```
rbaseFF2Err: * succeed in writing rm with rbase=0 from
error; * ff field. t = real val, rmx0=expected val.
```

```
rmx2 ← q;      * restore old value
q ← rmx4;      * save current value for "source" rm
t ← rmx4 ← cm1; * t ← "source rm" ← -1
rb4rm4 ← t-t; * zero "destination" rm
rb4rm4 ← rmx4; * "destin" rm (different rbase)← source rm
RBASE ← 4s;   * check the result. First fetch the value in
t ← rmx4, RBASE ← rbase[defaultRegion]; * the destination rm, then compare it to
t # (rmx4);   * the source rm. An error means we didn't
skpif[ALU=0];
```

```
rbaseFF4Err: * succeed in writing rm with rbase=0 from
error; * ff field. t = real val, rmx0=expected val.
```

```
rmx4 ← q;      * restore old value
q ← rmx10;     * save current value for "source" rm
t ← rmx10 ← cm1; * t ← "source rm" ← -1
rb10rm10 ← t-t; * zero "destination" rm
rb10rm10 ← rmx10;
RBASE ← 10s;  * check the result. First fetch the value in
t ← rmx10, RBASE ← rbase[defaultRegion]; * the destination rm, then compare it to
t # (rmx10); * the source rm. An error means we didn't
skpif[ALU=0];
```

```
rbaseFF10Err: * succeed in writing rm with rbase=0 from
error; * ff field. t = real val, rmx0=expected val.
```

```
%
Test RSTK destination function:
```

```
    FOR I IN [0..7] DO
        FOR J IN [0..7] DO
            RBASE[I]+RBASE[J];
            t←RBASE[I];
            IF T#RBASE[J] THEN error;
        ENDLOOP;
    ENDLOOP;
```

```
Of course, this code is "expanded" inline rather than in a loop
%
```

```
*    FOR I IN [0..7] DO RBASE[0] ← RBASE[I]; (EXCEPT FOR ←RBASE[0])
```

```
rstkTest0:
```

```
    rscr←3C;
    rscr2←4C;
```

```
    Q←r0;
    t←r0←r1;
    t#(Q);
    skpUnless[ALU=0], t←t#(r0);
    error;
    skpif[ALU=0];
    error;
```

```
    t←r0←rm1;
    t#(Q);
    skpUnless[ALU=0], t←t#(r0);
```

```
rstkTest02:
```

```
    error;
    skpif[ALU=0];
    error;
```

```
    t←r0←r01;
    t#(Q);
    skpUnless[ALU=0], t←t#(r0);
    error;
    skpif[ALU=0];
    error;
```

```
    t←r0←r10;
    t#(Q);
    skpUnless[ALU=0], t←t#(r0);
```

```
rstkTest04:
```

```
    error;
    skpif[ALU=0];
    error;
```

```
    t←r0←rhigh1;
    t#(Q);
    skpUnless[ALU=0], t←t#(r0);
    error;
    skpif[ALU=0];
    error;
```

```
    t←r0←rscr;
    t#(Q);
    skpUnless[ALU=0], t←t#(r0);
```

```
rstkTest06:
```

```
    error;
    skpif[ALU=0];
    error;
```

```
    t←r0←rscr2;
    t#(Q);
    skpUnless[ALU=0], t←t#(r0);
    error;
    skpif[ALU=0];
    error;
    r0←Q;
```

```
*    FOR I IN [0..7] DO RBASE[1] ← RBASE[I]; (EXCEPT FOR ←RBASE[1])
```

```
rstkTest1:
```

```
    Q←r1;
```

```

    t←r1←r0;
    t#(Q);
    skipUnless[ALU=0], t←t#(r1);
    error;
    skipif[ALU=0];
    error;

    t←r1←rm1;
    t#(Q);
    skipUnless[ALU=0], t←t#(r1);
rstkTest12:
    error;
    skipif[ALU=0];
    error;

    t←r1←r01;
    t#(Q);
    skipUnless[ALU=0], t←t#(r1);
    error;
    skipif[ALU=0];
    error;

    t←r1←r10;
    t#(Q);
    skipUnless[ALU=0], t←t#(r1);
rstkTest14:
    error;
    skipif[ALU=0];
    error;

    t←r1←rhigh1;
    t#(Q);
    skipUnless[ALU=0], t←t#(r1);
    error;
    skipif[ALU=0];
    error;

    t←r1←rscr;
    t#(Q);
    skipUnless[ALU=0], t←t#(r1);
rstkTest16:
    error;
    skipif[ALU=0];
    error;

    t←r1←rscr2;
    t#(Q);
    skipUnless[ALU=0], t←t#(r1);
    error;
    skipif[ALU=0];
    error;
    r1←Q;

*   FOR I IN [0..7] DO RBASE[2] ← RBASE[I]; (EXCEPT FOR ←RBASE[2])
rstkTest2:
    Q←rm1;
    t←rm1←r0;
    t#(Q);
    skipUnless[ALU=0], t←t#(rm1);
    error;
    skipif[ALU=0];
    error;

    t←rm1←r1;
    t#(Q);
    skipUnless[ALU=0], t←t#(rm1);
rstkTest22:
    error;
    skipif[ALU=0];
    error;

    t←rm1←r01;
    t#(Q);
    skipUnless[ALU=0], t←t#(rm1);
    error;
    skipif[ALU=0];

```

```
    error;

    t←rm1←r10;
    t#(Q);
    skpUnless[ALU=0], t←t#(rm1);
rstkTest24:
    error;
    skpif[ALU=0];
    error;

    t←rm1←rhigh1;
    t#(Q);
    skpUnless[ALU=0], t←t#(rm1);
    error;
    skpif[ALU=0];
    error;

    t←rm1←rscr;
    t#(Q);
    skpUnless[ALU=0], t←t#(rm1);
rstkTest26:
    error;
    skpif[ALU=0];
    error;

    t←rm1←rscr2;
    t#(Q);
    skpUnless[ALU=0], t←t#(rm1);
    error;
    skpif[ALU=0];
    error;
    rm1←Q;

*   FOR I IN [0..7] DO RBASE[3] ← RBASE[I]; (EXCEPT FOR ←RBASE[3])
rstkTest3:
    Q←r01;
    t←r01←r0;
    t#(Q);
    skpUnless[ALU=0], t←t#(r01);
    error;
    skpif[ALU=0];
    error;

    t←r01←r1;
    t#(Q);
    skpUnless[ALU=0], t←t#(r01);
rstkTest32:
    error;
    skpif[ALU=0];
    error;

    t←r01←rm1;
    t#(Q);
    skpUnless[ALU=0], t←t#(r01);
    error;
    skpif[ALU=0];
    error;

    t←r01←r10;
    t#(Q);
    skpUnless[ALU=0], t←t#(r01);
rstkTest34:
    error;
    skpif[ALU=0];
    error;

    t←r01←rhigh1;
    t#(Q);
    skpUnless[ALU=0], t←t#(r01);
    error;
    skpif[ALU=0];
    error;

    t←r01←rscr;
    t#(Q);
    skpUnless[ALU=0], t←t#(r01);
```



```

rstkTest36:
    error;
    skipif[ALU=0];
    error;

    t←r01←rscr2;
    t#(Q);
    skipUnless[ALU=0], t←t#(r01);
    error;
    skipif[ALU=0];
    error;
    r01←Q;

*   FOR I IN [0..7] DO RBASE[4] ← RBASE[I]; (EXCEPT FOR ←RBASE[4])
rstkTest4:
    Q←r10;
    t←r10←r0;
    t#(Q);
    skipUnless[ALU=0], t←t#(r10);
    error;
    skipif[ALU=0];
    error;

    t←r10←r1;
    t#(Q);
    skipUnless[ALU=0], t←t#(r10);
rstkTest42:
    error;
    skipif[ALU=0];
    error;

    t←r10←r01;
    t#(Q);
    skipUnless[ALU=0], t←t#(r10);
    error;
    skipif[ALU=0];
    error;

    t←r10←rm1;
    t#(Q);
    skipUnless[ALU=0], t←t#(r10);
rstkTest44:
    error;
    skipif[ALU=0];
    error;

    t←r10←rhigh1;
    t#(Q);
    skipUnless[ALU=0], t←t#(r10);
    error;
    skipif[ALU=0];
    error;

    t←r10←rscr;
    t#(Q);
    skipUnless[ALU=0], t←t#(r10);
rstkTest46:
    error;
    skipif[ALU=0];
    error;

    t←r10←rscr2;
    t#(Q);
    skipUnless[ALU=0], t←t#(r10);
    error;
    skipif[ALU=0];
    error;
    r10←Q;

*   FOR I IN [0..7] DO RBASE[5] ← RBASE[I]; (EXCEPT FOR ←RBASE[5])
rstkTest5:
    Q←rhigh1;
    t←rhigh1←r0;
    t#(Q);
    skipUnless[ALU=0], t←t#(rhigh1);
    error;

```

```

        skipif[ALU=0];
        error;

        t←rhigh1←r1;
        t#(Q);
        skipUnless[ALU=0], t←t#(rhigh1);
rstkTest52:
        error;
        skipif[ALU=0];
        error;

        t←rhigh1←r01;
        t#(Q);
        skipUnless[ALU=0], t←t#(rhigh1);
        error;
        skipif[ALU=0];
        error;

        t←rhigh1←r10;
        t#(Q);
        skipUnless[ALU=0], t←t#(rhigh1);
rstkTest54:
        error;
        skipif[ALU=0];
        error;

        t←rhigh1←rm1;
        t#(Q);
        skipUnless[ALU=0], t←t#(rhigh1);
        error;
        skipif[ALU=0];
        error;

        t←rhigh1←rscr;
        t#(Q);
        skipUnless[ALU=0], t←t#(rhigh1);
rstkTest56:
        error;
        skipif[ALU=0];
        error;

        t←rhigh1←rscr2;
        t#(Q);
        skipUnless[ALU=0], t←t#(rhigh1);
        error;
        skipif[ALU=0];
        error;
        rhigh1←Q;

*   FOR I IN [0..7] DO RBASE[6] ← RBASE[I]; (EXCEPT FOR ←RBASE[6])
rstkTest6:
        Q←rscr;
        t←rscr←r0;
        t#(Q);
        skipUnless[ALU=0], t←t#(rscr);
        error;
        skipif[ALU=0];
        error;

        t←rscr←r1;
        t#(Q);
        skipUnless[ALU=0], t←t#(rscr);
rstkTest62:
        error;
        skipif[ALU=0];
        error;

        t←rscr←r01;
        t#(Q);
        skipUnless[ALU=0], t←t#(rscr);
        error;
        skipif[ALU=0];
        error;

        t←rscr←r10;
        t#(Q);

```

```
    skipUnless[ALU=0], t+t#(rscr);
rstkTest64:
    error;
    skipif[ALU=0];
    error;

    t←rscr←rhigh1;
    t#(0);
    skipUnless[ALU=0], t+t#(rscr);
    error;
    skipif[ALU=0];
    error;

    t←rscr←rm1;
    t#(0);
    skipUnless[ALU=0], t+t#(rscr);
rstkTest66:
    error;
    skipif[ALU=0];
    error;

    t←rscr←rscr2;
    t#(0);
    skipUnless[ALU=0], t+t#(rscr);
    error;
    skipif[ALU=0];
    error;
    rscr←0;

*   FOR I IN [0..7] DO RBASE[7] ← RBASE[I]; (EXCEPT FOR ←RBASE[7])
rstkTest7:
    Q←rscr2;
    t←rscr2+r0;
    t#(0);
    skipUnless[ALU=0], t+t#(rscr2);
    error;
    skipif[ALU=0];
    error;

    t←rscr2+r1;
    t#(0);
    skipUnless[ALU=0], t+t#(rscr2);
rstkTest72:
    error;
    skipif[ALU=0];
    error;

    t←rscr2+r01;
    t#(0);
    skipUnless[ALU=0], t+t#(rscr2);
    error;
    skipif[ALU=0];
    error;

    t←rscr2+r10;
    t#(0);
    skipUnless[ALU=0], t+t#(rscr2);
rstkTest74:
    error;
    skipif[ALU=0];
    error;

    t←rscr2←rhigh1;
    t#(0);
    skipUnless[ALU=0], t+t#(rscr2);
    error;
    skipif[ALU=0];
    error;

    t←rscr2←rscr;
    t#(0);
    skipUnless[ALU=0], t+t#(rscr2);
rstkTest76:
    error;
    skipif[ALU=0];
    error;
```

```
t←rscr2←rm1;
t#(Q);
skpUnless[ALU=0], t←t#(rscr2);
error;
skpif[ALU=0];
error;
rscr2←Q;

goto[afterKernel12];
```

%

Page Numbers: Yes First Page: 1

Heading:

kerne13.mc November 13, 1978 10:43 AM%

* INSERT[D1ALU.MC];

* TITLE[KERNEL3];

top level;

beginKerne13: noop;

%

TEST CONTENTS

SHCtestRW Read and write SHC

Rlsh test RM shiftlmask, lsh r[i]; 0<=i<=15

Tlsh test T shiftlmask, lsh t[i]; 0<=i<=15

Rrsh test RM shiftlmask, rsh r[i]; 0<=i<=15

Trsh test T shiftlmask, rsh t[i]; 0<=i<=15

TRlcyTest test T R lcy[i]; 0<=i<=15. (cycle 0,,1 and 177777,,177776)

RTlcyTest test R T lcy[i]; 0<=i<=15. (cycle 0,,1 and 177777,,177776)

rcy16, lcy16 test 16 bit cycles with selected bit values

cycleTest Test 32 bit cycles by generating possible r,t, count values

RWFtest test RF← and WF←

aluRSH test alu right shift (ie., H3 ← ALU rightshift 1)

aluRCY test alu right cycle (ie., H3 ← ALU rightCycle 1)

aluARSH test alu arithmetic right shift (ie., H3 ← ALU rightshift 1, sign preserved)

aluLSH test alu left shift

aluLCY test alu left cycle

aluSHTEST exhaustive test of alu shifts

%

* August 9, 1977 12:33 PM

```
%
    TEST SHC: READ AND WRITE

    FOR I IN[0..177777B] DO
        SHC+I;
        T+SHC XOR I;
        IF T#0 THEN ERROR;
    ENDLLOOP;
```

Note: ShC is a 16 bit register AND the upper three bits [0..2] are not used by the shifter!

```
%
SHCtestRW:
    rscr+r0;
SHCRWL:
    SHC+rscr;
    t ← SHC;
    t←(rscr)#(t);
    branch[.+2,ALU=0];
    error;
    rscr←(rscr)+1;
    loopUntil[ALU=0,SHCRWL];
```

```
% TEST THE SHIFTER
MAKE SURE THAT ALL SHIFTS WORK PROPER AMOUNT
MAKE SURE ALL MASKS WORK
MAKE SURE SHIFTS AND MASKS WORK TOGETHER
These tests work by left (or right) shifting bit15 (bit 0) 0 thru
15 times. rscr or rscr2 holds the expected value. The result is XOR'd
with the expected value and those bits are placed in T. If t #0 there
has been an error.
```

```
test order:
R shift left
T shift left
R shift right
T shift right
```

```
R T cycle left
T R cycle left
R T cycle right
T R cycle right
```

Note: The cycle tests are duplicated with the bits inverted (eg., bit15 {bit0} is zero and all other bits are one.

%

```
R1sh:
    t←r1;
    rscr←t;
    t←B15;
    rscr ← lsh[rscr,0];
    t←t#(rscr);
    skipif[ALU=0];
    error;
```

```
t←r1;
rscr←t;
t←B14;
rscr ← lsh[rscr,1];
t←t#(rscr);
skipif[ALU=0];
error;
```

```
R1sh2:
    t←r1;
    rscr←t;
    t←B13;
    rscr ← lsh[rscr,2];
    t←t#(rscr);
    skipif[ALU=0];
    error;
```

```
t←r1;
rscr←t;
```

```
t←B12;
rscr ← lsh[rscr,3];
t←t#(rscr);
skpif[ALU=0];
error;
```

R1sh4:

```
t←r1;
rscr←t;
t←B11;
rscr ← lsh[rscr,4];
t←t#(rscr);
skpif[ALU=0];
error;
```

```
t←r1;
rscr←t;
t←B10;
rscr ← lsh[rscr,5];
t←t#(rscr);
skpif[ALU=0];
error;
```

R1sh6:

```
t←r1;
rscr←t;
t←B9;
rscr ← lsh[rscr,6];
t←t#(rscr);
skpif[ALU=0];
error;
```

```
t←r1;
rscr←t;
t←B8;
rscr ← lsh[rscr,7];
t←t#(rscr);
skpif[ALU=0];
error;
```

R1sh8:

```
t←r1;
rscr←t;
t←B7;
rscr ← lsh[rscr,10];
t←t#(rscr);
skpif[ALU=0];
error;
```

```
t←r1;
rscr←t;
t←B6;
rscr ← lsh[rscr,11];
t←t#(rscr);
skpif[ALU=0];
error;
```

R1sh10:

```
t←r1;
rscr←t;
t←B5;
rscr ← lsh[rscr,12];
t←t#(rscr);
skpif[ALU=0];
error;
```

```
t←r1;
rscr←t;
t←B4;
rscr ← lsh[rscr,13];
t←t#(rscr);
skpif[ALU=0];
error;
```

R1sh12:

```
t←r1;
```

```
rscr←t;
t←B3;
rscr ← lsh[rscr,14];
t←t#(rscr);
skpif[ALU=0];
error;
```

```
t←r1;
rscr←t;
t←B2;
rscr ← lsh[rscr,15];
t←t#(rscr);
skpif[ALU=0];
error;
```

R1sh14:

```
t←r1;
rscr←t;
t←B1;
rscr ← lsh[rscr,16];
t←t#(rscr);
skpif[ALU=0];
error;
```

```
t←r1;
rscr←t;
t←RHIGH1;
rscr ← lsh[rscr,17];
t←t#(rscr);
skpif[ALU=0];
error;
```


* October 20, 1978 10:32 AM

T1sh:

```
t←rscr+B15;
noop;
t←lsh[t,0];
t←t#(rscr);
skipif[ALU=0];
error;
```

```
t←r1;
rscr←B14;
t←lsh[t,1];
t←t#(rscr);
skipif[ALU=0];
error;
```

T1sh2:

```
t←r1;
rscr←B13;
t←lsh[t,2];
t←t#(rscr);
skipif[ALU=0];
error;
```

```
t←r1;
rscr←B12;
t←lsh[t,3];
t←t#(rscr);
skipif[ALU=0];
error;
```

T1sh4:

```
t←r1;
rscr←B11;
t←lsh[t,4];
t←t#(rscr);
skipif[ALU=0];
error;
```

```
t←r1;
rscr←B10;
t←lsh[t,5];
t←t#(rscr);
skipif[ALU=0];
error;
```

T1sh6:

```
t←r1;
rscr←B9;
t←lsh[t,6];
t←t#(rscr);
skipif[ALU=0];
error;
```

```
t←r1;
rscr←B8;
t←lsh[t,7];
t←t#(rscr);
skipif[ALU=0];
error;
```

T1sh8:

```
t←r1;
rscr←B7;
t←lsh[t,10];
t←t#(rscr);
skipif[ALU=0];
error;
```

```
t←r1;
rscr←B6;
t←lsh[t,11];
t←t#(rscr);
skipif[ALU=0];
```

```
error;
```

```
T1sh10:
```

```
t←r1;  
rscr←B5;  
t←lsh[t,12];  
t←t#(rscr);  
skpif[ALU=0];  
error;
```

```
t←r1;  
rscr←B4;  
t←lsh[t,13];  
t←t#(rscr);  
skpif[ALU=0];  
error;
```

```
T1sh12:
```

```
t←r1;  
rscr←B3;  
t←lsh[t,14];  
t←t#(rscr);  
skpif[ALU=0];  
error;
```

```
t←r1;  
rscr←B2;  
t←lsh[t,15];  
t←t#(rscr);  
skpif[ALU=0];  
error;
```

```
T1sh14:
```

```
t←r1;  
rscr←B1;  
t←lsh[t,16];  
t←t#(rscr);  
skpif[ALU=0];  
error;
```

```
t←r1;  
rscr←RHIGH1;  
t←lsh[t,17];  
t←t#(rscr);  
skpif[ALU=0];  
error;
```

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%

KEEP 100000 IN Q FOR THESE TESTS !!!

%

Rrsh:

```
Q←RHIGH1;
GOTO[Rrsh1];    * Temporary EXPEDIENT
rscr←Q;
t←RHIGH1;
rscr ← rsh[rscr,0];
t←t#(rscr);
skpif[ALU=0];
error;
```

Rrsh1:

```
rscr←Q;
t←B1;
rscr ← rsh[rscr,1];
t←t#(rscr);
skpif[ALU=0];
error;
```

Rrsh2:

```
rscr←Q;
t←B2;
rscr ← rsh[rscr,2];
t←t#(rscr);
skpif[ALU=0];
error;
```

```
rscr←Q;
t←B3;
rscr ← rsh[rscr,3];
t←t#(rscr);
skpif[ALU=0];
error;
```

Rrsh4:

```
rscr←Q;
t←B4;
rscr ← rsh[rscr,4];
t←t#(rscr);
skpif[ALU=0];
error;
```

```
rscr←Q;
t←B5;
rscr ← rsh[rscr,5];
t←t#(rscr);
skpif[ALU=0];
error;
```

Rrsh6:

```
rscr←Q;
t←B6;
rscr ← rsh[rscr,6];
t←t#(rscr);
skpif[ALU=0];
error;
```

```
rscr←Q;
t←B7;
rscr ← rsh[rscr,7];
t←t#(rscr);
skpif[ALU=0];
error;
```

Rrsh8:

```
rscr←Q;
t←B8;
rscr ← rsh[rscr,10];
t←t#(rscr);
skpif[ALU=0];
error;
```

```
rscr←0;
t←B9;
rscr ← rsh[rscr,11];
t←t#(rscr);
skpif[ALU=0];
error;
```

Rrsh10:

```
rscr←0;
t←B10;
rscr ← rsh[rscr,12];
t←t#(rscr);
skpif[ALU=0];
error;
```

```
rscr←0;
t←B11;
rscr ← rsh[rscr,13];
t←t#(rscr);
skpif[ALU=0];
error;
```

Rrsh12:

```
rscr←0;
t←B12;
rscr ← rsh[rscr,14];
t←t#(rscr);
skpif[ALU=0];
error;
```

```
rscr←0;
t←B13;
rscr ← rsh[rscr,15];
t←t#(rscr);
skpif[ALU=0];
error;
```

Rrsh14:

```
rscr←0;
t←B14;
rscr ← rsh[rscr,16];
t←t#(rscr);
skpif[ALU=0];
error;
```

```
rscr←0;
t←B15;
rscr ← rsh[rscr,17];
t←t#(rscr);
skpif[ALU=0];
error;
```

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Trsh:

```
GOTO[Trshift1]; * Temporary EXPEDIENT
t←rscr←RHIGH1;
NOOP;
T←rsh[t,0];
t←t#(rscr);
skpif[ALU=0];
error;
```

Trshift1:

```
t←rhigh1;
rscr←B1;
t←rsh[t,1];
t←t#(rscr);
skpif[ALU=0];
error;
```

Trsh2:

```
t←rhigh1;
rscr←B2;
t←rsh[t,2];
t←t#(rscr);
skpif[ALU=0];
error;
```

```
t←rhigh1;
rscr←B3;
t←rsh[t,3];
t←t#(rscr);
skpif[ALU=0];
error;
```

Trsh4:

```
t←rhigh1;
rscr←B4;
t←rsh[t,4];
t←t#(rscr);
skpif[ALU=0];
error;
```

```
t←rhigh1;
rscr←B5;
t←rsh[t,5];
t←t#(rscr);
skpif[ALU=0];
error;
```

Trsh6:

```
t←rhigh1;
rscr←B6;
t←rsh[t,6];
t←t#(rscr);
skpif[ALU=0];
error;
```

```
t←rhigh1;
rscr←B7;
t←rsh[t,7];
t←t#(rscr);
skpif[ALU=0];
error;
```

Trsh8:

```
t←rhigh1;
rscr←B8;
t←rsh[t,10];
t←t#(rscr);
skpif[ALU=0];
error;
```

```
t←rhigh1;
rscr←B9;
t←rsh[t,11];
```

```
t←t#(rscr);
skpif[ALU=0];
error;
```

Trsh10:

```
t←rhigh1;
rscr←B10;
t←rsh[t,12];
t←t#(rscr);
skpif[ALU=0];
error;
```

```
t←rhigh1;
rscr←B11;
t←rsh[t,13];
t←t#(rscr);
skpif[ALU=0];
error;
```

Trsh12:

```
t←rhigh1;
rscr←B12;
t←rsh[t,14];
t←t#(rscr);
skpif[ALU=0];
error;
```

```
t←rhigh1;
rscr←B13;
t←rsh[t,15];
t←t#(rscr);
skpif[ALU=0];
error;
```

Trsh14:

```
t←rhigh1;
rscr←B14;
t←rsh[t,16];
t←t#(rscr);
skpif[ALU=0];
error;
```

```
t←rhigh1;
rscr←B15;
t←rsh[t,17];
t←t#(rscr);
skpif[ALU=0];
error;
```

* October 20, 1978 10:14 AM

%

These tests work by cycling by 0, 1, ...17B. The predicted result is kept in RSCR2 and the actual result XOR's w/ predicted result is kept in T. Note that each test is done twice: once w/ one "1" bit and all the rest "0" bits, and once w/ one "0" bit and all the rest "1" bits.

FOR THESE TESTS WE WILL REDEFINE R01 TO BE RM2 (-2)!

%

TR1cyTest:

RM[rm2,IP[R01]];
 rm2 ← CM2;

t←r0;
 rscr2←B15; * RSCR2 ← PREDICTED RESULT
 t←lcy[t,r1,0];
 t←t#(rscr2);
 skipif[alu=0];
 error;
 t←rm1;
 rscr2←NB15; * RSCR2 ← PREDICTED RESULT
 t←lcy[t,rm2,0];
 t←t#(rscr2);
 skipif[alu=0];
 error;

t←r0;
 rscr2←B14; * RSCR2 ← PREDICTED RESULT
 t←lcy[t,r1,1];
 t←t#(rscr2);
 skipif[alu=0];
 error;
 t←rm1;
 rscr2←NB14; * RSCR2 ← PREDICTED RESULT
 t←lcy[t,rm2,1];
 t←t#(rscr2);
 skipif[alu=0];
 error;

TR1cy2:

t←r0;
 rscr2←B13; * RSCR2 ← PREDICTED RESULT
 t←lcy[t,r1,2];
 t←t#(rscr2);
 skipif[alu=0];
 error;
 t←rm1;
 rscr2←NB13; * RSCR2 ← PREDICTED RESULT
 t←lcy[t,rm2,2];
 t←t#(rscr2);
 skipif[alu=0];
 error;

t←r0;
 rscr2←B12; * RSCR2 ← PREDICTED RESULT
 t←lcy[t,r1,3];
 t←t#(rscr2);
 skipif[alu=0];
 error;
 t←rm1;
 rscr2←NB12; * RSCR2 ← PREDICTED RESULT
 t←lcy[t,rm2,3];
 t←t#(rscr2);
 skipif[alu=0];
 error;

TR1cy4:

t←r0;
 rscr2←B11; * RSCR2 ← PREDICTED RESULT
 t←lcy[t,r1,4];
 t←t#(rscr2);
 skipif[alu=0];
 error;
 t←rm1;
 rscr2←NB11; * RSCR2 ← PREDICTED RESULT

```

t←lcy[t,rm2,4];
t←t#(rscr2);
skpif[alu=0];
error;

```

```

t←r0;
rscr2←B10;      * RSCR2 ← PREDICTED RESULT
t←lcy[t,r1,5];
t←t#(rscr2);
skpif[alu=0];
error;
t←rm1;
rscr2←NB10;     * RSCR2 ← PREDICTED RESULT
t←lcy[t,rm2,5];
t←t#(rscr2);
skpif[alu=0];
error;

```

TR1cy6:

```

t←r0;
rscr2←B9;       * RSCR2 ← PREDICTED RESULT
t←lcy[t,r1,6];
t←t#(rscr2);
skpif[alu=0];
error;
t←rm1;
rscr2←NB9;     * RSCR2 ← PREDICTED RESULT
t←lcy[t,rm2,6];
t←t#(rscr2);
skpif[alu=0];
error;

```

```

t←r0;
rscr2←B8;       * RSCR2 ← PREDICTED RESULT
t←lcy[t,r1,7];
t←t#(rscr2);
skpif[alu=0];
error;
t←rm1;
rscr2←NB8;     * RSCR2 ← PREDICTED RESULT
t←lcy[t,rm2,7];
t←t#(rscr2);
skpif[alu=0];
error;

```

TR1cy8:

```

t←r0;
rscr2←B7;       * RSCR2 ← PREDICTED RESULT
t←lcy[t,r1,10];
t←t#(rscr2);
skpif[alu=0];
error;
t←rm1;
rscr2←NB7;     * RSCR2 ← PREDICTED RESULT
t←lcy[t,rm2,10];
t←t#(rscr2);
skpif[alu=0];
error;

```

```

t←r0;
rscr2←B6;       * RSCR2 ← PREDICTED RESULT
t←lcy[t,r1,11];
t←t#(rscr2);
skpif[alu=0];
error;
t←rm1;
rscr2←NB6;     * RSCR2 ← PREDICTED RESULT
t←lcy[t,rm2,11];
t←t#(rscr2);
skpif[alu=0];
error;

```

TR1cy10:

```

t←r0;
rscr2←B5;       * RSCR2 ← PREDICTED RESULT
t←lcy[t,r1,12];

```



```

t←t#(rscr2);
skpif[alu=0];
error;
t←rm1;
rscr2←NB5;      * RSCR2 ← PREDICTED RESULT
t←lcy[t,rm2,12];
t←t#(rscr2);
skpif[alu=0];
error;

```

```

t←r0;
rscr2←B4;      * RSCR2 ← PREDICTED RESULT
t←lcy[t,r1,13];
t←t#(rscr2);
skpif[alu=0];
error;
t←rm1;
rscr2←NB4;     * RSCR2 ← PREDICTED RESULT
t←lcy[t,rm2,13];
t←t#(rscr2);
skpif[alu=0];
error;

```

TR1cy12:

```

t←r0;
rscr2←B3;      * RSCR2 ← PREDICTED RESULT
t←lcy[t,r1,14];
t←t#(rscr2);
skpif[alu=0];
error;
t←rm1;
rscr2←NB3;     * RSCR2 ← PREDICTED RESULT
t←lcy[t,rm2,14];
t←t#(rscr2);
skpif[alu=0];
error;

```

```

t←r0;
rscr2←B2;      * RSCR2 ← PREDICTED RESULT
t←lcy[t,r1,15];
t←t#(rscr2);
skpif[alu=0];
error;
t←rm1;
rscr2←NB2;     * RSCR2 ← PREDICTED RESULT
t←lcy[t,rm2,15];
t←t#(rscr2);
skpif[alu=0];
error;

```

TR1cy14:

```

t←r0;
rscr2←B1;      * RSCR2 ← PREDICTED RESULT
t←lcy[t,r1,16];
t←t#(rscr2);
skpif[alu=0];
error;
t←rm1;
rscr2←NB1;     * RSCR2 ← PREDICTED RESULT
t←lcy[t,rm2,16];
t←t#(rscr2);
skpif[alu=0];
error;

```

```

t←r0;
rscr2←B0;      * RSCR2 ← PREDICTED RESULT
t←lcy[t,r1,17];
t←t#(rscr2);
skpif[alu=0];
error;
t←rm1;
rscr2←NB0;     * RSCR2 ← PREDICTED RESULT
t←lcy[t,rm2,17];
t←t#(rscr2);
skpif[alu=0];
error;

```



```
*
*      October 20, 1978  10:06 AM
RT1cyTest:      *RSCR2 HOLDS THE PREDICTED RESULT, T HOLDS ACTUAL RESULT
```

```
  t+r1;      * rscr2, T ← [0,1] LCY[0]
  rscr2←B15;      * rscr2 ← PREDICTED RESULT
  t+lcy[r0,t,0];
  t+t#(rscr2);
  skipif[alu=0];
  error;
  t+rm2;
  rscr2←NB15;      * rscr2 ← PREDICTED RESULT
  t+lcy[rm1,t,0];
  t+t#(rscr2);
  skipif[alu=0];
  error;
```

```
  t+r1;
  rscr2←B14;      * rscr2 ← PREDICTED RESULT
  t+lcy[r0,t,1];
  t+t#(rscr2);
  skipif[alu=0];
  error;
  t+rm2;
  rscr2←NB14;      * rscr2 ← PREDICTED RESULT
  t+lcy[rm1,t,1];
  t+t#(rscr2);
  skipif[alu=0];
  error;
```

```
RT1cy2:
```

```
  t+r1;
  rscr2←B13;      * rscr2 ← PREDICTED RESULT
  t+lcy[r0,t,2];
  t+t#(rscr2);
  skipif[alu=0];
  error;
  t+rm2;
  rscr2←NB13;      * rscr2 ← PREDICTED RESULT
  t+lcy[rm1,t,2];
  t+t#(rscr2);
  skipif[alu=0];
  error;
```

```
  t+r1;
  rscr2←B12;      * rscr2 ← PREDICTED RESULT
  t+lcy[r0,t,3];
  t+t#(rscr2);
  skipif[alu=0];
  error;
  t+rm2;
  rscr2←NB12;      * rscr2 ← PREDICTED RESULT
  t+lcy[rm1,t,3];
  t+t#(rscr2);
  skipif[alu=0];
  error;
```

```
RT1cy4:
```

```
  t+r1;
  rscr2←B11;      * rscr2 ← PREDICTED RESULT
  t+lcy[r0,t,4];
  t+t#(rscr2);
  skipif[alu=0];
  error;
  t+rm2;
  rscr2←NB11;      * rscr2 ← PREDICTED RESULT
  t+lcy[rm1,t,4];
  t+t#(rscr2);
  skipif[alu=0];
  error;
```

```
  t+r1;
  rscr2←B10;      * rscr2 ← PREDICTED RESULT
  t+lcy[r0,t,5];
  t+t#(rscr2);
```

```
skpif[alu=0];
error;
t←rm2;
rscr2←NB10;      * rscr2 ← PREDICTED RESULT
t←lcy[rm1,t,5];
t←t#(rscr2);
skpif[alu=0];
error;
```

RT1cy6:

```
t←r1;
rscr2←B9;      * rscr2 ← PREDICTED RESULT
t←lcy[r0,t,6];
t←t#(rscr2);
skpif[alu=0];
error;
t←rm2;
rscr2←NB9;      * rscr2 ← PREDICTED RESULT
t←lcy[rm1,t,6];
t←t#(rscr2);
skpif[alu=0];
error;
```

```
t←r1;
rscr2←B8;      * rscr2 ← PREDICTED RESULT
t←lcy[r0,t,7];
t←t#(rscr2);
skpif[alu=0];
error;
t←rm2;
rscr2←NB8;      * rscr2 ← PREDICTED RESULT
t←lcy[rm1,t,7];
t←t#(rscr2);
skpif[alu=0];
error;
```

RT1cy8:

```
t←r1;
rscr2←B7;      * rscr2 ← PREDICTED RESULT
t←lcy[r0,t,10];
t←t#(rscr2);
skpif[alu=0];
error;
t←rm2;
rscr2←NB7;      * rscr2 ← PREDICTED RESULT
t←lcy[rm1,t,10];
t←t#(rscr2);
skpif[alu=0];
error;
```

```
t←r1;
rscr2←B6;      * rscr2 ← PREDICTED RESULT
t←lcy[r0,t,11];
t←t#(rscr2);
skpif[alu=0];
error;
t←rm2;
rscr2←NB6;      * rscr2 ← PREDICTED RESULT
t←lcy[rm1,t,11];
t←t#(rscr2);
skpif[alu=0];
error;
```

RT1cy10:

```
t←r1;
rscr2←B5;      * rscr2 ← PREDICTED RESULT
t←lcy[r0,t,12];
t←t#(rscr2);
skpif[alu=0];
error;
t←rm2;
rscr2←NB5;      * rscr2 ← PREDICTED RESULT
t←lcy[rm1,t,12];
t←t#(rscr2);
skpif[alu=0];
error;
```

```

t←r1;
rscr2←B4;      * rscr2 ← PREDICTED RESULT
t←lcy[r0,t,13];
t←t#(rscr2);
skpif[alu=0];
error;
t←rm2;
rscr2←NB4;     * rscr2 ← PREDICTED RESULT
t←lcy[rm1,t,13];
t←t#(rscr2);
skpif[alu=0];
error;

```

RTlcy12:

```

t←r1;
rscr2←B3;      * rscr2 ← PREDICTED RESULT
t←lcy[r0,t,14];
t←t#(rscr2);
skpif[alu=0];
error;
t←rm2;
RSCR2←NB3;     * RSCR2 ← PREDICTED RESULT
t←lcy[rm1,t,14];
t←t#(rscr2);
skpif[alu=0];
error;

```

```

t←r1;
RSCR2←B2;      * RSCR2 ← PREDICTED RESULT
t←lcy[r0,t,15];
t←t#(rscr2);
skpif[alu=0];
error;
t←rm2;
RSCR2←NB2;     * RSCR2 ← PREDICTED RESULT
t←lcy[rm1,t,15];
t←t#(rscr2);
skpif[alu=0];
error;

```

RTlcy14:

```

t←r1;
RSCR2←B1;      * RSCR2 ← PREDICTED RESULT
t←lcy[r0,t,16];
t←t#(rscr2);
skpif[alu=0];
error;
t←rm2;
RSCR2←NB1;     * RSCR2 ← PREDICTED RESULT
t←lcy[rm1,t,16];
t←t#(rscr2);
skpif[alu=0];
error;

```

```

t←r1;
RSCR2←B0;      * RSCR2 ← PREDICTED RESULT
t←lcy[r0,t,17];
t←t#(rscr2);
skpif[alu=0];
error;
t←rm2;
RSCR2←NB0;     * RSCR2 ← PREDICTED RESULT
t←lcy[rm1,t,17];
t←t#(rscr2);
skpif[alu=0];
error;

```

RTlcyDone:

```

r01 ← NOT(r10); * REDEFINE r01 !!!!!!!

```

* November 3, 1978 6:43 PM

%

rcy16, lcy16 Test the 16 bit cycles with selected
bit values. This is not an exhaustive test.

%

rcyTest: * test 16 bit right cycle

t ← rcy[r01, r01, 1];

t # (r10);

skipif[ALU=0];

rcy16Err1: * r10 rcy 1 should be r01

error;

t ← r01;

t ← rcy[t, t, 1]; * try it again from t

t # (r10);

skipif[ALU=0];

rcy16Err2: * r10 rcy 1 should be r01. (done from

error; * t this time)

t ← rcy[r1, r1, 1];

t # (rhigh1);

skipif[ALU=0];

rcy16Err3: * 1 rcy 1 should be 100000B

error;

t ← r1;

t ← rcy[t, t, 1];

t # (rhigh1);

skipif[ALU=0];

rcy16Err4: * 1 rcy 1 should be 100000B. (done from t

error; * time).

t ← rcy[r10, r10, 1];

t # (r01);

skipif[ALU=0];

rcy16Err5: * r10 rcy 1 should be r01

error;

t ← r10;

t ← rcy[t, t, 1]; * t ← r10 rcy 1

t # (r01);

skipif[ALU=0];

rcy16Err6: * r10 rcy 1 should be r01

error; * done from t this time.

t ← rcy[r01, r01, 2];

t # (r01);

skipif[ALU=0];

rcy16Err7: * r01 rcy 2 should be r01

error;

t ← rcy[r01, r01, 3];

t # (r10);

skipif[ALU=0];

rcy16Err8: * r01 rcy 3 should be r10

error;

t ← rcy[r01, r01, 10];

t # (r01);

skipif[ALU=0];

rcy16Err9: * r01 rcy 10 should be r01

error;

lcyTest:

t ← lcy[r01, r01, 1];

t # (r10);

skipif[ALU=0];

lcy16Err1: * r01 lcy 1 should be r10

error;

t ← lcy[rhigh1, rhigh1, 1];

t # (r1);

skipif[ALU=0];

lcy16Err2: * 100000B lcyd 1 should be 1

error;

```
t ← lcy[r1, r1, 1];
t#(2c);
skpif[ALU=0];
lcy16Err3: * 1 lcy 1 should be 2
error;

t ← lcy[r10, r10, 1]; * t ← r10 lcy 1
t # (r01);
skpif[ALU=0];
lcy16Err4: * r10 lcy 1 should be r01
error;

t ← lcy[r10, r10, 2];
t # (r10);
skpif[ALU=0];
lcy16Err5: * r10 lcy 2 should be r10
error;

t ← lcy[r10, r10, 3];
t # (r01);
skpif[ALU=0];
lcy16Err6: * r10 lcy 3 should be r01
error;

t ← lcy[r10, r10, 4];
t # (r10);
skpif[ALU=0];
lcy16Err7: * r10 lcy 4 should be r10
error;

t ← lcy[r10, r10, 10];
t # (r10);
skpif[ALU=0];
lcy16Err8: * r10 lcy 10 should be r10
error;
```

* November 13, 1978 10:40 AM

%

cycleTest Test the cycle machinery by generating all possible values for the r, t, and count fields in Shc (6 bits). The 16-bit data patterns must test all possible starting positions (ie., bit 0, bit 1, ...). Furthermore, set the "other word" to all 1s when single one-bits are being tested, and set it to all zeros when single zero bits are being tested. Since we test cycling, we don't set any of the mask fields in Shc.

```
SHC: TYPE = MACHINE DEPENDENT RECORD[
ShifterIgnores: IN[0..3],      -- bits 0, 1
a: IN [0..1],    -- bit 2, shA select. 1 ==> "select T"
b: IN [0..1],    -- bit 3, shB select. 1 ==> "select T"
Count: IN [0..17B],  -- bits 4:7, shift count
RMask: IN [0..17B],  -- bits 8:11
LMask: IN [0..17B],  -- bits 12:15
];
shcVals: IN [0..77B];  -- iterate thru all possible counts, sha, shb
```

```
FOR pats IN NPats DO
FOR shcVal In SHCVals DO
Shc.a ← shcVals AND 40B;
Shc.b ← shcVals AND 20B;
Shc.count ← shcVals AND 17B;
r ← getPattern[pats];
t ← IF numberOfZeroBitsGr1[r] THEN -1 ELSE 0;
result ← doShift[];
expected ← simulateCycle[t,r,shcVal];
IF result ≠ expected THEN SIGNAL BadShift[result, expected, shcVals];
ENDLOOP;      -- end of shcVals loop
ENDLOOP;      -- end of pats loop
simulateCycle: PROCEDURE[t, r: WORD, shcVal: SHCVals] RETURNS [expected: WORD] =
BEGIN
tCycle: CARDINAL = 60B; -- 2 highest bits in shcVal are 1
rCycle: CARDINAL = 0;  -- 2 highest bits in shcVal are 0
trCycle: CARDINAL = 40B; -- highest bit in shcVal is 1
rtCycle: CARDINAL = 20BB; -- 2nd highest bit in ShcVal is 1
shAB ← shcVal AND 60B;
SELECT shAB INTO
tCycle=> BEGIN
left ← right<t; END,
rCycle=> BEGIN
left ← right ← r; END,
trCycle=> BEGIN
left ← t, right ← r; END,
rtCycle=> BEGIN
left ← r;
right ← t; END,
END;
shiftCount ← shcVal AND 17B;
saveMask ← SELECT shiftCount INTO
1=>100000;
2=>140000B;
3=>160000B;
4=>170000B;
5=>174000B;
6=>176000B;
7=>177000B;
8=>177400B;
9=>177600B;
10=>177700B;
11=>177740B;
12=>177760B;
13=>177770B;
14=>177774B;
15=>177776B;
0=>177777B,
END;
savedValue ← left AND saveMask;
right ← LeftShift[right, shiftCount];
savedValue ← RightShift[right, 16-shiftCount];
result ← savedValue OR left;
END;
numberOfZeroBitsGr1: PROCEDURE[ x: WORD] RETURNS[result: BOOLEAN] =
BEGIN
count ← 0;
```



```
FOR i IN [0..15] DO
IF (x AND 1) =0 THEN count ← count + 1;
x ← RightShift[x,1];
ENDLOOP;
result ← IF count >1 THEN TRUE ELSE FALSE;
END;
```

```
%
```

```
%
November 17, 1978 2:39 PM
```

```
TEST RF AND WF
```

```
ShC: TYPE = MACHINE DEPENDENT RECORD [
  IGNORE: TYPE = [0..7B]
  SHIFTCOUNT: TYPE = [0..37B]
  RMASK: TYPE = [0..17B]
  LMASK: TYPE = [0..17B]
]
```

```
MesaDescriptor: TYPE = MACHINE DEPENDENT RECORD[ -- this is the value stored w/ rf←, wf←
  IGNORE: TYPE = [0..377B] -- IGNORE FIRST BYTE
  POS: TYPE = [0..17B] -- RIGHT SHIFT OF POS WILL RIGHT JUSTIFY THE FIELD
  SIZE: TYPE = [0..17B] -- LENGTH OF FIELD IN BITS
]
```

```
THIS TEST PROCEEDS BY WRITING ShC W/ ALL POSSIBLE RF AND WF VALUES. THEN
ShC IS READ AND CHECKED TO MAKE SURE THAT IT WAS LOADED PROPERLY.
```

```
FOR I IN [0..377B] DO
  RF←I;
  RSCR←SHC;
  SIZE ← I AND 17B;
  POS ← BITSHIFT[I,-4] AND 17B;
  IF RSCR.LMASK # (16-SIZE-1) THEN ERROR; -- BAD LMASK
  IF RSCR.RMASK # 0 THEN ERROR; -- BAD RMASK
  IF RSCR.SHIFTCOUNT # (16+pos+size+1) THEN ERROR; -- BAD SHIFT COUNT
  (Actually this computation isn't quite right.
  * let count = 16+pos+size+1. realCount ← (count and 17b).
  * IF (realCount and 17b) #0 then realCount ← realCount OR 20B. This funny computation
  * accommodates hardware limitations associated w/ carry across boards.
```

```
  -- now test wf
  WF←I;
  RSCR←SHC;
  IF RSCR.RMASK # (16-POS-SIZE-1) THEN ERROR; -- BAD RMASK
  IF RSCR.LMASK NE POS THEN ERROR; -- BAD LMASK
  IF RSCR.SHIFTCOUNT # (16-pos-size-1) THEN ERROR; -- BAD SHIFT COUNT
```

```
ENDLOOP;
```

```
%
```

```
RM[r4BitMsk,IP[R01]];
r4BitMsk ← 17C; * RENAME R01 AS r4BitMsk !!!
RM[lastShC, IP[RSCR]]; * RENAME RSCR AS lastShC
```

```
RFWFtest:
  Q←R0; * Q WILL HOLD THE INDEX VARIABLE
  t←377C;
  CNT←t; * LOOP LIMIT
```

```
RFTESTL:
  t←Q;
  RF←t;
  lastShC←SHC;
```

```
* CHECK LMASK
  T← (r4BitMsk)AND (Q); * COMPUTE LMASK (= 16-SIZE-1) FROM INDEX VAR
  rscr2←t; * rscr2 ← size
  t←17C; * 16-1
  rscr2 ← t - (rscr2); * rscr2 ← expected Lmask = 16-size-1
  t ← (lastShC) and (17c); * t ← Lmask from ShC
  t # (rscr2);
  branch[.+2, ALU=0];
```

```
RFLMASK: * t = Lmask from ShC, rscr2 = expected Lmask
  error; * LMASK FIELD WRONG IN ShC
```

```
* CHECK RMASK
  t←(lastShC) and (360c); * t = isolated Rmask field of ShC
  skipif[ALU=0];
```

```
RFRMASK:
  error; * RMASK FIELD NOT 0
```

* CHECK SHIFT COUNT = 16+pos+size+1 (Actually this computation isn't quite right.
* let count = 16+pos+size+1. realCount ← (count and 17b).
* IF (realCount and 17b) #0 then realCount ← realCount OR 20B. This funny computation
* accommodates hardware limitations associated w/ carry across boards.

```
rscr2 ← (Q);  
t ← r4BitMsk;  
rscr2 ← rsh[rscr2,4];  
rscr2 ← t AND (rscr2); * rscr2 = POS  
t ← 21c; * 16 + 1  
t ← t + (rscr2); * 16 + 1 + pos  
rscr2 ← q;  
rscr2 ← (rscr2) and (17c); * isolate size  
rscr2 ← t + (rscr2); * rscr2 ← 16 + 1 + pos + size  
rscr2 ← (rscr2) and (17c); * isolate to 17 bits  
skpif[alu=0]; * see if bit 0 of count is one  
rscr2 ← (rscr2) or (20c); * set bit0 of count if count[1:4]#0  
rscr2 ← (rscr2) and (17c); * isolate result to 5 bits
```

```
t ← rsh[LastShC, 10];
```

RFSHIFTC:

```
t#(rscr2); * t=value from LastShC, rscr2 = computed value  
skpif[ALU=0];  
error; * BAD SHIFT COUNT
```

* June 28, 1978 5:06 PM
 * NOW TEST WF

WFTTEST:

t←0;
 WF←t;
 lastShC←SHC;

* CHECK LMASK: COMPUTE pos
 rscr2 ← q;
 noop;
 t ← rsh[rscr2,4];
 rscr2 ← t and (r4BitMsk); * isolate pos bits in rscr2

t←lastShC;
 t ← t AND (r4BitMsk); * T← LMASK

WFLMASK:

t#(rscr2); * t=LastShC's Lmask, rscr2 = computed value
 branch[.+2, ALU=0];
 error; * SHC'S LMASK # pos

* CHECK THAT RMASK = 16 - pos - size - 1

rscr2 ← q;
 t ← (r4BitMsk) and (q); * isolate size in t
 rscr2 ← rsh[rscr2,4]; * rscr2 ← pos
 rscr2 ← (rscr2) + t; * rscr2 ← pos + size
 t ← 17c; * t ← 16 - 1
 rscr2 ← t - (rscr2); * rscr2 ← 16 - pos - size - 1
 rscr2 ← (rscr2) and (17c); * isolate to 17 bits
 t ← rsh[lastShC,4]; * t = ShC's shift count
 t ← t and (r4BitMsk);

WFRMASK:

t#(rscr2); * t = ShC's shift count
 skipif[ALU=0]; * rscr2 = 16-pos-size-1
 error; * RMASK NE (16-POS-SIZE-1)

* CHECK SHIFT COUNT=16-pos-size-1

t ← rsh[lastShC,10]; * put ShC's shift count into t
 t ← t and (37c);
 t#(rscr2); * t = ShC's shift count
 skipif[ALU=0]; * rscr2 = 16-pos-size-1, as computed above

WFSHIFTC:

* for the Rmask check
 error; * SHC'S SHIFTCOUNT # POS

rscr2←(R1) + (0);
 loopUntil[CNT=0&-1,RFTESTL],Q←(rscr2);

RFXITL:

R01 ← NOT(R10); * RESET R01 !!!!!

* October 20, 1978 10:23 AM

%

TEST ALU SHIFT OPERATIONS

%

* TEST RESULT ← ALU RSH 1 (RESULT[0] ← 0)

aluRSH:

```
r01 ← not(r10); * RESET r01 !!!! INCASE WE SKIPPED RFXITL
t←(PD←(r1))rsh 1;
t←t#(r0);
skpif[ALU=0];
error; * 1 RSH[1] SHOULD BE 0
```

```
t←(PD←(rm1))rsh 1;
rscr←77777C;
t←t#(rscr);
skpif[ALU=0];
error; * -1 RSH[1] SHOULD BE 77777B
```

```
t←(PD←r10)rsh 1;
t←t#(r01);
skpif[ALU=0];
error; * (ALTERNATING 10) rsh 1 SHOULD BE (ALT. 01)
```

* TEST RESULT ← ALU RCY[1] (RESULT[0]←ALU[15])

aluRCY:

```
t←(PD←rm1)rcy 1;
t←t#(rm1);
skpif[ALU=0];
error; * -1 RCY[1] SHOULD BE -1
```

```
t←(PD←r0)rcy 1;
t←t#(r0);
skpif[ALU=0];
error; * 0 RCY[1] SHOULD BE 0
```

```
t←(PD←r10)rcy 1;
t←t#(r01);
skpif[ALU=0];
error; * (ALTERNATING 10) RCY[1] SHOULD BE (ALT 01)
```

```
t←(PD←r01)rcy 1;
t←t#(r10);
skpif[ALU=0];
error; * (ALT 01) RCY[1] SHOULD BE (ALT 10)
```

* REST RESULT ← ALU Arsh 1 (RESULT[0] ← ALU[0]) (SIGN PRESERVING)

aluARSH:

```
t←(PD←rm1)Arsh 1;
t←t#(rm1);
skpif[ALU=0];
error; * -1 ARSH SHOULD BE -1
```

```
t←(PD←r0)Arsh 1;
t←t#(r0);
skpif[ALU=0];
error; * 0 ARSH SHOULD BE 0
```

```
t←(PD←rhigh1)Arsh 1;
rscr←140000C;
t←t#(rscr);
skpif[ALU=0];
error; * 100000 ARSH SHOULD BE 140000
```

```
t←rhigh1;
rscr←t+(r01);
t←(PD←r10)Arsh 1;
t←t#(rscr);
skpif[ALU=0];
error; * (ALT. 10) ARSH SHOULD BE(ALT. 01+100000)
```

* TEST RESULT ← ALU lsh 1

aluLSH:

```
t←(PD←rhigh1)lsh 1;
t←t#(r0);
```

```
skpif[ALU=0];
error; * 100000 LSH SHOULD BE 0
```

```
aluLSHb:
```

```
t←(PD←r1)lsh 1;
rscr←(r1)+(r1);
t←t#(rscr);
skpif[ALU=0];
error; * 1 LSH SHOULD BE 2
```

```
aluLSHc:
```

```
t←(PD←r01)lsh 1;
t←t#(r10);
skpif[ALU=0];
error; * (ALT. 01) LSH SHOULD BE (ALT. 10)
```

```
aluLSHd:
```

```
t←(PD←rm1)lsh 1;
rscr←CM2;
t←t#(rscr);
skpif[ALU=0];
error; * -1 LSH SHOULD BE -2
```

```
* TEST RESULT ← ALU LCY1
```

```
aluLCY:
```

```
t←(PD←rm1)lcy 1;
t←t#(rm1);
skpif[ALU=0];
error; * -1 LCY SHOULD BE -1
```

```
aluLCYb:
```

```
t←(PD←r10)lcy 1;
t←t#(r01);
skpif[ALU=0];
error; * (ALT. 10) LCY SHOULD BE (ALT. 01)
```

```
aluLCYc:
```

```
t←(PD←r01)lcy 1;
t←t#(r10);
skpif[ALU=0];
error; * (ALT. 01) LCY SHOULD BE (ALT. 10)
```

```
aluLCYd:
```

```
t←(PD←r0)lcy 1;
t←t#(r0);
skpif[ALU=0];
error; * 0 LCY SHOULD BE 0
```

```
aluLCYe:
```

```
t←(PD←r1)lcy 1;
rscr←(r1)+(r1);
t←t#(rscr);
skpif[ALU=0];
error; * 1 LCY SHOULD BE 2
```

* October 20, 1978 10:24 AM

%

EXHAUSTIVE TEST OF ALU SHIFT FUNCTIONS

FOR Q IN[0..177777B] DO

rscr2←Q rsh 1;
t←predictedRSH[I];
IF t←(T XOR rscr2) THEN ERROR;

rscr2←Q)rcy 1;
t←predictedRCY[I];
IF t←(T XOR rscr2) THEN ERROR;

rscr2←Q Arsh 1;
t←predictedARSH[I];
IF t←(T XOR rscr2) THEN ERROR;

rscr2←Q lsh 1;
t←predictedLSH[I];
IF t←(T XOR rscr2) THEN ERROR;

rscr2←Q rsh 1;
t←predictedRSH[I];
IF t←(T XOR rscr2) THEN ERROR;

rscr2←Q lcy 1;
t←predictedLCY[I];
IF t←(T XOR rscr2) THEN ERROR;

ENDLOOP;

%

aluSHTEST:

Q←r0; * USE Q AS LOOP VARIABLE

aluSHL: * TOP OF LOOP

* RSH TEST

rscr←Q;
t←(PD←rscr)rsh 1;
rscr←rsh[rscr,1];
t←t#(rscr);
branch[.+2,ALU=0];

RSHER:

error; * PREDICTED RESULT DIFFERENT FROM REAL ONE

* RCY TEST

rscr←Q; * Q IS LOOP VARIABLE
t←(PD←rscr)rcy 1;
rscr2←t; * REAL RESULT IN rscr2
t←rsh[rscr,1];
skpif[R EVEN], (PD←rscr); * ADD HIGH BIT IF NECESSARY
t←t+(rhigh1); * PREDICTED RESULT IN T
t←t#(rscr2);
branch[.+2,ALU=0];

RCYER:

error; * T ← (PREDICTED T) XOR rscr2

* ARSH TEST

rscr←Q; * Q IS LOOP VARIABLE
t←(PD←rscr)Arsh 1;
rscr2 ← T; * rscr2 = ACTUAL RESULT
t←rsh[rscr,1];
PD←Q; * ADD SIGN BIT IF REQUIRED
skpUnless[ALU<0]; * ADD SIGN BIT IF REQUIRED
t←t+(rhigh1);

t←t#(rscr2);
skpif[ALU=0];

ARSHER:

error; * t←(PREDICTED T) XOR rscr2

* LSH TEST

rscr←Q; * Q IS LOOP VARIABLE
rscr2←(PD←Q)lsh 1; * rscr2 = ACTUAL RESULT
t←lsh[rscr,1];

```
    t←t#(rscr2);
    skipif[ALU=0];
LSHER:
    error; * t← (PREDICTED T) XOR rscr2
* LCY TEST
    t←rscr ← Q; * Q IS LOOP VARIABLE
    rscr2 ← (PD+t)lcy 1; * rscr2 = ACTUAL RESULT
    t←lsh[rscr,1]; * t← PREDICTED RESULT
    rscr←(rscr);
    skipUnless[ALU<0];
    t←t+(r1); * t← T+ 1 FOR CYCLED BIT 0
    t←t#(rscr2);
    skipif[ALU=0];
LCYER:
    error; * T ← (PREDICTED T) XOR rscr2

    t←(r1)+(Q);
    dblBranch[.+1,aluSHL,ALU=0],Q←t;
goto[afterkerne13];
```



```
%
Page Numbers: Yes First Page: 1
Heading:
kernel4.mc      July 14, 1979  4:53 PM  %
%
July 14, 1979  4:53 PM
      Fix bug in computation of correct value for stkp after performing stack+1←.
May 8, 1979  11:53 AM
      Add bypass checking to stack test.
January 25, 1979  1:12 PM
      Add call to checkTaskNum at beginKernel4 to skip the stack tests when we're not executing in ta
**sk 0.
%
top level;

%
      CONTENTS

TEST      DESCRIPTION

stkTest test all stack operations
carry20Test      tests CARRY20 function
xorCarryTest     test XORCARRY function (CIN to bit0, provided by ALUFM)
useSavedCarry    test function (use aluCarry from preeceding instr as CIN)
multiplyTest     test multiply step fcn (affects Q, result. its a slowbranch, too)
divide          test divide step fcn (affects Q, result)
cdivide         test divide step fcn (affects Q, result)
slowBR          tests 8-way slow dispatch
%
beginKernel4:
      call[checkTaskNum], t ← r0;
      skipif[ALU=0];
      branch[stkXitTopL];      * don't try task 0 tests.
```

* May 8, 1979 11:53 AM

%

TEST STKP PUSH AND POP OPERATIONS

-- I AND STKP SHOULD BE INCREMENTING TOGETHER.

-- notation: stack&+1[stkp] ← val : place val into stack[stkp], then increment stkp by 1

-- stack+1[stkp] ← val : increment stkp by one, then place val into stack[stkp]

-- The strategy for this test is to perform all the various stack manipulations (+1, +2, +3,

-- -1, -2, -3, -4, &+1, &+2, &+3, &-1, &-2, &-3, &-4) for every value of stkp that won't

-- cause a hardware error (underflow). The test knows what to expect in RM by setting

-- each rm location to its address (stack[i] ← i).

FOR top2StkpBits IN [0..3] DO

FOR index IN [0..stkpMax] DO -- check simple loading of stkp and stk

i ← index + LeftShift[top2StkpBits, 6];

stkp ← i;

IF POINTERS[].stkError THEN ERROR; -- underflow or overflow

IF POINTERS[].stkp # i THEN ERROR; -- stkp not the value we loaded

stk[stkp] ← 0;

IF stk[stkp] # 0 THEN ERROR; -- loaded zero, got back something different.

stk[stkp] ← -1;

IF stk[stkp] # -1 THEN ERROR; -- loaded -1, got back something different.

stk[stkp] ← Alternating01;

IF stk[stkp] # Alternating01 THEN ERROR; -- loaded Alternating01, got back something dif

**ferent.

stk[stkp] ← Alternating10;

IF stk[stkp] # Alternating10 THEN ERROR; -- loaded Alternating10, got back something dif

**ferent.

stk[stkp] ← i; -- init current stack to stk[i] ← i

IF stk[stkp] # i THEN ERROR;

ENDLOOP;

FOR index IN [0..STKPMAX] DO -- test other operations within that stack

i ← index + LeftShift[top2StkpBits, 6];

stkp ← i;

IF i # STKPMAX THEN -- check +, - 1 operations

BEGIN -- begin w/ operations that use current stkp for loading rm

stack&+1[stkp] ← -1;

IF POINTERS[].stack # i+1 THEN ERROR; -- auto increment of stkp failed

IF stack[stkp] # i+1 THEN ERROR; -- got wrong data after increment

stack&-1[stkp] ← i+1; -- rewrite current data, decrement stkp

IF POINTERS[].stack # i THEN ERROR; -- auto decrement of stkp failed

IF stack[stkp] # -1 THEN ERROR; -- got wrong data after decrement

stkp ← i+1; -- set stkp to check data

IF stack[stkp] # i+1 THEN ERROR; -- wrong data during auto decrement

stkp ← i; -- reset stkp

stack[stkp] ← 0;

IF stack[stkp] # 0 THEN ERROR; -- bypass error

stack[stkp] ← 1;

IF stack[stkp] = 0 THEN ERROR; -- bypass error

stack[stkp] ← i; -- reset data at "current" stkp

-- now test operations that modify stkp before loading rm

stack+1[stkp] ← -1; -- increment stkp, then load rm

IF POINTERS[].stack # i+1 THEN ERROR; -- stkp auto increment failed

IF stack[stkp] # -1 THEN ERROR; -- didn't get data we wrote

stack[stkp] ← i+1; -- fix clobbered location

stack-1[stkp] ← -1; -- decrement stkp, then load rm

IF POINTERS[].stack # i THEN ERROR; -- stkp auto decrement failed

IF stack[stkp] # -1 THEN ERROR; -- didn't get data we wrote

stack[stkp] ← i; -- fix clobbered location

END;

IF i+1 < stkpMax THEN -- check +, - 2 operations

BEGIN -- begin w/ operations that use current stkp for loading rm

stack&+2[stkp] ← -1;

IF POINTERS[].stack # i+2 THEN ERROR; -- auto increment of stkp failed

IF stack[stkp] # i+2 THEN ERROR; -- got wrong data after increment

stack&-2[stkp] ← i+2; -- rewrite current data, decrement stkp

IF POINTERS[].stack # i THEN ERROR; -- auto decrement of stkp failed

IF stack[stkp] # -1 THEN ERROR; -- got wrong data

stack[stkp] ← i; -- reset data at current stkp

stkp ← i+2; -- set stkp to check data

IF stack[stkp] # i+2 THEN ERROR; -- wrong data during auto decrement

```

    stkp ← i;          -- reset stkp

    -- now test operations that modify stkp before loading rm
    stack+2[stkp] ← -1;          -- increment stkp, then load rm
    IF POINTERS[.stack # i+2 THEN ERROR; -- stkp auto increment failed
    IF stack[stkp] # -1 THEN ERROR; -- didn't get data we wrote
    stack[stkp] ← i+2;          -- fix clobbered location
    stack-2[stkp] ← -1;          -- decrement stkp, then load rm
    IF POINTERS[.stack # i THEN ERROR; -- stkp auto decrement failed
    IF stack[stkp] # -1 THEN ERROR; -- didn't get data we wrote
    stack[stkp] ← i;          -- fix clobbered location
    END;
IF i+2 < stkpMAX THEN -- check +,- 3 operations
    BEGIN -- begin w/ operations that use current stkp for loading rm
    stack&+3[stkp] ← -1;
    IF POINTERS[.stack # i+3 THEN ERROR; -- auto increment of stkp failed
    IF stack[stkp] # i+3 THEN ERROR; -- stkp auto increment failed
    stack&-2[stkp] ← i+2; -- rewrite current data, decrement stkp
    IF POINTERS[.stack # i THEN ERROR; -- auto decrement of stkp failed
    IF stack[stkp] # -1 THEN ERROR; -- got wrong data
    stack[stkp] ← i; -- reset data at current stkp
    stkp ← i+3; -- set stkp to check data
    IF stack[stkp] # i+3 THEN ERROR; -- wrong data during auto decrement

    stkp ← i;          -- reset stkp

    -- now test operations that modify stkp before loading rm
    stack+3[stkp] ← -1;          -- increment stkp, then load rm
    IF POINTERS[.stack # i+3 THEN ERROR; -- stkp auto increment failed
    IF stack[stkp] # -1 THEN ERROR; -- stkp auto increment failed
    stack[stkp] ← i+3;          -- fix clobbered location
    stack-3[stkp] ← -1;          -- decrement stkp, then load rm
    IF POINTERS[.stack # i THEN ERROR; -- stkp auto decrement failed
    IF stack[stkp] # -1 THEN ERROR; -- didn't get data we wrote
    stack[stkp] ← i;          -- fix clobbered location
    END;
IF i>3 THEN -- check -4 operations
    BEGIN -- begin w/ operations that use current stkp for loading rm
    stack&-4[stkp] ← -1; -- decrement stkp, then load rm
    IF POINTERS[.stack # i-4 THEN ERROR; -- stkp auto decrement failed
    IF stack[stkp] # i-4 THEN ERROR; -- didn't get data we wrote
    stkp ← i;
    IF stack[stkp] # -1 THEN ERROR; -- didn't get data we wrote
    stack[stkp] ← i-4; -- fix clobbered location

    -- now test operations that modify stkp before loading rm
    stack-4[stkp] ← -1;          -- decrement stkp, then load rm
    IF POINTERS[.stack # i-4 THEN ERROR; -- stkp auto decrement failed
    IF stack[stkp] # -1 THEN ERROR; -- didn't get data we wrote
    stack[stkp] ← i-4          -- fix clobbered location
    stkp ← i;
    IF stack[stkp] # i THEN ERROR;
    ENDLOOP;
-- end of stkp loop
ENDLOOP;
-- end of top2StkpBits loop
%
```

* July 14, 1979 4:53 PM

%

stkTest Test the various stack operations

%

```
mc[stkPMaxXC, 77];
mc[pointers.stkOvf, b8];
mc[pointers.stkUnd, b9];
mc[pointers.stkErr, b8, b9];
stkTest:      * initialize the top2bits loop
              call[iTopStkBits];
stkTopL:      * top of "top 2 bits of stkp" loop
              call[nextTopStkBits];
              skipif[ALU#0];
              branch[stkXitTopL];
              noop;
```

* This code writes the current stack with the address (stack[stkP] ← stkP).

* It also checks that stkp←, ←stack work properly.

```
              call[iStkPAddr];      * initialize stack index [1..maxStkXC]
stkIL:
              call[nextStkPAddr];    * top of stk init loop. here we check stkp← and ←stack.
              skipif[ALU#0];
              branch[stkiXit];
              stkp ← t;              * load stkp
              call[chkStkErr];
              skipif[ALU=0];
stkiErr0:     * got stack underflow or overflow
              error;

              call[getRealStkAddr], rscr ← t;
              t # (rscr);           * compare real stkp with value we loaded
              skipif[ALU=0];
stkiErr1:     * t = tskp, rscr = value we loaded
              error;
```

* This is a limited test of the bits in the stack memory: write zero, -1, alternating 10, 01

```
              t ← stack ← t-t;
              t ← t #(Q←stack) ;
              skipif[ALU=0];
stkiErr2:     * wrote zero, got back something else
              error; * Q = value from stack

              t ← rm1;
              stack ← t;
              t ← t #(Q←stack) ;
              skipif[ALU=0];
stkiErr3:     * wrote -1 got back something else.
              error; * t = bad bits Q = value from stack

              t ← r01;
              stack ← t;
              t ← t #(Q←stack) ;
              skipif[ALU=0];
stkiErr4:     * wrote r01 got back something else.
              error; * t = bad bits. Q = value from stack

              t ← r10;
              stack ← t;
              t ← t #(Q←stack) ;
              skipif[ALU=0];
stkiErr5:     * wrote r10 got back something else.
              error; * t = bad bits. Q = value from stack

              t ← rscr;             * t ← current index
              stack←t;             * stack[i] ←i, then check it
              t # (Q←stack);
              skipif[ALU=0];
stkiErr6:     * wrote stkp from rscr. Q = value from stack
              error; * read it into t. they aren't the same
              branch[stkiL];
stkiXit:
```

```

* July 14, 1979 4:53 PM
* We have successfully written the stack using non incrementing and non decrementing
* operations. Now we test stack&+1←, stack&-1←, stack+1←, stack-1←

    call[iStkPAddr];      * init the main loop for the main test
stkTestL:  * top of main loop
    call[nextStkPAddr];  * get next stack index or exit loop
    skipif[alu#0];
    branch[stkTestXitL];
    stkp ← t;      * stackP ← i

    call[chkStkErr];
    skipif[ALU=0];
stkpErr10:  * got stack underflow or overflow
    error;

    rscr ← t and (77c);    * isolate the index (exclude top 2 bits)
    (rscr)-(stkPMaxXC);   * skip this test if it would cause overflow
    branch[afterStkTest1, ALU=0];

    t # (Q←stack); * see if stack[stkp] = stkp
    skipif[ALU=0]; * if not, an earlier execution of this loop clobbered
stkpErr1:  * the stack entry at location in t, or this is first time
    error; * thru, and the initialization didn't work properly.
    Q=value from stack

* stack&+1 stack&+1 stack&+1 stack&+1 stack&+1 stack&+1 stack&+1 stack&+1

    stack&+1 ← cm1; * stack[stackP] ← -1, then stackP ← stackP+1
    call[chkStkErr];
    skipif[ALU=0];
stk&+1Err0:  * got stack underflow or overflow
    error;

    call[getRealStkAddr], rscr←t+1; * compare stackPAddr from Pointers w/ expected val
    t # (rscr);
    skipif[ALU=0];
stkP1AddrErr:  * auto increment of StackP failed. rscr = expected value,
    error; * t = value from Pointers

    t ← t # (Q←stack);
    skipif[ALU=0];
stkP1ValErr:  * value at stackp is bad. Q = value from stack
    error; * t = expected val, rscr = stack's val from Pointers
    t ← rscr; * restore t

* stack&-1 stack&-1 stack&-1 stack&-1 stack&-1 stack&-1 stack&-1 stack&-1 stack&-1

    stack&-1 ← t; * stack["i+1"] ← i+1, stackp ← i.

    call[chkStkErr];
    skipif[ALU=0];
stkpErr12:  * got stack underflow or overflow
    error;

    call[getRealStkAddr], rscr ← t-1;
    t # (rscr); * compare expected stkP (rscr) with
    skipif[ALU=0]; * actual stkp (t)
stkM1AddrErr:  * auto decrement failed
    error;

    t ← cm1;
    t ← t # (Q←stack);
    skipif[ALU=0]; * see if original stack&+1 ← cm1 worked
stkP1ValErr2:  * stack&+1 seems to have clobbered the
    error; * (i+1)th value. t = bad bits. Q = value from stack

    t ← rscr; * restore t
    (stack)← t; * reset stk[stkp] to contain stkp

    call[chkStkErr];
    skipif[ALU=0];
stkpErr13:  * got stack underflow or overflow
    error;

```

```

    rscr ← t ← t+1;
    stkp ← t;      * check the data modified during "stack&-1" instruction
    t ← t # (Q←stack);      * compare tos with expected valu
    skipif[ALU=0];
stkM1ValErr:    * . Q = value from stack
error;    * t = bad bits, rscr = expected value

    t ← rscr ← (rscr)-1;    * t, rscr ← "i"
    stkp ← t;    * stkp is at i+1 now. Fix it.

    call[chkStkErr];
    skipif[ALU=0];
stkpErr14:    * got stack underflow or overflow
error;

    Q ← stack;    * save stack value
    stack ← t-t;
    PD←(stack);
    skipif[ALU=0];
stkByPassErr0: * didn't notice that we just zeroed the stack
error;
    t ← cm1;
    stack ← cm1;
    PD ← (stack) # t;
    skipif[ALU=0];
stkByPassErr1: * didn't notice that we just put all ones
error;    * in the stack.

    stack ← Q;    * restore stack

* stack+1 stack+1 stack+1 stack+1 stack+1 stack+1 stack+1 stack+1 stack+1 stack+1

    rscr ← (rscr)+1;    * compute expected stkp value
    t ← cm1;
    stack+1 ← t;    * stkp ← i+1, stack[stkp] ← -1

    call[chkStkErr];
    skipif[ALU=0];
stkpErr15:    * got stack underflow or overflow
error;

    call[getRealStkAddr];
    (rscr) # t;
    skipif[ALU=0];
stkP1AddrErr2: * expected Rscr, got stackp in t, they're different
error;

    t ← cm1;
    t ← t # (Q←stack);    * check that we loaded -1 into incremented stack location
    skipif[ALU=0];    * Q = value from stack
stkP1ValErr3:    * t = bad bits
error;
    t ← rscr;    * restore t

* stack-1 stack-1 stack-1 stack-1 stack-1 stack-1 stack-1 stack-1 stack-1 stack-1

    stack ← t;    * reset stack which was clobbered by "stack+1←cm1"

    call[chkStkErr];
    skipif[ALU=0];
stkpErr16:    * got stack underflow or overflow
error;

    rscr ← t-1;    * compute expected value of rscr
    t ← cm1;
    stack-1 ← t;    * (stack-1) ← "-1"

    call[chkStkErr];
    skipif[ALU=0];
stkpErr17:    * got stack underflow or overflow
error;

    call[getRealStkAddr];
    (rscr) # t;    * see if real stkp (t) matches expected stkp (rscr)
    skipif[ALU=0];
stkM1AddrErr2:

```

```
error;
t ← cm1;
t ← t # (Q←stack);      * compare tos with -1
skpif[ALU=0];
stkM1ValErr2: * Q = value from stack
error; * t = bad bits, expected -1

t ← rscr;      * restore t
(stack) ← t;   * restore addr as value in stack: stack[stkp]←stkp

call[chkStkErr];
skpif[ALU=0];
stkpErr18: * got stack underflow or overflow
error;
noop; * for placement

afterStkTest1:
```

```

* November 30, 1978 6:07 PM
%remember, don't execute if i=1, if i+2=77%
  call[getStkPAddr];
  rscr ← t and (77c);      * isolate the index (exclude top 2 bits)
  rscr ← (rscr)+1;
  (rscr)-(stkPMaxXC);     * skip this test if it would cause overflow
  branch[afterStkTest2, ALU>=0];

  t # (Q←stack); * see if stack[stkp] = stkp
  skipif[ALU=0]; * if not, an earlier execution of this loop clobbered
stkpErr21:      * the stack entry at location in t, or this is first time
error; * thru, and the initialization didn't work properly.
* Q=value from stack

* stack&+2 stack&+2 stack&+2 stack&+2 stack&+2 stack&+2 stack&+2 stack&+2

  stack&+2 ← cm1; * stack[stackP] ← -1, then stackP ← stackP+2
  call[chkStkErr];
  skipif[ALU=0];
stk&+2Err0:    * got stack underflow or overflow
error;

  call[getRealStkAddr], rscr←t+(2c);      * compare stackPAddr from Pointers w/ expected val
  t # (rscr);
  skipif[ALU=0];
stkP2AddrErr: * auto increment of StackP failed. rscr = expected value,
error; * t = value from Pointers

  t ← t # (Q←stack);
  skipif[ALU=0];
stkP2ValErr:  * value at stackp is bad. Q = value from stack
error; * t = expected val, rscr = stack's val from Pointers
  t ← rscr;      * restore t

* stack&-2 stack&-2 stack&-2 stack&-2 stack&-2 stack&-2 stack&-2 stack&-2 stack&-2

  stack&-2 ← t; * stack["i+2"] ← i+2, stackp ← i.

  call[chkStkErr];
  skipif[ALU=0];
stkpErr22:    * got stack underflow or overflow
error;

  call[getRealStkAddr], rscr ← t-(2c);
  t # (rscr); * compare expected stkP (rscr) with
  skipif[ALU=0]; * actual stkp (t)
stkM2AddrErr: * auto decrement failed
error;

  t ← cm1;
  t ← t # (Q←stack);
  skipif[ALU=0]; * see if original stack&+2 ← cm1 worked
stkP2ValErr2: * stack&+2 seems to have clobbered the
error; * (i+2)th value. t = bad bits. Q = value from stack

  t ← rscr;      * restore t
  (stack)← t;    * reset stk[stkp] to contain stkp

  call[chkStkErr];
  skipif[ALU=0];
stkpErr23:    * got stack underflow or overflow
error;

  rscr ← t ← t+(2C);
  stkp ← t;      * check the data modified during "stack&-2" instruction
  t ← t # (Q←stack); * compare tos with expected valu
  skipif[ALU=0];
stkM2ValErr:  *. Q = value from stack
error; * t = bad bits, rscr = expected value

* stack+2 stack+2 stack+2 stack+2 stack+2 stack+2 stack+2 stack+2 stack+2 stack+2

  t ← rscr ← (rscr)-(2c); * t, rscr ← "i"
  stkp ← t;      * stkp is at i+2c now. Fix it.

```



```

    call[chkStkErr];
    skipif[ALU=0];
stkErr24:    * got stack underflow or overflow
    error;

    rscr ← t+(2c); * compute expected stkp value
    t ← cm1;
    stack+2 ← t; * stkp ← i+2, stack[stkp] ← -1

    call[chkStkErr];
    skipif[ALU=0];
stkErr25:    * got stack underflow or overflow
    error;

    call[getRealStkAddr];
    (rscr) # t;
    skipif[ALU=0];
stkP2AddrErr2: * expected stackP t, got stackp in Rscr, they're different
    error;

    t ← cm1;
    t ← t # (Q←stack); * check that we loaded -1 into incremented stack location
    skipif[ALU=0]; * Q = value from stack
stkP2ValErr3: * t = bad bits
    error;
    t ← rscr; * restore t

* stack-2 stack-2 stack-2 stack-2 stack-2 stack-2 stack-2 stack-2 stack-2 stack-2
    stack ← t; * reset stack which was clobbered by "stack+2←cm1"

    call[chkStkErr];
    skipif[ALU=0];
stkErr26:    * got stack underflow or overflow
    error;

    rscr ← t-(2c); * compute expected value of rscr
    t ← cm1;
    stack-2 ← t; * (stack-2) ← "-1"

    call[chkStkErr];
    skipif[ALU=0];
stkErr27:    * got stack underflow or overflow
    error;

    call[getRealStkAddr];
    (rscr) # t; * see if real stkp (t) matches expected stkp (rscr)
    skipif[ALU=0];
stkM2AddrErr2:
    error;

    t ← cm1;
    t ← t # (Q←stack); * compare tos with -1
    skipif[ALU=0];
stkM2ValErr2: * Q = value from stack
    error; * t = bad bits, expected -1

    t ← rscr; * restore t
    (stack) ← t; * restore addr as value in stack: stack[stkp]←stkp

    call[chkStkErr];
    skipif[ALU=0];
stkErr228:    * got stack underflow or overflow
    error;
    noop; * for placement

afterStkTest2:

```

```

* December 1, 1978 3:19 PM
%remember, don't execute if i=1, if i+3>=77%
  noop; * placement for afterStkTest2 branch
  call[getStkPAddr];
  rscr ← t and (77c); * isolate the index (exclude top 2 bits)
  rscr ← (rscr)+(2c);
  (rscr)-(stkPMaxXC); * skip this test if it would cause overflow
  branch[afterStkTest3, ALU>=0];

  t # (Q←stack); * see if stack[stkp] = stkp
  skipif[ALU=0]; * if not, an earlier execution of this loop clobbered
stkpErr31: * the stack entry at location in t, or this is first time
error; * thru, and the initialization didn't work properly.
* Q=value from stack

* stack&+3 stack&+3 stack&+3 stack&+3 stack&+3 stack&+3 stack&+3 stack&+3

  stack&+3 ← cm1; * stack[stackP] ← -1, then stackP ← stackP+3
  call[chkStkErr];
  skipif[ALU=0];
stk&+3Err0: * got stack underflow or overflow
error;

  call[getRealStkAddr], rscr←t+(3c); * compare stackPAddr from Pointers w/ expected val
  t # (rscr);
  skipif[ALU=0];
stkP3AddrErr: * auto increment of StackP failed. rscr = expected value,
error; * t = value from Pointers

  t ← t # (Q←stack);
  skipif[ALU=0];
stkP3ValErr: * value at stackp is bad. Q = value from stack
error; * t = expected val, rscr = stack's val from Pointers
  t ← rscr; * restore t

* stack&-3 stack&-3 stack&-3 stack&-3 stack&-3 stack&-3 stack&-3 stack&-3 stack&-3

  stack&-3 ← t; * stack["i+3"] ← i+3, stackp ← i.

  call[chkStkErr];
  skipif[ALU=0];
stkpErr32: * got stack underflow or overflow
error;

  call[getRealStkAddr], rscr ← t-(3c);
  t # (rscr); * compare expected stkp (rscr) with
  skipif[ALU=0]; * actual stkp (t)
stkM3AddrErr: * auto decrement failed
error;

  t ← cm1;
  t ← t # (Q←stack);
  skipif[ALU=0]; * see if original stack&+3 ← cm1 worked
stkP3ValErr2: * stack&+3 seems to have clobbered the
error; * (i+3)th value. t = bad bits. Q = value from stack

  t ← rscr; * restore t
  (stack)← t; * reset stk[stkp] to contain stkp

  call[chkStkErr];
  skipif[ALU=0];
stkpErr33: * got stack underflow or overflow
error;

  rscr ← t ← t+(3C);
  stkp ← t; * check the data modified during "stack&-3" instruction
  t ← t # (Q←stack); * compare tos with expected valu
  skipif[ALU=0];
stkM3ValErr: * Q = value from stack
error; * t = bad bits, rscr = expected value

* stack+3 stack+3 stack+3 stack+3 stack+3 stack+3 stack+3 stack+3 stack+3 stack+3

  t ← rscr ← (rscr)-(3c); * t, rscr ← "i"
  stkp ← t; * stkp is at i+3c now. Fix it.

```

```

    call[chkStkErr];
    skipif[ALU=0];
stkpErr34:    * got stack underflow or overflow
    error;

    rscr ← t+(3c); * compute expected stkp value
    t ← cm1;
    stack+3 ← t; * stkp ← i+3, stack[stkp] ← -1

    call[chkStkErr];
    skipif[ALU=0];
stkpErr35:    * got stack underflow or overflow
    error;

    call[getRealStkAddr];
    (rscr) # t;
    skipif[ALU=0];
stkP3AddrErr2: * expected stackP t, got stackp in Rscr, they're different
    error;

    t ← cm1;
    t ← t # (Q←stack); * check that we loaded -1 into incremented stack location
    skipif[ALU=0]; * Q = value from stack
stkP3ValErr3: * t = bad bits
    error;
    t ← rscr; * restore t

* stack-3 stack-3 stack-3 stack-3 stack-3 stack-3 stack-3 stack-3 stack-3 stack-3
    stack ← t; * reset stack which was clobbered by "stack+3←cm1"

    call[chkStkErr];
    skipif[ALU=0];
stkpErr36:    * got stack underflow or overflow
    error;

    rscr ← t-(3c); * compute expected value of rscr
    t ← cm1;
    stack-3 ← t; * (stack-3) ← "-1"

    call[chkStkErr];
    skipif[ALU=0];
stkpErr37:    * got stack underflow or overflow
    error;

    call[getRealStkAddr];
    (rscr) # t; * see if real stkp (t) matches expected stkp (rscr)
    skipif[ALU=0];
stkM3AddrErr2:
    error;

    t ← cm1;
    t ← t # (Q←stack); * compare tos with -1
    skipif[ALU=0];
stkM3ValErr2: * Q = value from stack
    error; * t = bad bits, expected -1

    t ← rscr; * restore t
    (stack) ← t; * restore addr as value in stack: stack[stkp]←stkp

    call[chkStkErr];
    skipif[ALU=0];
stkpErr38:    * got stack underflow or overflow
    error;
    noop; * for placement

afterStkTest3:

```

```

* December 1, 1978 4:57 PM
%remember, don't execute if i=1, if i+4>=77%
  noop; * placement for the afterStkTest3 check
  call[getStkPAddr];
  rscr ← t and (77c); * isolate the index (exclude top 2 bits)
  rscr ← (rscr)+(3c);
  (rscr)-(stkPMaxXC); * skip this test if it would cause overflow
  branch[afterStkTest4, ALU>=0];

  t # (Q←stack); * see if stack[stkP] = stkP
  skipif[ALU=0]; * if not, an earlier execution of this loop clobbered
stkPErr41: * the stack entry at location in t, or this is first time
error; * thru, and the initialization didn't work properly.
* Q=value from stack

* Simulate stack&+4 -- hardware can perform stack&+3 as maximum increment

  stack&+3 ← cm1; * stack[stackP] ← -1, then stackP ← stackP+4
  stkP+1; * simulate +4
  call[chkStkErr];
  skipif[ALU=0];
stk&+4Err0: * got stack underflow or overflow
error;

  call[getRealStkAddr], rscr←t+(4c); * compare stackPAddr from Pointers w/ expected val
  t # (rscr);
  skipif[ALU=0];
stkP4AddrErr: * auto increment of StackP failed. rscr = expected value,
error; * t = value from Pointers

  t ← t # (Q←stack);
  skipif[ALU=0];
stkP4ValErr: * value at stackp is bad. Q = value from stack
error; * t = expected val, rscr = stack's val from Pointers
  t ← rscr; * restore t

* stack&-4 stack&-4 stack&-4 stack&-4 stack&-4 stack&-4 stack&-4 stack&-4 stack&-4

  stack&-4 ← t; * stack["i+4"] ← i+4, stackp ← i.

  call[chkStkErr];
  skipif[ALU=0];
stkPErr42: * got stack underflow or overflow
error;

  call[getRealStkAddr], rscr ← t-(4c);
  t # (rscr); * compare expected stkP (rscr) with
  skipif[ALU=0]; * actual stkP (t)
stkM4AddrErr: * auto decrement failed
error;

  t ← cm1;
  t ← t # (Q←stack);
  skipif[ALU=0]; * see if original stack&+4 ← cm1 worked
stkP4ValErr2: * stack&+4 seems to have clobbered the
error; * (i+4)th value. t = bad bits. Q = value from stack

  t ← rscr; * restore t
  (stack)← t; * reset stk[stkP] to contain stkP

  call[chkStkErr];
  skipif[ALU=0];
stkPErr43: * got stack underflow or overflow
error;

  rscr ← t ← t+(4C);
  stkP ← t; * check the data modified during "stack&-4" instruction
  t ← t # (Q←stack); * compare tos with expected valu
  skipif[ALU=0];
stkM4ValErr: * Q = value from stack
error; * t = bad bits, rscr = expected value

* stack+4 stack+4 stack+4 stack+4 stack+4 stack+4 stack+4 stack+4 stack+4 stack+4

  t ← rscr ← (rscr)-(4c); * t, rscr ← "i"

```

```

    stkp ← t;      * stkp is at i+4c now. Fix it.
    call[chkStkErr];
    skipif[ALU=0];
stkpErr44:      * got stack underflow or overflow
    error;

    rscr ← t+(4c); * compute expected stkp value
    t ← cm1;
    stkp+1; * simulate stack+4
    stack+3 ← t; * stkp ← i+4, stack[stkp] ← -1

    call[chkStkErr];
    skipif[ALU=0];
stkpErr45:      * got stack underflow or overflow
    error;

    call[getRealStkAddr];
    (rscr) # t;
    skipif[ALU=0];
stkP4AddrErr2: * expected stackP t, got stackp in Rscr, they're different
    error;

    t ← cm1;
    t ← t # (Q←stack); * check that we loaded -1 into incremented stack location
    skipif[ALU=0]; * Q = value from stack
stkP4ValErr3:  * t = bad bits
    error;
    t ← rscr; * restore t

* stack-4 stack-4 stack-4 stack-4 stack-4 stack-4 stack-4 stack-4 stack-4 stack-4
    stack ← t; * reset stack which was clobbered by "stack+4+cm1"

    call[chkStkErr];
    skipif[ALU=0];
stkpErr46:      * got stack underflow or overflow
    error;

    rscr ← t-(4c); * compute expected value of rscr
    t ← cm1;
    stack-4 ← t; * (stack-4) ← "-1"

    call[chkStkErr];
    skipif[ALU=0];
stkpErr47:      * got stack underflow or overflow
    error;

    call[getRealStkAddr];
    (rscr) # t; * see if real stkp (t) matches expected stkp (rscr)
    skipif[ALU=0];
stkM4AddrErr2:
    error;

    t ← cm1;
    t ← t # (Q←stack); * compare tos with -1
    skipif[ALU=0];
stkM4ValErr2:  * Q = value from stack
    error; * t = bad bits, expected -1

    t ← rscr; * restore t
    (stack) ← t; * restore addr as value in stack: stack[stkp]←stkp

    call[chkStkErr];
    skipif[ALU=0];
stkpErr48:      * got stack underflow or overflow
    error;
    noop; * for placement

afterStkTest4:
    branch[stkTestL];
stkTestXitL:
    branch[stkTopL];

```

* November 27, 1978 10:31 AM.

```

iStkPAddr: subroutine;
    return, stackPAddr ← t-t;      * first valid index is one.

nextStkPAddr: subroutine;      * stack indeces are 6 bits long.
*   Return (stackPAddr OR stackPtopBits) in T. It's an 8 bit address.
*   ALU#0 =>valid value.

    klink ← link;
    t ← stackPAddr ← (stackPAddr) + 1;      * increment the index
    t and (77c);      * check for 6 bit overflow
    skipif[ALU=0], t ← t + (stackPtopBits); * OR the top two bits into returned value
    skip, rscr ← 1c;      * indicate valid value
    rscr ← t-t;      * indicate invalid value
    returnAndBranch[klink, rscr];

getStkPAddr: subroutine;
    t ← stackPAddr;
    return, t ← t + (stackPtopBits);

iTopStkBits: subroutine;
    t ← (r0) - (100c);
    return, stackPtopBits ← t;      * first valid index is zero.

nextTopStkBits: subroutine;
    klink ← link;
    top level;
    t ← stackPtopBits ← (stackPtopBits)+(100c);
    t - (400c);
    skipif[ALU#0], rscr ← 1c;
    rscr ← t-t;
    returnAndBranch[klink, rscr];

getRealStkAddr: subroutine;
    t ← TIOA&Stkp;
    return, t ← t and (377c);

chkStkErr: subroutine; * rtn w/ ALU#0 ==> stk (underflow or overflow).
* Clobber rscr2 ← Pointers[]
    rscr2 ← Pointers[];
    return, rscr2 ← (rscr2) AND (pointers.stkErr);
    top level;

stkXitTopL:

```

* September 3, 1977 2:25 PM

%

TEST CARRY20 FUNCTION

This function causes a 1 to be or'd into the carry out bit that is used as input to bit 11 in the alu. Given that there was not already a carry, this function has the effect of adding 20B to the value in the alu.

%

carry20Test:

```
t←rscr+17C;
t←t+(r0),CARRY20;
rscr2 ← 37C;
t←t#(rscr2);
skipif[ALU=0];
error; * T NE 17B + 0 + CARRY20
```

```
t←rscr;
t←t+(r1),CARRY20; * t←17B+1+CARRY20
rscr2←20C;
t←t#(rscr2);
skipif[ALU=0];
error; * T NE 17B+1=20(=17B+1+CARRY20)
```

```
t←r0;
t←t+(r0),CARRY20; * t←0+0+CARRY20
rscr2←20C;
t←t#(rscr2);
skipif[ALU=0];
error; * T NE 20B=0+0+CARRY20
```

```
t←r0;
t←t-(rscr2); * t←-20B=(0-20B)
t←t+(r0),CARRY20;
skipif[ALU=0];
error;
```

* September 11, 1977 1:57 PM

%

TEST XORCARRY FUNCTION

XORCARRY causes the carry in bit for bit 15 of the alu to be xor'd. Normally this bit is 0. When the bit is one, alu arithmetic functions will see a carry into bit 0. For A-B the ALUFM is programmed to provide a one and XORCARRY will leave a result one less than expected.

%

xorCarryTest:

t←(r0)+(r0),XORCARRY;
t←t#(r1);
skpif[ALU=0];
error; * 1 = 0+0+XORCARRY

t←r1;
t←t+(r0),XORCARRY;
t←t#(2C);
skpif[ALU=0];

xorCarryb:

error; * 2= 0+1+XORCARRY

t←r1;
t←t+(r1),XORCARRY;
t←t#(3C);
skpif[ALU=0];

xorCarryc:

error; * 3= 1+1+XORCARRY

t←RM1;
t←t+(r0),XORCARRY;
skpif[ALU=0];

xorCarryd:

error; * 0= -1+XORCARRY

t←(r0)AND(r0),XORCARRY;
skpif[ALU=0];

xorCarrye:

error; * CIN SHOULD BE IGNORED ON LOGICAL OPS!

t←(r1)AND(r1),XORCARRY;
t←t#(r1);
skpif[ALU=0];

xorCarryf:

error; * SHOULD BE 1. CIN IGNORED ON LOGICAL OPS

t←(RM1)OR(RM1),XORCARRY;
t←t#(RM1);
skpif[ALU=0];

xorCarryg:

error; * SHOULD BE -1. CIN IGNORED ON LOGICAL OPS

t←(r1)-(r1),XORCARRY;
t←t#(RM1);
skpif[ALU=0];

xorCarryh:

error; * BWL SEZ THIS SHOULD BE -1. IE.,

* R1-R1 causes 1+1777777, but xorcarry causes the op to become,

* 1+177776 because "A-B" uses carryin = 1, and xorcarry causes it to be

* zero!

* September 12, 1977 9:52 AM

%

TEST USESAVEDCARRY

This function causes the alu carry bit from the last instruction to be used as the carry in bit to bit 15 during the current instruction. This bit is usually provided by the alufm and is usually zero (its the bit complemented by the xorcarry function).

%

% commented out

savedCarry:

T←(RHIGH1)+(RHIGH1); * T←0, CARRY←1
T←T+(R0),USESAVEDCARRY; * T←0+0+LAST CARRY
T←T#(R1);
SKPIF[ALU=0];
ERROR; * EXPECTED 1, USED LASTCARRY=1

T←(RM1)+(RM1); * T←-2, CARRY ←1
T←T+(R1),USESAVEDCARRY; * T←-2+1+LAST CARRY
SKPIF[ALU=0];

savedCarryB:

ERROR; * EXPECTED 0, USED LASTCARRY=1

T←(R0)+(R0); * T←0, CARRY←0
T←(R1)+(R1),USESAVEDCARRY; * T←1+1+LAST CARRY
T←T#(2C);
SKPIF[ALU=0];

savedCarryC:

ERROR; * EXPECTED 2, USED LASTCARRY=0

T←(R0)+(R0); * T←0, CARRY←0
T←(RM1)+(RM1),USESAVEDCARRY; * T←(-1)+(-1)+LAST CARRY
T←T#(177776C);
SKPIF[ALU=0];

savedCarryD:

ERROR; * EXPECTED -2, USED LASTCARRY=0

T←(RM1)+(RM1); * T←-2, CARRY←1
T←(R1)+(R1),USESAVEDCARRY; * T←1+1+LAST CARRY
T←T#(3C);
SKPIF[ALU=0];

savedCarryE:

ERROR; * EXPECTED 3, USED LASTCARRY=1
commented out%

* September 14, 1977 12:03 PM

%

TEST MULTIPLY STEP.

MULTIPLY works as follows:

H3[0:15] ← CARRY,,ALU[0:14] ==> CARRY,,ALU/2

Q[0:15] ← ALU[15],,Q[0:14] ==> ALU[15],,Q/2

Q[14] OR'D INTO TNIA[10] AS SLOW BRANCH! ==> ADDR OR'D 2 IF Q[14]=1

These tests invoke mulCheck, a subroutine that performs two services:

1) T←T+RSCR,MULTIPLY

2) RSCR2←1 IF Q[14] BRANCH IS TAKEN (ZERO OTHERWISE)

ERRORS are numbered 1 thru 3:

mulXerr1 ==> T value wrong (H3)

mulXerr2 ==> Q[14] branch wrong

mulXerr3 ==> Q value wrong

%

multiplyTest:

* Q[14]=0, CARRY=0, ALU15=0

t←Q←r0;

rscr←t;

call[mulCheck]; * t←t+rscr,MULTIPLY==>ALU←0, CARRY←0

skpif[R EVEN],rscr2←rscr2;

mulAerr1:

error; * TOOK Q[14]=1 BRANCH

t←t;

skpif[ALU=0];

mulAerr2:

error; * CARRY,,(0+0)/2 SHOULD BE ZERO

t←Q;

skpif[ALU=0];

mulAerr3:

error; * ALU15,,Q[0:14] SHOULD BE ZERO

* Q[14]=0, CARRY=1, ALU15=0

multiplyB: * Q=0; t←-1+1,MULTIPLY

Q←r0;

rscr←r1;

call[mulCheck],t←rm1; * t←t+rscr,MULTIPLY==>ALU←0, CARRY←1

skpif[R EVEN], rscr2←rscr2;

mulBerr1:

error; * TOOK Q[14]=1 BRANCH

t←t#(rhigh1); *-1+1 GENERATES CARRY BIT

skpif[ALU=0];

mulBerr2:

error; * CARRY,,(0+0)/2 SHOULD BE 100000

t←Q; * -1+1 WOULD LEAVE ALU15=0

skpif[ALU=0];

mulBerr3:

error; * ALU15,,Q[0:14] SHOULD BE ZERO

* Q[14]=0, CARRY=0, ALU15=1

multiplyC: * Q=0; t←0+1,MULTIPLY

t←Q←r0;

rscr←r1;

call[mulCheck]; * t←t+rscr,MULTIPLY==>ALU←1, CARRY←0

skpif[R EVEN], rscr2←rscr2;

mulCerr1:

error; * TOOK Q[14]=1 BRANCH

t←t; * 0+1==> CARRY←0

skpif[ALU=0];

mulCerr2:

error; * CARRY,,(0+1)/2 SHOULD BE 0

t←(rhigh1)#(Q); * 0+1 WOULD LEAVE ALU15=1

skpif[ALU=0];

mulCerr3:

error; * ALU15,,Q[0:14] SHOULD BE 100000

```

*
multiplyD:      Q[14]=0, CARRY=1, ALU15=1
                * Q=0; t←100001+100000,MULTIPLY
                Q←r0;
                t+rscr←rhigh1;
                call[mulCheck],t←t+(r1);      * t←t+rscr,MULTIPLY==>ALU←1, CARRY←1
                skipif[R EVEN], rscr2←rscr2;
mulDerr1:
  error; * TOOK Q[14]=1 BRANCH

                t←t#(rhigh1); * 1000001+100000==> CARRY←1
                skipif[ALU=0];
mulDerr2:
  error; * CARRY,,(1000001+100000)/2 SHOULD BE 100000

                t←(rhigh1)#(Q); * 1000001+100000 WOULD LEAVE ALU15=1
                skipif[ALU=0];
mulDerr3:
  error; * ALU15,,Q[0:14] SHOULD BE 100000

multiplyE:
*
                Q[14]=1, CARRY=0, ALU15=0
                t←(r1)+(r1);
                Q←t; * Q[14]←1
                t←r0;
                rscr←t;
                call[mulCheck]; * t←t+rscr,MULTIPLY==>ALU←0, CARRY←0
                skipif[R ODD],rscr2←rscr2;
mulEerr1:
  error; * DIDN'T TAKE Q[14]=1 BRANCH

                t←t;
                skipif[ALU=0];
mulEerr2:
  error; * CARRY,,(0+0)/2 SHOULD BE ZERO

                t←(r1)#(Q);
                skipif[ALU=0];
mulEerr3:
  error; * ALU15,,Q[0:14] SHOULD BE 1

*
multiplyF:      Q[14]=1, CARRY=1, ALU15=0
                * Q=1; t←-1+1,MULTIPLY
                t←(r1)+(r1);
                Q←t; * Q[14]←1
                rscr←r1;
                call[mulCheck],t←rm1; * t←t+rscr,MULTIPLY==>ALU←0, CARRY←1
                skipif[R ODD], rscr2←rscr2;
mulFerr1:
  error; * DIDN'T TAKE Q[14]=1 BRANCH

                t←t#(rhigh1); * -1+1 GENERATES CARRY BIT
                skipif[ALU=0];
mulFerr2:
  error; * CARRY,,(0+0)/2 SHOULD BE 100000

                t←(r1)#(Q); * -1+1 WOULD LEAVE ALU15=0
                skipif[ALU=0];
mulFerr3:
  error; * ALU15,,Q[0:14] SHOULD BE 1

*
multiplyG:      Q[14]=1, CARRY=0, ALU15=1
                * Q=1; t←0+1,MULTIPLY
                t←(r1)+(r1);
                Q←t; * Q[14]←1
                t←r0;
                rscr←r1;
                call[mulCheck]; * t←t+rscr,MULTIPLY==>ALU←1, CARRY←0
                skipif[R ODD], rscr2←rscr2;
mulGerr1:
  error; * DIDN'T TAKE Q[14]=1 BRANCH

                t←t; * 0+1==> CARRY←0
                skipif[ALU=0];
mulGerr2:
  error; * CARRY,,(0+1)/2 SHOULD BE 0

```

```

        t←(rhigh1)+1; * 0+1 WOULD LEAVE ALU15=1
        t←t#(Q);
        skipif[ALU=0];
mulGerr3:
        error; * ALU15,,Q[0:14] SHOULD BE 100001

*
        Q[14]=1, CARRY=1, ALU15=1
multiplyH: * Q=2; t←100001+100000,MULTIPLY
        t←(r1)+(r1);
        Q←t; * Q[14]←1
        t←rscr←rhigh1;
        call[mulCheck],t←t+(r1); * t←t+rscr,MULTIPLY==>ALU←1, CARRY←1
        skipif[R ODD], rscr2←rscr2;
mulHerr1:
        error; * DIDN'T TAKE Q[14]=1 BRANCH

        t←t#(rhigh1); * 100001+100000==> CARRY←1
        skipif[ALU=0];
mulHerr2:
        error; * CARRY,,(100001+100000)/2 SHOULD BE 100000

        t←(rhigh1)+1; * 100001+100000 WOULD LEAVE ALU15=1
        t←t#(Q);
        skipif[ALU=0];
mulHerr3:
        error; * ALU15,,Q[0:14] SHOULD BE 100001

multiplyJ:
*
        Q←r01=>Q[14]=0; CARRY=0,ALU15=0
        t←r01;
        Q←t; * Q[14]←0
        t←r0;
        rscr←t;
        call[mulCheck]; * t←t+rscr,MULTIPLY==>ALU←0, CARRY←0
        skipif[R EVEN],rscr2←rscr2;
mulJerr1:
        error; * TOOK Q[14]=1 BRANCH

        t←t;
        skipif[ALU=0];
mulJerr2:
        error; * CARRY,,(0+0)/2 SHOULD BE ZERO

        t←(r01) RSH 1;
        t←t#(Q);
        skipif[ALU=0];
mulJerr3:
        error; * ALU15,,Q[0:14] SHOULD BE (r01) RSH 1

multiplyK:
*
        Q←r10=>Q[14]=1; CARRY=0,ALU15=0
        t←r10;
        Q←t; * Q[14]←1
        t←r0;
        rscr←t;
        call[mulCheck]; * t←t+rscr,MULTIPLY==>ALU←0, CARRY←0
        skipif[R ODD],rscr2←rscr2;
mulKerr1:
        error; * DIDN'T TAKE Q[14]=1 BRANCH

        t←t;
        skipif[ALU=0];
mulKerr2:
        error; * CARRY,,(0+0)/2 SHOULD BE ZERO

        t←(r10) RSH 1;
        t←t#(Q);
        skipif[ALU=0];
mulKerr3:
        error; * ALU15,,Q[0:14] SHOULD BE (r10) RSH 1

mulDone: BRANCH[afterMul];

```

* September 11, 1977 2:46 PM

%

MULCHECK

This subroutine performs,
t←t+(rscr),MULTIPLY;

It sets rscr2=0 IF Q[14] branch DID NOT HAPPEN.
It sets rscr2=1 IF Q[14] branch DID HAPPEN

T and rscr2 ARE CLOBBERED! IT ASSUMES r0=0, r1=1.

%

SUBROUTINE;

mulCheck:

t←t+(rscr),MULTIPLY, GLOBAL, AT[700];

GOTO[.+1];

rscr2←r0,AT[701];

RETURN;

rscr2←r1,AT[703];

RETURN;

TOP LEVEL;

afterMul:

* September 14, 1977 12:03 PM

%

DIVIDE TEST

H3 ← ALU[1:15],,Q[0] ==> H3← 2*ALU,,Q[0]
Q ← Q[1:15], CARRY ==> Q ← 2*Q,,CARRY

%

divideTest:

```
*      Q0=0, CARRY=0
      Q←r0;
      t←(r1)+(r1),DIVIDE;      *t←1+1,DIVIDE==> CARRY←0,
      t←t#(4C);
      skipif[ALU=0];
      error; * 2*(1+1)=4, Q0=0

      t←Q;
      skipif[ALU=0];
      error; * CARRY WAS ZERO, Q SHOULD BE ZERO
```

divB:

```
*      Q0=0, CARRY=1
      Q←r0;
      t←rhigh1;
      t←t+(r1);      * T = 100001
      t←t+(rhigh1),DIVIDE; * t←100001+100000,DIVIDE==>ALU=1, CARRY=1
      t←t#(2C);
      skipif[ALU=0];
      error; * 2*(1+1)=4, Q0=0

      t←(r1)#(Q);
      skipif[ALU=0];
      error; * 2*0,,CARRY SHOULD BE 1
```

divC:

```
*      Q0=1, CARRY=0
      Q←rhigh1;
      t←(r1)+(r1),DIVIDE; * t←1+1,DIVIDE==>ALU=2, CARRY=0
      t←t#(5C);
      skipif[ALU=0];
      error; * 2*(2),,Q[0]=5

      t←(Q);
      skipif[ALU=0];
      error; * Q[1:15],,CARRY SHOULD BE ZERO
```

divD:

```
*      Q0=1, CARRY=1
      t←Q←rhigh1; * SET Q[0] TO ONE
      t←t+(r1); * T = 100001
      t←t+(rhigh1),DIVIDE; * t←100001+100000,DIVIDE==>ALU=1, CARRY=1
      t←t#(3C);
      skipif[ALU=0];
      error; * 2*(1),,Q[0]=3

      t←(r1)#(Q);
      skipif[ALU=0];
      error; * 2*0,,CARRY SHOULD BE 1
```

divE:

```
*      Q←r01=>Q0=0, CARRY=1
      Q←r01;
      t←(rhigh1)+1; * T = 100001
      t←t+(rhigh1),DIVIDE; * t←100001+100000,DIVIDE==>ALU=1, CARRY=1
      t←t#(2C);
      skipif[ALU=0];
      error; * 2*(1),,Q[0]=2

      t←(r01) LSH 1;
      t←t+(r1); * ADD ONE FOR CARRY
      t←t#(Q);
      skipif[ALU=0];
      error; * 2*r01,,CARRY SHOULD BE ((r01)LSH 1)+1
```

* May 1, 1909 1:09 PM

%

CDIVIDE TEST

H3 ← ALU[1:15],,Q[0] ==> H3← 2*ALU,,Q[0]
Q ← Q[1:15], CARRY ==> Q ← 2*Q,,CARRY' (COMPLEMENTED CARRY)

%

CdivideTest:

```
*          Q0=0, CARRY=0
Q←r0;
t←(r1)+(r1),CDIVIDE;    *t←1+1,DIVIDE==> CARRY←0
t←t#(4C);
skipif[ALU=0];
error; * 2*(1+1)=4, Q0=0
```

```
t←(r1)#(Q);
skipif[ALU=0];
error; * CARRY' WAS 1, Q SHOULD BE 1
```

CdivB:

```
*          Q0=0, CARRY=1
Q←r0;
t←rhigh1;
t←t+(r1);          * T = 100001
t←t+(rhigh1),CDIVIDE; * t←100001+100000,DIVIDE==>ALU=1, CARRY=1
t←t#(2C);
skipif[ALU=0];
error; * 2*(1+1)=4, Q0=0
```

```
t←(Q);
skipif[ALU=0];
error; * 2*0,,CARRY' SHOULD BE 0
```

CdivC:

```
*          Q0=1, CARRY=0
Q←rhigh1;
t←(r1)+(r1),CDIVIDE; * t←1+1,DIVIDE==>ALU=2, CARRY=0
t←t#(5C);
skipif[ALU=0];
error; * 2*(1),,Q[0]=3
```

```
t←(r1)#(Q);
skipif[ALU=0];
error; * Q[1:15],,CARRY' SHOULD BE 1
```

CdivD:

```
*          Q0=1, CARRY=0
T←Q←rhigh1;      * SET Q[0] TO ONE
T←T+(R1);        * T = 100001
T←T+(rhigh1),CDIVIDE; * T←100001+100000,DIVIDE==>ALU=1, CARRY=1
T←T#(3C);
skipif[ALU=0];
ERROR; * 2*(1),,Q[0]=3
```

```
T←(Q);
skipif[ALU=0];
ERROR; * 2*0,,CARRY' SHOULD BE 0
```

CdivE:

```
*          Q←R01=>Q0=0, CARRY=1
Q←R01;
T←(rhigh1)+1;    * T = 100001
T←T+(rhigh1),CDIVIDE; * T←100001+100000,CDIVIDE==>ALU=1, CARRY=1
T←T#(2C);
skipif[ALU=0];
ERROR; * 2*(1),,Q[0]=2
```

```
T←(R01) LSH 1;
T←T#(Q);
skipif[ALU=0];
ERROR; * 2*R01,,CARRY' SHOULD BE (R01)LSH 1
```

* September 9, 1977 5:09 PM

%

TEST 8 WAY SLOW B DISPATCH

Go thru the loop 16 times to make sure that only the low 3 bits are
or'd into next pc. keep counter in Q, loop control in CNT.

%

slowBr:

```
cnt←17s;
t←q←r0;
rscr2←7C;      * THIS WILL BE A 3 BIT MASK
```

slowBrL: * TOP OF LOOP

```
t←rm1;
rscr←t;
t←q;
BDISPATCH←t;  * DO DISPATCH W/ BITS IN T (=Q)
branch[BDTb1];
```

slowBRnoBr: * should have branched and didn't
error;

bdConverge:

```
t←(rscr2)AND(q);      * MASK DISPATCH VALUE
t←t#(rscr);
branch[. +2,ALU=0];
```

slowBRbadBr:

error; * DIDN'T GO WHERE WE EXPECTED

```
t←(r1)+(q);      * GET NEXT VALUE FOR DISPATCH
loopChk[slowBrL,CNT#0&-1],q←t;
```

afterBD:

goto[afterKerne14];

SET[BDTb1Loc,5110];

BDTb1:

```
goto[bdConverge], rscr←r0, AT[BDTb1Loc];
goto[bdConverge], rscr←r1, AT[BDTb1Loc,1];
goto[bdConverge], rscr←2C, AT[BDTb1Loc,2];
goto[bdConverge], rscr←3C, AT[BDTb1Loc,3];
goto[bdConverge], rscr←4C, AT[BDTb1Loc,4];
goto[bdConverge], rscr←5C, AT[BDTb1Loc,5];
goto[bdConverge], rscr←6C, AT[BDTb1Loc,6];
goto[bdConverge], rscr←7C, AT[BDTb1Loc,7];
goto[bdConverge], rscr←10C, AT[BDTb1Loc,10];      * shouldn't be here
goto[bdConverge], rscr←11C, AT[BDTb1Loc,11];      * shouldn't be here
goto[bdConverge], rscr←12C, AT[BDTb1Loc,12];      * shouldn't be here
goto[bdConverge], rscr←13C, AT[BDTb1Loc,13];      * shouldn't be here
goto[bdConverge], rscr←14C, AT[BDTb1Loc,14];      * shouldn't be here
goto[bdConverge], rscr←15C, AT[BDTb1Loc,15];      * shouldn't be here
goto[bdConverge], rscr←16C, AT[BDTb1Loc,16];      * shouldn't be here
goto[bdConverge], rscr←17C, AT[BDTb1Loc,17];      * shouldn't be here
```


%
Page Numbers: Yes First Page: 1
Heading:
kernelAlu.mc October 20, 1978 10:38 AM %
TITLE[KernelALU-mode11.8/30/78];

%
October 19, 1978 1:13 PM
PD← for LU←

This file defines ALU operations for the program. Definitions for all possible ALU operations are given here with one line/definition. The lines for the 20 operations enabled should have the "n" replaced by consecutive values from 0 to 17. 16 should be "NOT A" for the shifter and 17 is normally used as a variable by BitBlt (i.e., not defined here), but can be used otherwise, if desired. Other lines should be commented out.

Note that the letter "E" added as an argument where noted below converts the operation to emulator-only. This is intended for locations used as variables by BitBlt, which must restore the smashed ALUFM locations to the proper value before exit.

The "A" and "B" operations must both be defined because many macros optionally route using either A or B for RB, T, MD, and Q (B path is preferred). Since ALUF[0] is the default for microinstructions, it should contain a non-arithmetic operation to preserve CARRY and OVERFLOW branch conditions.

Also, when running Midas the "B" operation is stored in ALUFM 0 and the NOT A operation in ALUFM 16, so these should be loaded the same way by the microprogram to avoid confusion, unless there is some good reason to do otherwise.

Note that an important choice must be made between the arithmetic and logical versions of ALU←A. The arithmetic version allows XORCARRY with ALU←A but smashes OVERFLOW and CARRY branch conditions on ALU←A.

%

```

ABOPB[PD+,0,25];      *ALU+B (use n=0)
ABOPA[PD+,1,0];      *ALU+A
                      *(arithmetic--so XORCARRY ok. Requires no more
                      *time than boolean "A" in absence of carryin)
*      ABOPA[PD+,n,37];      *ALU+A (logical, so XORCARRY illegal but OVERFLOW
                      *and CARRY branch conditions not smashed).

*"NOT B" and "NOT A" must be both defined also.
ABOPB[NOT,,15,13];   *NOT B
ABOPA[NOT,,16,1];    *NOT A (use n=16 for shifter)

*Operations of no args (4th arg = letter E makes emulator-only)
ZALUOP[A0,2,31];     *A0 [= all zeroes]
ZALUOP[A1,3,7];      *A1 [= all ones]

*Arithmetic operations of two args [do not accept "slow" (external)
*B sources] (5th arg = letter E makes emulator-only)
SALUOP[+,4,,14];     *A+B
*      SALUOP[+,+1,n,,214];   *A+B+1
SALUOP[-,,5,,222];   *A-B
*      SALUOP[-,-1,n,,22];    *A-B-1

*Boolean operations of two args (5th arg = letter E makes emulator-only)
**Note how synonyms are defined for # (XOR) and = (XNOR, EQV)
ALUOP[AND,,6,,35];   *A AND B
ALUOP[OR,,7,,27];    *A OR B
ALUOP[#,,10,,23];    XALUOP[XOR,,n];      *A # B (A XOR B)
*      ALUOP[,,n,,15];        XALUOP[XNOR,,n];    XALUOP[EQV,,n];
ALUOP[AND NOT,,11,,33]; *A AND NOT B [= A AND (NOT B)]
ALUOP[OR NOT,,12,,17]; *A OR NOT B [= A OR (NOT B)]
*      ALUOP[NOT,AND,,n,,21]; *NOT A AND B [= (NOT A) AND B]
*      ALUOP[NOT,OR,,n,,5];   *NOT A OR B [= (NOT A) OR B]
*      ALUOP[NOT,AND NOT,,n,,11]; *NOT A AND NOT B [= (NOT A) AND (NOT B)]
*      ALUOP[NOT,OR NOT,,n,,3]; *NOT A OR NOT B [= (NOT A) OR (NOT B)]

*Operations of one arg (5th arg = letter E makes emulator-only)
*      AOP[2,,n,6];          *2A
*      AOP[2,+1,n,206];     *2A+1
AOP[+,1,13,200];     *A+1
AOP[-,-1,14,36];     *A-1

```

```

%
Page Numbers: Yes First Page: 1
Heading:
postamble.mc February 1, 1980 6:24 PM %
TITLE[POSTAMBLE];
TOP LEVEL;
%
February 1, 1980 6:24 PM
    Fix goto[preBegin], described below, into goto[restartDiagnostic]. Postamble already defines a
**nd uses preBegin.
February 1, 1980 11:52 AM
    Fix restart to goto[preBegin]. This allows each diagnostic to perform whatever initialization
**it wants.
September 19, 1979 9:18 PM
    Fix another editing bug in chkSimulating, used the wrong bit to check for flags.conditionOK --
**just did it wrong.
September 19, 1979 9:08 PM
    Fix bug in chkSimulating wherein an edit lost a carriage return and a statment became part of a
** comment. Unfortunately, automatic line breaks made the statement look as if it were still there rat
**her than making it look like part of the comment line.
September 19, 1979 4:23 PM
    Fix placement errors associated with bumming locations from makeholdvalue and from checksimulat
**ing.
September 19, 1979 3:48 PM
    Bum locations to fit postamble with current os/microD: reallyDone, checkFlags global, make che
**ckFlags callers exploit FF, eliminate noCirculate label, make others shorter..
September 19, 1979 10:41 AM
    change callers of getIM*, putIM* to use FF field when calling them.
September 19, 1979 10:18 AM
    Create zeroHoldTRscr which loops to zero hold-- called by routines that invoke resetHold when t
**he hold simulator may be functioning. Make getIM*, putIM* routines global.
September 16, 1979 1:27 PM
    Bum code to fix storage full problem that occurs because OS 16/6 is bigger than OS 15/5: remove
** kernel specific patch locations (patch*).
August 1, 1979 3:28 PM
    Add scopeTrigger.
June 17, 1979 4:48 PM
    Move IM data locations around to accommodate Ifu entry points
April 26, 1979 11:03 AM
    Make justReturn global.
April 19, 1979 5:03 PM
    Remove calls to incTask/HoldFreq from enable/disableConditionalTask.
April 18, 1979 3:24 PM
    Remove DisplayOff from postamble.
April 18, 1979 11:11 AM
    Rename chkTaskSim, chkHoldSim, simControl to incTaskFreq, incHoldFreq, makeHoldValue; clean up
**setHold.
April 17, 1979 10:51 PM
    SimControl now masks holdFreq and taskFreq & shifts them w/ constants defined in Postamble.
April 11, 1979 3:49 PM
    Add breakpoint to "done", and fix, again, a bug associated with task simulation. Set defaultFl
**agsP (when postamble defines it) to force taskSim and holdSim.
March 7, 1979 11:42 PM
    Set RBASE to defaultRegion upon entry to postamble. thnx to Roger.
February 16, 1979 2:54 PM
    Modify routines that read IM to invert the value returned in link if b1 from that value =1 (thi
**s implies the whole value was inverted).
January 25, 1979 10:41 AM
    Change taskCirculate code to accommodate taskSim wakeups for task 10D, 12B
January 18, 1979 5:13 PM
    Modify checkTaskNum to use the RM value, currentTaskNum, and modify taskCircInc to keep the cop
**y in currentTaskNum.
January 15, 1979 1:25 PM
    add justReturn, a subroutine that just returns
January 9, 1979 12:07 PM
    breakpoint on xorTaskSimXit to avoid midas bug
%
%*****
TABLE of CONTENTS, by order of Occurence
done Location where diagnostics go when they are finished --gives control to postamble code the incr
**ements iterations, implements hold and task simulation and task circulation.
reallyDone Location where postamble inits 2 rm locations then performs "GoTo BEGIN"
restart Reinit diagnostic state, then restart the diagnostic.
incTaskFreq Increment the task frequency counter
incHoldFreq Increment the hold frequency counter
makeHoldValue Counstruct the "Hold&TaskSim" value from holdFreq and taskFreq, given that each is enab

```

```

**led in Flags
chkRunSimulators      Cause Hold or Task sim to happen, if required
chkSimulating         Return ALU#0 if some sort of simulating occurring
taskCirculate         Implement task circulation
incIterations         Increment iterations counter (>16 bits)
resetHold             Reset Hold&TaskSim to its previous value.
setHold              Set hold&task sim, notifying task simulator to do it.
simInit              Entry point for initialization in task simulator code
testTaskSim          Subroutine that tests task simulator
fixSim               Run Hold&TaskSim given current holdFreq and taskFreq
readByte3            Return byte 3 of an IM location
getIMRH              Return right half of an IM location
getIMLH              Return left half of an IM location
putIMRH              Write Right half of an IM location
putIMLH              Write Left half aof an IM location
checkFlags           Return Alu result & t based on entry mask & current flags
checkTaskNum         Return "currentTaskNum" # expectedTaskNum
notifyTask           Awaken take in T
topLvlPostRtn       Code that returns through mainPostRtn
scopeTrigger         Global subroutine that performs TIOA<0,TIOA<17777
justReturn           global subroutine that returns only
random               random numbers, used w/ getRand[] macro.
saveRandState        Save random number generator's state
restoreRandState     Restore old state to random number generator
getRandV             Part of random number linkage
xorFlags             Xor Flags w/ t
xorTaskCirc          toggle flags.taskCirc
xorHoldSim           toggle flags.holdSim
xorTaskSim           toggle flags.taskSim
disableConditionalTask  disable conditional tasking
enableConditionalTask  enable conditional tasking
ERR                  global label where ERROR macro gives control
IMdata               beginning of Postamble's FLAGS, et c.
%*****
IM[ILC,0];
%*****
    This code presumes R0=0 and uses RSCR, RSCR2, T, and Q.  It uses a number of other registers in
** a different RM region.
    When Postamble gets control of the processor at "Done", bits in "Flags", a word in IM determine
** which of Postamble's functions will occur when it runs.  At the least, Postamble increments at 32 bit
** number in IM called Iterations.  If flags.taskSim is true, the task simulator started.  The task si
** mulator awakens after a software controllable number of clocks has occurred.  The microcode that wakes
** up must reset the task simulator before it (the microcode) blocks to cause a task wakeup to occur ag
** ain.  The first time a program runs (ie., the time before it gives control to "done") the task simula
** tor and the hold simulator (discussed below) are inactive.  Running the task simulator forces task sp
** ecific hardware functions to effect the state of the machine.
When flags.holdSim is set, Postamble sets the hold simulator to a non-zero value.  The 8 bit "hold valu
** e" enters a circulating shift register where occurrence of a "1" bit at b[0] causes an external hold.
** This exercises the hold hardware.
The body of postamble contains a number of procedures for user programs, including routines to read and
** write IM, a routine to return a random number, and routines to initialize a task's pc and to notify
** it.
%*****
done:

```

```

* June 17, 1979  4:49 PM          POSTAMBLE CONSTANTS

    set[randomTloc, 620]; * random number generator may have to
* be moved to "global" call location if extensively used!

    set[flagsLoc, 1000];   mc[flagsLocC, flagsLoc];
    set[taskFreqLoc, 1400]; mc[taskFreqLocC, taskFreqLoc];
    set[holdFreqLoc, 2000]; mc[holdFreqLocC, holdFreqLoc];
    set[nextTaskLoc, 2400]; mc[nextTaskLocC, nextTaskLoc];
    set[itrsLoc, 3000];   mc[itrsLocC, itrsLoc];
*   holdValueLoc defined in preamble!
    set[preBeginLoc, 4000]; mc[preBeginLocC, preBeginLoc];
    set[initTloc, 4400];   mc[initTlocC, initTloc];

    ifdef[simInitLocC, mc[simInitLocC, initTloc] ]; * define the bmux constant for the
*   address of the task simulator code. If its already been defined, leave it as is.

*   flags.taskSim defined in preamble!
*   flags.holdSim defined in preamble!
*   flags.simulating defined in preamble!
mc[flags.testTasks, b13]; * than 8 flags (since READIM rtns a BYTE)
mc[flags.conditional, flags.conditionalP]; * allow simulating iff flags.simulating
* AND flags.conditionOK
mc[flags.conditionOK, flags.conditionOKp]; * enable conditional simulating

%*****
This portion of the kernel code encapsulates the microdiagnostic with an outer loop. This outer
** loop has several features that it implements:
    task simulation
    hold simulation
    task switching

Task simulation refers to the taskSim register in the hardware. It is 4 bits wide; taskSim[0] enables t
**he task simulator and taskSim[1:3] form a counter that determines the number of cycles before a task
**wakeup occurs.

Hold simulation is similar: holdSim is an 8-bit recirculating shift register in which the presence of a
** 1 in bit 7 causes HOLD two instructions later.

Task switching determines which task will run the microdiagnostics.

These features are controlled by the flags word in IM. If the appropriate bits are set to one, the asso
**ciated feature will function. The bits are defined above (flags.taskSim, flags.holdSim, flags.testTas
**ks).
%*****

    rmRegion[rmForKernelRtn];
    knowRbase[rmForKernelRtn];

    rv[setHoldRtn, 0];
    rv[oldt, 0]; * save t, rscr, rscr2, rtn link for resetHold
    rv[oldrscr, 0];
    rv[oldrscr2, 0];
    rv[resetHoldRtn, 0];
    rv[xorFlagsRtn, 0];
    rv[flagSubrsRtn, 0];
    rv[mainPostRtn, 0];

    knowRbase[rm2ForKernelRtn]; * defined in preamble because of macros
* that reference randV, randX

    knowRbase[defaultRegion];

```

```

* February 1, 1980 6:24 PM POSTAMBLE CONTROL CODE
D-O-E:
RBASE ← rbase[defaultRegion], breakpoint;      * set RBASE incase user's is different.
call[incTaskFreq];
call[incHoldFreq];
call[makeHoldValue];
call[taskCirculate];
call[incIterations];

call[checkFlags]t←flags.testTasks;      * bookkeeping is done. switch tasks if required
skipif[ALU#0];
branch[preBegin];      * xit if not running other tasks

taskCircInit:  * now that bookkeeping is done, switch tasks if required
noop;      * for placement.
call[checkTaskNum];
rscr ← t;      * rscr ← nextTask

t ← preBeginLocC;      * link ← t ← preBeginLoc
subroutine;

taskCirc:
zeroHold[rscr2];      * turn off hold-task sim during LdTpc←, wakeup
link ← t;
t ← rscr;
top level;
ldTPC ← t;      * tpc[nextTask] ← preBeginLoc
call[notifyTask];      * wakeup nextTask: task num in t
set[xtask, 1];
block;
set[xtask, 0];

preBegin: noop, at[preBeginLoc];
call[chkRunSimulators]; * check for simulator conditions and run if required

reallyDone:      * LOOP TO BEGIN
t←RSCR←a1;
goto[begin], RSCR2←t;

restart:      * restart diagnostics from "initial" state
randV ← t-t;      * restart random number generator
randX ← t-t;

rscr ← t-t;      * restart hold/task simulator stuff
call[putIMRH], t ← holdFreqLocC;
rscr ← t-t;
call[putIMRH], t←taskFreqLocC;
rscr ← t-t;
call[putIMRH], t ← holdValueLocC;

rscr ← t-t;      * restart iterations count
call[putIMRH], t ← itrsLocC;

branch[restartDiagnostic];      * special entry point so each diagnostic
* can perform whatever special initialization that it wants to perform

```

* January 18, 1978 1:51 PM

This code sets the taskSim value with the next value if flags.testTasks is true. Otherwise 0 is used.

```

IF flags.taskSim THEN
BEGIN -- when hardware counts to 17 it awakens
taskFreq ← (taskFreq + 1) or 10b;      -- simTask
IF taskFreq > 15 THEN taskFreq ← 12;  -- always wait min=2 cycles
END
      ELSE taskFreq ← 0;
IF flags.holdSim THEN
BEGIN
holdFreq ← holdFreq+1;
IF holdFreq >376 THEN holdFreq ← 0;
END;
      ELSE holdFreq ← 0;

```

```

incTaskFreq: subroutine;
  t ← link;
  mainPostRtn ← t;
  top level;

  call[checkFlags], t←flags.taskSim;      * see if taskSim enabled
  branch[writeTaskSim, alu=0],t←r0;      * use 0 if not enabled
  t←taskFreqLocC; * increment next taskSim
  call[readByte3];
  t←(r1)+(t);
  t-(156c);      * Use [1..156]. 156 => max wait,
  skipif[alu<0]; * 1 = > min wait. Beware infinite hold!
  t←r1;      * see discussion at simInit, simSet code
  noop;

```

```

writeTaskSim:
  rscr ← t;
  call[putIMRH], t←taskFreqLocC; * update IM location

```

```

taskSimRtn:
  goto[topLvlPostRtn];

```

```

incHoldFreq: subroutine;      * see if holdSim enabled
  t ← link;
  mainPostRtn ← t;
  top level;

  call[checkFlags], t←flags.holdSim;
  branch[noHoldSim, alu=0],t←r0; * use zero if hold not enabled
  t←holdFreqLocC;
  call[readByte3];
  t←t+(r1);
  t-(377c);      * IF holdFreq >376
  skipif[alu<0];
  t←r1;      * THEN holdFreq ← 1;
  noop;      * here for placement

```

```

noHoldSim:
  rscr ← t;      * rewrite IM
  call[putIMRH], t ← holdFreqLocC;

```

```

holdSimRtn:
  goto[topLvlPostRtn];

```

* April 17, 1979 10:51 PM

This code actually controls the task and hold loading. It is responsible for initializing T for the task at simTaskLevel, and it is responsible for initializing HOLD.

The code proceeds by constructing the current value to be loaded into hold and placing it in IM at holdValueLoc. Kernel loads HOLD as its last act before looping to BEGIN.

hold&tasksim requires hold value in left byte, task counter value in right byte.

makeHoldValue: subroutine; * construct holdValue
saveReturn[mainPostRtn];

call[chkSimulating];
skpif[alu#0];
branch[simCtr10];

t←holdFreqLocC; * rscr2 ← holdFreq
call[readByte3], t←holdFreqLocC;
rscr2 ← t;

call[readByte3], t←taskFreqLocC; * t ← taskFreq

t←lsh[t, sim.taskShift]; * position hold and task values
t←t and (sim.taskMask);
rscr2 ← lsh[rscr2, sim.holdShift];
rscr2 ← (rscr2) and (sim.holdMask);
rscr2 ← (rscr2) and (377c);

rscr2 ← (t) + (rscr2); * taskFreq, holdFreq
rscr ← rscr2;

%

now, save combined taskSim, holdSim values in IM. Last thing done before exiting postamble is to set H **OLD if simulating.

%

simCtr1WHold: * may branch here from simCtr10
call[putIMRH], t ← holdValueLocC; * write holdValue into holdValueLoc
branch[simCtr1Rtn];

simCtr10:
branch[simCtr1WHold], rscr ← t-t; * write zero into holdValueLoc

simCtr1Rtn:
goto[topLv1PostRtn];

* September 19, 1979 9:09 PM

```
%*+++++
IF chkSimulating[] THEN fixSimulator[];
* cause hold or task simulator to run, if required
%*+++++
```

```
chkRunSimulators: subroutine;
    saveReturn[chkRunSimRtn];
    call[chkSimulating];
    dblBranch[chkRunSimXit, chkRunSimDoIt, alu=0];
chkRunSimDoIt: * run the simulator
    noop;
    call[fixSim];
    noop;
```

```
chkRunSimXit:
    returnUsing[chkRunSimRtn];
```

* September 19, 1979 9:19 PM

```
%*+++++
chkSimulating: PROCEDURE RETURNS[weAreSimulating: BOOLEAN] =
BEGIN
weAreSimulating ← FALSE;
IF flags.Simulating THEN
IF ~(flags.Conditional) OR (flags.Conditional AND flags.ConditionOK) THEN
weAreSimulating ← TRUE;
END;
```

```
%*+++++
chkSimulating: subroutine;
    saveReturn[chkSimulatingRtn];
    call[checkFlags], t←flags.simulating; * check for taskSim OR holdSim
    branch[chkSimNo, alu=0];
    t ← flags.conditional; * We're simulating. check
    call[checkFlags], t ← flags.conditional;
    dblbranch[chkSimYes, chkSimCond, alu=0];
chkSimCond: * conditional simulation. check for ok
    t ← flags.conditionOK;
    call[checkFlags];
    skipif[alu=0];
    branch[chkSimYes]; * conditionOK is set, do it!
    branch[chkSimNo];
chkSimYes: * run the simulator
    t ← (r0)+1; * rtn w/ alu#0
chkSimRtn:
    returnAndBranch[chkSimulatingRtn, t];
chkSimNo:
    branch[chkSimRtn], t ← r0; * rtn w/ alu=0
```

* January 25, 1979 10:44 AM

%

This code controls task circulation for the diagnostics: when flags.testTasks is set, postamble ** causes successive tasks to execute the diagnostic code when the current task has completed. If flag **s.taskSim is true the diagnostic is using the taskSimulator to periodically awaken the simulator task **; consequently, that task (simTaskLevel) must not execute the diagnostic -- otherwise the advantage o **f the simulator for testing the effects of task switching will be lost.

IF ~flags.testTasks THEN RETURN;

temp ← getTaskNum[] + 1; -- increment the current number

IF flags.taskSim THEN

IF temp = simTaskLevel THEN temp ← temp+1;

IF temp > maxTaskLevel THEN temp ← 0;

putTaskNum[temp]; -- remember it in IM

%

taskCirculate: subroutine;

saveReturn[mainPostRtn];

call[checkFlags], t ← flags.testTasks;

branch[taskCircRtn, ALU=0]; * Don't bother if not task circulating.

noop;

call[checkTaskNum]; * Increment the current task number.

t ← t + (r1); * Current value came back in t.

q ← t; * Remember incremented value in Q.

call[checkFlags], t ← flags.taskSim;

skpif[ALU#0], t ← q; * Now, see if using task simulator.

branch[taskCircChk]; * If not task simulating, check for max size.

t - (simTaskLevelC); * Since we're task simulating, avoid

skpif[ALU#0]; * we must avoid simTaskLevel.

t ← t+1; * Increment over simTask if required.

noop;

taskCircChk:

t - (20C); * See if tasknum is too big.

skpif[ALU#0];

t ← t-t; * We wraparound to zero.

currentTaskNum ← t; * keep it in both RM and IM

rscr ← t;

call[putIMRH], t ← nextTaskLocC;

noop; * for placement

taskCircRtn:

goto[topLv1PostRtn];

* January 18, 1978 1:57 PM

```
incIterations: subroutine;      * maintain double precision count at incItrsLoc
    t ← link;
    mainPostRtn ← t;
    top level;

    call[getIMRH], t ← itrsLocC;  * increment iteration count at tableloc+1
    rscr ← (t)+1;

    rscr2 ← rscr;  * copy new itrs
    rscr2 ← (t) #(rscr);  * see if new b0 # old b0

    rscr2 ← (rscr2) AND (b0);
    skipif[alu#0];
    branch[incItrs2], q ← r0;      * new b0 = old b0. remember in q and write
    t and (B0);  * see if b0 went from 0 to 1 or 1 to 0 (carry)
    skipif[alu=0], q←r0;  * skipif old b0 = 0
    q ← r1;

incItrs2:

    call[putIMRH], t ← itrsLocC;  * T = addr, rscr = value
    rscr2 ← q;
    branch[incItrsRtn,alu=0];    * goto incItrsRtn if no carry

incItrsHi16:
    link ← t;  * read hi byte of hi 16 bits
    call[getIMLH];
    rscr ← (t)+1;
    call[putIMLH], t ← itrsLocC;  * T = addr, rscr = value
    noop;  * help the instruction placer.

incItrsRtn:
    goto[topLv1PostRtn];
```

* March 20, 1978 1:51 PM KERNEL - COMMON SUBROUTINES

resetHold: subroutine; * special subroutine called by IM manipulating
 * code. This subr saves t, rscr, rscr2 and causes hold to be initialized to the value in
 * holdValueLocC. It restores the RM and T values before returning.

```

oldT ← t;
t ← link;
resetHoldRtn ← t;
top level;
t ← rscr;
oldrscr ← t;
t ← rscr2;
oldrscr2 ← t; * link, t, rscr, and rscr2 are now saved.

```

```

t ← HoldValueLocC; * READ RIGHT HALF, HoldValueLocC
subroutine;
link ← t;
top level;
readim[3]; * read low order byte
subroutine;
t ← link; * low byte in t
t and (b1); * see if the data is inverted. If so, b1 will
skpif[ALU=0]; * 1, and we must reinvert the data.
t ← not(t);
t ← t and (getIMmask); * isolate the byte

```

```

rscr ← HoldValueLocC;
subroutine;
link ← rscr;
top level;
readim[2]; * read hi order byte
subroutine;
rscr2 ← link; * hi byte in rscr2
(rscr2) and (b1); * see if the data is inverted. If so, b1 will
skpif[ALU=0]; * 1, and we must reinvert the data.
rscr2 ← not(rscr2);
rscr2 ← (rscr2) and (getIMmask); * isolate the byte

```

```

top level;
noop;
rscr2 ← lsh[rscr2, 10]; * left shift hi byte
t ← t and (377C); * isolate low byte
t ← t OR (rscr2); * add hi byte
call[setHold];

```

```

knowRbase[rmForKernelRtn]; * restore link, t, rscr, rscr2, then return
RBASE ← rbase[rmForKernelRtn];
t ← oldrscr;
rscr ← t;
t ← oldrscr2;
rscr2 ← t;
subroutine;
link ← resetHoldRtn;
return, t ← oldt, RBASE ← rbase[defaultRegion];

```

* June 23, 1978 10:22 AM

setHold: subroutine; * ENTER w/ T = HOLD value
 * clobber t, rscr, rscr2

```

zeroHold[rscr2]; * kill hold-task sim before polyphas instrs xqt
rscr2 ← q; * SAVE Q
q ← t; * save hold value, then save rtn link
t ← link;
setHoldRtn ← t;

```

```

taskingon;
t ← simInitLocC; * defined w/ postamble constants OR in
* some user specific code (eg., memSubrsA where RM values are defined). This
* convention allows users to specify their own code to run when the simulator task runs.
link ← t; * cause task taskSimLevel to put
top level;
ldTPC ← simTaskLevelC; * proper hold value in T for refresh
notify[simTaskLevel]; * after task switch occurs. Remember
* taskSim is a counter. refresh it!
noop; * wakeup will happen soon
noop;
rbase ← RBASE[rmForKernelRtn];
t ← setHoldRtn, rbase ← RBASE[defaultRegion];
Q ← rscr2; * restore Q
subroutine;
link ← t;
return;

```

* This code actually causes T to be set properly and branches to the code that sets HOLD.

```

set[xtask, 1];

```

simInit:

```

t ← q, at[initTloc];

```

simSet:

```

hold&tasksim←t; * T init'd at simInit
noop; * this noop doesn't cause hold to count

```

simBlock:

```

branch[simSet], block; * count hold, block

```

%

Note: if t = 14, then hold = 16 when the simulator blocks. The preempted task will execute one instruct
 **ion, then the task simulator will waken the simulator task.

%

```

set[xtask, 0];

```

* November 6, 1978 12:07 PM MIDAS SUBROUTINE for testing the task simulator

testTaskSim: subroutine;

```

rscr ← link; * save return in case we later want it
top level;
t ← lsh[t, 10]; * ENTER w/ T = task sim val NOT shifted
q ← t; * simInit expects q = hold value

```

```

subroutine;
t ← initTlocC;
link ← t;
top level;
LDTPC ← simTaskLevelC;
notify[simTaskLevel];

```

```

noop;
t ← t - t; * t ← 0
branch[., alu=0], t+t; * this shouldn't change

```

testTaskErr:

```

branch[.], breakpoint;

```

```

subroutine;
link ← rscr;
return;

```

fixSim: subroutine;

```

t←link; * save return in fixSimRtn
fixSimRtn ← t;
top level;

```

```
    call[makeHoldValue]; * compose holdValue and set hardware
    call[getIMRH], t ← holdValueLocC;
    call[setHold];

    returnUsing[fixSimRtn];

zeroHoldTRscr: subroutine;
    t←4c;
    rscr←a0;
zeroHoldTRscrL:
    Hold&TaskSim←rscr;
    t←t-1, Hold&TaskSim←rscr;
    loopUntil[alu<0, zeroHoldTRscrL];
    return;
```

* January 18, 1979 5:18 PM * READ/WRITE IM
%

The subroutines that read and write IM turn OFF hold simulator before touching IM. Before they **return to the caller, they invoke "resetHold" to reset the hold register to the contents of "holdValue **Loc". By convention, the current value of the two simulator registers is kept in "holdValueLoc" for t **his express purpose. Zeroing and resetting hold is done because of hardware restrictions: hold and po **lyphase instructions don't mix.

ReadIM[] instructions are followed by a mask operation with getIMmask because of the interactio **n between DWATCH (Midas facility) and LINK[0].
%

readByte3: subroutine; * CLOBBER T, RSCR!
zeroHold[rscr];

rscr ← link; * this routine assumes t points to IM
link ← t; * it reads the least significant byte in IM

top level; * read byte 3
readim[3];
subroutine;
t←link; * t = byte3
t and (b1); * see if the data is inverted. If so, b1 will
skpif[ALU=0]; * 1, and we must reinvert the data.
t ← not(t);
t ← t and (getIMmask); * isolate the byte

top level; * reset value of hold and return
call[resetHold];
subroutine;
link ← rscr;
return; * return w/ byte in t

getIMRH: subroutine, global; * CLOBBER T, RSCR, RSCR2!
zeroHold[rscr]; * disable task/hold Sim before touching IM

rscr ← link; * ENTER w/ T pointing to IM location
link ← t;

top level; * read hi byte of right half
readim[2];
subroutine;
rscr2 ← link; * rscr2 = high byte
(rscr2) and (b1); * see if the data is inverted. If so, b1 will
skpif[ALU=0]; * 1, and we must reinvert the data.
rscr2 ← not(rscr2);
rscr2 ← (rscr2) and (getIMmask); * isolate the byte

link ← t; * read low byte of right half
top level;
readim[3];
subroutine;
t ← link; * t = low byte, rscr2 = hi byte
t and (b1); * see if the data is inverted. If so, b1 will
skpif[ALU=0]; * 1, and we must reinvert the data.
t ← not(t);
t ← t and (getIMmask); * isolate the byte

rscr2 ← lsh[rscr2, 10];
t ← t + (rscr2); * RETURN w/ T = IMRH

top level;
call[resetHold];
subroutine;
link ← rscr;
return;

getIMLH: subroutine, global; * CLOBBER T, RSCR, RSCR2!
zeroHold[rscr]; * disable task/hold Sim before touching IM

rscr ← link; * ENTER w/ T pointing to IM location
link ← t;

top level; * read hi byte of left half
readim[0];

```

subroutine;
rscr2 ← link; * rscr2 = hi byte
(rscr2) and (b1); * see if the data is inverted. If so, b1 will
skipif[ALU=0]; * 1, and we must reinvert the data.
rscr2 ← not(rscr2);
rscr2 ← (rscr2) and (getIMmask); * isolate the byte

link ← t; * read low byte of left half
top level;
readim[1];
subroutine; * CLOBBER T, RSCR, RSCR2!
t ← link; * t = low byte, rscr2 = hi byte
t and (b1); * see if the data is inverted. If so, b1 will
skipif[ALU=0]; * 1, and we must reinvert the data.
t ← not(t);
t ← t and (getIMmask); * isolate the byte

rscr2 ← lsh[rscr2, 10];
t ← t + (rscr2); * RETURN w/ T = IMLH

top level;
call[resetHold];
subroutine;
link ← rscr;
return;

putIMRH: subroutine, global; * T = addr, RSCR = value, clobberr RSCR2
rscr2 ← link;
link ← t;

zeroHold[t]; * disable task/hold Sim before touching IM

top level;
t ← rscr;
IMRHB'POK ← t;

call[resetHold];
subroutine;
link ← rscr2;
return;

putIMLH: subroutine, global; * T = addr, RSCR = value, Clobber RSCR2
rscr2 ← rscr;
rscr ← link;
link ← t;

zeroHold[t]; * disable task/hold Sim before touching IM

top level;
t ← rscr2;
IMLHR0'POK ← t;

call[resetHold];
subroutine;
link ← rscr;
return;

checkFlags: subroutine, global; * CLOBBER T, RSCR, RSCR2
rscr ← link; * this routine assumes t has a bit mask

zeroHold[rscr2]; * disable task/hold Sim before touching IM

rscr2 ← flagsLocC; * it reads the flags word in IM
link ← rscr2; * and performs t←tANDflag
top level;
readim[3];
subroutine;
rscr2 ← link;
(rscr2) and (b1); * see if the data is inverted. If so, b1 will
skipif[ALU=0]; * 1, and we must reinvert the data.
rscr2 ← not(rscr2);
rscr2 ← (rscr2) and (getIMmask); * isolate the byte

top level;
call[resetHold];
subroutine;

```



```
link ← rscr;
return, t ← t AND(rscr2);      * returnee can do alu=0 fast branch

checkTaskNum: subroutine;      * enter: T=expected task num,
rscr←t, RBASE ← rbase[currentTaskNum]; * return: T=current task num, branch condition
t ← currentTaskNum, RBASE ← rbase[defaultRegion];      * clobber rscr, rscr2

return, t#(rscr);      * rtn w/ branch condition, t=current task
* number, rscr = expected task number.
```

```

* August 1, 1979 3:30 PM      other, miscellaneous subroutines:
notifyTask: subroutine;
    rscr← link;
    bigBDispatch←t;
    top level;
    branch[dispatchTbl];
set[nloc, 6600];
dispatchTbl:
    branch[nxit], notify[0],      at[nloc,0];
    branch[nxit], notify[1],      at[nloc,1];
    branch[nxit], notify[2],      at[nloc,2];
    branch[nxit], notify[3],      at[nloc,3];
    branch[nxit], notify[4],      at[nloc,4];
    branch[nxit], notify[5],      at[nloc,5];
    branch[nxit], notify[6],      at[nloc,6];
    branch[nxit], notify[7],      at[nloc,7];
    branch[nxit], notify[10],     at[nloc,10];
    branch[nxit], notify[11],     at[nloc,11];
    branch[nxit], notify[12],     at[nloc,12];
    branch[nxit], notify[13],     at[nloc,13];
    branch[nxit], notify[14],     at[nloc,14];
    branch[nxit], notify[15],     at[nloc,15];
    branch[nxit], notify[16],     at[nloc,16];
    branch[nxit], notify[17],     at[nloc,17];
    branch[.], breakpoint, at[nloc,20];
    branch[.], breakpoint, at[nloc,21];

subroutine;
nxit:
    link ← rscr;
    return;

topLvlPostRtn:
    RBASE ← rbase[mainPostRtn];
    link ← mainPostRtn;
    return, RBASE ← rbase[defaultRegion];

scopeTrigger: subroutine;
    t ← a0, global;
    TIOA ← t, T←a1;
    return, TIOA←t;

justReturn: * this subroutine ONLY RETURNS. Calling justReturn forces the instruction
            return, global; * (logically) after the call to occur in the physically
* next location after the call. This is a way of reserving a noop that can ALWAYS be
* safely patched with a "call".

```

* April 24, 1978 6:51 PM

```
knowRbase[randomRM];
random:
  goto[random1], t ← rndm0, RBASE ← rbase[randV], at[randomTloc,0]; knowRbase[randomRM];
  goto[random1], t ← rndm1, RBASE ← rbase[randV], at[randomTloc,1]; knowRbase[randomRM];
  goto[random1], t ← rndm2, RBASE ← rbase[randV], at[randomTloc,2]; knowRbase[randomRM];
  goto[random1], t ← rndm3, RBASE ← rbase[randV], at[randomTloc,3]; knowRbase[randomRM];
  goto[random1], t ← rndm4, RBASE ← rbase[randV], at[randomTloc,4]; knowRbase[randomRM];
  goto[random1], t ← rndm5, RBASE ← rbase[randV], at[randomTloc,5]; knowRbase[randomRM];
  goto[random1], t ← rndm6, RBASE ← rbase[randV], at[randomTloc,6]; knowRbase[randomRM];
  goto[random1], t ← rndm7, RBASE ← rbase[randV], at[randomTloc,7];

random1:
  return, t ← randV ← (randV)+t;
  knowRbase[defaultRegion];

saveRandState: subroutine;      * remember random number seed
  RBASE ← rbase[randV];
  oldRandV ← randV;
  oldRandX ← randX;
  return, RBASE ← rbase[defaultRegion];

restoreRandState: subroutine;   * restore remembered random number seed
  RBASE ← rbase[randV];
  randV ← oldRandV;
  randX ← oldRandX;
  return, RBASE ← rbase[defaultRegion];

getRandV: subroutine;
  RBASE ← rbase[randV];
  RETURN, t ← randV, RBASE ← rbase[defaultRegion];
```

* January 20, 1978 3:04 PM 'FLAGS' manipulating code

```
xorFlags: subroutine; * T = value to XOR into flags
* CLOBBER RSCR, RSCR2, T
  rscr2 ← t; * save bits
  t ← link;
  xorFlagsRtn ← t;
  top level;

  t ← flagsLocC;
  call[readByte3];
  t ← t # (rscr2); * xor new bits

  rscr ← t; * put new value back into IM
  call[putIMRH], t ← flagsLocC;

  returnUsing[xorFlagsRtn];

xorTaskCirc: subroutine; * xor the flags.testTasks bit in FLAGS
* CLOBBER RSCR, RSCR2, T
  saveReturn[flagSubrsRtn];
  t ← flags.testTasks;
  call[xorFlags];
  noop;

  returnUsing[flagSubrsRtn];

xorHoldSim: subroutine; * xor the flags.holdSim bit in FLAGS
  saveReturn[flagSubrsRtn];
  t ← flags.holdSim;
  call[xorFlags];

  rscr←a0; * whether off or on, clear holdFreqLoc
  call[putIMRH], t ← holdFreqLocC; * holdFreq ← 0

  call[fixSim];
xorHoldSimXit:
  noop, breakpoint;

  returnUsing[flagSubrsRtn];

xorTaskSim: subroutine; * xor the flags.taskSim bit in FLAGS
  saveReturn[flagSubrsRtn];
  t ← flags.taskSim;
  call[xorFlags];

  rscr←a0; * whether off or on, clear taskFreqLoc
  call[putIMRH], t ← taskFreqLocC; * taskFreq ← 0

  call[fixSim]; * fix the holdValueLoc, set hardware

xorTaskSimXit:
  breakpoint, noop;
  returnUsing[flagSubrsRtn];
  top level;
```

* June 22, 1978 10:15 AM

%

This code supports the conditional simulation mechanism. disableConditionalTask is a subroutine ** that requires no parameters. It clears flags.conditionOK and sets flags.conditional. It also turns on ***ff the hold simulator.

enableConditionalTask sets flags.conditionOK and flags.conditional, then it calls makeHoldValue to force **e the hold simulator into working.

%

```
disableConditionalTask: subroutine;
  saveReturn[flagSubrsRtn];
  call[checkFlags], t ← (r0)-1; * use mask = -1 to force a read of all the bits
  rscr ← not (flags.conditionOK);
  rscr ← t and (rscr);
  rscr ← (rscr) or (flags.conditional);
  rscr ← (rscr) and (377C); * isolate lower byte
  call[putIMRH], t ← flagsLocC;

  call[makeHoldValue]; * compose a new hold value from task and
```

```

        * hold simulator sub values
    call[zeroHoldTRscr]; * stop hold
    call[resetHold]; * jam the hold register w/ holdValue
    returnUsing[flagSubrsRtn];

enableConditionalTask: subroutine;
    saveReturn[flagSubrsRtn];
    call[checkFlags], t ← (r0)-1; * use mask = -1 to force a read of all the bits
    noop; * make placement easier
    rscr ← t or (flags.conditionOK);
    rscr ← (rscr) or (flags.conditional);
    noop; * make placement easier
    call[putIMRH], t ← flagsLocC; * write the new value

    call[makeHoldValue]; * compose a new hold value from task and
        * hold simulator sub values
    call[zeroHoldTRscr]; * stop hold
    call[resetHold]; * jam the hold register w/ holdValue
    returnUsing[flagSubrsRtn];
    top level;

* ERRORS come here!
    branch[err];
SET[ERRLOC,400];
ERR:
    BREAKPOINT,GLOBAL, AT[ERRLOC];
    GOTO[.],BREAKPOINT, AT[ERRLOC,1];
    GOTO[.], AT[ERRLOC,2];

* DATA HELD IN IM

IMdata:
    ifdef[defaultFlagsP,,set[defaultFlagsP,add[flags.taskSim!, flags.holdSim!]]; * define default
**t flags if undefined

    data[(Flags: lh[0] rh[defaultFlagsP], at[flagsLoc]); * CONTROL FLAGS
    data[(taskFreq: lh[0] rh[0], at[taskFreqLoc]); * task sim value
    data[(holdFreq: lh[0] rh[0], at[holdFreqLoc]); * hold sim value
    data[(nextTask: lh[0] rh[0], at[nextTaskLoc]); * next task value
    data[(holdValue: lh[0] rh[0], at[holdValueLoc]); * current hold value
    data[(iterations: lh[0] rh[0], at[itrsLoc]); * iteration count

postDone: noop;

```